

32

SH7205 Group

User's Manual: Hardware

Renesas 32-Bit RISC Microcomputer SuperH™ RISC engine Family / SH7200 Series

SH7205 R5S72050W200BG

Rev.3.00 Mar 2014

Notice

- 1. Descriptions of circuits, software and other related information in this document are provided only to illustrate the operation of semiconductor products and application examples. You are fully responsible for the incorporation of these circuits, software, and information in the design of your equipment. Renesas Electronics assumes no responsibility for any losses incurred by you or third parties arising from the use of these circuits, software, or information.
- 2. Renesas Electronics has used reasonable care in preparing the information included in this document, but Renesas Electronics does not warrant that such information is error free. Renesas Electronics assumes no liability whatsoever for any damages incurred by you resulting from errors in or omissions from the information included herein.
- 3. Renesas Electronics does not assume any liability for infringement of patents, copyrights, or other intellectual property rights of third parties by or arising from the use of Renesas Electronics products or technical information described in this document. No license, express, implied or otherwise, is granted hereby under any patents, copyrights or other intellectual property rights of Renesas Electronics or others.
- 4. You should not alter, modify, copy, or otherwise misappropriate any Renesas Electronics product, whether in whole or in part. Renesas Electronics assumes no responsibility for any losses incurred by you or third parties arising from such alteration, modification, copy or otherwise misappropriation of Renesas Electronics product.
- 5. Renesas Electronics products are classified according to the following two quality grades: "Standard" and "High Quality". The recommended applications for each Renesas Electronics product depends on the product's quality grade, as indicated below.
	- "Standard": Computers; office equipment; communications equipment; test and measurement equipment; audio and visual equipment; home electronic appliances; machine tools; personal electronic equipment; and industrial robots etc.
	- "High Quality": Transportation equipment (automobiles, trains, ships, etc.); traffic control systems; anti-disaster systems; anti-crime systems; and safety equipment etc.

Renesas Electronics products are neither intended nor authorized for use in products or systems that may pose a direct threat to human life or bodily injury (artificial life support devices or systems, surgical implantations etc.), or may cause serious property damages (nuclear reactor control systems, military equipment etc.). You must check the quality grade of each Renesas
Electronics product before using it in a particular application. You may not use any Renesas Electronics application for which it is not intended. Renesas Electronics shall not be in any way liable for any damages or losses incurred by you or third parties arising from the use of any Renesas Electronics product for which the product is not intended by Renesas Electronics.

- 6. You should use the Renesas Electronics products described in this document within the range specified by Renesas Electronics, especially with respect to the maximum rating, operating supply voltage range, movement power voltage range, heat radiation characteristics, installation and other product characteristics. Renesas Electronics shall have no liability for malfunctions or damages arising out of the use of Renesas Electronics products beyond such specified ranges.
- 7. Although Renesas Electronics endeavors to improve the quality and reliability of its products, semiconductor products have specific characteristics such as the occurrence of failure at a certain rate and malfunctions under certain use conditions. Further, Renesas Electronics products are not subject to radiation resistance design. Please be sure to implement safety measures to guard them against the possibility of physical injury, and injury or damage caused by fire in the event of the failure of a Renesas Electronics product, such as safety design for hardware and software including but not limited to redundancy, fire control and malfunction prevention, appropriate treatment for aging degradation or any other appropriate measures. Because the evaluation of microcomputer software alone is very difficult, please evaluate the safety of the final products or systems manufactured by you.
- 8. Please contact a Renesas Electronics sales office for details as to environmental matters such as the environmental compatibility of each Renesas Electronics product. Please use Renesas Electronics products in compliance with all applicable laws and regulations that regulate the inclusion or use of controlled substances, including without limitation, the EU RoHS Directive. Renesas Electronics assumes no liability for damages or losses occurring as a result of your noncompliance with applicable laws and regulations.
- 9. Renesas Electronics products and technology may not be used for or incorporated into any products or systems whose manufacture, use, or sale is prohibited under any applicable domestic or foreign laws or regulations. You should not use Renesas Electronics products or technology described in this document for any purpose relating to military applications or use by the military, including but not limited to the development of weapons of mass destruction. When exporting the Renesas Electronics products or technology described in this document, you should comply with the applicable export control laws and regulations and follow the procedures required by such laws and regulations.
- 10. It is the responsibility of the buyer or distributor of Renesas Electronics products, who distributes, disposes of, or otherwise places the product with a third party, to notify such third party in advance of the contents and conditions set forth in this document, Renesas Electronics assumes no responsibility for any losses incurred by you or third parties as a result of unauthorized use of Renesas Electronics products.
- 11. This document may not be reproduced or duplicated in any form, in whole or in part, without prior written consent of Renesas Electronics.
- 12. Please contact a Renesas Electronics sales office if you have any questions regarding the information contained in this document or Renesas Electronics products, or if you have any other inquiries.
- (Note 1) "Renesas Electronics" as used in this document means Renesas Electronics Corporation and also includes its majority-owned subsidiaries.
- (Note 2) "Renesas Electronics product(s)" means any product developed or manufactured by or for Renesas Electronics.

General Precautions in the Handling of MPU/MCU Products

The following usage notes are applicable to all MPU/MCU products from Renesas. For detailed usage notes on the products covered by this document, refer to the relevant sections of the document as well as any technical updates that have been issued for the products.

1. Handling of Unused Pins

Handle unused pins in accord with the directions given under Handling of Unused Pins in the manual.

- ⎯ The input pins of CMOS products are generally in the high-impedance state. In operation with unused pin in the open-circuit state, extra electromagnetic noise is induced in the vicinity of LSI, an associated shoot-through current flows internally, and malfunctions occur due to the false recognition of the pin state as an input signal become possible. Unused pins should be handled as described under Handling of Unused Pins in the manual.
- 2. Processing at Power-on

The state of the product is undefined at the moment when power is supplied.

⎯ The states of internal circuits in the LSI are indeterminate and the states of register settings and pins are undefined at the moment when power is supplied. In a finished product where the reset signal is applied to the external reset pin, the states of pins are not guaranteed from the moment when power is supplied until the reset process is completed.

In a similar way, the states of pins in a product that is reset by an on-chip power-on reset function are not guaranteed from the moment when power is supplied until the power reaches the level at which resetting has been specified.

3. Prohibition of Access to Reserved Addresses

Access to reserved addresses is prohibited.

- ⎯ The reserved addresses are provided for the possible future expansion of functions. Do not access these addresses; the correct operation of LSI is not guaranteed if they are accessed.
- 4. Clock Signals

After applying a reset, only release the reset line after the operating clock signal has become stable. When switching the clock signal during program execution, wait until the target clock signal has stabilized.

- ⎯ When the clock signal is generated with an external resonator (or from an external oscillator) during a reset, ensure that the reset line is only released after full stabilization of the clock signal. Moreover, when switching to a clock signal produced with an external resonator (or by an external oscillator) while program execution is in progress, wait until the target clock signal is stable.
- 5. Differences between Products

Before changing from one product to another, i.e. to a product with a different part number, confirm that the change will not lead to problems.

 $-$ The characteristics of an MPU or MCU in the same group but having a different part number may differ in terms of the internal memory capacity, layout pattern, and other factors, which can affect the ranges of electrical characteristics, such as characteristic values, operating margins, immunity to noise, and amount of radiated noise. When changing to a product with a different part number, implement a system-evaluation test for the given product.

Configuration of This Manual

This manual comprises the following items:

- 1. General Precautions in the Handling of MPU/MCU Products
- 2. Configuration of This Manual
- 3. Preface
- 4. Contents
- 5. Overview
- 6. Description of Functional Modules
	- CPU and System-Control Modules
	- ⎯ On-Chip Peripheral Modules

The configuration of the functional description of each module differs according to the module. However, the generic style includes the following items:

- i) Feature
- ii) Input/Output Pin
- iii) Register Description
- iv) Operation
- v) Usage Note

When designing an application system that includes this LSI, take notes into account. Each section includes notes in relation to the descriptions given, and usage notes are given, as required, as the final part of each section.

- 7. List of Registers
- 8. Electrical Characteristics
- 9. Appendix
	- Product Type, Package Dimensions, etc.

10. Main Revisions for This Edition (only for revised versions)

The list of revisions is a summary of points that have been revised or added to earlier versions. This does not include all of the revised contents. For details, see the actual locations in this manual.

11. Index

Preface

This LSI is an RISC (Reduced Instruction Set Computer) microcomputer which includes a Renesas-original RISC CPU as its core, and the peripheral functions required to configure a system.

- Target Users: This manual was written for users who will be using this LSI in the design of application systems. Target users are expected to understand the fundamentals of electrical circuits, logical circuits, and microcomputers.
- Objective: This manual was written to explain the hardware functions and electrical characteristics of this LSI to the target users. Refer to the SH-2A, SH2A-FPU Software Manual for a detailed description of the instruction set.

Notes on reading this manual:

- In order to understand the overall functions of the chip Read the manual according to the contents. This manual can be roughly categorized into parts on the CPU, system control functions, peripheral functions and electrical characteristics.
- In order to understand the details of the CPU's functions Read the SH-2A, SH2A-FPU Software Manual.
- In order to understand the details of a register when its name is known Read the index that is the final part of the manual to find the page number of the entry on the register. The addresses, bits, and initial values of the registers are summarized in section 32, List of Registers.

• Description of Numbers and Symbols

Aspects of the notations for register names, bit names, numbers, and symbolic names in this manual are explained below.

• Description of Registers

Each register description includes a bit chart, illustrating the arrangement of bits, and a table of bits, describing the meanings of the bit settings. The standard format and notation for bit charts and tables are described below.

All trademarks and registered trademarks are the property of their respective owners.

Contents

Section 1 Overview

1.1 SH7205 Features

This LSI is a single-chip RISC (reduced instruction set computer) microcontroller that includes two Renesas-original RISC CPUs as its cores, and the peripheral functions required to configure a system.

This LSI features a multi-processor architecture, that is, a dual-core architecture that includes two units of SH-2A CPU, which provides upward compatibility for SH-1, SH-2, and SH-2E CPUs at object code level. The SH-2A CPU has a RISC-type instruction set and uses a superscalar architecture and a Harvard architecture, which greatly improves instruction execution speed. In addition, the 32-bit internal-bus architecture that is independent from the direct memory access controller (DMAC) enhances data processing power. This CPU brings the user the ability to set up high-performance systems with strong functionality at less expense than was achievable with previous microcontrollers, and is even able to handle realtime control applications requiring highspeed characteristics.

This LSI includes a floating-point unit (FPU) and cache for each of the CPU cores (CPU0 and CPU1). In addition, this LSI has on-chip peripheral functions necessary for system configuration: 64-Kbyte (CPU0) and 32-Kbyte (CPU1) RAM for high-speed operation, 16-Kbyte RAM for data retention, an interrupt controller (INTC), a multi-function timer pulse unit 2 (MTU2), a compare match timer (CMT), a realtime clock (RTC), a serial communication interface with FIFO (SCIF), a synchronous serial communication unit (SSU), an $I²C$ bus interface 3 (IIC3), a serial sound interface with FIFO (SSIF), a controller area network (RCAN-TL1), an A/D converter (ADC), a D/A converter (DAC), an AND/NAND flash memory controller (FLCTL), a USB2.0 host/function module supporting two ports (USB), an AT attachment packet interface (ATAPI), a 2D engine (2DG), and I/O ports.

This LSI also provides an external memory access support function to enable direct connection to various memory devices or peripheral LSIs. These on-chip functions significantly reduce costs of designing and manufacturing application systems. Furthermore, I/O pins in this LSI have weak keeper circuits that prevent the pin voltage from entering an intermediate potential range. Therefore, no external circuits for fixing the input level are required, which reduces the number of parts considerably.

The features of this LSI are listed in table 1.1.

Table 1.1 SH7205 Features

1.2 Product Lineup

Table 1.2 Product Lineup

1.3 Block Diagram

Figure 1.1 Block Diagram

1.4 Pin Assignment

Figure 1.2 Pin Assignment

Table 1.3 Pin Numbers and Corresponding Pin Names

Section 1 Overview SH7205 Group

Section 1 Overview SH7205 Group

Function 1 Function 2 Function 3 Function 4 Pin

 $Y10$ NAF7 $I(s)/O$ $-$ AUDATA1 O $-$ Yes Figure 1.3 (14) $Y11$ $-$ AUDATA3 O $-$ Yes Figure 1.3 (14) $Y12$ FCDE 0 Y es Figure 1.3 (10) $Y13$ TCLKB $I(s)$ $-$ Figure 1.3 (6)

[Legend]

(s): Schmitt

(a): Analog

(o): Open drain

Figure 1.3 (1) Simplified Circuit Diagram (Schmitt Input Buffer)

Figure 1.3 (2) Simplified Circuit Diagram (Schmitt AND Input Buffer with Pull-Up)

Figure 1.3 (3) Simplified Circuit Diagram (TTL AND Input Buffer with Pull-Up)

Figure 1.3 (4) Simplified Circuit Diagram (TTL OR Input and A/D Input Buffer)

Figure 1.3 (6) Simplified Circuit Diagram (Schmitt OR Input and A/D Input Buffer)

Figure 1.3 (7) Simplified Circuit Diagram (Output Buffer with Enable, with Latch)

Figure 1.3 (8) Simplified Circuit Diagram (Output Buffer with Enable, with Latch and Weak Keeper)

Figure 1.3 (9) Simplified Circuit Diagram (Bidirectional Buffer, TTL AND Input, with Latch)

Figure 1.3 (10) Simplified Circuit Diagram (Bidirectional Buffer, TTL AND Input, with Latch and Weak Keeper)

Figure 1.3 (11) Simplified Circuit Diagram (Bidirectional Buffer, TTL AND Input, with Latch and Pull-up)

Figure 1.3 (12) Simplified Circuit Diagram (Bidirectional Buffer, Schmitt AND Input, with Latch and Weak Keeper)

Figure 1.3 (13) Simplified Circuit Diagram (Bidirectional Buffer, Schmitt AND Input, with Latch and Pull-up)

Figure 1.3 (14) Simplified Circuit Diagram

(Bidirectional Buffer, TTL AND Input, Schmitt AND Input, with Latch and Weak Keeper)

Figure 1.3 (15) Simplified Circuit Diagram (Open Drain Output and Schmitt OR Input Buffer)

Figure 1.3 (16) Simplified Circuit Diagram (Open Drain/Normal Output and Schmitt OR Input Buffer, with Latch)

Figure 1.3 (17) Simplified Circuit Diagram (Oscillation Buffer 1)

Figure 1.3 (18) Simplified Circuit Diagram (Oscillation Buffer 2)

1.5 Pin Functions

Table 1.4 Pin Functions

1.6 Bus Structure

The bus structure of this LSI largely consists of CPU buses, internal buses, and peripheral buses.

The bus master of the CPU bus is a CPU. Each of the CPUs (CPU0 and CPU1) is provided with a CPU bus for its own use, allowing both CPUs to run independently. A CPU bus actually consists of two buses: an instruction-fetch bus and a memory-access bus (Harvard architecture).

The circuit has multiple (four) internal buses. The master modules of the internal bus are the two CPUs and the DMAC. CPU0 and CPU1 are connected to the internal bus via the CPU bus and a bus bridge. The read port and write port of the DMAC act as master modules for the corresponding buses. That is, CPU0, CPU1, the DMA read port, and the DMA write port are individually connected to the corresponding internal buses. This allows each of the master modules to occupy its own internal bus without bus arbitration.

The slave modules of the internal buses are multiple peripheral buses (including the external bus and high-speed on-chip RAM access bus). On each internal bus, arbitration for bus mastership is performed between internal buses (master modules), after which access to the individual peripheral bus proceeds. In this LSI, internal modules called bus interface units (BIUs) perform this bus mastership arbitration. Since the BIUs perform arbitration per slave module, multiple accesses can proceed in parallel as long as access by each master module is to a different BIU. However, if more than one attempt at access to a given BIU is made at the same time, arbitration between the master modules is performed. The master module that failed to gain bus mastership is kept waiting until it succeeds, and thus the multiple accesses are executed one after another. The order of priority in bus-mastership arbitration is as follows: DMA write port > DMA read port > CPU. The priority order of CPU0 and CPU1 alternates in a round-robin manner.

The peripheral buses are used for the connections with the on-chip peripheral modules.

Section 2 CPU

2.1 Register Configuration

The register set consists of sixteen 32-bit general registers, four 32-bit control registers, and four 32-bit system registers.

2.1.1 General Registers

Figure 2.1 shows the general registers.

The sixteen 32-bit general registers are numbered R0 to R15. General registers are used for data processing and address calculation. R0 is also used as an index register. Several instructions have R0 fixed as their only usable register. R15 is used as the hardware stack pointer (SP). Saving and restoring the status register (SR) and program counter (PC) in exception handling is accomplished by referencing the stack using R15.

Figure 2.1 General Registers

2.1.2 Control Registers

The control registers consist of four 32-bit registers: the status register (SR), the global base register (GBR), the vector base register (VBR), and the jump table base register (TBR).

The status register indicates instruction processing states.

The global base register functions as a base address for the GBR indirect addressing mode to transfer data to the registers of on-chip peripheral modules.

The vector base register functions as the base address of the exception handling vector area (including interrupts).

The jump table base register functions as the base address of the function table area.

Figure 2.2 Control Registers

(1) Status Register (SR)

should always be 0.

instruction.

These bits are always read as 0. The write value

Specifies a saturation operation for a MAC

True/false condition or carry/borrow bit

(2) Global Base Register (GBR)

1 S — R/W S Bit

0 T — R/W T Bit

GBR is referenced as the base address in a GBR-referencing MOV instruction.

(3) Vector Base Register (VBR)

VBR is referenced as the branch destination base address in the event of an exception or an interrupt.

(4) Jump Table Base Register (TBR)

TBR is referenced as the start address of a function table located in memory in a JSR/N@@(disp8,TBR) table-referencing subroutine call instruction.

2.1.3 System Registers

The system registers consist of four 32-bit registers: the high and low multiply and accumulate registers (MACH and MACL), the procedure register (PR), and the program counter (PC). MACH and MACL store the results of multiply or multiply and accumulate operations. PR stores the return address from a subroutine procedure. PC indicates the four bytes ahead of the current instruction being executed.

Figure 2.3 System Registers

(1) Multiply and Accumulate Register High (MACH) and Multiply and Accumulate Register Low (MACL)

MACH and MACL are used as the addition value in a MAC instruction, and store the result of a MAC or MUL instruction.

(2) Procedure Register (PR)

PR stores the return address of a subroutine call using a BSR, BSRF, or JSR instruction, and is referenced by a subroutine return instruction (RTS).

(3) Program Counter (PC)

PC indicates the four bytes ahead of the current instruction being executed.

2.1.4 Register Banks

For the nineteen 32-bit registers comprising general registers R0 to R14, control register GBR, and system registers MACH, MACL, and PR, high-speed register saving and restoration can be carried out using a register bank. The register contents are automatically saved in the bank after the CPU accepts an interrupt that uses a register bank. Restoration from the bank is executed by issuing a RESBANK instruction in an interrupt processing routine.

The LSI has two sets of 15 banks, one set for each CPU core. For details, see the SH-2A, SH2A-FPU Software Manual and section 7.8, Register Banks.

2.1.5 Initial Values of Registers

Table 2.1 lists the values of the registers after a reset.

Table 2.1 Initial Values of Registers

2.2 Data Formats

2.2.1 Data Format in Registers

Register operands are always longwords (32 bits). If the size of memory operand is a byte (8 bits) or a word (16 bits), it is changed into a longword by expanding the sign-part when loaded into a register.

Figure 2.4 Data Format in Registers

2.2.2 Data Formats in Memory

Memory data formats are classified into bytes, words, and longwords. Memory can be accessed in 8-bit bytes, 16-bit words, or 32-bit longwords. A memory operand of fewer than 32 bits is stored in a register in sign-extended or zero-extended form.

A word operand should be accessed at a word boundary (an even address of multiple of two bytes: address 2n), and a longword operand at a longword boundary (an even address of multiple of four bytes: address 4n). Otherwise, an address error will occur. A byte operand can be accessed at any address.

Only big-endian byte order can be selected for the data format.

Data formats in memory are shown in figure 2.5.

Figure 2.5 Data Formats in Memory

2.2.3 Immediate Data Format

Byte (8-bit) immediate data is located in an instruction code. Immediate data accessed by the MOV, ADD, and CMP/EQ instructions is sign-extended and handled in registers as longword data. Immediate data accessed by the TST, AND, OR, and XOR instructions is zero-extended and handled as longword data. Consequently, AND instructions with immediate data always clear the upper 24 bits of the destination register.

20-bit immediate data is located in the code of a MOVI20 or MOVI20S 32-bit transfer instruction. The MOVI20 instruction stores immediate data in the destination register in sign-extended form. The MOVI20S instruction shifts immediate data by eight bits in the upper direction, and stores it in the destination register in sign-extended form.

Word or longword immediate data is not located in the instruction code, but rather is stored in a memory table. The memory table is accessed by an immediate data transfer instruction (MOV) using the PC relative addressing mode with displacement.

See examples given in section 2.3.1 (10), Immediate Data.

2.3 Instruction Features

2.3.1 RISC-Type Instruction Set

Instructions are RISC type. This section details their functions.

(1) 16-Bit Fixed-Length Instructions

Basic instructions have a fixed length of 16 bits, improving program code efficiency.

(2) 32-Bit Fixed-Length Instructions

The SH-2A additionally features 32-bit fixed-length instructions, improving performance and ease of use.

(3) One Instruction per State

Each basic instruction can be executed in one cycle using the pipeline system.

(4) Data Length

Longword is the standard data length for all operations. Memory can be accessed in bytes, words, or longwords. Byte or word data in memory is sign-extended and handled as longword data. Immediate data is sign-extended for arithmetic operations or zero-extended for logic operations. It is also handled as longword data.

Table 2.2 Sign Extension of Word Data

Note: @(disp, PC) accesses the immediate data.

(5) Load-Store Architecture

Basic operations are executed between registers. For operations that involve memory access, data is loaded to the registers and executed (load-store architecture). Instructions such as AND that manipulate bits, however, are executed directly in memory.

(6) Delayed Branch Instructions

With the exception of some instructions, unconditional branch instructions, etc., are executed as delayed branch instructions. With a delayed branch instruction, the branch is taken after execution of the instruction immediately following the delayed branch instruction. This reduces disturbance of the pipeline control when a branch is taken.

In a delayed branch, the actual branch operation occurs after execution of the slot instruction. However, instruction execution such as register updating excluding the actual branch operation, is performed in the order of delayed branch instruction \rightarrow delay slot instruction. For example, even though the contents of the register holding the branch destination address are changed in the delay slot, the branch destination address remains as the register contents prior to the change.

(7) Unconditional Branch Instructions with No Delay Slot

The SH-2A additionally features unconditional branch instructions in which a delay slot instruction is not executed. This eliminates unnecessary NOP instructions, and so reduces the code size.

(8) Multiply/Multiply-and-Accumulate Operations

-bit \times 16-bit \rightarrow 32-bit multiply operations are executed in one to two cycles. 16-bit \times 16-bit + -bit \rightarrow 64-bit multiply-and-accumulate operations are executed in two to three cycles. 32-bit \times -bit \rightarrow 64-bit multiply and 32 -bit \times 32-bit + 64-bit \rightarrow 64-bit multiply-and-accumulate operations are executed in two to four cycles.

(9) T Bit

The T bit in the status register (SR) changes according to the result of the comparison. Whether a conditional branch is taken or not taken depends upon the T bit condition (true/false). The number of instructions that change the T bit is kept to a minimum to improve the processing speed.

Table 2.4 T Bit

(10) Immediate Data

Byte immediate data is located in an instruction code. Word or longword immediate data is not located in instruction codes but in a memory table. The memory table is accessed by an immediate data transfer instruction (MOV) using the PC relative addressing mode with displacement.

With the SH-2A, 17- to 28-bit immediate data can be located in an instruction code. However, for 21- to 28-bit immediate data, an OR instruction must be executed after the data is transferred to a register.

Table 2.5 Immediate Data Accessing

Note: @(disp, PC) accesses the immediate data.

(11) Absolute Address

When data is accessed by an absolute address, the absolute address value should be placed in the memory table in advance. That value is transferred to the register by loading the immediate data during the execution of the instruction, and the data is accessed in register indirect addressing mode.

With the SH-2A, when data is referenced using an absolute address not exceeding 28 bits, it is also possible to transfer immediate data located in the instruction code to a register and to reference the data in register indirect addressing mode. However, when referencing data using an absolute address of 21 to 28 bits, an OR instruction must be used after the data is transferred to a register.

(12) 16-Bit/32-Bit Displacement

When data is accessed by 16-bit or 32-bit displacement, the displacement value should be placed in the memory table in advance. That value is transferred to the register by loading the immediate data during the execution of the instruction, and the data is accessed in the indexed indirect register addressing mode.

2.3.2 Addressing Modes

Addressing modes and effective address calculation are as follows:

Table 2.8 Addressing Modes and Effective Addresses

2.3.3 Instruction Format

The instruction formats and the meaning of source and destination operands are described below. The meaning of the operand depends on the instruction code. The symbols used are as follows:

- xxxx: Instruction code
- mmmm: Source register
- nnnn: Destination register
- iiii: Immediate data
- dddd: Displacement

Table 2.9 Instruction Formats

Note: $*$ In multiply-and-accumulate instructions, nnnn is the source register.

2.4 Instruction Set

2.4.1 Instruction Set by Classification

Table 2.10 lists the instructions according to their classification.

Table 2.10 Classification of Instructions

The table below shows the format of instruction codes, operation, and execution states. They are described by using this format according to their classification.

Notes: 1. Instruction execution cycles: The execution cycles shown in the table are minimums. In practice, the number of instruction execution states will be increased in cases such as the following:

- a. When there is a conflict between an instruction fetch and a data access
- b. When the destination register of a load instruction (memory \rightarrow register) is the same as the register used by the next instruction.
- 2. Depending on the operand size, displacement is scaled by \times 1, \times 2, or \times 4. For details, refer to the SH-2A, SH2A-FPU Software Manual.

2.4.2 Data Transfer Instructions

Table 2.11 Data Transfer Instructions

R01UH0473EJ0300 Rev. 3.00
Mar. 27, 2014 **RENESAS** Mar. 27, 2014

2.4.3 Arithmetic Operation Instructions

Table 2.12 Arithmetic Operation Instructions

2.4.4 Logic Operation Instructions

Table 2.13 Logic Operation Instructions

2.4.5 Shift Instructions

Table 2.14 Shift Instructions

2.4.6 Branch Instructions

Table 2.15 Branch Instructions

Note: * One cycle when the program does not branch.

2.4.7 System Control Instructions

Table 2.16 System Control Instructions

Notes: Instruction execution cycles: The execution cycles shown in the table are minimums. In practice, the number of instruction execution states in cases such as the following:

a. When there is a conflict between an instruction fetch and a data access

b. When the destination register of a load instruction (memory \rightarrow register) is the same as the register used by the next instruction.

* In the event of bank overflow, the number of cycles is 19.

2.4.8 Floating-Point Operation Instructions

Table 2.17 Floating-Point Operation Instructions

Compatibility

Compatibility

2.4.9 FPU-Related CPU Instructions

Table 2.18 FPU-Related CPU Instructions

2.4.10 Bit Manipulation Instructions

Table 2.19 Bit Manipulation Instructions

2.5 Processing States

The LSI has four CPU processing states: reset, dual-processor active, single-processor active, and power-down. Figure 2.6 shows the transitions between the states.

Figure 2.6 Transitions between Processing States

(1) Reset State

In the reset state, the CPU is reset. There are two kinds of reset, power-on reset and manual reset.

(2) Dual-Processor Active State

In this state, CPU0 and CPU1 sequentially execute their own programs.

(3) Single-Processor Active State

In this state, either CPU0 or CPU1 operates. In single-processor 0 mode, CPU0 is active and CPU1 is in the sleep state. In single-processor 1 mode, CPU1 is active and CPU0 is in the sleep state.

(4) Power-Down State

In the power-down state, both CPUs stop operating to reduce power consumption.

Section 3 Floating-Point Unit (FPU)

3.1 Features

The FPU has the following features.

- Conforms to IEEE754 standard
- 16 single-precision floating-point registers (can also be referenced as eight double-precision registers)
- Two rounding modes: Round to nearest and round to zero
- Denormalization modes: Flush to zero
- Five exception sources: Invalid operation, divide by zero, overflow, underflow, and inexact
- Comprehensive instructions: Single-precision, double-precision, and system control

3.2 Data Formats

3.2.1 Floating-Point Format

A floating-point number consists of the following three fields:

- $Sign(s)$
- Exponent (e)
- Fraction (f)

This LSI can handle single-precision and double-precision floating-point numbers, using the formats shown in figures 3.1 and 3.2.

Figure 3.1 Format of Single-Precision Floating-Point Number

The exponent is expressed in biased form, as follows:

 $e = E + bias$

The range of unbiased exponent E is $E_{min} - 1$ to $E_{max} + 1$. The two values $E_{min} - 1$ and $E_{max} + 1$ are distinguished as follows. $E_{min} - 1$ indicates zero (both positive and negative sign) and a denormalized number, and $E_{\text{max}} + 1$ indicates positive or negative infinity or a non-number (NaN). Table 3.1 shows E_{min} and E_{max} values.

Parameter	Single-Precision	Double-Precision	
Total bit width	32 bits	64 bits	
Sign bit	1 bit	1 bit	
Exponent field	8 bits	11 bits	
Fraction field	23 bits	52 bits	
Precision	24 bits	53 bits	
Bias	$+127$	$+1023$	
$\mathsf{E}_{\scriptscriptstyle \sf max}$	$+127$	$+1023$	
E_{\min}	-126	-1022	

Table 3.1 Floating-Point Number Formats and Parameters

Floating-point number value v is determined as follows:

If $E = E_{max} + 1$ and $f \ne 0$, v is a non-number (NaN) irrespective of sign s If $E = E_{\text{max}} + 1$ and $f = 0$, $v = (-1)^s$ (infinity) [positive or negative infinity] If $E_{\text{min}} \leq E \leq E_{\text{max}}$, $v = (-1)^{s} 2^{E} (1.f)$ [normalized number] If $E = E_{min} - 1$ and $f \neq 0$, $v = (-1)^{5} 2^{Emin} (0.f)$ [denormalized number] If $E = E_{min} - 1$ and $f = 0$, $v = (-1)^{s}0$ [positive or negative zero]

Table 3.2 shows the ranges of the various numbers in hexadecimal notation.

Table 3.2 Floating-Point Ranges

3.2.2 Non-Numbers (NaN)

Figure 3.3 shows the bit pattern of a non-number (NaN). A value is NaN in the following case:

- Sign bit: Don't care
- Exponent field: All bits are 1
- Fraction field: At least one bit is 1

The NaN is a signaling NaN (sNaN) if the MSB of the fraction field is 1, and a quiet NaN (qNaN) if the MSB is 0.

Figure 3.3 Single-Precision NaN Bit Pattern

An sNaN is input in an operation, except copy, FABS, and FNEG, that generates a floating-point value.

- When the EN.V bit in FPSCR is 0, the operation result (output) is a qNaN.
- When the EN.V bit in FPSCR is 1, an invalid operation exception will generate FPU exception processing. In this case, the contents of the operation destination register are unchanged.

If a qNaN is input in an operation that generates a floating-point value, and an sNaN has not been input in that operation, the output will always be a qNaN irrespective of the setting of the EN.V bit in FPSCR. An exception will not be generated in this case.

The qNAN values as operation results are as follows:

- Single-precision qNaN: H'7FBF FFFF
- Double-precision qNaN: H'7FF7 FFFF FFFF FFFF

See the individual instruction descriptions for details of floating-point operations when a nonnumber (NaN) is input.

3.2.3 Denormalized Numbers

For a denormalized number floating-point value, the exponent field is expressed as 0, and the fraction field as a non-zero value.

In the SH2A-FPU, the DN bit in the status register FPSCR is always set to 1, therefore a denormalized number (source operand or operation result) is always flushed to 0 in a floatingpoint operation that generates a value (an operation other than copy, FNEG, or FABS).

When the DN bit in FPSCR is 0, a denormalized number (source operation operation result) is processed as it is. See the individual instruction descriptions for details of floating-point operations when a denormalized number is input.

3.3 Register Descriptions

3.3.1 Floating-Point Registers

Figure 3.4 shows the floating-point register configuration. There are sixteen 32-bit floating-point registers FPR0 to FPR15, referenced by specifying FR0 to FR15, DR0/2/4/6/8/10/12/14. The correspondence between FRPn and the reference name is determined by the PR and SZ bits in FPSCR. Refer to figure 3.4.

- 1. Floating-point registers, FPRi (16 registers) FPR0 to FPR15
- 2. Single-precision floating-point registers, FRi (16 registers) FR0 to FR15 indicate FPR0 to FPR15
- 3. Double-precision floating-point registers or single-precision floating-point vector registers in pairs, DRi (8 registers)

A DR register comprises two FR registers.

 $DR0 = {FR0, FR1}, DR2 = {FR2, FR3}, DR4 = {FR4, FR5}, DR6 = {FR6, FR7},$ $DR8 = {FR8, FR9}, DR10 = {FR10, FR11}, DR12 = {FR12, FR13}, DR14 = {FR14, FR15}$

	Reference name		Register name
Transfer instruction case: Operation instruction case: $FPSCR.PR = 0$ $FPSCR.PR = 1$	$FPSCR.SZ = 0$ $FPSCR.SZ = 1$		
	FR ₀	DR ₀	FPR ₀
		FR ₁	FPR ₁
	FR ₂ DR ₂	FPR ₂	
	FR ₃		FPR ₃
	FR4	DR4	FPR4
	FR ₅		FPR ₅
	FR ₆	DR ₆	FPR6
	FR ₇		FPR7
	FR ₈	DR8	FPR8
	FR ₉		FPR9
	FR10	DR10	FPR ₁₀
	FR11		FPR ₁₁
	FR12	DR12	FPR ₁₂
	FR13		FPR ₁₃
	FR14		FPR14
	FR15	DR14	FPR ₁₅

Figure 3.4 Floating-Point Registers

3.3.2 Floating-Point Status/Control Register (FPSCR)

FPSCR is a 32-bit register that controls floating-point instructions, sets FPU exceptions, and selects the rounding mode.

Table 3.3 Bit Allocation for FPU Exception Handling

Note: No FPU error occurs in the SH2A-FPU.

3.3.3 Floating-Point Communication Register (FPUL)

Information is transferred between the FPU and CPU via FPUL. FPUL is a 32-bit system register that is accessed from the CPU side by means of LDS and STS instructions. For example, to convert the integer stored in general register R1 to a single-precision floating-point number, the processing flow is as follows:

 $\mathsf{R1} \rightarrow (\mathsf{LDS}\text{ instruction}) \rightarrow \mathsf{FPUL} \rightarrow (\text{single-precision}\text{ FLOAT}\text{ instruction}) \rightarrow \mathsf{FR1}$

3.4 Rounding

In a floating-point instruction, rounding is performed when generating the final operation result from the intermediate result. Therefore, the result of combination instructions such as FMAC will differ from the result when using a basic instruction such as FADD, FSUB, or FMUL. Rounding is performed once in FMAC, but twice in FADD, FSUB, and FMUL.

Which of the two rounding methods is to be used is determined by the RM bits in FPSCR.

 $FPSCR.RM[1:0] = 00$: Round to Nearest $FPSCR.RM[1:0] = 01$: Round to Zero

(1) Round to Nearest

The operation result is rounded to the nearest expressible value. If there are two nearest expressible values, the one with an LSB of 0 is selected.

If the unrounded value is $2^{Emax} (2 - 2^{-p})$ or more, the result will be infinity with the same sign as the unrounded value. The values of Emax and P, respectively, are 127 and 24 for single-precision, and 1023 and 53 for double-precision.

(2) Round to Zero

The digits below the round bit of the unrounded value are discarded.

If the unrounded value is larger than the maximum expressible absolute value, the value will become the maximum expressible absolute value.

3.5 FPU Exceptions

3.5.1 FPU Exception Sources

FPU exceptions may occur on floating-point operation instruction and the exception sources are as follows:

- FPU error (E): When FPSCR.DN = 0 and a denormalized number is input (No error occurs in the SH2A-FPU)
- Invalid operation (V): In case of an invalid operation, such as NaN input
- Division by zero (Z): Division with a zero divisor
- Overflow (O): When the operation result overflows
- Underflow (U): When the operation result underflows
- Inexact exception (I): When overflow, underflow, or rounding occurs

The FPU exception cause field in FPSCR contains bits corresponding to all of above sources E, V, Z, O, U, and I, and the FPU exception flag and enable fields in FPSCR contain bits corresponding to sources V, Z, O, U, and I, but not E. Thus, FPU errors cannot be disabled.

When an FPU exception occurs, the corresponding bit in the FPU exception cause field is set to 1, and 1 is added to the corresponding bit in the FPU exception flag field. When an FPU exception does not occur, the corresponding bit in the FPU exception cause field is cleared to 0, but the corresponding bit in the FPU exception flag field remains unchanged.

3.5.2 FPU Exception Handling

FPU exception handling is initiated in the following cases:

- FPU error (E): FPSCR.DN = 0 and a denormalized number is input (No error occurs in the SH2A-FPU)
- Invalid operation (V) : FPSCR. Enable. $V = 1$ and invalid operation
- Division by zero (Z): FPSCR. Enable. $Z = 1$ and division with a zero divisor
- Overflow (O): FPSCR. Enable. $O = 1$ and instruction with possibility of operation result overflow
- Underflow (U): FPSCR. Enable. $U = 1$ and instruction with possibility of operation result underflow
- Inexact exception (I): $FPSCR$. Enable. I = 1 and instruction with possibility of inexact operation result

These possibilities of each exceptional handling on floating-point operation are shown in the individual instruction descriptions. All exception events that originate in the floating-point operation are assigned as the same FPU exceptional handling event. The meaning of an exception generated by floating-point operation is determined by software by reading from FPSCR and interpreting the information it contains. Also, the destination register is not changed when FPU exception handling operation occurs.

Except for the above, the FPU disables exception handling. In every processing, the bit corresponding to source V, Z , O , U , or I is set to 1, and a default value is generated as the operation result.

- Invalid operation (V): qNaN is generated as the result.
- Division by zero (Z) : Infinity with the same sign as the unrounded value is generated.
- Overflow (O):

When rounding mode $= RZ$, the maximum normalized number, with the same sign as the unrounded value, is generated.

When rounding mode = RN, infinity with the same sign as the unrounded value is generated.

- Underflow (U): Zero with the same sign as the unrounded value is generated.
- Inexact exception (I): An inexact result is generated.

Section 4 Multi-Core Processor

This LSI includes two SH2A CPUs (CPU0 and CPU1). Dual CPUs provide this LSI with strong levels of performance (through distribution of load) and functionality (through distribution of functions) in processing, which cannot be achieved by a single CPU.

4.1 Features

- Synchronization control between CPUs Inter-processor interrupts support control of synchronization between the two CPUs (see section 7, Interrupt Controller (INTC)).
- Exclusive control for shared resources Semaphore control registers support exclusive control for shared resources.
- Floating-point unit (FPU), cache memory, and high-speed on-chip RAM are provided for each CPU

Each CPU has its own FPU, cache memory, and high-speed on-chip RAM. The high-speed on-chip RAM can be configured as shared RAM space or as CPU-specific RAM space by enabling or disabling access from the other CPU (see section 29, On-Chip RAM).

• Power-down modes (see section 30, Power-Down Modes)

To reduce power consumption, this LSI can move between dual-processor mode, where both CPUs are operating, single-processor mode, where one CPU is in the sleep state, and dualsleep mode, where both CPUs are in the sleep state. By making transitions between these modes in accordance with the load, power consumption can be reduced while high performance is maintained.

• Multiple-internal-bus structure (see section 1, Overview)

To prevent deterioration in performance due to both CPUs and the DMAC not being able to get bus mastership, multiple (four) internal buses are provided.

4.2 Register Descriptions

The following registers are provided for control of the multi-core processor.

Table 4.1 Register Configuration

Notes: 1. The values H'10111000 and H'50110800, respectively, are read out in response to reading by CPU0 and CPU1.

2. After being read, the register is cleared to H'00.

4.2.1 CPU ID Register (CPUIDR)

The CPU ID register indicates the CPU number (CPU0 or CPU1).

Note: * Overall values of H'10111000 and H'50110800, respectively, are read out in response to reading by CPU0 and CPU1.

4.2.2 Semaphore Registers 0 to 31 (SEMR0 to SEMR31)

Semaphore registers 0 to 31 (SEMR0 to SEMR31) support exclusive control for resource access by the two CPUs.

Access to SEMR0 to SEMR31 by a given CPU does not interfere with the operation of the other CPU or the DMAC.

4.3 Operation

4.3.1 Initializing This LSI

Use the following procedure to initialize this LSI. A sample program for the procedure is given in figure 4.1.

- 1. After exit from the power-on reset state, the two CPUs execute power-on reset exception handling. The CPUs should execute the same exception-handling routine.
- 2. In the power-on reset exception handling routine, each CPU should identify itself as CPU0 or CPU1 by reading out CPUIDR and testing the value of the ID bit. The value read from the ID bit is 0 for CPU0 and 1 for CPU1.
- 3. Each CPU then branches to the corresponding processing routine.

; In the power-on reset exception handling routine, ; read ID bit in CPUIDR and check the value. MOVI20 #H'FFFC1404, R0 MOV.L MOV.L AND CMP/EQ R2, R1 BF BRA NOP ; Processing routine for CPU0 CPU0_ROUTINE: : : ; Processing routine for CPU1 CPU1_ROUTINE: : : @R0, R1 #H'40000000, R2 R2, R1 CPU0_ROUTINE CPU1_ROUTINE

Figure 4.1 Example of a Program for Initialization of This LSI

4.3.2 Exclusive Control for CPUs

(1) Using the Semaphore Registers for Exclusive Control of CPU Access to Resources

A procedure for exclusive control of resource access by the two CPUs is given below. A sample program for this procedure is shown in figure 4.2.

- 1. In the initialization routine for either of the CPUs, set all of the SEMF bits in SEMR0 to SEMR31 to 1 (this indicates that all resources are free).
- 2. For example, assume that SEMR0 is used for semaphore control of resource A and that CPU0 wants to use resource A. In this case, CPU0 should read the SEMF bit in SEMR0 repeatedly until the bit is read as 1
- 3. CPU0 recognizes that it has read 1 from the SEMF bit in SEMR0. This clears the SEMF bit to 0.
- 4. CPU0 then uses resource A. While resource A is in use by CPU0, CPU1 can only read 0 (resource A is in use) from the SEMF bit of SEMR0, and thus cannot use resource A.
- 5. After CPU0 has finished using resource A, it sets the SEMF bit in SEMR0 to 1 (resource A is free).

; Initialization routine Make initial settings ; 1. Initialize SEMR0 to SEMR31 MOVI20 #H'FFFC1E00, R0 MOV MOV.B R1,@(H'000, R0) ; SEMR0.SEMF = 1 MOV.B MOV.B MOV.B : : MOV.B ; 2. Read SEMR0 ; 3. Make sure that 1 has been read from SEMF bit in SEMR0 MOVI20 #H'FFFC1E00, R0 LOOP: BLD.B **BF** ; 4. Use resource A : : ; 5. Set SEMF bit in SEMR0 to 1 MOVI20 MOV MOV.B #H'01, R1 R1,@(H'004, R0) ; SEMR1.SEMF = 1 $R1, \omega$ (H'008, R0) ; SEMR2 SEMF = 1 R1,@(H'00C, R0) ; SEMR3.SEMF = 1 R1,@(H'07C, R0) ; SEMR31.SEMF = 1 #0,@(H'000, R0) LOOP #H'FFFC1E00, R0 #H'01, R1 R1,@(H'000, R0)

Figure 4.2 Example of a Program for Exclusive Control

(2) Notes

As a general precaution in exclusive control, pay attention to ensuring that the system does not enter a deadlock. For example, the system enters a deadlock in the following case.

- 1. CPU0 is to use currently available resource A and thus reads 1 from the SEMF bit in SEMR0 (this operation clears the SEMF bit in SEMR0 to 0).
- 2. CPU1 is to use currently available resource B and thus reads 1 from the SEMF bit in SEMR1 (this operation clears the SEMF bit in SEMR1 to 0).
- 3. CPU0 is to use resource B and thus reads the SEMF bit in SEMR1, but it keeps reading the bit as 0 since it was cleared in step 2.
- 4. CPU1 is to use resource A and thus reads the SEMF bit in SEMR0, but it keeps reading the bit as 0 since it was cleared in step 1.

Section 5 Clock Pulse Generator (CPG)

This LSI has a clock pulse generator (CPG) that generates CPU0 clock $(I_0 \phi)$, CPU1 clock $(I_1 \phi)$, a peripheral clock (Pφ), and a bus clock (Bφ). The CPG consists of a crystal oscillator, PLL circuits, and divider circuits.

5.1 Features

• Four clock operating modes

The mode is selected from among the four clock operating modes based on the frequency range to be used and the input clock type: the clock from the crystal resonator, the external clock or the clock for USB.

• Three clocks generated independently

CPU0 clock $(I_0 \phi)$ and CPU1 clock $(I_1 \phi)$ for the CPU and cache, a peripheral clock $(P\phi)$ for peripheral modules, and a bus clock ($B\phi = CKIO$) for the external bus interface can be generated independently.

• Frequency change function

CPU0 clock, CPU1 clock, and peripheral clock frequencies can be changed independently using the PLL (phase locked loop) circuits and divider circuits within the CPG. Frequencies are changed by software using the settings of the frequency control registers 0 and 1 (FRQCR0 and FROCR1).

• Power-down mode control

The clock can be stopped and specific modules can be stopped using the module standby function in power-down modes. For details on clock control in power-down mode, see section 30, Power-Down Modes.

Figure 5.1 Block Diagram of Clock Pulse Generator

The clock pulse generator blocks function as follows:

(1) Crystal Oscillator

The crystal oscillator is used in which the crystal resonator is connected to the XTAL/EXTAL pin or USB_X1/USB_X2 pin. One of them is selected according to the clock operating mode.

(2) Divider 1

Divider 1 divides the output from the crystal oscillator or the external clock input. The division ratio depends on the clock operating mode.

(3) PLL Circuit

PLL circuit multiplies the frequency of the output from the divider 1. The multiplication ratio is set by the frequency control register.

(4) Divider 2

Divider 2 generates a clock signal whose operating frequency can be used for the CPU0 clock, CPU1 clock, peripheral clock, and bus clock. The division ratio of the CPU0 clock, CPU1 clock, and peripheral clock are set by the frequency control register. The division ratio of the bus clock is determined by the clock operating mode and the PLL multiplication ratio.

(5) Clock Frequency Control Circuit

The clock frequency control circuit controls the clock frequency using the MD_CLK0 and MD_CLK1 pins and the frequency control registers 0 and 1 (FRQCR0 and FRQCR1).

(6) Standby Control Circuit

The standby control circuit controls the states of the on-chip oscillation circuit and other modules in power-down modes.

In addition, the standby control register is provided to control power-down mode of other modules. For details on the standby control register, see section 30, Power-Down Modes.

(7) Frequency Control Register 0 (FRQCR0)

The frequency control register 0 (FRQCR0) has control bits assigned for the following functions: clock output/non-output from the CKIO pin during software standby mode, the frequency multiplication rate of the PLL circuit, and the frequency division ratio of the CPU0 clock $(I_0 \phi)$ and the peripheral clock $(P\phi)$.

(8) Frequency Control Register 1 (FRQCR1)

The frequency control register 1 (FRQCR1) has control bits assigned for the following function: frequency division ratio of the CPU1 clock $(I_1\phi)$.

5.2 Input/Output Pins

Table 5.1 lists the clock pulse generator pins and their functions.

Table 5.1 Pin Configuration and Functions of the Clock Pulse Generator

5.3 Clock Operating Modes

Table 5.2 shows the relationship between the combinations of the mode control pins (MD_CLK1 and MD_CLK0) and the clock operating modes. Table 5.3 shows the usable frequency ranges in clock operating modes.

Table 5.2 Clock Operating Modes

Mode 0

In mode 0, a clock is input from the EXTAL pin or the crystal oscillator. The PLL circuit shapes waveforms and the frequency is multiplied according to the frequency control register setting before the clock is supplied to the LSI. The oscillating frequency for the crystal resonator and EXTAL pin input clock ranges from 10 to 16.67 MHz. The frequency range of CKIO is from 40 to 66.66 MHz. The internal clock frequency is the EXTAL pin frequency multiplied by the frequency multiplication rate of the PLL circuit and the division ratio of the divider 2. To reduce current, fix the USB_X1 pin (pull up, pull down, connect to power supply, or connect to ground) and open the USB_X2 pin when USB is not used.

Mode 1

In mode 1, a clock is input from the EXTAL pin or the crystal oscillator. The PLL circuit shapes waveforms and the frequency is multiplied according to the frequency control register setting before the clock is supplied to the LSI. The oscillating frequency for the crystal resonator and EXTAL pin input clock ranges from 20 to 33.33 MHz. The frequency range of CKIO is from 40 to 66.66 MHz. The internal clock frequency is half the EXTAL pin frequency multiplied by the frequency multiplication rate of the PLL circuit and the division ratio of the divider 2. To reduce current, fix the USB_X1 pin (pull up, pull down, connect to power supply, or connect to ground) and open the USB_X2 pin when USB is not used.

Mode 2

In mode 2, the CKIO pin functions as an input pin and draws an external clock signal. The PLL circuit shapes waveforms and the frequency is multiplied according to the frequency control register setting before the clock is supplied to the LSI. The frequency range of CKIO is from 40 to 66.66 MHz. The internal clock frequency is quarter the CKIO pin frequency multiplied by the frequency multiplication rate of the PLL circuit and the division ratio of the divider 2. To reduce current, fix the EXTAL pin (pull up, pull down, connect to power supply, or connect to ground) and open the XTAL pin when the LSI is used in mode 2. When USB is not used, fix the USB_X1 pin (pull up, pull down, connect to power supply, or connect to ground) and open the USB_X2 pin.

Mode 3

In mode 3, a clock is input from the USB X1 pin or the crystal oscillator. The external clock is input through this pin and a waveform is shaped in the PLL circuit. Then the frequency is multiplied according to the frequency control register setting before the clock is supplied to the LSI. The frequency of CKIO is 48 MHz (USB_X1/crystal resonator). The internal clock frequency is quarter the USB_X1 pin frequency multiplied by the frequency multiplication rate of the PLL circuit and the division ratio of the divider 2. To reduce current, fix the EXTAL pin (pull up, pull down, connect to power supply, or connect to ground) and open the XTAL pin when the LSI is used in mode 3. When the USB crystal resonator is not used, open the USB_X2 pin.

Table 5.3 Relationship between Clock Operating Mode and Frequency Range

PLL

PLL

2. The ratio of clock frequencies, where the input clock frequency is assumed to be 1.

 3. In mode 0 or 1, the frequency of the EXTAL pin input clock or the crystal resonator In mode 2, the frequency of the CKIO pin input clock. In mode 3, the frequency of the USB_X1 pin input clock or the crystal resonator

Cautions: Do not use this LSI for frequency settings other than those in table 5.3.

5.4 Register Descriptions

The clock pulse generator has the following registers.

Table 5.4 Register Configuration

5.4.1 Frequency Control Registers 0 and 1 (FRQCR0 and FRQCR1)

(1) FRQCR0

FRQCR0 is a 16-bit readable/writable register used to specify whether a clock is output from the CKIO pin during normal operation mode, changes in the multiplication rate of the PLL circuit, software standby mode, and standby mode cancellation. The register also specifies the frequency multiplication rate of the PLL circuit and the frequency division ratio for the CPU0 clock $(I_0 \phi)$ and peripheral clock (Pφ). The FRQCR0 register should be changed only from CPU0.

Note: * The initial value depends on clock mode.

Table 5.5 CKOEN[1:0] Settings

(2) FRQCR1

FRQCR1 is a 16-bit readable/writable register used to specify the frequency division ratio of the CPU1 clock $(I_1\phi)$. The FRQCR1 register should be changed only from CPU1.

Note: * The initial value depends on clock mode.

5.5 Changing the Frequency

The frequencies of the CPU0 clock (I_{ϕ}) , CPU1 clock (I_{ϕ}) , and peripheral clock (P ϕ) can be changed either by changing the multiplication rate of the PLL circuit or by changing the division ratio of the divider. All of these are controlled by software through the frequency control registers 0 and 1 (FRQCR0 and FRQCR1). The methods are described below.

The multiplication rate and division ratio must be changed so that the register values satisfy the conditions shown in table 5.3. Otherwise, the operation is not guaranteed.

5.5.1 Changing the Multiplication Rate

An oscillation stabilization time is required when the multiplication rate of the PLL circuit is changed. On-chip WDT0 counts the stabilization time. The oscillation stabilization time becomes the same time as that of recovery from the software standby mode.

When changing the multiplication rate, after setting IFC1 and IFC0 in FRQCR1 to B'00 and specifying CPU1 not to be interrupted, execute the SLEEP command from CPU1, confirm that single processor 0 mode (CPU1 is in sleep mode) (for details, see section 30, Power-Down Modes), and perform the following procedure from CPU0.

- 1. In the initial state, the multiplication rate of the PLL circuit is 12 times in clock modes 0, 1, and 2 or 16 times in clock mode 3.
- 2. Set a value that will produce the specified oscillation stabilization time in WDT0 for CPU0 and stop WDT0. The following must be set:

 $WTCSR0. TME = 0$: WDT stops

WTCSR0.CKS[2:0]: Division ratio for WDT counter clock

WTCNT0 counter: Initial counter value

(The WDT0 count is incremented using the clock after the setting.)

- 3. Set the desired value in the STC[1:0] bits of FRQCR0. The division ratios can also be set in the IFC[1:0] and PFC[2:0] of FRQCR0.
- 4. This LSI pauses temporarily and WDT0 starts to increment. A clock is supplied to the WDT only, and the other internal clocks all stop. The clock will continue to be output at the CKIO pin. Low level output can also be selected by setting CKOEN2 of FRQCR0. This state is the same as software standby mode. Whether or not registers are initialized depends on the module. For details, see section 32.3, Register States in Each Operating Mode.
- 5. Supply of the clock that has been set begins at WDT0 count overflow, and CPU0 of this LSI begins to operate again. WDT0 stops after it overflows. At this time, WOVF of WRCR0 is not set. The counter (WTCNT0) stops at H'00.
- 6. To change the WTCNT0 value after executing a frequency changing instruction, read WTCNT0 to make sure that it holds a value of H'00.
- 7. Since the CPU1 stays in the sleep state after execution of the frequency changing instruction, wake it up by using an interrupt or other means before using it.

Figure 5.2 shows a sample procedure for changing the multiplication rate of the PLL circuit from \times 12 to \times 16 in clock mode 0.

Figure 5.2 Sample Procedure for Changing the Multiplication Rate of the PLL Circuit from ×**12 to** ×**16**

5.5.2 Changing the Division Ratio

Counting by WDT0 does not proceed if the frequency division ratio is changed but the multiplication rate is not. However, when changing the division ratio of a peripheral clock, after specifying CPU1 not to be interrupted, place CPU1 in the sleep state and perform the operation from CPU0. When changing the division ratio of only the CPU1 clock, it is not necessary to place CPU1 in the sleep state.

- When the CPU0 clock is changed from CPU0
	- 1. The initial state depends on clock mode. See table 5.4.
	- 2. Set the desired values in the IFC1 and IFC0 bits and the PFC2 to PFC0 bits of FRQCR0. The values that can be set are limited by the clock operating mode and the multiplication rate of the PLL circuit. Note that if the wrong value is set, this LSI will malfunction.
	- 3. After the register bits (IFC[1:0] and PFC[2:0] of FRQCR0) have been set, the clock generated with the new division ratio is supplied.
- When the CPU1 clock is changed from CPU1
	- 1. The initial state depends on clock mode. See table 5.4.
	- 2. Set the desired value in IFC[1:0] of FRQCR1. The values that can be set are limited by clock operating mode and the multiplication rate of the PLL circuit. Note that if the wrong value is set, this LSI will malfunction.
	- 3. After the register bits (IFC[1:0] of FRQCR1) have been set, the clock generated with the new division ratio is supplied.
- Note: When executing the SLEEP instruction after changing the frequency, read the frequency control register 0 (FRQCR0) or frequency control register 1 (FRQCR1) three times and then execute the SLEEP instruction.

5.5.3 Notes on Changing the Multiplication Rate and Division Ratio

- 1. When the division ratio for the CPU1 clock is changed, if IFC[1:0] of FRQCR1 are changed while CPU1 is in the sleep state, the change is not reflected. To prevent malfunction, always change the FRQCR1 register from CPU1.
- 2. When the multiplication rate or division ratio is changed through the frequency control registers 0 and 1 (FRQCR0 and FRQCR1) while the DMAC is transferring data, the DMA transfer is not guaranteed because the frequency is changed without waiting for the completion of the DMA transfer. Therefore, to change the multiplication rate or division ratio through the frequency control registers 0 and 1 (FRQCR0 and FRQCR1), wait for the completion of the DMA transfer or stop the DMA transfer and then change the frequency control registers 0 and 1 (FRQCR0 and FRQCR1).

5.6 Usage of the Clock Pins

For the connection of a crystal resonator or the input of a clock signal, this LSI circuit has the pins listed in table 5.6. With regard to these pins, take care on the following points. Furthermore, Xin pin and Xout pin are used in this section to refer to the pins listed in the table.

Table 5.6 Clock Pins

5.6.1 In the Case of Inputting an External Clock

An example of the connection of an external clock is shown in figure 5.3. In cases where the Xout pin is left open state, take the parasitic capacitance as less than 10 pF.

Figure 5.3 Example of the Connection of an External Clock

5.6.2 In the Case of Using a Crystal Resonator

An example of the connection of crystal resonator is shown in figure 5.4.

Place the crystal resonator and capacitors (CL1 and CL2) as close to pins Xin and Xout as possible. Furthermore, to avoid inductance so that oscillation is correct, use the points where the capacitors are connected to the crystal resonator in common and do not place wiring patterns close to these components.

Since the design of the user board is closely connected with the effective characteristics of the crystal resonator, refer to the example of connection of the crystal resonator that is introduced in this section and perform thorough evaluation on the user side as well. The rated value of the crystal resonator will vary with the floating capacitances and so on of the crystal resonator and mounted circuit, so proceed with decisions on the basis of full discussions with the maker of the crystal resonator. Ensure that voltages applied to the clock pins do not exceed the maximum rated values.

Although the feedback resistor is included in this LSI, an external feedback resistor may be required in some cases. This depends on the characteristics of the crystal resonator.

Set the parameters (of resistors and capacitors) with thorough evaluation on the user side.

5.6.3 In the Case of Not Using the Clock Pin

In cases where the pins are not in use, fix the level on the Xin pin (pull it up or down, or connect it to the power-supply or ground level), and leave the Xout pin open state.

5.7 Oscillation Stabilizing Time

5.7.1 Oscillation Stabilizing Time of the On-chip Crystal Oscillator

In the case of using a crystal resonator, please wait longer than the oscillation stabilizing time at the following cases, to keep the oscillation stabilizing time of the on-chip crystal oscillator (In the case of inputting an external clock input, it is not necessary).

- Power on
- Canceling software standby mode or deep standby mode by using the RES or MRES pin
- Changing from halting oscillation to running oscillation by power-on reset or register setting (AUDIO_X1, RTC_X1)

5.7.2 Oscillation Stabilizing Time of the PLL circuit

In clock mode 0 or 1 the clock input on EXTAL, in clock mode 2 the clock input on CKIO, and in clock mode 3 the clock input on USB_X1 is supplied to the PLL circuit. So, regardless of whether using a crystal resonator or inputting an external clock from EXTAL (clock mode 0 and 1) or USB_X1 (clock mode 3), please wait longer than the oscillation stabilizing time at the following cases, to keep the oscillation stabilizing time of the PLL circuit.

- Power on (in the case of using the crystal resonator)/start inputting external clock (in the case of inputting the external clock)
- Canceling software standby mode or deep standby mode by using the \overline{RES} or \overline{MRES} pin
- Changing the multiplication ratio of the PLL circuit by power-on reset from RES pin

[Remarks]

The oscillation stabilizing time is kept by the counter running in the LSI at the following cases.

• Canceling software standby mode or deep standby mode by using the RES or MRES pin

5.8 Notes on Board Design

5.8.1 Note on Using a PLL Oscillation Circuit

In the PLLVcc and PLLVss connection pattern for the PLL, signal lines from the board power supply pins must be as short as possible and pattern width must be as wide as possible to reduce inductive interference.

Since the analog power supply pins of the PLL are sensitive to the noise, the system may malfunction due to inductive interference at the other power supply pins. To prevent such malfunction, the analog power supply pin Vcc and digital power supply pin PVcc should not supply the same resources on the board if at all possible.

Ensure that PLLVcc has the same electric potential as Vcc.

Figure 5.5 Note on Using a PLL Oscillation Circuit

Section 6 Exception Handling

6.1 Overview

6.1.1 Types of Exception Handling and Priority

Exception handling is started by sources, such as resets, address errors, register bank errors, interrupts, and instructions as shown in table 6.1. Since the exception sources have their priorities as shown in table 6.1, when multiple exception source events coincide, they are processed according to the priority shown below.

Type	Exception Handling	Priority
Reset	Power-on reset	High
	Manual reset	
	Address error CPU address error	
Instruction	FPU exception	
	Integer division exception (division by zero)	
	Integer division exception (overflow)	
error	Register bank Bank underflow	
	Bank overflow	
Sleep error	Sleep error	
Interrupt	NMI	
	User break	
	H-UDI	
	Inter-processor	
	IRQ	
	PINT	
	On-chip peripheral modules	Low

Table 6.1 Types of Exception Handling and Priority Order

- 2. Instructions that rewrite the PC: JMP, JSR, BRA, BSR, RTS, RTE, BT, BF, TRAPA, BF/S, BT/S, BSRF, BRAF, JSR/N, RTV/N
- 3. 32-bit instructions: BAND.B, BANDNOT.B, BCLR.B, BLD.B, BLDNOT.B, BOR.B, BORNOT.B, BSET.B, BST.B, BXOR.B, MOV.B@disp12, MOV.W@disp12, MOV.L@disp12, MOVI20, MOVI20S, MOVU.B, MOVU.W

6.1.2 Exception Handling Operations

Exception sources are detected and their corresponding processing is started with the timing shown in table 6.2.

When exception handling starts, the CPU operates as follows:

(1) Exception Handling Triggered by Reset

The initial values of the program counter (PC) and stack pointer (SP) are fetched from the exception handling vector table (PC and SP are respectively at the H'00000000 and H'00000004 addresses for power-on resets and the H'00000008 and H'0000000C addresses for manual resets). See section 6.1.3, Exception Handling Vector Table, for more information. The vector base register (VBR) is then initialized to H'00000000, the interrupt mask level bits (I3 to I0) of the status register (SR) are initialized to H'F (B'1111), and the BO and CS bits are initialized to 0. The BN bit in IBNR of the interrupt controller (INTC) is also initialized to 0. The program begins running from the PC address fetched from the exception handling vector table.

(2) Exception Handling Triggered by Address Errors, Sleep Errors, Register Bank Errors, Interrupts, and Instructions

SR and PC are saved to the stack indicated by R15. In the case of interrupt exception handling other than NMI or UBC with usage of the register banks enabled, general registers R0 to R14, control register GBR, system registers MACH, MACL, and PR, and the vector table address offset of the interrupt exception handling to be executed are saved to the register banks. In the case of exception handling due to an address error, sleep error, register bank error, NMI interrupt, UBC interrupt, or instruction, saving to a register bank is not performed. When saving is performed to all register banks, automatic saving to the stack is performed instead of register bank saving. In this case, an interrupt controller setting must have been made so that register bank overflow exceptions are not accepted (the BOVE bit in IBNR of the INTC is 0). If a setting to accept register bank overflow exceptions has been made (the BOVE bit in IBNR of the INTC is 1), register bank overflow exception will be generated. In the case of interrupt exception handling, the interrupt priority level is written to the I3 to I0 bits in SR. In the case of exception handling due to an address error, sleep error, register bank error, or instruction, the I3 to I0 bits are not affected. The start address is then fetched from the exception handling vector table and the program begins running from that address.

6.1.3 Exception Handling Vector Table

Before exception handling begins running, the exception handling vector table must be set in memory. The exception handling vector table stores the start addresses of exception service routines. (The reset exception handling table holds the initial values of PC and SP.)

All exception sources are given different vector numbers and vector table address offsets, from which the vector table addresses are calculated. During exception handling, the start addresses of the exception service routines are fetched from the exception handling vector table, which is indicated by this vector table address.

Table 6.3 shows the vector numbers and vector table address offsets. Table 6.4 shows how vector table addresses are calculated.

Table 6.3 Exception Handling Vector Table

Notes: 1. Both CPUs: Indicates that exception handling is performed on both CPU0 and CPU1. Each CPU: Indicates that exception handling is performed on CPU0 when the exception source occurred on CPU0 or performed on CPU1 when the exception source occurred on CPU1.

> User selection A: Indicates that only CPU0, only CPU1, or both CPU0 and CPU1 can be selected to perform exception handling.

 User selection B: Indicates that only CPU0 or only CPU1 can be selected to perform exception handling.

 2. The vector numbers and vector table address offsets for each external interrupt and onchip peripheral module interrupt are given in table 7.8.

Table 6.4 Calculating Exception Handling Vector Table Addresses

Notes: 1. Vector table address offset: See table 6.3.

2. Vector number: See table 6.3.

6.2 Resets

6.2.1 Input/Output Pins

Table 6.5 shows the reset-related pin configuration.

Table 6.5 Pin Configuration

6.2.2 Types of Reset

A reset is the highest-priority exception source. There are two kinds of reset, power-on and manual. As shown in table 6.6, the CPU state is initialized in both a power-on reset and a manual reset. On-chip peripheral module registers are also initialized by a power-on reset, but not by a manual reset.

Table 6.6 Reset States

Notes: 1. See section 32.3, Register States in Each Operating Mode.

2. Data are retained when the setting of either the RAME or RAMWE bit is disabled.

6.2.3 Power-On Reset

(1) Power-On Reset by Means of RES **Pin**

When the RES pin is driven low, this LSI enters the power-on reset state. To reliably reset this LSI, the RES pin should be kept at the low level for the duration of the oscillation settling time at power-on or in software standby mode (when the clock is halted), or at least 20-tcyc when the clock is running. In the power-on reset state, the internal state of the CPU and all the on-chip peripheral module registers are initialized. See appendix A, Pin States, for the status of individual pins during the power-on reset state.

In the power-on reset state, power-on reset exception handling starts when the RES pin is first driven low for a fixed period and then returned to high. The CPU operates as follows:

- 1. The initial value (execution start address) of the program counter (PC) is fetched from the exception handling vector table.
- 2. The initial value of the stack pointer (SP) is fetched from the exception handling vector table.
- 3. The vector base register (VBR) is cleared to H'00000000, the interrupt mask level bits (I3 to I0) of the status register (SR) are initialized to H'F (B'1111), and the BO and CS bits are initialized to 0. The BN bit in IBNR of the INTC is also initialized to 0.
- 4. The values fetched from the exception handling vector table are set in the PC and SP, and the program begins executing.

Be certain to always perform power-on reset processing when turning the system power on. For a recommended flow of power-on reset processing, see section 4, Multi-Core Processor.

(2) Power-On Reset by Means of H-UDI Reset Assert Command

When the H-UDI reset assert command is set, this LSI enters the power-on reset state. Power-on reset by means of an H-UDI reset assert command is equivalent to power-on reset by means of the RES pin. Setting the H-UDI reset negate command cancels the power-on reset state. The time required between an H-UDI reset assert command and H-UDI reset negate command is the same as the time to keep the RES pin low to initiate a power-on reset. In the power-on reset state generated by an H-UDI reset assert command, setting the H-UDI reset negate command starts power-on reset exception handling. The CPU operates in the same way as when a power-on reset was caused by the RES pin.

(3) Power-On Reset Initiated by WDT

Each CPU has a watchdog timer (WDT).

When either or both of the WDTs are set so that a power-on reset occurs in watchdog timer mode, and the WTCNT (or WTCNTs) of the WDT (or WDTs) overflows, this LSI enters the power-on reset state.

In this case, WRCSR of the WDT and FRQCR of the CPG are not initialized by the reset signal generated by the WDT.

If a reset caused by the RES pin or the H-UDI reset assert command occurs simultaneously with a reset caused by WDT overflow, the reset caused by the RES pin or the H-UDI reset assert command has priority, and the WOVF bit in WRCSR is cleared to 0. When power-on reset exception handling is started by the WDT, the CPU operates in the same way as when a power-on reset was caused by the RES pin.

6.2.4 Manual Reset

(1) Manual Reset by Means of MRES **Pin**

When the MRES pin is driven low, this LSI enters the manual reset state. To reset this LSI without fail, the MRES pin should be kept at the low level for at least 20-tcyc. In the manual reset state, the CPU's internal state is initialized, but the on-chip peripheral module registers are not initialized. In the manual reset state, manual reset exception handling starts when the MRES pin is first driven low for a fixed period and then returned to high. The CPU operates as follows:

- 1. The initial value (execution start address) of the program counter (PC) is fetched from the exception handling vector table.
- 2. The initial value of the stack pointer (SP) is fetched from the exception handling vector table.
- 3. The vector base register (VBR) is cleared to H'00000000, the interrupt mask level bits (I3 to I0) of the status register (SR) are initialized to H'F (B'1111), and the BO and CS bits are initialized to 0. The BN bit in IBNR of the INTC is also initialized to 0.
- 4. The values fetched from the exception handling vector table are set in the PC and SP, and the program begins executing.

(2) Manual Reset Initiated by WDT

Each CPU has a watchdog timer (WDT).

When either or both of the WDTs are set so that a manual reset occurs in watchdog timer mode, and the WTCNT (or WTCNTs) of the WDT (or WDTs) overflows, this LSI enters the manual reset state.

When manual reset exception handling is started by the WDT, the CPU operates in the same way as when a manual reset was caused by the MRES pin.

(3) Note in Manual Reset

When a manual reset is generated, the bus cycle is retained, but if a manual reset occurs during DMAC burst transfer, manual reset exception handling will be deferred until the CPU acquires the bus mastership. The CPU and the BN bit in IBNR of the INTC are initialized by a manual reset. The FPU and other modules are not initialized.

6.3 Address Errors

6.3.1 Address Error Sources

Address errors occur when instructions are fetched or data read or written, as shown in table 6.7.

Table 6.7 Bus Cycles and Address Errors

Bus Cycle

see section 10, Bus State Controller (BSC).

6.3.2 Address Error Exception Handling

When an address error occurs, the bus cycle in which the address error occurred ends*, the executing instruction finishes, and address error exception handling starts. The CPU operates as follows:

- 1. The exception service routine start address which corresponds to the address error that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the last executed instruction.
- 4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.
- Note: * In the case of address error related to data read/write. In the case of address error related to instruction fetch, if the bus cycle in which the address error occurred doesn't end until the entire three above-mentioned operations end, the CPU will start address error exception handling again until the bus cycle in which the address error occurred ends.

6.4 Register Bank Errors

6.4.1 Register Bank Error Sources

(1) Bank Overflow

In the state where saving has already been performed to all register bank areas, bank overflow occurs when acceptance of register bank overflow exception has been set by the interrupt controller (the BOVE bit in IBNR of the INTC is set to 1) and an interrupt that uses a register bank has occurred and been accepted by the CPU.

(2) Bank Underflow

Bank underflow occurs when an attempt is made to execute an RESBANK instruction while saving has not been performed to register banks.

6.4.2 Register Bank Error Exception Handling

When a register bank error occurs, register bank error exception handling starts. The CPU operates as follows:

- 1. The exception service routine start address which corresponds to the register bank error that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the last executed instruction for a bank overflow, and the start address of the executed RESBANK instruction for a bank underflow.

To prevent multiple interrupts from occurring at a bank overflow, the interrupt priority level that caused the bank overflow is written to the interrupt mask level bits (I3 to I0) of the status register (SR).

4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.

6.5 Sleep Errors

6.5.1 Sleep Error Source

A sleep error occurs if issuance of the sleep instruction by CPU0 is detected when the sleep error occurrence notification for CPU0 is set and the sleep error enable bit (SLPERE) of the standby control register 1 (STBCR1) is 1. For details, see section 30, Power-Down Modes.

6.5.2 Sleep Error Exception Handling

When a sleep error occurs, sleep error exception handling starts after the bus cycle in which the sleep error occurred ends and the execution of the current instruction is completed. The CPU operates as follows:

- 1. The exception service routine start address which corresponds to the sleep error that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the last executed instruction.
- 4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.
- 5. Clear the sleep error enable bit (SLPERE) of the standby control register 1 (STBCR1) to 0 in the sleep error exception handling routine.

To detect a sleep error again, set the sleep error enable bit of the standby control register 1 to 1 after the corresponding sleep instruction of CPU0.

6.6 Interrupts

6.6.1 Interrupt Sources

The interrupt sources that trigger interrupt exception handling are NMI, user breaks, H-UDI, interprocessor interrupts, IRQ, PINT, and on-chip peripheral modules.

Each interrupt source is assigned a different vector number and vector table offset. See table 7.8, for more information on vector numbers and vector table address offsets.

6.6.2 Interrupt Priority Level

The interrupt priority order is predetermined. When multiple interrupts occur simultaneously (overlap), the interrupt controller (INTC) determines their relative priorities and starts exception handling according to the results.

The priority order of interrupts is expressed as priority levels 0 to 16, with priority 0 the lowest and priority 16 the highest. The NMI interrupt has priority 16 and cannot be masked, so it is always accepted. The user break interrupt and H-UDI interrupt priority level is 15. The interprocessor interrupt has an interrupt priority level of 15 to 8 according to the interrupt source. Priority levels of IRQ interrupts, PINT interrupts, and on-chip peripheral module interrupts can be set freely using the interrupt priority registers of the INTC (table 6.8). The priority levels that can be set are 0 to 15. Level 16 cannot be set. See section 7.3.1, Interrupt Priority Registers 01, 02, 05 to 21 (C0IPR01, C0IPR02, C0IPR05 to C0IPR21, C1IPR01, C1IPR02, C1IPR05 to C1IPR21), for details of the interrupt priority registers.

Table 6.8 Interrupt Priority Order

6.6.3 Interrupt Exception Handling

When an interrupt occurs, its priority level is determined by the interrupt controller (INTC). NMI is always accepted, but other interrupts are only accepted if they have a priority level higher than the priority level set in the interrupt mask level bits (I3 to I0) of the status register (SR).

When an interrupt is accepted, interrupt exception handling begins. In interrupt exception handling, the CPU fetches the exception service routine start address which corresponds to the accepted interrupt from the exception handling vector table, and saves the SR and the program counter (PC) to the stack. In the case of interrupt exception handling other than NMI or UBC with usage of the register banks enabled, general registers R0 to R14, control register GBR, system registers MACH, MACL, and PR, and the vector table address offset of the interrupt exception handling to be executed are saved in the register banks. In the case of exception handling due to an address error, NMI interrupt, UBC interrupt, or instruction, saving is not performed to the register banks. If saving has been performed to all register banks (0 to 14), automatic saving to the stack is performed instead of register bank saving. In this case, an interrupt controller setting must have been made so that register bank overflow exceptions are not accepted (the BOVE bit in IBNR of the INTC is 0). If a setting to accept register bank overflow exceptions has been made (the BOVE bit in IBNR of the INTC is 1), register bank overflow exception occurs. Next, the priority level value of the accepted interrupt is written to the I3 to I0 bits in SR. For NMI, however, the priority level is 16, but the value set in the I3 to I0 bits is H'F (level 15). Then, after jumping to the start address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch. See section 7.6, Operation, for further details of interrupt exception handling.

6.7 Exceptions Triggered by Instructions

6.7.1 Types of Exceptions Triggered by Instructions

Exception handling can be triggered by the trap instruction, slot illegal instructions, general illegal instructions, integer division exceptions, and FPU exceptions, as shown in table 6.9.

Table 6.9 Types of Exceptions Triggered by Instructions

6.7.2 Trap Instruction

When a TRAPA instruction is executed, trap instruction exception handling starts. The CPU operates as follows:

- 1. The exception service routine start address which corresponds to the vector number specified in the TRAPA instruction is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the TRAPA instruction.
- 4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.

6.7.3 Slot Illegal Instructions

An instruction placed immediately after a delayed branch instruction is said to be placed in a delay slot. When the instruction placed in the delay slot is undefined code, an instruction that rewrites the PC, a 32-bit instruction, an RESBANK instruction, a DIVS instruction, or a DIVU instruction, slot illegal exception handling starts when such kind of instruction is decoded. In slot illegal instruction exception handling, the CPU operates as follows:

- 1. The exception service routine start address is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the jump address of the delayed branch instruction immediately before the undefined code, the instruction that rewrites the PC, the 32-bit instruction, the RESBANK instruction, the DIVS instruction, or the DIVU instruction.
- 4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.

6.7.4 General Illegal Instructions

When an undefined code placed anywhere other than immediately after a delayed branch instruction, i.e., in a delay slot, is decoded, general illegal instruction exception handling starts. In general illegal instruction exception handling, the CPU handles general illegal instructions in the same way as slot illegal instructions. Unlike processing of slot illegal instruction exception handling, however, the program counter value stored is the start address of the undefined code.

6.7.5 Integer Division Exceptions

When an integer division instruction performs division by zero or the result of integer division overflows, integer division instruction exception handling starts. The instructions that may become the source of division-by-zero exception are DIVU and DIVS. The only source instruction of overflow exception is DIVS, and overflow exception occurs only when the negative maximum value is divided by −1. The CPU operates as follows:

- 1. The exception service routine start address which corresponds to the integer division exception that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the integer division instruction at which the exception occurred.
- 4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.

6.7.6 FPU Exceptions

An FPU exception handling is generated when the V, Z, O, U or I bit in the FPU enable field (Enable) of the floating point status register (FPSCR) is set. This indicates the occurrence of an invalid operation exception defined by the IEEE standard 754, a division-by-zero exception, overflow (in the case of an instruction for which this is possible), underflow (in the case of an instruction for which this is possible), or inexact exception (in the case of an instruction for which this is possible).

The floating-point operation instructions that may cause generation of an FPU exception handling are FADD, FSUB, FMUL, FDIV, FMAC, FCMP/EQ, FCMP/GT, FLOAT, FTRC, FCNVDS, FCNVSD, and FSQRT.

An FPU exception handling is generated only when the corresponding FPU exception enable bit (enabled) is set. When the FPU detects an exception source by a floating-point operation, FPU operation is halted and FPU exception handling generation is reported to the CPU. When exception handling is started, the CPU operations are as follows.

- 1. The start address of the exception service routine corresponding to the FPU exception handling that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the last executed instruction.

4. After jumping to the address fetched from the exception handling vector table, program execution starts. This jump is not a delayed branch.

The FPU exception flag field (Flag) of FPSCR is always updated regardless of whether or not an FPU exception handling has been accepted, and remains set until explicitly cleared by the user through an instruction. The FPU exception source field (Cause) of FPSCR changes each time a floating-point operation is executed. When the V bit in the FPU exception enable field (Enable) of FPSCR is set and the QIS bit in FPSCR is also set, FPU exception handling is generated when qNaN or ±∞ is input to a floating point operation instruction source.

6.8 When Exception Sources Are Not Accepted

When an address error, sleep error, FPU exception, register bank error (overflow), or interrupt is generated immediately after a delayed branch instruction, it is sometimes not accepted immediately but stored instead, as shown in table 6.10. When this happens, it will be accepted when an instruction that can accept the exception is decoded.

Table 6.10 Exception Source Generation Immediately after Delayed Branch Instruction

Note: * Delayed branch instructions: JMP, JSR, BRA, BSR, RTS, RTE, BF/S, BT/S, BSRF, BRAF

6.9 Stack Status after Exception Handling Ends

The status of the stack after exception handling ends is as shown in table 6.11.

Table 6.11 Stack Status after Exception Handling Ends

6.10 Usage Notes

6.10.1 Value of Stack Pointer (SP)

The value of the stack pointer must always be a multiple of four. If it is not, an address error will occur when the stack is accessed during exception handling.

6.10.2 Value of Vector Base Register (VBR)

The value of the vector base register must always be a multiple of four. If it is not, an address error will occur when the stack is accessed during exception handling.

6.10.3 Address Errors Caused by Stacking of Address Error Exception Handling

When the stack pointer is not a multiple of four, an address error will occur during stacking of the exception handling (interrupts, etc.) and address error exception handling will start up as soon as the first exception handling is ended. Address errors will then also occur in the stacking for this address error exception handling. To ensure that address error exception handling does not go into an endless loop, no address errors are accepted at that point. This allows program control to be shifted to the address error exception service routine and enables error processing.

When an address error occurs during exception handling stacking, the stacking bus cycle (write) is executed. During stacking of the status register (SR) and program counter (PC), the SP is decremented by 4 for both, so the value of SP will not be a multiple of four after the stacking either. The address value output during stacking is the SP value, so the address where the error occurred itself is output. This means the write data stacked will be undefined.

6.10.4 Interrupt Control via Modification of Interrupt Mask Bits

When enabling interrupts by changing the Interrupt Mask bits (I3-I0) of the Status Register (SR) using the LDC or LDC.L instructions, interrupts might not be accepted during the execution of the 5 instructions immediately after the LDC/LDC.L instruction.

Therefore, when enabling/disabling interrupts by changing the Interrupt Mask bits (I3-I0) of the Status Register (SR) using LDC/LDC.L instructions, please place at least 5 instructions between the interrupt-enable instruction and the interrupt-disable instruction.

6.10.5 Note before Exception Handling Begins Running

Before exception handling begins running, the exception handling vector table must be stored in a memory, and the CPU must be able to access the memory. So, if the exception handling is generated

- Ex. 1: when the exception handling vector table is stored in an external address space, but the settings of bus state controller and general I/O ports to access the external address space have been not completed yet, or
- Ex. 2: when the exception handling vector table is stored in the on-chip RAM, but the vector base register (VBR) has been not changed to the on-chip RAM address yet,

the CPU fetches an unintended value as the execution start address, and starts executing programs from unintended address.

Section 7 Interrupt Controller (INTC)

The interrupt controller (INTC) identifies the priorities of interrupt sources and controls interrupt requests to the CPU. The INTC has registers used to set interrupt priorities; interrupt requests are processed according to the priorities set in these registers by the user.

7.1 Features

• 16 levels of interrupt priority can be set.

By setting 19 interrupt priority registers, the priorities of IRQ interrupts, PINT interrupts, and on-chip peripheral module interrupts can be selected from 16 levels for individual request sources.

• NMI noise canceller function

An NMI input-level bit indicates the NMI pin state. The pin state can be checked by reading this bit in the interrupt exception service routine, and this LSI can be used for the noise canceller function.

• Register banks

The LSI has two sets of register banks, one set for each CPU core, that enable register contents to be saved and restoration processing to be performed at high speed for the interrupt processing.

Inter-processor interrupts

By configuring the inter-processor interrupt control registers, inter-processor interrupts can be generated with programmed priority levels of 15 to 8.

Figure 7.1 shows a block diagram of the INTC.

Figure 7.1 Block Diagram of INTC

7.2 Input/Output Pins

Table 7.1 shows the pin configuration of the INTC.

Table 7.1 Pin Configuration

7.3 Register Descriptions

The INTC has the following registers. These registers are used to set the interrupt priorities and control detection of the external interrupt input signals. The registers are classified as the following: CPU0-dedicated, CPU1-dedicated, and shared.

(1) CPU0-Dedicated Registers

Table 7.2 CPU0-Dedicated Register Configuration

Notes: 1. The initial value is either H'8000 when the NMI pin is high, or H'0000 when the NMI pin is low.

2. To clear the flag, only 0 can be written after 1 is read.

(2) CPU1-Dedicated Registers

Table 7.3 CPU1-Dedicated Register Configuration

Notes: 1. The initial value is either H'8000 when the NMI pin is high, or H'0000 when the NMI pin is low.

2. To clear the flag, only 0 can be written after 1 is read.

(3) Shared Registers

Table 7.4 Shared Register Configuration

7.3.1 Interrupt Priority Registers 01, 02, 05 to 21 (C0IPR01, C0IPR02, C0IPR05 to C0IPR21, C1IPR01, C1IPR02, C1IPR05 to C1IPR21)

C0IPR01, C0IPR02, and C0IPR05 to C0IPR21 and C1IPR01, C1IPR02, and C1IPR05 to C0IPR21 are 16-bit readable/writable registers in which priority levels from 0 to 15 are set for IRQ interrupts, PINT interrupts, and on-chip peripheral module interrupts. Table 7.5 shows the correspondence between the interrupt request sources and the bits in C0IPR01, C0IPR02, and C0IPR05 to C0IPR21 and C1IPR01, C1IPR02, and C1IPR05 to C0IPR21.

Table 7.5 Interrupt Request Sources and C0IPR01, C0IPR02, C0IPR05 to C0IPR21, C1IPR01, C1IPR02, and C1IPR05 to C0IPR21

As shown in table 7.5, by setting the 4-bit groups (bits 15 to 12, bits 11 to 8, bits 7 to 4, and bits 3 to 0) with values from H'0 (0000) to H'F (1111), the priority of each corresponding interrupt is set. Setting of H'0 means priority level 0 (the lowest level) and H'F means priority level 15 (the highest level).

7.3.2 Interrupt Control Registers 0 (C0ICR0, C1ICR0)

C0ICR0 and C1ICR0 are 16-bit registers that set the input signal detection mode for the external interrupt input pin NMI, and indicate the input level at the NMI pin.

Note: $*$ The initial value is either 1 when the NMI pin is high, or 0 when the NMI pin is low.

7.3.3 Interrupt Control Registers 1 (C0ICR1, C1ICR1)

C0ICR1 and C1ICR1 are 16-bit registers that specify the detection mode for external interrupt input pins IRQ7 to IRQ0 individually: falling edge, rising edge, both edges, or low level.

7.3.4 Interrupt Control Registers 2 (C0ICR2, C1ICR2)

C0ICR2 and C1ICR2 are 16-bit registers that specify the detection mode for external interrupt input pins PINT7 to PINT0 individually: low level or high level.

[Legend]

7.3.5 IRQ Interrupt Request Registers (C0IRQRR, C1IRQRR)

C0IRQRR and C1IRQRR are 16-bit registers that indicate interrupt requests from external interrupt input pins IRQ7 to IRQ0. If edge detection is set for the IRQ7 to IRQ0 interrupts, the retained interrupt requests can be cancelled by reading 1 from the IRQ7F to IRQ0F bits and then writing 0 to these bits. However, this register is enabled only when C0IRQER and C1IRQER accept an interrupt request input. When an interrupt request input is disabled, this register always becomes 0.

7.3.6 PINT Interrupt Enable Registers (C0PINTER, C1PINTER)

C0PINTER and C1PINTER are 16-bit registers that enable interrupt request inputs to external interrupt input pins PINT7 to PINT0.

[Legend]

7.3.7 PINT Interrupt Request Registers (C0PIRR, C1PIRR)

C0PIRR and C1PIRR are 16-bit registers that indicate interrupt requests from external interrupt input pins PINT7 to PINT0.

However, this register is enabled only when C0PINTER and C1PINTER accept an interrupt request input. When an interrupt request input is disabled, this register always becomes 0.

[Legend]

7.3.8 Bank Control Registers (C0IBCR, C1IBCR)

C0IBCR and C1IBCR are 16-bit registers that enable or disable the use of register banks for each interrupt priority level.

7.3.9 Bank Number Registers (C0IBNR, C1IBNR)

C0IBNR and C1IBNR are 16-bit registers that enable or disable the use of register banks and register bank overflow exception. In bits BN3 to BN0, C0IBNR and C1IBNR indicate the number of the bank to which saving is performed next.

7.3.10 Inter-Processor Interrupt Control Registers 15 to 08 (C0IPCR15 to C0IPCR08, C1IPCR15 to C1IPCR08)

C0IPCR15 to C0IPCR08 and C1IPCR15 to C1IPCR08 are 16-bit registers that generate interprocessor interrupts when 1 is written to any of the CI bits. Each CI bit remains 1 until the target processor accepts interrupt processing, and the bit is cleared to 0 upon completion of the acceptance.

An inter-processor interrupt request made from CPU0 is set in one of C1IPCR15 to C1IPCR08. An inter-processor interrupt request made from CPU1 to CPU0 is set in one of C0IPCR15 to C0IPCR08.

Note: Although 0 can be written to the CI bit, an inter-processor interrupt request is held pending internally, and cannot be cleared.

The interrupt priorities for the registers are predefined as shown below.

Table 7.6 Interrupt Priorities for Registers

7.3.11 Inter-processor Interrupt Enable Registers (C0IPER, C1IPER)

C0IPER and C1IPER are 16-bit registers that enable or disable inter-processor interrupts of each interrupt priority level. The interrupt controller decides whether to accept interrupts, according to the inter-processor interrupt enable settings. C0IPER enables or disables interrupts to CPU0, while C1IPER enables or disables interrupts to CPU1.

7.3.12 Interrupt Enable Control Registers (C0INTER, C1INTER)

C0INTER and C1INTER are 16-bit registers that control whether to enable or disable acceptance of interrupt requests by processors CPU0 and CPU1. If the same bits in both registers C0INTER and C1INTER are set to 0, the acceptance by CPU0 is enabled.

Note: * The initial value is 1 for C0INTER and 0 for C1INTER.

7.3.13 IRQ Interrupt Enable Control Registers (C0IRQER, C1IRQER)

C0IRQER and C1IRQER are 16-bit registers that control whether to enable or disable acceptance of IRQ interrupt requests by processors CPU0 and CPU1. If the same bits in both registers C0IRQER and C1IRQER are set to 0, the acceptance by CPU0 is enabled.

[Legend]

 $n = 7$ to 0

Note: * The initial value is 1 for C0IRQER and 0 for C1IRQER.

7.3.14 Interrupt Detect Control Registers (IDCNT6 to IDCNT139)

IDCNT6 to IDCNT139 (except IDCNT65 and 127 to 129) are 16-bit registers that control whether to enable interrupt requests from on-chip peripheral modules and also control which CPU should accept the requests.

Table 7.7 shows the correspondence between the sources of on-chip peripheral module interrupt requests and the IDCNT registers.

Table 7.7 Correspondence between Sources of On-Chip Peripheral Module Interrupt Requests and IDCNT Registers

7.3.15 DMA Transfer Request Enable Registers 0 to 8 (DREQER0 to DREQER8)

DREQER0 to DREQER8 are 8-bit readable/writable registers that enable or disable DMA transfer requests from on-chip peripheral modules, and enable or disable CPU interrupts.

(1) DREQER0

(2) DREQER1

(3) DREQER2

(4) DREQER3

(5) DREQER4

(6) DREQER5

(7) DREQER6

(8) DREQER7

(9) DREQER8

7.4 Interrupt Sources

There are six types of interrupt sources: NMI, user break, H-UDI, IRQ, PINT, and on-chip peripheral modules. Each interrupt has a priority level (0 to 16), with 0 the lowest and 16 the highest. When set to level 0, that interrupt is masked at all times.

7.4.1 NMI Interrupts

An NMI interrupt has a priority level of 16 and is accepted at all times. NMI pin inputs are edgedetected, and the NMI sense select bits (NMIS) in interrupt control registers 0 (C0ICR0 and C1ICR0) select whether interrupt requests are detected on the rising edge or falling edge. The CPU that should accept the NMI interrupts can be selected by the interrupt enable control registers (C0INTER and C1INTER).

Though the priority level of the NMI interrupt is 16, the NMI interrupt exception handling sets the interrupt mask bits (I3 to I0) in the status register (SR) to level 15.

7.4.2 User Break Interrupts

A user break interrupt has a priority level of 15, and occurs when a break condition set in the user break controller (UBC) is satisfied. User break interrupt requests are edge-detected and retained until they are accepted. The user break exception handling sets the I3 to I0 bits in the SR to level 15. For user break interrupts, see section 8, User Break Controller (UBC).

7.4.3 H-UDI Interrupts

A user debugging interface (H-UDI) interrupt has a priority level of 15, and occurs at serial input of an H-UDI interrupt instruction. H-UDI interrupt requests are edge-detected and retained until they are accepted. The CPU that should accept the H-UDI interrupts can be selected by the interrupt enable control registers (C0INTER and C1INTER). The H-UDI exception handling sets the I3 to I0 bits in the SR to level 15. For H-UDI interrupts, see section 31, User Debugging Interface (H-UDI).

7.4.4 IRQ Interrupts

IRQ interrupts are input from pins IRQ7 to IRQ0. For an explanation of how to configure pins IRQ7 to IRQ0, see section 27, Pin Function Controller (PFC). For IRQ7 to IRQ0, low-level, falling-edge, rising-edge, or both-edge detection can be selected individually for each pin by the IRQ sense select bits (IRQ71S to IRQ01S and IRQ70S to IRQ00S) in interrupt control registers 1 (C0ICR1 and C1ICR1). The priority level can be set individually in a range from 0 to 15 for each pin by interrupt priority registers 01 and 02 (C0IPR01, C0IPR02, C1IPR01, and C1IPR02). The CPU that should accept the IRQ interrupts can be selected by the IRQ interrupt enable control registers (C0IRQER and C1IRQER).

When low-level detection is used for IRQ interrupts, an interrupt request signal is sent to the INTC while the IRQ7 to IRQ0 pins are low. An interrupt request signal is no longer sent to the INTC when the IRQ7 to IRQ0 pins are driven high. The status of the interrupt requests can be checked by reading the IRQ interrupt request bits (IRQ7F to IRQ0F) in the IRQ interrupt request registers (C0IRQRR and C1IRQRR).

When edge-detection is used for IRQ interrupts, an interrupt request is detected due to change of the IRQ7 to IRQ0 pin states, and an interrupt request signal is sent to the INTC. The result of IRQ interrupt request detection is retained until that interrupt request is accepted. Whether IRQ interrupt requests have been detected or not can be checked by reading IRQ7F to IRQ0F in C0IRQRR and C1IRQRR. The result of IRQ interrupt request detection can be cleared by reading 1 from these bits and then writing 0 to them.

The IRQ interrupt exception handling sets the I3 to I0 bits in the SR to the priority level of the accepted IRQ interrupt.

When returning from the IRQ interrupt exception service routine, execute the RTE instruction after using C0IRQRR and C1IRQRR to ensure that the interrupt request has been cleared, so as not to accidentally receive the interrupt request again.

7.4.5 PINT Interrupts

PINT interrupts are input from pins PINT7 to PINT0. For an explanation of how to configure pins PINT7 to PINT0, see section 27, Pin Function Controller (PFC). Input of the interrupt requests is enabled by the PINT enable bits (PINT7E to PINT0E) in the PINT interrupt enable registers (C0PINTER and C1PINTER). For PINT7 to PINT0, low-level or high-level detection can be selected individually for each pin by the PINT sense select bits (PINT7S to PINT0S) in interrupt control registers 2 (C0ICR2 and C1ICR2). A single priority level in a range from 0 to 15 can be set for all PINT7 to PINT0 interrupts by bits 15 to 12 in interrupt priority registers 05 (C0IPR05 and C1IPR05).

When low-level detection is used for the PINT7 to PINT0 interrupts, an interrupt request signal is sent to the INTC while the PINT7 to PINT0 pins are low. An interrupt request signal is no longer sent to the INTC when the PINT7 to PINT0 pins are driven high. The status of the interrupt requests can be checked by reading the PINT interrupt request bits (PINT7R to PINT0R) in the PINT interrupt request registers (C0PIRR and C1PIRR). The above description also applies to a case in which high-level detection is used, except for the polarity being reversed. The PINT interrupt exception handling sets the I3 to I0 bits in the SR to the priority level of the PINT interrupt.

When returning from the PINT interrupt exception service routine, execute the RTE instruction after using C0PIRR and C1PIRR to ensure that the interrupt request has been cleared, so as not to accidentally receive the interrupt request again.

7.4.6 On-Chip Peripheral Module Interrupts

On-chip peripheral module interrupts are generated by the following on-chip peripheral modules:

- Direct memory access controller (DMAC)
- USB2.0 host/function module (USB)
- Compare match timer (CMT)
- Watchdog timer (WDT)
- Multi-function timer pulse unit 2 (MTU2)
- I²C bus interface 3 (IIC3)
- Serial communications interface with FIFO (SCIF)
- Serial sound interface with FIFO (SSIF)
- Synchronous serial communications unit (SSU)
- A/D converter (ADC)
- 2D engine (2DG)
- AT attachment packet interface (ATAPI)
- AND/NAND flash memory controller (FLCTL)
- Realtime clock (RTC)
- Controller area network (RCAN-TL1)

As every source is assigned a different interrupt vector, the source does not need to be identified in the exception service routine. A priority level in a range from 0 to 15 can be set for each module by interrupt priority registers 06 to 21 (C0IPR06 to C0IPR21 and C1IPR06 to C1IPR21). The onchip peripheral module interrupt exception handling sets the I3 to I0 bits in the SR to the priority level of the accepted on-chip peripheral module interrupt.

7.4.7 Inter-Processor Interrupts

Inter-processor interrupts are generated by setting the inter-processor interrupt control registers (C0IPCR15 to C0IPCR08 and C1IPCR15 to C1IPCR08). Interrupts can be generated from CPU0 to CPU1 and vise versa.

When the inter-processor interrupt enable registers (C0IPER and C1IPER) are set, interrupt requests from the inter-processor interrupt control registers are enabled and sent to the CPU.

7.5 Interrupt Exception Handling Vector Tables and Priorities

Table 7.8 lists interrupt sources and their vector numbers, vector table address offsets, and interrupt priorities.

Each interrupt source is allocated a different vector number and vector table address offset. Vector table addresses are calculated from the vector numbers and vector table address offsets. In interrupt exception handling, the exception service routine start address is fetched from the vector table indicated by the vector table address. For details of calculation of the vector table address, see table 6.4.

The priorities of IRQ interrupts, PINT interrupts, and on-chip peripheral module interrupts can be set freely between 0 and 15 for each pin or module by setting interrupt priority registers 01, 02, and 05 to 21 (C0IPR01, C0IPR02, and C0IPR05 to C0IPR21 and C1IPR01, C1IPR02, and C1IPR05 to C1IPR21). However, if two or more interrupts specified by the same IPR setting among C0IPR05 to C0IPR21 and C1IPR05 to C1IPR21 occur, the priorities are defined as shown in the default priorities in table 7.8, and the priorities cannot be changed. A power-on reset assigns priority level 0 to IRQ interrupts, PINT interrupts, and on-chip peripheral module interrupts. If the same priority level is assigned to two or more interrupt sources and interrupts from those sources occur simultaneously, they are processed according to the default priorities indicated in table 7.8.

Table 7.8 Interrupt Exception Vectors and Priorities

7.6 Operation

7.6.1 Interrupt Operation Sequence

The interrupt operation sequence is described below. Figure 7.2 shows the operation flow.

- 1. The interrupt request sources send interrupt request signals to the interrupt controller.
- 2. The interrupt controller selects the highest-priority interrupt from the sent interrupt requests, according to the priority levels set in interrupt priority registers 01, 02, and 05 to 21 (C0IPR01, C0IPR02, and C0IPR05 to C0IPR21 and C1IPR01, C1IPR02, and C1IPR05 to C1IPR21). Lower priority interrupts are ignored*. If two of more interrupts have the same priority level or if two or more interrupts specified by the same IPR setting occur, the interrupt with the highest priority is selected, according to the default priorities shown in table 7.8.
- 3. The priority level of the interrupt selected by the interrupt controller is compared with the interrupt mask level bits (I3 to I0) in the status register (SR) of the CPU. If the priority level is equal to or lower than the level set in bits I3 to I0, the interrupt is ignored. Only when the priority level is higher than the level in bits I3 to I0, the interrupt controller accepts the interrupt and sends an interrupt request signal to the CPU.
- 4. The CPU detects the interrupt request sent from the interrupt controller when the CPU decodes the instruction to be executed. Instead of executing the decoded instruction, the CPU starts interrupt exception handling (figure 7.4).
- 5. The exception service routine start address is fetched from the exception handling vector table corresponding to the accepted interrupt.
- 6. The status register (SR) is saved onto the stack, and the priority level of the accepted interrupt is written to bits I3 to I0 in the SR.
- 7. The program counter (PC) is saved onto the stack.
- 8. The CPU jumps to the fetched exception service routine start address and starts executing the program. The jump that occurs is not a delayed branch.
- Notes: The interrupt source flag should be cleared in the interrupt handler. After clearing the interrupt source flag, "time from occurrence of interrupt request until interrupt controller identifies priority, compares it with mask bits in the SR, and sends interrupt request signal to CPU" shown in table 7.9 is required before the interrupt source sent to the CPU is actually cancelled. To ensure that an interrupt request that should have been cleared is not inadvertently accepted again, read the interrupt source flag after it has been cleared, and then execute an RTE instruction.

 * Interrupt requests that are set for edge-detection are held pending until they are accepted. IRQ interrupts, however, can be cancelled by accessing the IRQ interrupt request registers (C0IRQRR and C1IRQRR). For details, see section 7.4.4, IRQ Interrupts.

Interrupts held pending due to edge-detection are cleared by a power-on reset.

7.6.2 Stack Status after Interrupt Exception Handling

Figure 7.3 shows the stack status after interrupt exception handling.

Figure 7.3 Stack Status after Interrupt Exception Handling

7.7 Interrupt Response Time

Table 7.9 lists the interrupt response time, which is the time from the occurrence of an interrupt request until the interrupt exception handling starts and fetching of the first instruction in the exception service routine begins. The interrupt processing operations differ in the cases when banking is disabled, when banking is enabled without register bank overflow, and when banking is enabled with register bank overflow. Figures 7.4 and 7.5 show examples of pipeline operation when banking is disabled. Figures 7.6 and 7.7 show examples of pipeline operation when banking is enabled without register bank overflow. Figures 7.8 and 7.9 show examples of pipeline operation when banking is enabled with register bank overflow.

Table 7.9 Interrupt Response Time

Number of States*¹

Notes: m1 to m4 are the number of states needed for the following memory accesses.

- m1: Vector address read (longword read)
- m2: SR save (longword write)
- m3: PC save (longword write)
- m4: Restoration of banked registers (R0 to R14, GBR, MACH, MACL, and PR) from the stack
- 1. n in Incyc indicates the number (0 or 1) of the CPU to which an interrupt request is sent.
- 2. Case where $In\phi: B\phi: P\phi = 200 \text{ MHz}: 66 \text{ MHz}: 33 \text{ MHz}$ and $m1 = m2 = m3 = m4 = 1 \text{lncyc}$

Figure 7.4 Example of Pipeline Operation when IRQ Interrupt is Accepted (No Register Banking)

Figure 7.6 Example of Pipeline Operation when IRQ Interrupt is Accepted (Register Banking without Register Bank Overflow)

Figure 7.7 Example of Pipeline Operation when Interrupt is Accepted during RESBANK Instruction Execution (Register Banking without Register Bank Overflow)

Figure 7.8 Example of Pipeline Operation when IRQ Interrupt is Accepted (Register Banking with Register Bank Overflow)

Figure 7.9 Example of Pipeline Operation when Interrupt is Accepted during RESBANK Instruction Execution (Register Banking with Register Bank Overflow)

7.8 Register Banks

The LSI has two sets of 15 register banks, one set for each CPU core, to enable saving and restoring of register data at high speed during interrupt handling. Figure 7.10 shows the register bank configuration.

Figure 7.10 Overview of Register Bank Configuration

7.9 Register Banks and Bank Control Registers

(1) Banked Registers

The general registers (R0 to R14), global base register (GBR), multiply-and-accumulate registers (MACH and MACL), procedure register (PR), and the vector table address offset are banked.

(2) Input/Output of Banks

This LSI has 15 register banks, bank 0 to bank 14. Register banks are stacked in first-in last-out (FILO) sequence. Saving takes place in order, beginning from bank 0, and restoration takes place in the reverse order, beginning from the last bank saved to.

7.9.1 Bank Save and Restore Operations

(1) Saving to Bank

Figure 7.11 shows register bank save operation. The following operation is performed when the CPU accepts an interrupt and the use of register banks is enabled for that interrupt.

- a. Assume that the values of the bank number bits (BN) in the bank number registers (C0IBNR and C1IBNR) are i before the interrupt is generated.
- b. The values in registers R0 to R14, GBR, MACH, MACL, and PR, and the vector table address offset (VTO) of the accepted interrupt are saved to bank i indicated by BN.
- c. The BN value is incremented by 1.

Figure 7.11 Bank Save Operation

Figure 7.12 shows the timing for saving to a register bank. Saving to a register bank takes place between the start of interrupt exception handling and the start of fetching the first instruction in the exception service routine.

(2) Restoration from Bank

The RESBANK (restore from register bank) instruction is used to restore data saved in a register bank. After restoring data from the register banks with the RESBANK instruction at the end of the interrupt service routine, execute the RTE instruction to return from the exception handling.

7.9.2 Save and Restore Operations after Saving to All Banks

If the CPU accepts an interrupt and the use of the register banks is enabled for that interrupt when saving to all register banks has been performed, automatic saving to the stack is performed instead of register bank saving if the BOVE bits in the bank number registers (C0IBNR and C1IBNR) are cleared to 0. If the BOVE bits in C0IBNR and C1IBNR are set to 1, a register bank overflow exception occurs and data is not saved to the stack.

Saving to the stack and restoration from the stack take place as described below:

(1) Saving to Stack

- 1. The status register (SR) and program counter (PC) are saved to the stack during interrupt exception handling.
- 2. The values in the banked registers (R0 to R14, GBR, MACH, MACL, and PR) are saved to the stack. The register values are saved to the stack in the order of MACL, MACH, GBR, PR, R14, R13, …, R1, and R0.
- 3. The register bank overflow bit (BO) in the SR is set to 1.
- 4. The bank number bits (BN) in the bank number registers (C0IBNR and C1IBNR) remain set to the maximum value of 15.

(2) Restoration from Stack

When the RESBANK (restore from register bank) instruction is executed with the register bank overflow bit (BO) in the SR set to 1, the following operation is performed:

- 1. The values in the banked registers (R0 to R14, GBR, MACH, MACL, and PR) are restored from the stack. The register values are restored from the stack in the order of R0, R1, …, R13, R14, PR, GBR, MACH, and MACL.
- 2. The bank number bits (BN) in the bank number registers (C0IBNR and C1IBNR) remain set to the maximum value of 15.

7.9.3 Register Bank Exceptions

There are two types of register bank exceptions (register bank errors): register bank overflow and register bank underflow.

(1) Register Bank Overflow

This exception occurs if, after data has been saved to all of the register banks, the CPU accepts an interrupt and the use of the register banks is enabled for that interrupt, and the BOVE bits in the bank number registers (C0IBNR and C1IBNR) are set to 1. In this case, the bank number bits (BN) in the bank number registers (C0IBNR and C1IBNR) remain set to the bank count of 15 and saving to the register bank is not performed.

(2) Register Bank Underflow

This exception occurs if the RESBANK (restore from register bank) instruction is executed when no data has been saved to the register banks. In this case, the values of R0 to R14, GBR, MACH, MACL, and PR do not change. In addition, the bank number bits (BN) in the bank number registers (C0IBNR and C1IBNR) remain set to 0.

7.10 Register Bank Error Exception Handling

When a register bank error occurs, register bank error exception handling starts. When this happens, the CPU operates as follows:

- 1. The exception service routine start address is fetched from the exception handling vector table corresponding to the register bank error that has occurred.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. For a register bank overflow, the saved PC value is the start address of the instruction to be executed after the last executed instruction. For a register bank underflow, the saved PC value is the start address of the executed RESBANK instruction. To prevent multiple interrupts from occurring at a register bank overflow, the priority level of the interrupt that caused the register bank overflow is written to the interrupt mask level bits (I3 to I0) of the status register (SR).
- 4. Program execution starts from the exception service routine start address.

7.11 Data Transfer with Interrupt Request Signals

Interrupt request signals can be used to activate the DMAC and transfer data.

DMA transfer request enable registers 0 to 8 (DREQER0 to DREQER8) are used to specify whether the interrupt request signals start interrupt exception handling or activate the DMAC. When the bits corresponding to on-chip peripheral modules are set to 1, DMA transfer requests are generated; when these bits are set to 0, CPU interrupt requests are generated.

7.12 Usage Note

7.12.1 Timing to Clear an Interrupt Source

The interrupt source flags should be cleared in the interrupt handler. After clearing the interrupt source flag, "time from occurrence of interrupt request until interrupt controller identifies priority, compares it with mask bits in the SR, and sends interrupt request signal to CPU" shown in table 7.9 is required before the interrupt source sent to the CPU is actually cancelled. To ensure that an interrupt request that should have been cleared is not inadvertently accepted again, read the interrupt source flag after it has been cleared, and then execute an RTE instruction.

Section 8 User Break Controller (UBC)

The user break controller (UBC) provides functions that simplify program debugging. These functions make it easy to design a self-monitoring debugger, enabling this LSI chip to debug programs without using an in-circuit emulator. Instruction fetch or data read/write of CPU, data size, data contents, address value, and stop timing in the case of instruction fetch are break conditions that can be set in the UBC. Since this LSI uses a Harvard architecture, instruction fetch on the CPU bus (C bus) is performed by issuing bus cycles on the instruction fetch bus (F bus), and data access on the C bus is performed by issuing bus cycles on the memory access bus (M bus). The UBC monitors the C bus and internal bus (I bus).

There are two UBCs: UBC0, which monitors the operation of CPU0, and UBC1, which monitors the operation of CPU1. These UBCs are quite the same. The control registers of UBC0 and UBC1 are mapped to the same addresses, but the registers for UBC0 are accessed when access from CPU0 is made and the registers for UBC1 are accessed when access from CPU1 is made. In this section, UBC0 and UBC1 are collectively called UBC.

8.1 Features

1. The following break comparison conditions can be set. Number of break channels: two channels (channels 0 and 1)

User break can be requested as the independent condition on channels 0 and 1.

— Address

Comparison of the 32-bit address is maskable in 1-bit units.

One of the three address buses (F address bus (FAB), M address bus (MAB), and I address bus (IAB)) can be selected.

— Data

Comparison of the 32-bit data is maskable in 1-bit units.

One of the two data buses (M data bus (MDB) and I data bus (IDB)) can be selected.

— Bus cycle

Instruction fetch (only when C bus is selected) or data access

- Read/write
- Operand size
	- Byte, word, and longword
- 2. In an instruction fetch cycle, it can be selected whether break is set before or after execution of an instruction.
- 3. When a break condition is satisfied, a trigger signal can be output from the \overline{UBCTRG} pin.

Figure 8.1 shows a block diagram of the UBC.

Figure 8.1 Block Diagram of UBC (for One CPU)

8.2 Input/Output Pin

Table 8.1 shows the pin configuration of the UBC.

Table 8.1 Pin Configuration

8.3 Register Descriptions

The UBC has the following registers: five registers for each channel and a control register common to channels 0 and 1. These registers are provided for each of UBC0 and UBC1.

The channel of the UBC registers is indicated as follows: for example, BAR_0 represents the BAR register for channel 0.

8.3.1 Break Address Register (BAR)

BAR is a 32-bit readable/writable register. BAR specifies the address used as a break condition in each channel. Control bits CD1 and CD0 in the break bus cycle register (BBR) select one of the three address buses for a break condition.

Note: When setting the instruction fetch cycle as a break condition, clear the LSB in BAR to 0.

8.3.2 Break Address Mask Register (BAMR)

BAMR is a 32-bit readable/writable register. BAMR specifies the bits to be masked of the break address bits specified by BAR.

8.3.3 Break Data Register (BDR)

BDR is a 32-bit readable/writable register. Control bits CD1 and CD0 in the break bus cycle register (BBR) select one of the two data buses for a break condition.

Notes: 1. Set the operand size when specifying a value on a data bus as the break condition.

 2. When the byte size is selected as a break condition, the same byte data must be set in bits 31 to 24, 23 to 16, 15 to 8, and 7 to 0 in BDR as the break data. Similarly, when the word size is selected, the same word data must be set in bits 31 to 16 and 15 to 0.

8.3.4 Break Data Mask Register (BDMR)

BDMR is a 32-bit readable/writable register. BDMR specifies the bits to be masked of the break data bits specified by BDR.

Notes: 1. Set the operand size when specifying a value on a data bus as the break condition.

 2. When the byte size is selected as a break condition, the same byte data must be set in bits 31 to 24, 23 to 16, 15 to 8, and 7 to 0 in BDMR as the break mask data. Similarly, when the word size is selected, the same word data must be set in bits 31 to 16 and 15 to 0.

8.3.5 Break Bus Cycle Register (BBR)

BBR is a 16-bit readable/writable register, which specifies (1) disabling or enabling of user break interrupts, (2) including or excluding of the data bus value, (3) C bus cycle or I bus cycle, (4) instruction fetch or data access, (5) read or write, and (6) operand size as the break conditions.

8.3.6 Break Control Register (BRCR)

BRCR sets the following conditions:

- 1. Specifies whether a start of user break interrupt exception handling by instruction fetch cycle is set before or after instruction execution.
- 2. Specifies the pulse width of the UBCTRG output when a break condition is satisfied.

BRCR is a 32-bit readable/writable register that has break condition match flags and bits for setting other break conditions. For the condition match flags of bits 15 to 12, writing 1 is invalid (previous values are retained) and writing 0 is only possible. To clear the flag, write 0 to the flag bit to be cleared and 1 to all other flag bits.

8.4 Operation

8.4.1 Flow of the User Break Operation

The flow from setting of break conditions to user break exception handling is described below:

- 1. The break address is set in the break address register (BAR). The masked address bits are set in the break address mask register (BAMR). The break data is set in the break data register (BDR). The masked data bits are set in the break data mask register (BDMR). The bus break conditions are set in the break bus cycle register (BBR). No user break will be generated if any one of the three control bit pairs in BBR (C bus cycle/I bus cycle select, instruction fetch/data access select, and read/write select) is set to 00. The break control settings are made in the bits of the break control register (BRCR). Make sure to set all registers related to breaks before setting BBR, and branch after reading from the last written register. The newly written register values become valid from the instruction at the branch destination.
- 2. If a the break condition is satisfied, UBC0 (UBC1) sends a user break request to CPU0 (CPU1) through the INTC, sets the C bus condition match flag (SCMFC) or I bus condition match flag (SCMFD) for the appropriate channel, and outputs a pulse to the UBCTRG pin with the width set by the CKS[1:0] bits. Setting the UBID bit in BBR to 1 enables external monitoring of the trigger output without requesting user break interrupts.
- 3. On receiving a user break interrupt request signal, the INTC determines its priority. Since the user break interrupt has a priority level of 15, it is accepted when the priority level set in the interrupt mask level bits (I3 to I0) of the status register (SR) is 14 or lower. If the I3 to I0 bits are set to a priority level of 15, the user break interrupt is not accepted, but the conditions are checked, and condition match flags are set if the conditions match. For details on ascertaining the priority, see section 7, Interrupt Controller (INTC).
- 4. Condition match flags (SCMFC and SCMFD) can be used to check which condition has been satisfied. Clear the condition match flags during the user break interrupt exception handling routine. The interrupt occurs again if this operation is not performed.
- 5. There is a possibility that the break set in channel 0 and the break set in channel 1 occur around the same time. In this case, there will be only one user break request to the INTC, but these two break channel match flags may both be set.
- 6. When selecting the I bus as the break condition, note as follows:
	- Whether or not the access the CPU issued on the C bus is issued on the I bus depends on the setting of the cache. As regard to the I bus operation that depends on cache conditions, see table 9.8.
	- When a break condition is specified for the I bus, only the data access cycle is monitored. The instruction fetch cycle (including cache update cycle) is not monitored.

⎯ If a break condition is specified for the I bus, even when the condition matches in an I bus cycle resulting from an instruction executed by the CPU, at which instruction the user break interrupt request is to be accepted cannot be clearly defined.

8.4.2 Break on Instruction Fetch Cycle

- 1. When C bus/instruction fetch/read/word or longword is set in the break bus cycle register (BBR), the break condition is the FAB bus instruction fetch cycle. Whether a break is set before or after the execution of the instruction can be selected with the PCB0 or PCB1 bit in the break control register (BRCR) for the appropriate channel. If an instruction fetch cycle is set as a break condition, clear the BA0 bit in the break address register (BAR) to 0. A break cannot be generated as long as this bit is set to 1.
- 2. A break for instruction fetch which is set as a break before instruction execution occurs when it is confirmed that the instruction has been fetched and will be executed. This means a break does not occur for instructions fetched by overrun (instructions fetched at a branch or during an interrupt transition, but not to be executed). When this kind of break is set for the delay slot of a delayed branch instruction, the user break interrupt request is not received until the execution of the first instruction at the branch destination.
- Note: If a branch does not occur at a delayed branch instruction, the subsequent instruction is not recognized as a delay slot.
- 3. When setting a break condition for break after instruction execution, the instruction that matched the break condition is executed and then the break is generated prior to execution of the next instruction. As with pre-execution breaks, a break does not occur with overrun fetch instructions. When this kind of break is set for a delayed branch instruction and its delay slot, the user break interrupt request is not received until the first instruction at the branch destination.
- 4. When an instruction fetch cycle is set, the break data register (BDR) is ignored. Therefore, break data cannot be set for the break of the instruction fetch cycle.
- 5. If the I bus is set for a break of an instruction fetch cycle, the setting is invalidated.

8.4.3 Break on Data Access Cycle

- 1. If the C bus is specified as a break condition for data access break, condition comparison is performed for the logical addresses (and data) accessed by the executed instructions, and a break occurs if the condition is satisfied. If the I bus is specified as a break condition, condition comparison is performed for the addresses (and data) of the data access cycles, and a break occurs if the condition is satisfied. For details on the CPU bus cycles issued on the I bus, see paragraph 6 in section 8.4.1, Flow of the User Break Operation.
- 2. The relationship between the data access cycle address and the comparison condition for each operand size is listed in table 8.3.

Table 8.3 Data Access Cycle Addresses and Operand Size Comparison Conditions

This means that when address H'00001003 is set in the break address register (BAR), for example, the bus cycle in which the break condition is satisfied is as follows (where other conditions are met).

Longword access at H'00001000

Word access at H'00001002

Byte access at H'00001003

3. When the data value is included in the break conditions:

When the data value is included in the break conditions, either longword, word, or byte is specified as the operand size in the break bus cycle register (BBR). When data values are included in break conditions, a break is generated when the address conditions and data conditions both match. To specify byte data for this case, set the same data in the four bytes at bits 31 to 24, 23 to 16, 15 to 8, and 7 to 0 of the break data register (BDR) and break data mask register (BDMR). To specify word data for this case, set the same data in the two words at bits 31 to 16 and 15 to 0.

- 4. Access by a PREF instruction is handled as read access in longword units without access data. Therefore, if data is included in the break condition, of PREF instruction, a break will not occur.
- 5. If the data access cycle is selected, the instruction at which the break will occur cannot be determined.

8.4.4 Value of Saved Program Counter

When a user break interrupt request is received, the address of the instruction from where execution is to be resumed is saved to the stack, and the exception handling state is entered. If the C bus (FAB)/instruction fetch cycle is specified as a break condition, the instruction at which the break should occur can be uniquely determined. If the C bus/data access cycle or I bus/data access cycle is specified as a break condition, the instruction at which the break should occur cannot be uniquely determined.

1. When C bus (FAB)/instruction fetch (before instruction execution) is specified as a break condition:

The address of the instruction that matched the break condition is saved to the stack. The instruction that matched the condition is not executed, and the break occurs before it. However, when a delay slot instruction matches the condition, the instruction is executed, and the branch destination address is saved to the stack.

2. When C bus (FAB)/instruction fetch (after instruction execution) is specified as a break condition:

The address of the instruction following the instruction that matched the break condition is saved to the stack. The instruction that matches the condition is executed, and the break occurs before the next instruction is executed. However, when a delayed branch instruction or delay slot matches the condition, the instruction is executed, and the branch destination address is saved to the stack.

3. When C bus/data access cycle or I bus/data access cycle is specified as a break condition: The address after executing several instructions of the instruction that matched the break condition is saved to the stack.

8.4.5 Usage Examples

(1) Break Condition Specified for C Bus Instruction Fetch Cycle

(Example 1-1)

• Register specifications

BAR_0 = H'00000404, BAMR_0 = H'00000000, BBR_0 = H'0054, BAR_1 = H'00008010, BAMR $1 = H'00000006$, BBR $1 = H'0054$, BDR_1 = H'00000000, BDMR_1 = H'00000000, BRCR = H'00000020 <Channel 0> Address: H'00000404, Address mask: H'00000000 Bus cycle: C bus/instruction fetch (after instruction execution)/read (operand size is not included in the condition) <Channel 1> Address: H'00008010, Address mask: H'00000006 Data: H'00000000, Data mask: H'00000000 Bus cycle: C bus/instruction fetch (before instruction execution)/read (operand size is not included in the condition)

A user break occurs after an instruction of address H'00000404 is executed or before instructions of addresses H'00008010 to H'00008016 are executed.

(Example 1-2)

• Register specifications

BAR_0 = H'00027128, BAMR_0 = H'00000000, BBR_0 = H'005A, BAR_1= H'00031415, BAMR $1 = H'00000000$, BBR $1 = H'0054$, BDR_1 = H'00000000, BDMR_1 = H'00000000, BRCR = H'00000000 <Channel 0> Address: H'00027128, Address mask: H'00000000 Bus cycle: C bus/instruction fetch (before instruction execution)/write/word <Channel 1> Address: H'00031415, Address mask: H'00000000 Data: H'00000000, Data mask: H'00000000 Bus cycle: C bus/instruction fetch (before instruction execution)/read (operand size is not included in the condition)

On channel 0, a user break does not occur since instruction fetch is not a write cycle. On channel 1, a user break does not occur since instruction fetch is performed for an even address. (Example 1-3)

• Register specifications

BAR_0 = H'00008404, BAMR_0 = H'00000FFF, BBR_0 = H'0054, BAR_1= H'00008010, BAMR $1 = H'00000006$, BBR $1 = H'0054$, BDR_1 = H'00000000, BDMR_1 = H'00000000, BRCR = H'00000020 <Channel 0> Address: H'00008404, Address mask: H'00000FFF Bus cycle: C bus/instruction fetch (after instruction execution)/read (operand size is not included in the condition) <Channel 1> Address: H'00008010, Address mask: H'00000006 Data: H'00000000, Data mask: H'00000000 Bus cycle: C bus/instruction fetch (before instruction execution)/read (operand size is not included in the condition)

A user break occurs after an instruction with addresses H'00008000 to H'00008FFE is executed or before an instruction with addresses H'00008010 to H'00008016 are executed.

(2) Break Condition Specified for C Bus Data Access Cycle

(Example 2-1)

• Register specifications

BAR_0 = H'00123456, BAMR_0 = H'00000000, BBR_0 = H'0064, BAR_1= H'000ABCDE, BAMR $1 = H'000000$ FF, BBR $1 = H'106$ A, BDR_1 = H'A512A512, BDMR_1 = H'00000000, BRCR = H'00000000 <Channel 0> Address: H'00123456, Address mask: H'00000000 Bus cycle: C bus/data access/read (operand size is not included in the condition) <Channel 1> Address: H'000ABCDE, Address mask: H'000000FF Data: H'0000A512, Data mask: H'00000000 Bus cycle: C bus/data access/write/word On channel 0, a user break occurs with longword read from address H'00123454, word read from address H'00123456, or byte read from address H'00123456. On channel 1, a user break occurs when word H'A512 is written in addresses H'000ABC00 to H'000ABCFE.

(3) Break Condition Specified for I Bus Data Access Cycle

(Example 3-1)

• Register specifications

BAR_0 = H'00314156, BAMR_0 = H'00000000, BBR_0 = H'0094, BAR_1= H'00055555, BAMR $1 = H'00000000$, BBR $1 = H'11A9$, BDR_1 = H'78787878, BDMR_1 = H'0F0F0F0F, BRCR = H'00000000 <Channel 0> Address: H'00314156, Address mask: H'00000000 Bus cycle: I bus/instruction fetch/read (operand size is not included in the condition) <Channel 1> Address: H'00055555, Address mask: H'00000000 Data: H'00000078, Data mask: H'0000000F Bus cycle: I bus/data access/write/byte On channel 0, the setting of I bus/instruction fetch is ignored. On channel 1, a user break

occurs when the CPU writes byte data H'7x in address H'00055555 on the I bus.

8.5 Usage Notes

- 1. The CPU can read from or write to the UBC registers via the I bus. Accordingly, during the period from executing an instruction to rewrite the UBC register till the new value is actually rewritten, the desired break may not occur. In order to know the timing when the UBC register is changed, read from the last written register. Instructions after then are valid for the newly written register value.
- 2. The UBC cannot monitor access to the C bus and I bus cycles in the same channel.
- 3. When a user break and another exception source occur at the same instruction, which has higher priority is determined according to the priority levels defined in table 6.1. If an exception source with higher priority occurs, the user break interrupt request is not received.
- 4. Note the following when a break occurs in a delay slot. If a pre-execution break is set at a delay slot instruction, the break does not occur before execution of the branch destination.
- 5. User breaks are disabled during UBC module standby mode. Do not read from or write to the UBC registers during UBC module standby mode; the values are not guaranteed.
- 6. Do not set an address within an interrupt exception handling routine whose interrupt priority level is at least 15 (including user break interrupts) as a break address.
- 7. Do not set break after instruction execution for the SLEEP instruction or for the delayed branch instruction where the SLEEP instruction is placed at its delay slot.
- 8. When setting a break address for a 32-bit instruction, set the address where the upper 16 bits are placed. If the address of the lower 16 bits is set and a break before instruction execution is set as a break condition, the break is handled as a break after instruction execution.
- 9. Do not set a pre-execution break for the instruction that comes after the DIVU or DIVS instruction. If a pre-execution break is set for the instruction that comes after the DIVU or DIVS instruction and an exception or interrupt occurs during execution of the DIVU or DIVS instruction, a pre-execution break occurs even though execution of the DIVU or DIVS instruction is halted.

Section 9 Cache

9.1 Features

• Capacity

Instruction cache: 8 Kbytes × 2 cores (CPU0/CPU1)

Operand cache: 8 Kbytes × 2 cores (CPU0/CPU1)

- Structure: Instructions/data separated, 4-way set associative
- Way lock function (only for operand cache): Way 2 and way 3 are lockable
- Line size: 16 bytes
- Number of entries: 128 entries/way
- Write system: Write-back/write-through selectable
- Replacement method: Least-recently-used (LRU) algorithm

9.1.1 Cache Structure

The cache separates data and instructions and uses a 4-way set associative system. It is composed of four ways (banks), each of which is divided into an address section and a data section.

The address and data sections per way are divided into 128 entries. The data section of the entry is called a line. Each line consists of 16 bytes (4 bytes \times 4). The data capacity per way is 2 Kbytes (16 bytes × 128 entries), which makes a total of 8 Kbytes as a whole cache (four ways).

There are two caches: cache 0 is incorporated in CPU0 and cache 1 is incorporated in CPU1. The two have the same functions.

Although the control registers for cache 0 and cache 1 are allocated to the same address, access from CUP0 will be to cache 0 and access from CPU1 will be to cache 1.

In this section, "cache" inclusively refers to both cache 0 and cache 1.

Figure 9.1 shows the operand cache structure. The instruction cache structure is the same as the operand cache structure except for not having the U bit.

Figure 9.1 Operand Cache Structure

(1) Address Array

The V bit indicates whether the entry data is valid. When the V bit is 1, data is valid; when 0, data is not valid.

The U bit (only for operand cache) indicates whether the entry has been written to in write-back mode. When the U bit is 1, the entry has been written to; when 0, it has not.

The tag address holds the physical address used in the external memory access. It consists of 21 bits (address bits 31 to 11) used for comparison during cache searches. In this LSI, the addresses of the cache-enabled space are H'00000000 to H'1FFFFFFF, and therefore the upper three bits of the tag address are cleared to 0.

The V and U bits are initialized to 0 by a power-on reset but not initialized by a manual reset or in standby mode. The tag address is not initialized by a power-on reset, manual reset, or in standby mode.

(2) Data Array

Holds a 16-byte instruction or data. Entries are registered in the cache in line units (16 bytes).

The data array is not initialized by a power-on reset, manual reset, or in standby mode.
(3) LRU

With the 4-way set associative system, up to four instructions or data with the same entry address can be registered in the cache. When an entry is registered, LRU shows which of the four ways it is recorded in. There are six LRU bits, controlled by hardware. A least-recently-used (LRU) algorithm is used to select the way that has been least recently accessed.

Six LRU bits indicate the way to be replaced in case of a cache miss. The relationship between LRU and way replacement is shown in table 9.1 when the cache lock function (only for operand cache) is not used (concerning the case where the cache lock function is used, see section 9.2.2, Cache Control Register 2 (CCR2)). If a bit pattern other than those listed in table 9.1 is set in the LRU bits by software, the cache will not function correctly. When modifying the LRU bits by software, set one of the patterns listed in table 9.1.

The LRU bits are initialized to B'000000 by a power-on reset but not initialized by a manual reset or in standby mode.

LRU (Bits 5 to 0)	Way to Be Replaced
000000, 000100, 010100, 100000, 110000, 110100	
000001, 000011, 001011, 100001, 101001, 101011	
000110, 000111, 001111, 010110, 011110, 011111	
111000, 111001, 111011, 111100, 111110, 111111	

Table 9.1 LRU Bits and Way Replacement (Cache Lock Function Not Used)

9.2 Register Descriptions

The cache has the following registers.

Table 9.2 Register Configuration

9.2.1 Cache Control Register 1 (CCR1)

The instruction cache is enabled or disabled using the ICE bit. The ICF bit controls disabling of all instruction cache entries. The operand cache is enabled or disabled using the OCE bit. The OCF bit controls disabling of all operand cache entries. The WT bit selects either write-through mode or write-back mode for operand cache.

Programs that change the contents of CCR1 should be placed in a cache-disabled space, and a cache-enabled space should be accessed after reading the contents of CCR1.

9.2.2 Cache Control Register 2 (CCR2)

CCR2 is used to enable or disable the cache locking function for operand cache and is valid only in cache locking mode. In cache locking mode, the lock enable bit (LE bit) in CCR2 is set to 1. In non-cache-locking mode, the cache locking function is invalid.

When a cache miss occurs in cache locking mode by executing the prefetch instruction (PREF @Rn), the line of data pointed to by Rn is loaded into the cache according to bits 9 and 8 (the W3LOAD and W3LOCK bits) and bits 1 and 0 (the W2LOAD and W2LOCK bits) in CCR2. The relationship between the setting of each bit and a way, to be replaced when the prefetch instruction is executed, are listed in table 9.3. On the other hand, when the prefetch instruction is executed and a cache hit occurs, new data is not fetched and the entry which is already enabled is held. For example, when the prefetch instruction is executed with W3LOAD = 1 and W3LOCK = 1 specified in cache locking mode while one-line data already exists in way 0 which is specified by Rn, a cache hit occurs and data is not fetched to way 3.

In the cache access other than the prefetch instruction in cache locking mode, ways to be replaced by bits W3LOCK and W2LOCK are restricted. The relationship between the setting of each bit in CCR2 and ways to be replaced are listed in table 9.4.

Programs that change the contents of CCR2 should be placed in a cache-disabled space, and a cache-enabled space should be accessed after reading the contents of CCR2.

Note: $*$ The W3LOAD and W2LOAD bits should not be set to 1 at the same time.

SH7205 Group Section 9 Cache S

Note: * The W3LOAD and W2LOAD bits should not be set to 1 at the same time.

LE	W3LOAD*	W3LOCK	W2LOAD*	W2LOCK	Way to Be Replaced
Ω	\times	\times	\times	\times	Decided by LRU (table 9.1)
	\times		\times		Decided by LRU (table 9.1)
	\times				Decided by LRU (table 9.5)
			\times		Decided by LRU (table 9.6)
	O				Decided by LRU (table 9.7)
		\times			Way 2
				\times	Way 3

Table 9.3 Way to Be Replaced When a Cache Miss Occurs in PREF Instruction

[Legend] x: Don't care

Note: * The W3LOAD and W2LOAD bits should not be set to 1 at the same time.

Table 9.4 Way to Be Replaced When a Cache Miss Occurs in Other Than PREF Instruction

[Legend] x: Don't care

Note: * The W3LOAD and W2LOAD bits should not be set to 1 at the same time.

Table 9.5 LRU and Way Replacement (When W2LOCK = 1 and W3LOCK = 0)

Table 9.6 LRU and Way Replacement (When W2LOCK = 0 and W3LOCK = 1)

Table 9.7 LRU and Way Replacement (When W2LOCK = 1 and W3LOCK = 1)

9.3 Operation

Operations for the operand cache are described here. Operations for the instruction cache are similar to those for the operand cache except for the address array not having the U bit, and there being no prefetch operation or write operation, or a write-back buffer.

9.3.1 Searching Cache

If the operand cache is enabled (OCE bit in CCR1 is 1), whenever data in a cache-enabled space is accessed, the cache will be searched to see if the desired data is in the cache. Figure 9.2 illustrates the method by which the cache is searched.

Entries are selected using bits 10 to 4 of the address used to access memory and the tag address of that entry is read. At this time, the upper three bits of the tag address are always cleared to 0. Bits 31 to 11 of the address used to access memory are compared with the read tag address. The address comparison uses all four ways. When the comparison shows a match and the selected entry is valid $(V = 1)$, a cache hit occurs. When the comparison does not show a match or the selected entry is not valid $(V = 0)$, a cache miss occurs. Figure 9.2 shows a hit on way 1.

Figure 9.2 Cache Search Scheme

9.3.2 Read Access

(1) Read Hit

In a read access, data is transferred from the cache to the CPU. LRU is updated so that the hit way is the latest.

(2) Read Miss

An external bus cycle starts and the entry is updated. The way replaced follows table 9.4. Entries are updated in 16-byte units. When the desired data that caused the miss is loaded from external memory to the cache, the data is transferred to the CPU in parallel with being loaded to the cache. When it is loaded in the cache, the V bit is set to 1, and LRU is updated so that the replaced way becomes the latest. In operand cache, the U bit is additionally cleared to 0. When the U bit of the entry to be replaced by updating the entry in write-back mode is 1, the cache update cycle starts after the entry is transferred to the write-back buffer. After the cache completes its update cycle, the write-back buffer writes the entry back to the memory. The write-back unit is 16 bytes. The update of cache and write-back to memory are performed in wrap around method. For example, the lower four bits of the address at which a read miss occurs indicate H'4, the update of cache and write-back to memory are performed in the order of H'4, H'8, H'C, H'0, which are the lower four bits of the address.

9.3.3 Prefetch Operation (Only for Operand Cache)

(1) Prefetch Hit

LRU is updated so that the hit way becomes the latest. The contents in other caches are not modified. No data is transferred to the CPU.

(2) Prefetch Miss

No data is transferred to the CPU. The way to be replaced follows table 9.3, Other operations are the same in case of read miss.

9.3.4 Write Operation (Only for Operand Cache)

(1) Write Hit

In a write access in write-back mode, the data is written to the cache and no external memory write cycle is issued. The U bit of the entry written is set to 1 and LRU is updated so that the hit way becomes the latest.

In write-through mode, the data is written to the cache and an external memory write cycle is issued. The U bit of the written entry is not updated and LRU is updated so that the hit way becomes the latest.

(2) Write Miss

In write-back mode, an external bus cycle starts when a write miss occurs, and the entry is updated. The way to be replaced follows table 9.4. When the U bit of the entry to be replaced is 1, the cache update cycle starts after the entry is transferred to the write-back buffer. Data is written to the cache, the U bit is set to 1, and the V bit is set to 1. LRU is updated so that the replaced way becomes the latest. After the cache completes its update cycle, the write-back buffer writes the entry back to the memory. The write-back unit is 16 bytes. The update of cache and write-back to memory are performed in wrap around method. For example, the lower four bits of the address at which a write miss occurs indicate H^{'4}, the update of cache and write-back to memory are performed in the order of H'4, H'8, H'C, H'0, which are the lower four bits of the address.

In write-through mode, no write to cache occurs in a write miss; the write is only to the external memory.

9.3.5 Write-Back Buffer (Only for Operand Cache)

When the U bit of the entry to be replaced in write-back mode is 1, it must be written back to the external memory. To increase performance, the entry to be replaced is first transferred to the write-back buffer and fetching of new entries to the cache takes priority over writing back to the external memory. After the cache completes to fetch the new entry, the write-back buffer writes the entry back to external memory. During the write-back cycles, the cache can be accessed.

The write-back buffer can hold one line of cache data (16 bytes) and its physical address. Figure 9.3 shows the configuration of the write-back buffer.

Figure 9.3 Write-Back Buffer Configuration

Table 9.8 summarizes the above operations in sections 9.3.2, Read Access to 9.3.5, Write-Back Buffer (Only for Operand Cache).

Table 9.8 Cache Operations

[Legend]

x: Don't care.

Notes: Cache renewal cycle: 16-byte read access

Write-back cycle in write-back buffer: 16-byte write access

∗ Neither LRU renewed. LRU is renewed in all other cases.

9.3.6 Coherency of Cache and External Memory

Use software to ensure coherency between the cache and the external memory.

When memory shared by this LSI and another device is allocated in the cache-enabled space, operate the memory-allocated cache to invalidate and write back as required. Do the same operation for memory shared by the CPU on this LSI and the DMAC.

9.4 Memory-Allocated Cache

To allow software management of the cache, cache contents can be read and written by means of MOV instructions. The instruction cache address array is allocated onto addresses H'F0000000 to H'F07FFFFF, and the data array onto addresses H'F1000000 to H'F17FFFFF. The operand cache address array is allocated onto addresses H'F0800000 to H'F0FFFFFF, and the data array onto addresses H'F1800000 to H'F1FFFFFF. Only longword can be used as the access size for the address array and data array, and instruction fetches cannot be performed.

9.4.1 Address Array

To access an address array, the 32-bit address field (for read/write accesses) and 32-bit data field (for write accesses) must be specified.

In the address field, specify the entry address selecting the entry, The W bit for selecting the way, and the A bit for specifying the existence of associative operation. In the W bit, B'00 is way 0, B'01 is way 1, B'10 is way 2, and B'11 is way 3. Since the access size of the address array is fixed at longword, specify B'00 for bits 1 and 0 of the address.

The tag address, LRU bits, U bit (only for operand cache), and V bit are specified as data. Always specify 0 for the upper three bits (bits 31 to 29) of the tag address.

Refer to figure 9.4 regarding the address and data format.

The following three operations are possible for the address array.

(1) Address Array Read

The tag address, LRU bits, U bit (only for operand cache), and V bit are read from the entry address specified by the address and the entry corresponding to the way. For the read operation, associative operation is not performed regardless of whether the associative bit (A bit) specified by the address is 1 or 0.

(2) Address-Array Write (Non-Associative Operation)

When writing 0 to the associative bit (A bit) in the address field, the tag address, LRU bits, U bit (only for operand cache), and V bit, specified by the data field, are written to the entry address specified by the address and the entry corresponding to the way. When writing to a cache line for which the U bit = 1 and the V bit = 1 in the operand cache address array, write the contents of the cache line back to memory, then write the tag address, LRU bits, U bit, and V bit specified by the data field. When 0 is written to the V bit, 0 must also be written to the U bit of that entry. The write-back to memory is performed in the order of H'0, H'4, H'8, H'C, which are the lower four bits of the address.

(3) Address-Array Write (Associative Operation)

When writing with the associative bit (A bit) in the address field set to 1, the addresses in the four ways for the entry specified by the address field are compared with the tag address that is specified by the data field. Write the U bit (only for operand cache) and the V bit specified by the data field to the entry of the way that has a hit. However, the tag address and LRU bits remain unchanged. When there is no way that has a hit, nothing is written and there is no operation. This function is used to invalidate a specific entry in the cache.

When the U bit of the entry that has had a hit is 1 in the operand cache, writing back should be performed. However, when 0 is written to the V bit, 0 must also be written to the U bit of that entry. The write-back to memory is performed in the order of H'0, H'4, H'8, H'C, which are the lower four bits of the address.

9.4.2 Data Array

To access a data array, the 32-bit address field (for read/write accesses) and 32-bit data field (for write accesses) must be specified. The address field specifies information for selecting the entry to be accessed; the data field specifies the longword data to be written to the data array.

Specify the entry address for selecting the entry, the L bit indicating the longword position within the (16-byte) line, and the W bit for selecting the way. In the L bit, B'00 is longword 0, B'01 is longword 1, B'10 is longword 2, and B'11 is longword 3. In the W bit, B'00 is way 0, B'01 is way 1, B'10 is way 2, and B'11 is way 3. Since the access size of the data array is fixed at longword, specify B'00 for bits 1 and 0 of the address.

Refer to figure 9.4 regarding the address and data format.

The following two operations are possible for the data array. Information in the address array is not modified by this operation.

(1) Data Array Read

The data specified by the L bit in the address is read from the entry address specified by the address and the entry corresponding to the way.

(2) Data Array Write

The longword data specified by the data is written to the position specified by the L bit in the address from the entry address specified by the address and the entry corresponding to the way.

9.4.3 Usage Examples

(1) Invalidating Specific Entries

Specific cache entries can be invalidated by writing 0 to the entry's V bit in the memory allocating cache access. When the A bit is 1, the tag address specified by the write data is compared to the tag address within the cache selected by the entry address, and data is written to the bits V and U specified by the write data when a match is found. If no match is found, there is no operation. When the V bit of an entry in the address array is set to 0, the entry is written back if the entry's U bit is 1.

An example when the write data is specified in R0 and an address is specified in R1 is shown below.

; R0=H'0110 0010; tag address(28 to 11)=B'0 0001 0001 0000 0000 0, U=0, V=0 ; R1=H'F080 0088; operand cache address array access, entry=B'000 1000, A=1 ; MOV.L R0,@R1

(2) Reading the Data of a Specific Entry

The data section of a specific cache entry can be read by the memory allocating cache access. The longword indicated in the data field of the data array in figure 9.4 is read into the register.

An example when an address is specified in R0 and data is read in R1 is shown below.

```
; R0=H'F100 004C; instruction cache data array access, entry=B'000 0100, Way=0, 
longword address=3 
; 
 MOV.L @R0,R1
```


9.4.4 Notes

- 1. Programs that access memory-mapped cache of the operand cache should be placed in a cachedisabled space. Programs that access memory-mapped cache of the instruction cache should be placed in a cache-disabled space, and in each of the beginning and the end of that, two or more read accesses to on-chip peripheral modules or external address space (cache-disabled address) should be executed.
- 2. Rewriting the address array contents so that two or more ways are hit simultaneously is prohibited. Operation is not guaranteed if the address array contents are changed so that two or more ways are hit simultaneously.
- 3. Only the CPU can access memory-allocated cache; the DMAC cannot access it.

Section 10 Bus State Controller (BSC)

The bus state controller (BSC) outputs control signals for external devices and various types of memory that is connected to the external address space. This enables the LSI to connect directly with SRAM, SDRAM, and other memory storage devices, and external devices.

10.1 Features

- 1. External address space
	- Maximum of 64 Mbytes for the SDRAM and each for areas CS0 to CS5
	- ⎯ Ability to select the data bus width (8, 16, or 32 bits) independently for each address space
- 2. Normal space interface
	- ⎯ Supports an interface for direct connection to SRAM
	- ⎯ Cycle wait function: Maximum of 31 wait states (maximum of seven wait states for page access cycles)
	- Wait control
		- Ability to select the assert/negate timing for chip select signals
		- Ability to select the assert/negate timing for the read strobe and write strobe signals
		- Ability to select the data output start/end timing
		- Ability to select the delay for chip select signals
	- ⎯ Write access modes: One-write strobe and byte-write strobe modes
	- Page access mode: Support for page read and page write (64-bit, 128-bit, and 256-bit page units)
- 3. SDRAM interface
	- Ability to set SDRAM in up to two areas
	- Refresh functions

Auto-refresh (on-chip programmable refresh counter)

- Self-refresh
- Ability to select the access timing (support for row-column latency, column latency, and row-active interval settings)
- ⎯ Initialization sequencer function, power-down function, deep-power-down function, and mode register setting function implemented on-chip

Figure 10.1 shows a block diagram of the BSC. The BSC consists of an area controller (CSC), an access controller, and an SDRAM controller (SDRAMC). The CSC controls accessing normal space in the external address space (see table 10.2). The SDRAMC controls accesses to the SDRAM space. The access controller controls operations common to both the above-mentioned normal space and SDRAM space.

Figure 10.1 Block Diagram of BSC

10.2 Input/Output Pins

Table 10.1 shows the pin configuration of the BSC.

Table 10.1 Pin Configuration

10.3 Area Overview

10.3.1 Address Map

In the architecture, this LSI has a 32-bit address space, which is divided into cache-enabled, cache-disabled, and on-chip spaces (on-chip high-speed RAM, on-chip RAM for data retention, on-chip peripheral modules, and reserved areas) according to the upper bits of the address.

External address spaces CS5 to CS0, SDRAM0, and SDRAM1 are cache-enabled when internal address $A29 = 0$ and cache-disabled when $A29 = 1$.

The kind of memory to be connected and the data bus width are specified independently for each partial space. The address map for the external address space is shown below.

Notes: 1. The term BIU stands for Bus Interface Unit. BIUs are internal modules through which the CPU and DMAC accesses each address space. Described below are on-chip BIUs, and address spaces and internal buses connected to these BIUs.

BIU E: External address spaces (normal and SDRAM spaces)

BIU_PB0: peripheral bus 0 (internal to the LSI)

BIU_PB1: peripheral bus 1 (internal to the LSI)

BIU_PB2: peripheral bus 2 (internal to the LSI)

BIU_PB3: peripheral bus 3 (internal to the LSI)

 Pipelined DMA transfer is not available for transfer from a BIU to the same BIU. For details, see section 11, Direct Memory Access Controller (DMAC).

 2. Cache address array space and some on-chip peripheral modules are not allocated to any BIU. These devices are accessed directly from individual CPUs without using the internal bus. The DMAC cannot access any of these devices.

10.3.2 Data Bus Width and Related Pin Settings for Each Area

The data bus width of area 0 can be set to either 16 or 32 bits by means of an external pin, but it cannot be changed by a program after startup. The initial data bus width of SDRAM0, SDRAM1, and areas 1 to 5 can be changed by a program by means of register settings.

Immediately after a power-on reset, certain address and data bus settings, and functions of CS0 and RD, are automatically selected as initial functions required to read the area 0 ROM. The other pins are set as general ports as the initial function setting, and they cannot be used until their functions are set by a program. Do not perform other than read access to area 0 until the pin function settings have been completed.

Table10.3 lists the external pin settings and the initial state of each area.

Note that sample access waveforms are shown in this section for pins such as $RD_$ WR and WEn, and these functions are available after they are selected through the pin function controller. For example, after startup with the 32-bit bus width, changing the bus width of the areas other than area 0 to 16 bits requires setting the function of the A1 pin, and changing the bus width to 8 bits requires setting the functions of the A1 and A0 pins.

For details on pin function settings, see section 27, Pin Function Controller (PFC).

function are also both assigned to a single pin (PB6).

10.4 Register Descriptions

The BSC has the following registers. All registers are initialized by a power-on reset or in deep standby mode.

Do not access spaces other than area 0 until settings are completed for the connected memory interface.

Table 10.4 Register Configuration

Note: * Depends on the MD0 pin setting made at start-up.

10.4.1 CSn Control Register (CSnCNT) (n = 0 to 5)

CSnCNT selects the width of the external bus and controls the operation of the CSC interface.

Notes: 1. The initial value of the BSIZE bits in CS0 differs depending on the MD0 pin setting.

2. The initial value of the EXENB bit in CS0 is 1.

To disable $(EXENB = 0)$ the operation for each channel, forcibly write out data tentatively stored in internal write buffer. The procedure is as follows:

- 1. Execute read access to the channel whose operation is to be disabled.
- 2. Then, write 0 to the EXENB bit (operation disabled).

10.4.2 CSn Recovery Cycle Setting Register (CSnREC) (n = 0 to 5)

CSnREC specifies the number of data recovery cycles to be inserted after read or write accesses.

- Consequently, there is no data recovery cycle setting for SDRAM. (The value is fixed at 0 cycles.)
	- 2. Writing to the CSn recovery cycle setting register (CSnREC) must be done while the CSC for the corresponding channel is disabled ($EXENB = 0$). Only channel 0 (CS0) is allowed for writing to the register without disabling the CSC ($EXENB = 1$). To write to CS0REC with CSC enabled, satisfy all of the following conditions:
		- Stop the DMAC.
		- Keep the CPU other than the one that is going to rewrite the register from accessing CS0 (including access for instruction fetch). For example, if CPU0 is going to rewrite the register, make CPU1 stay looping by a program copied to on-chip memory, or put CPU1 in a sleep state.
		- Do not perform data write access to CS0 after a reset is released but before the register is rewritten.

10.4.3 SDRAMCm Control Register (SDCmCNT) (m = 0, 1)

To disable the operation (EXENB = 0) for each channel, forcibly write out data tentatively stored in internal write buffer. The procedure is as follows:

- 1. Execute read access to the channel whose operation is to be disabled.
- 2. Then, write 0 to the EXENB bit (operation disabled).

10.4.4 CSn Mode Register (CSMODn) (n = 0 to 5)

CSMODn selects the mode for page read access and the bit boundary for page access, enables page read/write access and external wait, and selects the mode for write access.

Writing to the CSn mode register (CSMODn) must be done while the CSC for the corresponding channel is disabled ($EXENB = 0$). Only channel 0 (CS0) is allowed for writing to the register without disabling the CSC (EXENB = 1). To write to CSMOD0 with CSC enabled, satisfy all of the following conditions:

- 1. Stop the DMAC.
- 2. Keep the CPU other than the one that is going to rewrite the register from accessing CS0 (including access for instruction fetch). For example, if CPU0 is going to rewrite the register, make CPU1 stay looping by a program copied to on-chip memory, or put CPU1 in a sleep state.
- 3. Do not perform data write access to CS0 after a reset is released but before the register is updated.

10.4.5 CSn Wait Control Register 1 (CS1WCNTn) (n = 0 to 5)

CS1WCNTn specifies the number of wait states to be inserted into the read/write cycle or page read/page write cycle.

according to the configuration of your system.

- 2. Writing to the CSn wait control register 1 (CS1WCNTn) must be done while the CSC for the corresponding channel is disabled (EXENB = 0). Only channel 0 (CS0) is allowed for writing to the register without disabling the CSC (EXENB = 1). To write to CS1WCNT0 with CSC enabled, satisfy all of the following conditions:
	- Stop the DMAC.
	- Keep the CPU other than the one that is going to rewrite the register from accessing CS0 (including access for instruction fetch). For example, if CPU0 is going to rewrite the register, make CPU1 stay looping by a program copied to on-chip memory, or put CPU1 in a sleep state.
	- Do not perform data write access to CS0 after a reset is released but before the register is rewritten.

10.4.6 CSn Wait Control Register 2 (CS2WCNTn) (n = 0 to 5)

CS2WCNTn specifies the number of wait states and the number of delay cycles.

 CSON ≤ min. (CSPRWAIT, CSPWWAIT) WRON ≤ CSPWWAIT, RDON ≤ CSPRWAIT WDON ≤ CSPWWAIT

10.4.7 SDRAM Refresh Control Register 0 (SDRFCNT0)

SDRFCNT0 controls self-refresh operation.

10.4.8 SDRAM Refresh Control Register 1 (SDRFCNT1)

SDRFCNT1 controls auto-refresh operation.

Note: Auto-refresh requests are not accepted while multiple read or write accesses are in progress, or during a transfer using DMAC, so the auto-refresh interval may become extended in some cases. Set the DRFC bits to an auto-refresh request interval value that satisfies the auto-refresh interval specification of the SDRAM being used. Furthermore, make sure to set the auto-refresh request interval to a duration longer than the auto-refresh cycle.

(a) Auto-Refresh Request Interval and DRFC Set Value

SDRAMC includes a 12-bit refresh counter that generates auto-refresh requests at fixed intervals. The following equation is used to calculate the set value for the DRFC bits from the auto-refresh request interval.

DRFC = (Auto-refresh request interval / Bus clock cycle) – 1

Auto-refresh requests are not accepted while SDRAM is being accessed; they must wait until the access completes. However, the counter value is updated regardless of whether or not the request was accepted. Note that if two or more auto-refresh requests are generated while SDRAM is being accessed, the second and subsequent requests are ignored.

10.4.9 SDRAM Initialization Register 0 (SDIR0)

SDIR0 specifies the SDRAM initialization sequence timing.

Note: Make settings that satisfy the specifications of the connected SDRAM before starting the initialization sequence.

10.4.10 SDRAM Initialization Register 1 (SDIR1)

SDIR1 controls activation of the SDRAM initialization sequence.

10.4.11 SDRAM Power-Down Control Register (SDPWDCNT)

SDPWDCNT controls transition to and recovery from power-down mode.

10.4.12 SDRAM Deep-Power-Down Control Register (SDDPWDCNT)

SDDPWDCNT controls transition to and recovery from deep-power-down mode.

10.4.13 SDRAMm Address Register (SDmADR) (m = 0, 1)

SDmADR specifies the data bus width and the channel size for SDRAM.

10.4.14 SDRAMm Timing Register (SDmTR) (m = 0, 1)

SDmTR specifies the timing for read and write accesses to SDRAM.

[Legend] x: Don't care

10.4.15 SDRAMm Mode Register (SDmMOD) (m = 0, 1)

SDmMOD specifies the values to be written to the SDRAM mode register or extended mode register. Writing to this register causes a mode register set command or extended mode register set command to be issued automatically to SDRAM.

Notes: The following points should be kept in mind regarding SDRAMm mode register settings.

- 1. Make sure to set a burst length of 1 for SDRAM. Operation cannot be guaranteed with settings other than a burst length of 1.
- 2. The SDRAM column latency must match the setting of the SDRAM controller column latency setting bits (DCL) in SDRAMC. Operation cannot be guaranteed if the latency settings do not agree.
- 3. Make sure the status bits (DSRFST, DPWDST, DDPDST, and DMRSST) in the SDRAM status register (SDSTR) are all cleared to 0.

10.4.16 SDRAM Status Register (SDSTR)

SDSTR consists of the status flags that indicate the status of operation during self-refresh, initialization sequences, power-down mode, deep-power-down mode, and mode register setting.

"Transition or recovery in progress" refers to the interval from the point at which the bits listed in table 10.5 are written to until the corresponding commands are issued.

Note: Execution of a self-refresh, a transition to or recovery from power-down or deep-powerdown mode, an initialization sequence, or mode register setting may only be performed when all status bits are cleared to 0. Do not rewrite the registers (bits) listed in table 10.5 when any of the status bits (DSRFST,

DINIST, DPWDST, DDPDST, DMRSST) is set to 1.

10.4.17 SDRAM Clock Stop Control Signal Setting Register (SDCKSCNT)

SDCKSCNT enables or disables the clock stop control signal (signal in the chip) and specifies the number of assert cycles.

10.5 Operation

10.5.1 Accessing CS Space

(1) Normal Access

Normal read/write operation is used for all bus accesses when page read/write access is disabled $(PRENB = 0, PWENB = 0)$. Even when page read/write access is enabled $(PRENB = 1, PWENB)$ = 1), normal read/write operation is employed in cases where page access cannot be used. Figure 10.2 shows the basic operation of the external bus control signals for read and write operations in byte-write strobe mode. Figure 10.3 shows the basic operation of these signals for read and write operations in one-write strobe mode. In these figures, \overline{DACTn} is a DMA active output signal. For details, see section 11, Direct Memory Access Controller (DMAC).

Figure 10.2 Basic Bus Timing (Read Operation in Byte-Write Strobe Mode) (a)

Figure 10.2 Basic Bus Timing (Write Operation in Byte-Write Strobe Mode) (b)

Figure 10.3 Basic Bus Timing (Read Operation in One-Write Strobe Mode) (a)

Figure 10.3 Basic Bus Timing (Write Operation in One-Write Strobe Mode) (b)

1. Ts (Internal Bus Access Start)

This is a bus access request cycle initiated by the internal bus master to the external bus as the target. \overline{CSn} is always high during this cycle. In the next cycle, A25 to A0, BCn, and the write data change.

2. Tw1 to Twn (Read Cycle Wait, Write Cycle Wait)

These are the cycles between internal bus access start and the wait end cycle. A duration of from 0 to 31 clocks may be selected. During this interval the \overline{CSn} , \overline{RD} , \overline{WEn} , and \overline{WE} signals are asserted (low level) in accordance with the wait settings. The assert timing can be controlled using the CS assert wait, RD assert wait, WR assert wait, and write data output wait bits in CSn control registers 1 and 2. The number of wait cycles can be set to from 0 to 7 clocks, as counted from the cycle following internal bus access start (Ts). The number of clocks selected must be no greater than the number of read/write cycle wait cycles. The RD_WR signal operates with the same timing as for the $\overline{\text{CSn}}$ signal.

3. Tend (Wait End Cycle)

This is the final cycle in a series of read cycle wait or write cycle wait cycles. The RD, WEn, or WE signal is negated (high level) in the next cycle.

4. Tn1 to Tnm (CS Delay Cycle)

These are the cycles between the wait end cycle and when CSn is negated (high level). The negation timing can be controlled using write data output delay cycles. The number of cycles is counted beginning from the wait end cycle. In write access or if the number of CS delay cycles during a read is other than 0 or 1, the succeeding bus access can start from the cycle following the CS delay cycle end. If the number of CS delay cycles is 0 or 1 in read access, the succeeding bus access can start after the end of the read data sample cycle (see below).

5. Trd (Read Data Sample Cycle)

This is the sample cycle for read data.

(2) Page Access

Page read and write operation is employed for bus accesses for which page access can be used if page write access enable ($PWENB = 1$) and page read access enable ($PRENB = 1$) have been selected. Page access is used in the following cases.

- 1. CPU cache replace (cache filling and write-back). Each time a transfer occurs, the address is incremented by the number of transferred bytes.
- 2. When longword (32-bit) access to an 8-bit or 16-bit external data bus has been performed.
- 3. When word (16-bit) access to an 8-bit external data bus has been performed.

Figure 10.4 shows the basic operation of the external bus control signals for page read operation, and figure 10.5 shows the basic operation of these signals for page write operation. If the single page access bit boundary setting made by the PBCNT bits in the mode register is smaller than the cache line size, a single cache replacement will trigger multiple page accesses. When the address exceeds the page boundary, page access stops once (the CSn signal is negated) and starts again. If the PBCNT bit setting is not smaller than the cache line size, a cache replacement is processed in a single page access.

Figure 10.4 Basic Bus Timing (Page Read Operation in Byte-Write Strobe Mode)

Figure 10.5 Basic Bus Timing (Page Write Operation in Byte-Write Strobe Mode)

1. Ts (Internal Bus Access Start)

This is a bus access request cycle initiated by the internal bus master to the external bus as the target. \overline{CSn} is always high during this cycle. In the next cycle, A25 to A0, \overline{BCn} , and the write data change.

2. Tw1 to Twn (Read Cycle Wait, Write Cycle Wait)

For the first bus access in a page access, the control of the wait operation from internal bus access start to the wait end cycle is the same as in normal access.

3. Tend (First Wait End Cycle)

This is the final cycle in the first series of read cycle wait or write cycle wait cycles. In write access, the second and subsequent page accesses start from the next cycle, unless a write data output delay cycle has been specified (with a value other than 0). The RD, WEn, or WE signal is negated (high level) in the next cycle if the RD assert wait or WR assert wait setting is other than 0. If the RD assert wait or WR assert wait setting is 0, the RD, WEn, or WE signal continues to be asserted (low level). The CSn signal is not negated and continues to be asserted (low level). The RD_WR signal operates with the same timing as for the CSn signal. In page read access, the succeeding bus access starts without waiting for the read data sample cycle (Trd).

4. Tdw1 to Tdwn (Write Data Output Delay Cycle)

In write access, write data output delay cycles are inserted between the wait end cycle and the following page access if the write data output delay wait setting is other than 0. Assertion of the address and output data is extended for the duration of this interval. Also, the WEn and WE signals are negated (high level).

5. Tpw1 to Tpwn (Page Read Cycle Wait, Page Write Cycle Wait)

For the second and subsequent bus cycles in a page access, the page read cycle wait and page write cycle wait settings are used in place of the read cycle wait and write cycle wait settings. The WR assert wait setting works in the same way as for the first bus cycle. The RD assert wait setting operates differently depending on the page read access mode (PRMOD) setting value.

 $PRMOD = 0$: RD assert wait setting operates in the same way as for the first bus cycle. PRMOD = 1: RD assert wait setting is invalid.

Operation is the same as for RD assert wait setting of 0.

- 6. Tend/Tdw1 to Tdwn (Wait End Cycle/Write Data Output Delay Cycle) The operation is the same as for the first access (3 and 4 above).
- 7. Tn1 to Tnm (CS Delay Cycle)

These are the cycles between the final wait end cycle and when CSn is negated (high level). The number of CS delay cycles is counted beginning from the wait end cycle.

8. Trd (Final Read Data Sample Cycle) This is the final sample cycle for read data.

(3) External Wait Function

The external wait signal (\overline{WAIT}) can be used to extend the wait cycle duration beyond the value specified by the cycle wait (CSRWAIT, CSWWAIT) or page access cycle wait (CSPRWAIT, CSPWWAIT) settings in the CSn wait control register (CSWCNTn). If external wait enable $(EWENB = 1)$ has been selected, wait cycles are inserted for as long as the WAIT signal remains low level. The \overline{WAIT} signal is disabled if external wait disable (EWENB = 0) has been selected.

Note that the wait cycles specified by the settings of the CSn wait control register (CSWCNTn) are inserted regardless of the state of the \overline{WAIT} signal.

(a) Normal Read/Write Operation

The WAIT signal is sampled all the time and its result is reflected two cycles later. Thus, when the WAIT signal is low two cycles before the end of the wait cycles, external cycles are inserted. After the \overline{WAIT} signal has gone high, the wait cycles end two cycles later.

(b) Page Access Operation

The initial data read/write operation is the same as a normal read/write operation. That is, when the WAIT signal is low two cycles before the end of the wait cycles (Tend), external wait cycles are inserted. After the WAIT signal has gone high, the wait cycles end (Tend) two cycles later.

In the second and subsequent read accesses, the page wait cycle is extended if the WAIT signal is low two cycles before the end of the page access wait cycle (Tend), and the page wait cycles end (Tend) two cycles after the \overline{WAIT} signal has gone high.

Figures 10.6 and 10.7 show examples of external wait timing for page read access by longword (32-bit) access to a 16-bit channel. Figure 10.6 is an example in which one or more cycles of cycle wait state or page cycle wait state has been set. Figure 10.7 is an example in which no cycle wait or page cycle wait state has been set. Note that the value of the WAIT signal before the beginning of the bus cycle is reflected if there are only a few cycles of cycle wait states.

Figure 10.6 External Wait Timing Example (Page Read Access to 16-Bit Channel)

Figure 10.7 External Wait Timing Example (Page Read Access to 16-Bit Channel)

(4) Access Type and Data Alignment

(a) 32-Bit Bus Channel

If a 32-bit bus is selected by the external bus width select bits in the CSn control register, A25 to A2 are enabled as address signals for longword units and A1 and A0 are disabled (fixed at low level). Table 10.6 lists the data alignment corresponding to byte addresses for different data sizes.

If byte strobe mode (WRMOD = 0) is selected, the $\overline{WE3}$ to \overline{WEO} signals indicate the bits to be accessed. For read access, however, all bits are access targets regardless of the state of the $\overline{WE3}$ to WE0 signals.

If one-write strobe mode (WRMOD = 1) is selected, the $\overline{BC3}$ to $\overline{BC0}$ signals indicate access targets for both read and write accesses. For write access, the write strobe signal WE is also asserted.

Table 10.6 Data Alignment (32-Bit Bus Channel)

Note: The valid bits on the data bus for each data size are indicated by circles (O). Crosses (x) indicate bus data bits that are undefined.

(b) 16-Bit Bus Channel

If a 16-bit bus is selected by the external bus width select bits in the CSn control register, A25 to A1 are enabled as address signals for word units, and A0 is disabled (fixed at low level). Table 10.7 lists the data alignment corresponding to byte addresses for different data sizes.

If byte strobe mode (WRMOD = 0) is selected, the $\overline{WE1}$ and $\overline{WE0}$ signals indicate the bits to be accessed. For read access, however, all bits are access targets regardless of the state of the $\overline{WE1}$ and WE0 signals.

If one-write strobe mode (WRMOD = 1) is selected, the $\overline{BC1}$ to $\overline{BC0}$ signals indicate access targets for both read and write. For write access, the write strobe signal \overline{WE} is also asserted.

Table 10.7 Data Alignment (16-Bit Bus Channel)

Note: The valid bits on the data bus for each data size are indicated by circles (O). Crosses (×) indicate bus data bits that are undefined. Asterisks (*) indicate write/byte control bits that are disabled (fixed at high level).

(c) 8-Bit Bus Channel

If an 8-bit bus is selected by the external bus width select bits in the CSn control register, A25 to A0 are enabled as address signals for byte units. Table 10.8 lists the data alignment corresponding to byte addresses for different data sizes.

If byte strobe mode (WRMOD = 0) is selected, the \overline{WEO} signal is asserted only for write access; it is not asserted for read access.

If one-write strobe mode (WRMOD = 1) is selected, the \overline{BCO} signal is asserted for both read and write accesses. For write access, the write strobe signal \overline{WE} is also asserted.

Table 10.8 Data Alignment (8-Bit Bus Channel)

Note: The valid bits on the data bus for each data size are indicated by circles (O). Crosses (x) indicate bus data bits that are undefined. Asterisks (*) indicate write/byte control bits that are disabled (fixed at high level).

10.5.2 Accessing SDRAM

A description is provided here of the SDRAM controller (SDRAMC) operation enable and SDRAM bus width settings as well as operations involving SDRAM (read, write, auto-refresh, self-refresh, initialization sequence, and mode register settings).

(1) SDRAM Access Enable/Disable and SDRAM Bus Width Settings

Enabling and disabling SDRAM access is performed by making settings in the individual SDRAMCm control registers to enable or prohibit SDRAMC operation. SDRAM bus width settings are also performed by means of the SDRAMCm control registers.

Even if the SDRAMC control register is set to disable SDRAMC operation, refresh operation will still take place if self-refresh or auto-refresh operation is set as enabled.

(2) SDRAM Commands

SDRAMC controls SDRAM by issuing commands each bus cycle. These commands are defined by combinations of RAS, CAS, WE, CKE, CS, etc.

Table 10.9 lists the commands issued by the SDRAMC.

Table 10.9 SDRAMC Commands

[Legend]

H: High Level. L: Low Level. V: Valid. X: Don't Care.

(3) SDRAMC Register Setting Conditions

Rewriting of SDRAMC registers should only be performed when all of the conditions listed in table 10.10 are satisfied.

Table 10.10 Register Rewrite Conditions

Notes: 1. After writing 0 to EXENB, make sure that the EXENB bit has been cleared to 0. 2. Before rewriting this bit, make sure that all status bits in the SDRAM status register

(SDSTR) have been cleared to 0.

(4) Self-Refresh

Settings in SDRAM refresh control register 0 (SDRFCNT0) control transition to and recovery from self-refresh mode. A transition to and recovery from self-refresh mode takes place simultaneously for all channels. After settings for self-refresh mode have been made, this LSI continues in the self-refresh state even when it is placed on software standby or deep standby. The self-refreshing state is also maintained after interrupt-initiated recovery from standby state. However, the setting for the HIZBCS bit in the HIZCR register must be 0, and the CKE and other pins must be driven even in standby mode. With regard to the HIZCE register, refer to section 30, Power-Down Modes.

An auto-refresh cycle operation takes place immediately before transition to self-refresh mode. The CKE signal is at the low level during periods in self-refresh mode. Immediately after recovery from self-refresh mode, an auto-refresh cycle is triggered.

Figure 10.8 shows the timing of transition to self-refresh mode, and figure 10.9 shows the timing of recovery from self-refresh mode.

Figure 10.9 Example of Timing of Recovery from Self-Refresh Mode (DREFW Bits = 0010)

(5) Auto-Refresh

An auto-refresh cycle starts when the auto-refresh operation enable bit (DRFEN) in SDRAM refresh control register 1 (SDRFCNT1) is set to 1. After that, refresh requests are issued at fixed intervals according to the refresh counter, activating auto-refresh cycles. However, the activation of auto-refresh cycles may sometimes be delayed because refresh requests are not accepted during read or write accesses.

A refresh request is issued immediately when the auto-refresh operation enable bit (DRFEN) in SDRAM refresh control register 1 (SDRFCNT1) is set to 1 if auto-refresh is enabled.

The refresh counter is halted in self-refresh or deep-power-down mode. After recovery from selfrefresh or deep-power-down mode, an auto-refresh cycle is activated, after which the counter value is reset and the counter begins operating again.

Make auto-refresh settings in SDRAM refresh control register 1 (SDRFCNT1). Note that refresh cycles affect SDRAM for all channels. Figure 10.10 shows an auto-refresh cycle timing example.

Figure 10.10 Auto-Refresh Cycle Timing Example (DREFW Bits = 0010)

(6) Initialization Sequencer

SDRAMC is provided with a sequencer for issuing the commands for SDRAM initialization. The initialization sequence should always be initiated a single time only following a reset (all channels) or following recovery from deep-power-down mode (individual channels). No normal operation is guaranteed if the initialization sequence is not performed or is performed more than once.

The SDRAM initialization sequencer issues the precharge-all-banks command and $n (n = 1 to 15)$ auto-refresh commands in the stated order. Make timing settings for the initialization sequencer, using SDRAM initialization register 0 (SDIR0). To initiate initialization sequences, use SDRAM initialization register 1 (SDIR1).

Note that an initialization sequence for SDRAM0 and SDRAM1 is initiated simultaneously, using the DINIRQ bit.

Figure 10.11 shows a timing example for the initialization sequence. Setting DARFC to 2 or greater causes multiple initialization auto-refresh cycles to be performed.

Figure 10.11 Initialization Sequence Timing Example (DPC Bits = 001, DARFI Bits = 0001, DARFC Bits = 001)

(7) Power-Down Mode

SDRAMC supports power-down mode for SDRAM. In power-down mode, it drives the CKE signal low. In power-down mode, auto-refresh operations are performed at the intervals specified by the auto-refresh request interval setting (DRFC) bits in SDRAM refresh control register 1 (SDRFCNT1). The CKE signal goes high only when an auto-refresh command is issued.

To perform transition to and recovery from power-down mode, use the SDRAM power-down control register (SDPWDCNT).

Setting the DPWD bit to 1 causes SDRAMC to enter power-down mode. Clearing the DPWD bit to 0 causes SDRAMC to exit power-down mode.

SDRAMC drives the CKE signal high after recovery from power-down mode.

Figure 10.12 SDRAMC Power-Down Mode

Figure 10.13 Auto-Refresh Operation in SDRAMC Power-Down Mode

(8) Deep-Power-Down Mode

SDRAMC supports deep-power-down mode for SDRAM. In deep-power-down mode, it issues a deep-power-down command to drive the CKE signal low.

To perform transition to and recovery from deep-power-down mode, use the SDRAM deeppower-down control register (SDDPDCNT).

Setting the DDPD bit to 1 causes SDRAM0 and SDRAM1 in SDRAM to enter deep-power-down mode. Clearing the DDPD bit to 0 causes SDRAMC to recover from deep-power-down mode.

After recovery from deep-power-down mode, SDRAMC issues a deep-power-down exit command to drive the CKE signal high.

After waiting for the duration designated for the SDRAM being used after recovery from deeppower-down mode, execute an initialization sequence.

Figure 10.14 SDRAMC Deep-Power-Down Mode

(9) Read/Write Access

The following two types of read/write access are supported.

- Multiple reads/multiple writes
- Single read/single write

Multiple reads/multiple writes occur in the following cases.

- 1. CPU burst access (cache replace)
- 2. Longword (32-bit) access to the SDRAM data bus having 8-bit or 16-bit width
- 3. Word (16-bit) access to the SDRAM data bus having 8-bit width
- 4. Multiple data transfers by DMA pipeline transfer

The access timing can be set separately for each channel, using the SDRAMm timing register (SDmTR). Access timing examples are described below.

(a) Multiple Read/Multiple Write Accesses

Figure 10.15 shows a timing example for multiple reads of 4 units of data, and figure 10.16, for multiple writes of 4 units of data.

The number of DMA transfers performed will vary depending on factors such as the number of transfers and the transfer data size per operand and the SDRAM bus width. Read commands or write commands may or may not be issued consecutively in response to an access request from the bus master. If read commands or write commands are not issued consecutively, a deselect command is issued between them.

Furthermore, deactivation and activation are performed automatically when the SDRAM row address changes during a DMA transfer operation.

Figure 10.17 shows a timing example for multiple reads of 4 units of data, and figure 10.18, for multiple writes of 4 units of data, unless read/write commands are issued consecutively. Figure 10.19 shows a timing example for multiple writes with a row address change.

The access timing varies according to the settings in the SDRAMm timing register (SDmTR).

Figure 10.16 Multiple Write Timing Example (Multiple Writes of 4 Data Units, Shortest Timing Settings) Consecutive Write Commands Issued

Figure 10.17 Multiple Read Timing Example (Multiple Reads of 4 Data Units, Shortest Timing Settings) Non-Consecutive Read Commands Issued

Figure 10.18 Multiple Write Timing Example (Multiple Writes of 4 Data Units, Shortest Timing Settings) Non-Consecutive Write Commands Issued

Figure 10.19 Multiple Write Timing Example (Multiple Writes of 4 Data Units, Shortest Timing Settings) Access Spanning Rows

(b) Single Read/Single Write Access

Figure 10.20 shows a timing example for single read operation and figure 10.21 for single write operation. The access timing is modified by means of settings in the SDRAMm timing register (SDmTR).

Figure 10.20 Single Read Timing Example (Shortest Timing Settings)

Figure 10.21 Single Write Timing Example (Shortest Timing Settings)

(10) Mode Register Setting

Writing to the SDRAMm mode register (SDmMOD) causes mode register set and extended mode register set commands to be issued to SDRAM for individual channels. Settings the SDRAMm mode register (SDmMOD) should be made separately for each channel.

Figure 10.22 shows the operation timing for mode register setting.

Figure 10.22 Operation Timing for Mode Register Setting

(11) Clock Stop Control Signal

The SDRAMC is provided with a clock stop control function, which can stop CKIO in deeppower-down mode. The function can be enabled or disabled using the DCKSEN bit in the SDRAM clock stop control signal setting register (SDCKSCNT).

CKIO runs continuously if the clock stop control function is disabled.

If the clock stop control function is enabled, CKIO stops or restarts operation in synchronization with transition to or from deep-power-down mode.

In transition to deep-power-down mode, CKIO is stopped (low level) after the deep-power-down entry command is issued. In recovery from deep-power-down mode, a deep-power-down exit command is issued when the clearing of the DDPD bit to 0 is accepted by SDRAMC, and CKIO restarts operating.

DCKSC, the period from the issuance of deep power-down entry (or exit) command until CKIO stops (or restarts) operating, can be set using the SDRAM clock stop control signal setting register.

Figures 10.23 and 10.24 show the operation timing of the clock stop control function.

Figure 10.23 Clock Stop Control Function Operation Timing (Transition to Deep-Power-Down Mode)

Figure 10.24 Clock Stop Control Function Operation Timing (Recovery from Deep-Power-Down Mode)

(12) SDRAMC Setting Examples

The SDRAMC setting procedure, timing register setting examples, and the procedure for transitioning to and recovering from self-refresh mode, power-down mode, and deep-power-down mode are described below.

(a) SDRAMC Setting Procedure

Figure 10.25 shows the SDRAMC setting procedure.

Note that the specifications such as the power-up sequence may vary with the SDRAM in use.

Study the SDRAM specifications carefully before designing your system. For example, when the SDRAM in use requires that the DQM pin be held "H" during the initialization sequence, set the SDRAM according to the procedure shown in figure 10.25 (b). Since the initialization sequence adopted for this LSI is compliant with the JEDEC standard, the value of DQM pin is not guaranteed from the power-up is supplied and through the initialization sequence.

Figure 10.25 (a) SDRAMC Setting Procedure (Basic Setting Example)

Figure 10.25 (b) SDRAMC Setting Procedure (when DQM Pin is Need to Keep "H" in the Initialization)

(b) Procedure for Transition to and Recovery from Self-Refresh Mode

Figure 10.26 shows the procedure for transitioning to and recovering from self-refresh mode.

Figure 10.26 Procedure for Transition to and Recovery from Self-Refresh Mode

Notes: Before transitioning to or recovering from self-refresh mode it is necessary to halt SDRAM access to the affected channel. Consequently, it is not possible to transition to or recover from self-refresh mode while programs or DMA operations that access SDRAM are in progress. Pay attention to the following points when writing programs.

- 1. Before transitioning to self-refresh mode, disable any DMA channel transfers that access the SDRAM area of the channel.
- 2. If programs are to be executed during transition to self-refresh mode, in self-refresh mode, or during recovery from self-refresh mode, design them in such a way that they will not include operands accessing or fetching (including pre-fetching) instructions stored in the SDRAM area.

(c) Procedure for Transition to and Recovery from Deep-Power-Down Mode

Figure 10.27 shows the procedure for transitioning to deep-power-down mode.

Figure 10.27 Procedure for Transition to Deep-Power-Down Mode

Figure 10.28 shows the procedure for recovering from deep-power-down mode.

Figure 10.28 Procedure for Recovery from Deep-Power-Down Mode

- Notes: Before transitioning to or recovering from deep-power-down mode it is necessary to halt SDRAM access to the affected channels. Consequently, it is not possible to transition to or recover from deep-power-down mode while programs or DMA operations that access SDRAM are in progress. Pay attention to the following points when writing programs.
	- 1. Before transitioning to deep-power-down mode, prohibit any DMA channel transfers that access the SDRAM area of the affected channels.
	- 2. If programs are to be executed during transition to deep-power-down mode, in deeppower-down mode, or during recovery from deep-power-down mode, design them in such a way that they will not include operands accessing or fetching (including prefetching) instructions stored in the SDRAM area.

(d) Timing Register Set Values and Access Timing

The correspondence between the SDRAMm timing register (SDmTR) set values and the read and write access timing is described below.

• Multiple Read Timing Setting Examples

Figures 10.29 to 10.31 show the correspondence between the timing of multiple read operations involving 4 data units and the set values of the SDRAMm timing register (SDmTR). Table 10.11 lists the SDRAMm timing register (SDmTR) set values for each figure.

Table 10.11 SDITR Set Value Correspondence Table (Multiple Read Timing)

Figure 10.30 Multiple Read Timing Example 2

Figure 10.31 Multiple Read Timing Example 3

• Multiple Write Timing Setting Examples

Figures 10.32 to 10.34 show the correspondence between the timing of multiple write operations involving 4 data units and the set values of the SDRAMm timing register (SDmTR). Table 10.12 lists the SDRAMm timing register (SDmTR) set values for each figure.

Table 10.12 SDITR Set Value Correspondence Table (Multiple Write Timing)

Figure	DRAS	DRCD	DPCG	DWR
Figure 10.32	010	00	001	
Figure 10.33	000		001	
Figure 10.34	000		001	

Figure 10.32 Multiple Write Timing Example 1

Figure 10.33 Multiple Write Timing Example 2

Figure 10.34 Multiple Write Timing Example 3

• Single Read Timing Setting Examples

Figures 10.35 to 10.37 show the correspondence between the timing of single read operations and the set values of the SDRAMm timing register (SDmTR). Table 10.13 lists the SDRAMm timing register (SDmTR) set values for each figure.

Figure 10.35 Single Read Timing Example 1

Figure 10.36 Single Read Timing Example 2

Figure 10.37 Single Read Timing Example 3

• Single Write Timing Setting Examples

Figures 10.38 to 10.40 show the correspondence between the timing of single write operations and the set values of the SDRAMm timing register (SDmTR). Table 10.14 lists the SDRAMm timing register (SDmTR) set values for each figure.

Figure 10.38 Single Write Timing Example 1

Figure 10.39 Single Write Timing Example 2

Figure 10.40 Single Write Timing Example 3

(13) External Address/SDRAM Address Signal Multiplexing

(a) Address Multiplexing

Either of addresses used for accessing an external device or SDRAM is output through external address pins.

Table 10.15 External Address/SDRAM Address Pins

(14) Address Register Setting Values

(a) Supported SDRAM Configurations

Tables 10.16 to 10.21 list the SDRAM configurations supported for bus widths of 8, 16, and 32 bits. These tables are intended to help understand the relationships between the supported SDRAM configurations and address multiplexing.

addr[25:0] is the logical address used by the CPU and DMAC in access to the SDRAM. The table below lists how the settings of DSZ and DDBW determine which signals are output on the SDRAM-access pins.

Table 10.16 Case for 8-Bit External Data Bus Width $(BSIZE^* = (1, 0))$

Notes: 1. BSIZE represents the BSIZE bit in the SDCmCNT control register.

- 2. DSZ represents the DSZ bit in the SDRAMm address register.
- 3. DDBW represents the DDBW bit in the SDRAMm address register.
- 4. ba[1:0] and ma[12:0] represent, respectively, the SDRAM bank address and SDRAM address.
- 5. When the RD, WR or PRA command is issued, this carries the pre-charge option signal.

Notes: 1. BSIZE represents the BSIZE bit in the SDCmCNT control register.

2. DSZ represents the DSZ bit in the SDRAMm address register.

3. DDBW represents the DDBW bit in the SDRAMm address register.

 4. ba[1:0] and ma[12:0] represent, respectively, the SDRAM bank address and SDRAM address.

Table 10.18 Case for 16-Bit External Data Bus Width $(BSIZE^{*1} = (0, 0))$ **(2)**

Notes: 1. BSIZE represents the BSIZE bit in the SDCmCNT control register.

2. DSZ represents the DSZ bit in the SDRAMm address register.

3. DDBW represents the DDBW bit in the SDRAMm address register.

 4. ba[1:0] and ma[12:0] represent, respectively, the SDRAM bank address and SDRAM address.

Table 10.19 Case for 32-Bit External Data Bus Width $(BSIZE^{*1} = (0, 1))$ **(1)**

Notes: 1. BSIZE represents the BSIZE bit in the SDCmCNT control register.

2. DSZ represents the DSZ bit in the SDRAMm address register.

3. DDBW represents the DDBW bit in the SDRAMm address register.

 4. ba[1:0] and ma[12:0] represent, respectively, the SDRAM bank address and SDRAM address.

Notes: 1. BSIZE represents the BSIZE bit in the SDCmCNT control register.

- 2. DSZ represents the DSZ bit in the SDRAMm address register.
- 3. DDBW represents the DDBW bit in the SDRAMm address register.
- 4. ba[1:0] and ma[12:0] represent, respectively, the SDRAM bank address and SDRAM address.
- 5. When the RD, WR or PRA command is issued, this carries the pre-charge option signal.

Table 10.21 Case for 32-Bit External Data Bus Width $(BSIZE^{*1} = (0, 1))$ **(3)**

Notes: 1. BSIZE represents the BSIZE bit in the SDCmCNT control register.

2. DSZ represents the DSZ bit in the SDRAMm address register.

3. DDBW represents the DDBW bit in the SDRAMm address register.

 4. ba[1:0] and ma[12:0] represent, respectively, the SDRAM bank address and SDRAM address.

10.6 Connection Examples

The following figures show examples of connecting SRAM or SDRAM to this LSI.

Figure 10.42 Example of Connecting a 16-Bit Data-Width SRAM

Figure 10.43 Example of Connecting an 8-Bit Data-Width SRAM

Figure 10.44 Example of Connecting a 32-Bit Data-Width SRAM (with Byte Control)

Figure 10.45 Example of Connecting a 16-Bit Data-Width SRAM (with Byte Control)

Figure 10.46 Example of Connecting a 32-Bit Data-Width SDRAM

Figure 10.47 Example of Connecting a 16-Bit Data-Width SDRAM

10.7 Usage Notes

10.7.1 Write Buffer

In writing to locations in normal space or in SDRAM space, data to be written are held once in the internal write buffer of the BSC, after which writing to the device connected in external space or SDRAM space (external device) proceeds. Since writing of data from the write buffer to the external device is automatic, no processing by software is required.

At the same time, care is required on the following point. Writing by the CPU or DMAC appears to be complete at the point where the data are stored in the write buffer mentioned above. In other words, at the point where writing by the CPU or DMAC is completed, writing to the external device may in some cases not have been completed. To confirm completion of writing to the external device, execute dummy reading from the normal space or SDRAM space. Completion of dummy reading guarantees the completion of writing to the external device for previous writeaccess. The target device for dummy reading need not be the same as the target for writing. Furthermore, the space need not be the same.

10.7.2 Point for Caution at the Time of a Transition to Software Standby or deep Standby Mode

In cases where a transition to software standby or deep standby mode follows the execution of writing to locations in normal space or in SDRAM space, data may remain within the internal write buffer of the BSC. To ensure that data do not remain within the write buffer, execute dummy reading from an external device in the same way as described above.

Section 11 Direct Memory Access Controller (DMAC)

The DMA controller (DMAC) is a module that handles high-speed data transfer without CPU intervention in response to requests from software, on-chip peripheral modules, or external pins (external modules). The DMAC itself does not distinguish between requests from on-chip peripheral modules and those from external pins (external modules). The DMA supports data transfer between memories, between memory and on-chip peripheral modules, and between onchip peripheral modules.

11.1 Features

- Number of channels: 14 channels (four channels can accept external requests; two-dimensional addressing supported on eight channels)
- Transfer requests: Software trigger and requests from on-chip peripheral modules (48 sources) and external pins (4 sources)
- Maximum transfer bytes: 64 Mbytes
- Address space: 4 Gbytes
- Transfer data sizes:
	- Single data transfer: 8, 16, and 32 bits
	- \sim Single operand transfer: 1, 2, 4, 8, 16, 32, 64, and 128 data units
	- Non-stop transfer: Until the byte counter reaches "0"
- Transfer mode:
	- Cycle-stealing transfer
	- Piepelined transfer
- Maximum transfer speed
	- ⎯ Cycle-stealing transfer: Minimum of three bus clock cycles per unit data transfer
	- Pipeline transfer: Minimum of one bus clock cycle per unit data transfer
- Transfer conditions
	- The following transfer method can be selected.
	- Unit operand transfer: Transfers data of one operand per DMA request. Arbitrates channels per transfer of one operand.
		- Requires request trigger per transfer of one operand.
	- ⎯ Sequential operand transfer: Repeats transfer of one operand per DMA request until the byte count reaches "0".
		- Arbitrates channels per transfer of one operand.
		- Requires only the first request trigger.
- ⎯ Non-stop transfer: Transfers data until the byte count reaches "0" per DMA request. Does not arbitrate channels.
	- Requires only the first request trigger.
- Channel priorities: Channel $0 >$ channel $1 > ... >$ channel $12 >$ channel 13 (this priority order is fixed)
- Interrupt request: Two types of interrupt requests (generated when the byte counter reaches "0")
	- Interrupt request signal for each channel
	- Interrupt request signal common to multiple channels
- Reload function: Reloads the source address, destination address, and byte count.
- Rotate function: The address rotate function can be set.
- Two-dimensional addressing: This can be specified in channels 0 to 7.
- The DMAC suspend/restart/stop function can be set.

Note: Terminologies in this section are defined as follows: Single data transfer: Transfer in one read cycle or one write cycle by the DMAC Single operand transfer: Continuous data transfer by the DMAC on one channel (amount of data to be transferred is set in a register) Single DMA transfer: Transfer of data by the number of data set in the byte count register from the start address to the end address Channel number: $n = 0$ to 13 Two-dimensional addressing-supported channel number: $m = 0$ to 7 Request source number: $k = 0$ to 52

Figure 11.1 is a block diagram of the DMAC.

Figure 11.1 DMAC Block Diagram

11.2 Input/Output Pins

Table 11.1 shows DMAC pin functions.

Table 11.1 Pin Configuration

11.3 Register Descriptions

The DMAC has the registers shown in tables 11.2 and 11.3. All these registers are initialized by a power-on reset or in deep standby mode and the previous settings are lost.

Table 11.2 Register Configuration (Registers Not Related to Two-Dimensional Addressing)

Table 11.3 Register Configuration (Registers Related to Two-Dimensional Addressing)

11.3.1 DMA Current Source Address Registers (DMCSADRn)

DMCSADRn is a register used to specify the start address of the transfer source. The value in this register is transferred to the working source-address register when DMA transfer starts. The contents of the working source-address register are returned to this register when an operand transfer is completed. If the rotate setting $(SAMOD = 011)$ is made for the source address, however, the contents of the working source-address register are not returned. If the sourceaddress reload function is enabled, the contents stored in the DMA reload source address register (DMRSADRn) are returned to this register when DMA transfer is completed. This register must be set regardless of whether the reload function is enabled or disabled.

Notes: 1. Set this register so that DMA transfer is performed for the following selected transfer data sizes within the correctly arranged address boundaries:

- When the transfer data size is set to 16 bits (SZSEL = 001): (bit 0) = 0
- When the transfer data size is set to 32 bits (SZSEL = 010): (bit 1, bit 0) = $(0, 0)$
- 2. Only write data to this register when the corresponding channel is not undergoing single operand transfer (the DASTS bit of the corresponding channel in the DMA arbitration status register (DMASTS) is 0) and DMA transfer is disabled (the DMST bit of the DMA activation control register (DMSCNT) is 0 or the DEN bit of DMA control register B (DMCNTBn) is 0). In other cases, operation is not guaranteed when data is written to this register.

11.3.2 DMA Current Destination Address Registers (DMCDADRn)

DMCDADRn is a register used to specify the start address of the transfer destination. The value in this register is transferred to the working destination-address register when DMA transfer is started. The contents of the working destination-address register are returned to this register when an operand transfer is completed. If the rotate setting $(DAMOD = 011)$ is made for the destination address, however, the contents of the working destination-address register are not returned. If the destination-address reload function is enabled, the contents stored in the DMA reload destination address register (DMRDADRn) are returned to this register when DMA transfer is completed. This register must be set regardless of whether the reload function is enabled or disabled.

Notes: 1. Set this register so that DMA transfer is performed for the following selected transfer data sizes within the correctly arranged address boundaries:

- When the transfer data size is set to 16 bits (SZSEL = 001): (bit 0) = 0
- When the transfer data size is set to 32 bits (SZSEL = 010): (bit 1, bit 0) = $(0, 0)$
- 2. Only write data to this register when the corresponding channel is not undergoing single operand transfer (the DASTS bit of the corresponding channel in the DMA arbitration status register (DMASTS) is 0) and DMA transfer is disabled (the DMST bit of the DMA activation control register (DMSCNT) is 0 or the DEN bit of DMA control register B (DMCNTBn) is 0). In other cases, operation is not guaranteed when data is written to this register.

11.3.3 DMA Current Byte Count Register (DMCBCTn)

DMCBCTn is a register used to specify the number of bytes to be transferred by DMA. The value in this register is transferred to the working byte-count register when DMA transfer is started and decremented the number of bytes to be transferred per single data transfer. How much this value is decremented depends on the transfer data size as follows:

- When the transfer data size is set to 8 bits $(SZSEL = 000)$: -1
- When the transfer data size is set to 16 bits $(SZSEL = 001)$: -2
- When the transfer data size is set to 32 bits ($SZSEL = 010$): -4

When the value in the working byte count register reaches H'000 0000, DMA transfer ends (transfer end when the byte count reaches "0"). The corresponding bit of the DMA transfer end detection register (DMEDET) is set to 1. If the byte count reload function is disabled, the contents of the working byte count register are returned to this register when the channel for DMA transfer switches or DMA transfer ends. If the byte count reload function is enabled, the contents of the DMA reload byte count register (DMRBCTn) are returned to this register when DMA transfer is completed. This register must be set regardless of whether the reload function is enabled or disabled.

this register.

11.3.4 DMA Reload Source Address Register (DMRSADRn)

DMRSADRn is a register used to set an address to be reloaded to the DMA current source address register (DMCSADRn).

To enable the reload function, set the DMA source address reload function enable bit (SRLOD) in DMA control register A (DMCNTAn) to 1. In this case, set both the DMA current source address register (DMCSADRn) and DMA reload source address register (DMRSADRn).

Note: Set this register so that DMA transfer is performed for the following selected transfer data sizes within the correctly arranged address boundaries:

• When the transfer data size is set to 16 bits (SZSEL = 001): (bit 0) = 0

• When the transfer data size is set to 32 bits (SZSEL = 010): (bit 1, bit 0) = $(0, 0)$

11.3.5 DMA Reload Destination Address Register (DMRDADRn)

DMRDADRn is a register used to set an address to be reloaded to the DMA current destination address register (DMCDADRn).

To enable the reload function, set the DMA destination address reload function enable bit (DRLOD) in DMA control register A (DMCNTAn) to 1. In this case, set both the DMA current destination address register (DMCDADRn) and DMA reload destination address register (DMRDADRn).

Note: Set this register so that DMA transfer is performed for the following selected transfer data sizes within the correctly arranged address boundaries:

• When the transfer data size is set to 16 bits (SZSEL = 001): (bit 0) = 0

• When the transfer data size is set to 32 bits (SZSEL = 010): (bit 1, bit 0) = $(0, 0)$

11.3.6 DMA Reload Byte Count Register (DMRBCTn)

DMRBCTn is a register used to set the byte count to be reloaded to the DMA current byte count register (DMCBCTn). To enable the reload function, set the DMA byte count reload function enable bit (BRLOD) in DMA control register A (DMCNTAn) to 1. In this case, set both the DMA current byte count register (DMCBTn) and DMA reload byte count register (DMRBCTn).

25 to 0 RBC[25:0] Undefined R/W Number of DMA transfer bytes for reloading

Note: Set this register so that the byte count becomes 0 as follows when the final data is sent in a DMA transfer:

• When the transfer data size is set to 16 bits (SZSEL = 001): (bit 0) = 0

• When the transfer data size is set to 32 bits (SZSEL = 010): (bit 1, bit 0) = $(0, 0)$

11.3.7 DMA Mode Register (DMMODn)

DMMODn controls the amount of data, unit data size selection, address direction, and various signal outputs.

Note: Only write data to this register when the corresponding channel is not engaged in single operand transfer (the corresponding DASTS bit in the DMA arbitration status register (DMASTS) is 0) and DMA transfer is disabled (the DMST bit in the DMA activation control register (DMSCNT) is 0 or the DEN bit in DMA control register B (DMCNTBn) is 0). In other cases, operation is not guaranteed when data is written to this register.

Table 11.4 shows the counter increment and decrement of DMA source/destination address registers (for details on the "rotation" addressing mode, see section 11.12, Rotate Function). If two-dimensional addressing is specified in these bits (SAMOD and DAMOD), the settings of registers related to two-dimensional addressing (section 11.3.16, DMA Two-Dimensional Addressing Column Setting Register (DM2DCLMm), and after) become valid. When performing pipelined transfer to or from external devices and modules that support burst access, make sure to set the direction bits to select address incrementation (001), rotation (011), or two dimensions (100).

Table 11.4 Counter Increment/Decrement for DMA Source/Destination Address Registers

Table 11.5 shows the relationship between DMA request sources and the DMA-active signal output control bit for the source. If the DREQ0 to DREQ3 pins are selected as the DMA request source, select "0: Stop" or "1: Output" as required. The signal corresponding to this setting is output to the DACT0 to DACT3 external pins (see section 11.9, DMA Acknowledge Signal Output and DMA-Active Signal Output). If the software trigger is selected, setting of this bit has no effect, so either 0 or 1 can be set. If other DMA request sources are selected, be sure to set "1: Output".

Table 11.5 Relationship between DMA Request Sources and DMA-active signal Output Control Bit for Source

[Legend]

Ο: Can be set

×: Setting prohibited

 \equiv : Setting ignored

Table 11.6 shows the relationship between DMA request sources and the DMA-active signal output control bit for the destination. If the DREQ0 to DREQ3 pins are selected as the DMA request source, select "0: Stop" or "1: Output" as required. The signal corresponding to this setting is output to the $\overline{DACT0}$ to $\overline{DACT3}$ external pins (see section 11.9, DMA Acknowledge Signal Output and DMA-Active Signal Output). If the software trigger is selected, setting of this bit has no effect, so either 0 or 1 can be set. If other DMA request sources are selected, be sure to set "1: Output".

Table 11.6 Relationship between DMA Request Sources and DMA-active signal Output Control Bit for Destination

Ο: Can be set

×: Setting prohibited

 \equiv : Setting ignored

Table 11.7 shows the relationship between DMA request sources and the DMA end signal output control bit. If the DREQ0 to DREQ3 pins are selected as the DMA request source, select 00, 01, 10, or 11 as required. The signal corresponding to this setting is output to the TEND0 to TENDT3 external pins (see section 11.5.3, DMA End Signal Output). If USB_0, USB_1, or a 2DG-related source is selected, be sure to select 11. If the software trigger or other DMA request source is selected, setting of this bit has no effect, so either 0 or 1 can be set.

Table 11.7 Relationship between DMA Request Sources and DMA End Signal Output Control Bit

[Legend]

Ο: Can be set

×: Setting prohibited

-: Setting ignored

11.3.8 DMA Control Register A (DMCNTAn)

DMCNTAn is used to select transfer modes, transfer conditions, and DMA sources and control various reload functions.

Note: Modify the settings of bits of this register other than the reload function enable bits (BRLOD, SRLOD, and DRLOD) only when the corresponding channel is not undergoing single operand transfer (the DASTS bit of the DMA arbitration status register (DMASTS) is 0) and DMA transfer is disabled (the DMST bit of the DMA activation control register (DMSCNT) is 0 or the DEN bit of DMA control register B (DMCNTBn) is 0). In other cases, operation is not guaranteed when data is written to this register.

Table 11.8 Relationship between DMA Request Sources and Input Sense Modes

[Legend]

Ο: Can be set

×: Setting prohibited

Note: The input sense modes of other DMA request sources may be changed in the future because they are for the preliminary version.

11.3.9 DMA Control Register B (DMCNTBn)

DMCNTBn controls whether to enable or disable DMA transfer, transfer enable clearing, and internal status clearing. This register can also reference the DMA request status.

Note: If the software trigger is selected as the DMA request source, the DMA request bit (DREQ) can be set to 1 regardless of the settings of the DMA transfer enable bit (DEN) and DMAC module activate bit (DMST) and the operand transfer status. However, even if the software trigger is selected as the DMA request source, clear the DMA request bit (DREQ) to 0 or write data to the DMAC internal state clear bit (DSCLR) only when the corresponding channel is not engaged in single operand transfer (the corresponding DASTS bit in the DMA arbitration status register (DMASTS) is 0) and DMA transfer is disabled (the DMST bit in the DMA activation control register (DMSCNT) is 0 or the DEN bit in DMA control register B (DMCNTBn) is 0). In other cases, operation is not guaranteed when data is written to this register.

11.3.10 DMA Activation Control Register (DMSCNT)

DMSCNT controls the entire DMAC operation.

11.3.11 DMA Interrupt Control Register (DMICNT)

DMICNT controls DMA interrupt for each channel.

Note: Bits 31 to 18 correspond to channels 0 to 13.

11.3.12 DMA Common Interrupt Control Register (DMICNTA)

DMICNTA controls DMA interrupts for each channel.

Note: Bits 31 to 18 correspond to channels 0 to 13.

11.3.13 DMA Interrupt Status Register (DMISTS)

DMISTS consists of DMA interrupt request status bits.

not affect this setting.

• Condition for clearing these bits to 0

Clearing the DMA transfer end condition detect bit (DEDET) of the DMA transfer end detection register (DMEDET) corresponding to the channel where the interrupt occurred

Notes: 1. This register is read-only.

2. Bits 31 to 18 correspond to channels 0 to 13.

11.3.14 DMA Transfer End Detection Register (DMEDET)

DMEDET references the DMA transfer end detection status of each channel. Writing 0 to the DEDET bits is invalid and 1s written to these bits are not retained.

Note: Bits 31 to 18 correspond to channels 0 to 13.

11.3.15 DMA Arbitration Status Register (DMASTS)

DMASTS is used to reference the DMA transfer status of each channel. Writing to this register is invalid.

Note: Bits 31 to 18 correspond to channels 0 to 13.

11.3.16 DMA Two-Dimensional Addressing Column Setting Register (DM2DCLMm)

DM2DCLMm is a register used to set the number of data columns in one block in twodimensional addressing.

Figure 11.2 Specifying Two-Dimensional Blocks

11.3.17 DMA Two-Dimensional Addressing Row Setting Register (DM2DROWm)

DM2DROWm is a register used to set the number of rows in one block in two-dimensional addressing.

11.3.18 DMA Two-Dimensional Addressing Block Setting Register (DM2DBLKm)

DM2DBLKm is a register used to set the number of blocks per line in two-dimensional addressing.

11.3.19 DMA Two-Dimensional Addressing Next Row Offset Register (DM2DNROSTm)

DM2DNROSTm is a register used to set the offset for calculating the start address of the next row in two-dimensional addressing.

11.3.20 DMA Two-Dimensional Addressing Next Block Offset Register (DM2DNBOSTm)

DM2DMBOSTm is a register used to set the offset for calculating the start address of the next block in two-dimensional addressing.

11.3.21 DMA Two-Dimensional Addressing Next Line Offset Register (DM2DNLOSTm)

DM2DNLOSTm is a register used to set the offset for calculating the start address of the next line in two-dimensional addressing.

11.3.22 DMA Reload Two-Dimensional Addressing Column Setting Register (DMR2DCLMm)

DMR2DCLMm is a register used to set the number of data columns to be reloaded to the DMA two-dimensional addressing column setting register (DM2DCLMm). To enable the reload function, set the two-dimensional reload function enable bit (2DRLOD) of DMA control register A (DMCNTAm) to 1. When enabled, it is necessary to set both the DMA two-dimensional addressing column setting register (DM2DCLMm) and DMA reload two-dimensional addressing column setting register (DMR2DCLMm).

11.3.23 DMA Reload Two-Dimensional Addressing Row Setting Register (DMR2DROWm)

DMR2DROWm is a register used to set the number of rows to be reloaded to the DMA twodimensional addressing row setting register (DM2DROWm). To enable the reload function, set the two-dimensional reload function enable bit (2DRLOD) of DMA control register A (DMCNTAm) to 1. When enabled, it is necessary to set both the DMA two-dimensional addressing row setting register (DM2DROWm) and DMA reload two-dimensional addressing row setting register (DMR2DROWm).

11.3.24 DMA Reload Two-Dimensional Addressing Block Setting Register (DMR2DBLKm)

DMR2DBLKm is a register used to set the number of blocks to be reloaded to the DMA twodimensional addressing block setting register (DM2DBLKm). To enable the reload function, set the two-dimensional reload function enable bit (2DRLOD) of DMA control register A (DMCNTAm) to 1. When enabled, it is necessary to set both the DMA two-dimensional addressing block setting register (DM2DBLKn) and DMA reload two-dimensional addressing block setting register (DMR2DBLKm).

11.3.25 DMA Reload Two-Dimensional Addressing Next Row Offset Register (DMR2DNROSTm)

DMR2DNROSTm is a register used to set the offset to be reloaded to the DMA two-dimensional row offset register (DM2DNROSTm). To enable the reload function, set the two-dimensional reload function enable bit (2DRLOD) of DMA control register A (DMCNTAm) to 1. When enabled, it is necessary to set both the DMA two-dimensional addressing row offset register (DM2DNROSTm) and DMA reload two-dimensional addressing row offset register (DMR2DNROSTm).

11.3.26 DMA Reload Two-Dimensional Addressing Next Block Offset Register (DMR2DNBOSTm)

DMR2DNBOSTm is a register used to set the offset to be reloaded to the DMA two-dimensional addressing next block offset register (DM2DNBOSTm). To enable the reload function, set the two-dimensional reload function enable bit (2DRLOD) of DMA control register A (DMCNTAm) to 1. When enabled, it is necessary to set both the DMA two-dimensional addressing next block offset register (DM2DNBOSTm) and DMA reload two-dimensional addressing next block offset register (DMR2DNBOSTm).

11.3.27 DMA Reload Two-Dimensional Addressing Next Line Offset Register (DMR2DNLOSTm)

DMR2DNLOSTm is a register used to set the offset to be reloaded to the DMA two-dimensional addressing next line offset register (DM2DNLOSTm). To enable the reload function, set the twodimensional reload function enable bit (2DRLOD) of DMA control register A (DMCNTAm) to 1. When enabled, it is necessary to set both the DMA two-dimensional addressing next line offset register (DM2DNLOSTm) and DMA reload two-dimensional addressing next line offset register (DMR2DNLOSTm).

11.4 Operation

11.4.1 DMA Transfer Mode

Two DMA transfer modes are available: cycle-stealing transfer mode and pipelined transfer mode. These modes can be selected by using the DMA transfer mode select bits (MDSEL) of DMA control register A (DMCNTAn).

Figure 11.3 shows how bus mastership alternates between the DMAC and CPU in DMA transfer modes.

(1) Cycle-Stealing Transfer Mode

Setting the DMA transfer mode select bits (MDSEL) to 00 selects cycle-stealing transfer mode.

In cycle-stealing transfer mode, the DMAC operates, leaving at least one cycle between the read and write access cycles (activations) of each single data transfer. For this reason, access from the CPU is possible during this interval (the CPU can access the BIU part of the source or destination target).

(2) Piepelined Transfer Mode

Setting the DMA transfer mode select bits (MDSEL) to 01 selects pipelined transfer mode.

In pipelined transfer mode, the DMAC consecutively accesses the bus for read access or write access or both. For this reason, access from the CPU is not accepted till the current single operand transfer ends (the CPU cannot access the BIU part of the source or destination target).

Piepelined transfer through a single BIU is also not possible.

Figure 11.3 Example of Bus Mastership Alternation between DMAC and CPU in Various DMA Transfer Modes

11.4.2 DMA Transfer Conditions

Three DMA transfer conditions are available: unit operand transfer, sequential operand transfer, and non-stop transfer. These conditions can be selected by using the DMA transfer condition select bits (DSEL) of DMA control register A (DMCNTAn). Each DMA transfer condition is described below. Table 11.9 lists DMA transfer conditions, which are illustrated in figure 11.4.

(1) Unit Operand Transfer

Setting the DMA transfer condition select bits (DSEL) to 00 selects unit operand transfer.

A single DMA request transfers data by the amount specified in the single operand transfer data count select bits (OPSEL) of the DMA mode register (DMMODn).

Each time a DMA transfer request is made, the DMAC repeats single operand transfer and ends single DMA transfer when the byte count reaches 0.

In the case that the DMA transfer condition is the unit operand transfer and the input sense mode of DMA request is the level sense, there is the mask period of the DMA request in the channel arbitration period after one operand transfer end (please refer to section 11.7.3, Sense Mode for DMA Requests for details). Therefore, in the channel arbitration period after one operand transfer end, in the case that there is no DMA request of the higher-priority channel than the transferring channel and there is the DMA request of the lower-priority channel than the transferring channel, the DMA transfer of the low-priority channel starts. To execute the DMA transfer of the highpriority channel in succession, please set the DMA transfer condition to the sequential operand transfer or the non-stop transfer.

(2) Sequential Operand Transfer

Setting the DMA transfer condition select bits (DSEL) to 01 selects sequential operand transfer.

A single DMA request transfers data in units of the number of data items set in the single operand transfer data count select bits (OPSEL) (single operand transfer). This data transfer is continued till a single DMA transfer ends (i.e., till the byte count reaches 0). Channel arbitration is performed each time the single operand transfer ends. Data transfer on the current channel is automatically continued if there is no DMA request from a higher-priority channel.

In the case that the DMA transfer condition is the sequential operand transfer, even if the input sense mode of DMA request is the level sense, there is no mask period before the byte count becomes 0. Therefore, the DMA transfer of the low-priority channel than the transferring channel cannot start.

(3) Non-Stop Transfer

Setting the DMA transfer condition select bits (DSEL) to 11 selects non-stop transfer.

A single DMA request continuously transfers data till a single DMA transfer ends (i.e., till the byte count reaches 0). During this interval, DMA requests from higher-priority channels are not accepted because channel arbitration is not performed.

In non-stop transfer, the settings of the single operand transfer data count select bits (OPSEL) are invalid and setting of two-dimensional addressing is prohibited.

Table 11.9 List of DMA Transfer Conditions

Figure 11.4 DMA Transfer Conditions

Table 11.10 shows the combinations of DMA transfer modes and DMA transfer conditions.

Table 11.10 Combinations of DMA Transfer Modes and DMA Transfer Conditions

Notes: 1. Non-stop transfer to BIU_E in pipelined transfer mode cannot be set.

2. Piepelined transfer to the same BIU is prohibited.

11.4.3 DMA Activation

(1) Initial setting of DMAC

Make the initial setting of each register before setting the DMA transfer enable bit (DEN) to 1. Once data transfer has been started, these settings cannot be changed.

The following shows an example of DMAC initial setting.

- 1. DMA mode register (DMMODn)
- 2. DMA control register A (DMCNTAn)
- 3. DMA control register B (DMCNTBn)
- 4. DMA current source address register (DMCSADRn)
- 5. DMA reload source address register (DMRSADRn)
	- ⎯ When the reload function is used
- 6. DMA current destination address register (DMCDADRn)
- 7. DMA reload destination address register (DMRDADRn)
	- ⎯ When the reload function is used
- 8. DMA current byte count register (DMCBCTn)
- 9. DMA reload byte count register (DMRBCTn)
	- ⎯ When the reload function is used

- 10. DMA interrupt control register (DMICNT)
	- When an interrupt is used
- 11. DMA common interrupt control register (DMICNTA)
	- When an interrupt is used
- 12. DMA transfer enable bit (DEN) of DMA control register B (DMCNTBn)
- 13. DMA activation control register (DMSCNT)

(2) DMA Activation

To enable DMA transfer on a channel, set the DMA transfer enable bit (DEN) of DMA control register B (DMCNTBn) and the DMAC module activate bit (DMST) of the DMA activation control register (DMSCNT) corresponding to the channel to 1.

If multiple DMA transfer requests are present, channel priorities are judged, the DMA request corresponding to the highest-priority channel is accepted, and DMA transfer on the channel is started.

Whether DMA requests are present can be checked from the DMA request bit (DREQ) of DMA control register B (DMCNTBn).

When a DMA request is accepted and DMA transfer is started, the DMA arbitration status bit (DASTS) of the channel corresponding to the DMA arbitration status register (DMASTS) is set to 1.

(3) DMA Activation Sources and Restrictions

When the DMAC is activated by an on-chip peripheral module, the transfer source, transfer destination, operand size, data size, transfer conditions, and transfer mode (whether pipelined transfer is available) may be fixed. Table 11.11 shows the transfer methods that can be selected for each DMA request source.

Notes: 1. Words in the Transfer Condition (DSEL) have the following meanings: Unit: Only unit operand transfer can be specified.

> Unit, sequential: Unit operand transfer or sequential operand transfer can be specified. No restriction: Unit operand transfer, sequential operand transfer, or non-stop transfer can be specified.

2. For single data access mode: Operand size = 1, data size = 8, 16, 32 For 16-byte sequential access mode: Set so that operand size \times data size = 16 bytes. For 32-byte sequential access mode: Set so that operand size \times data size = 32 bytes.

11.5 DMA Transfer End and Interrupts

11.5.1 DMA Transfer End

When the value in the DMA current byte count register (DMCBCTn) becomes H'000 0000 (transfer end of all data), the DMA transfer end condition is satisfied and one DMA transfer ends.

The following describes the operations performed when the DMA transfer end condition is detected.

• DMA transfer end detection

The DMA transfer end condition detect bit (DEDET) of the channel corresponding to the DMA transfer end detection register (DMEDET) is set to 1.

• Interrupt request generation

An interrupt request is generated for the interrupt controller according to the settings of the DMA interrupt control register (DMICNT) and the DMA common interrupt control register (DMICNTA).

• DMA end signal output

The DMA end signal (DMATC_N) is output according to the setting of the DMA end signal output control bit (DTCM) of the DMA mode register (DMMODn).

• DMA transfer enable bit (DEN) clearing

If the DMA transfer enable clear bit (ECLR) of DMA control register B (DMCNTBn) is set to 1, the DEN bit of DMA control register B (DMCNTBn) is cleared to 0 and the subsequent DMA transfer of the channel is suspended.

If the DMA transfer enable clear bit (ECLR) is cleared to 0, the DEN bit is not cleared.

• Source address register reloading

If the DMA source address reload function enable bit (SRLOD) of DMA control register A (DMCNTAn) is set to 1, the value in the DMA reload source address register (DMRSADRn) is reloaded to the DMA current source address register (DMCSADRn).

- Destination address register reloading If the DMA destination address reload function enable bit (DRLOD) of DMA control register A (DMCNTAn) is set to 1, the value in the DMA reload destination address register (DMRDADRn) is reloaded to the DMA current destination address register (DMCDADRn).
- Byte count register reloading

If the DMA byte count reload function enable bit (BRLOD) of DMA control register A (DMCNTAn) is set to 1, the value in the DMA reload byte count register (DMRBCTn) is reloaded to the DMA current byte count register (DMCBCTn).

Note: If registers are not reloaded, set $ECLR = 1$ so that the DEN bit is cleared.

11.5.2 DMA Interrupt Requests

The DMAC provides two types of interrupt request signals for the interrupt controller: interrupt request signals per channel (DMINTn_N, $n = 0$ to 13) and a common interrupt request signal (DMINTA_N) that is a collection of interrupt requests per channel.

Figure 11.5 is a block diagram of generating interrupt request signals per channel and a common interrupt request signal.

If the DMA interrupt control bit (DINTM) of the channel corresponding to the DMA interrupt control register (DMICNT) is set to 1 when a DMA transfer ends, an interrupt request for the corresponding channel is generated.

Only those channels for which the DMA common interrupt request signal control bit (DINTA) of the DMA common interrupt control register (DMICNTA) is set to 1 are collected into one and output as the common interrupt request signal.

The generated interrupt request can be cleared to 0 by writing 1 to the DMA transfer end condition detect bit (DEDET) of the corresponding channel.

Figure 11.5 Block Diagram of Generating Interrupt Request Signal per Channel and Common Interrupt Request Signal

11.5.3 DMA End Signal Output

(1) LSI Internal Signal (DMATC_N)

When DMA transfer for the amount of data set in the DMA current byte count register (DMCBCTn) ends, the DMAC outputs the DMA end signal (DMATC_N) in the LSI. DMATC_N is the LSI internal signal, so it cannot be monitored from outside the LSI.

DMATC_N output depends on the setting of the DMA end signal output control bits (DTCM) in the DMA mode register (DMMODn) for the corresponding channel.

- If DTCM is set to 00, DMATC N is not output and fixed at a high level even when DMA transfer ends.
- If DTCM is set to 01, DMATC_N is output in the read cycle immediately before the end of DMA transfer (read cycle of the last unit data transfer).
- If DTCM is set to 10, DMATC N is output in the write cycle immediately before the end of DMA transfer (write cycle of the last unit data transfer).
- If DTCM is set to 11, the low pulse signal of one clock cycle is output as DMATC_N at the same timing as the DMA transfer end interrupt.

Figure 11.6 shows the output timing of the DMA end signal.

Figure 11.6 Output Timing of DMA End Signal

(2) DMA End Output Signal (TENDi**, i = 0 to 3)**

If DMA transfer by the DREQi request from an external pin is selected, the read or write destination is normal space (CS0 to CS5), and transfer mode is cycle-stealing transfer mode, the DMA end output signal (TENDI) can be output to outside the LSI. For access to LSI interior or SDRAM space or when transfer mode is pipelined transfer mode, the TENDi signal is not output.

- To read from normal space by DMA, set the DMA end signal output control bits (DTCM) of the DMA mode register (DMMODn) to select that an end signal is output in the last read cycle $(DTCM = 01)$. TENDi is output when the last one data unit is read by DMA. The TENDi output timing is the same as the DMA-active signal (DACTi) timing (see section 11.9, DMA Acknowledge Signal Output and DMA-Active Signal Output, and section 10, Bus State Controller (BSC)). If DTCM is set to other than 01, TENDi is not output.
- To write to normal space by DMA, set the DTCM bits of DMMODn to select that an end signal is output in the last write cycle ($DTCM = 10$). TENDi is output when the last one data unit is written by DMA. The TENDi output timing is the same as the DMA-active signal (DACTi) timing (see section 11.9, DMA Acknowledge Signal Output and DMA-Active Signal Output, and section 10, Bus State Controller (BSC)). If DTCM is set to other than 10, TENDi is not output.

11.6 Suspending, Restarting, and Stopping of DMA Transfer

11.6.1 Suspending and Restarting of DMA Transfer

Clearing the DMAC module activate bit (DMST) of the DMA activation control register (DMSCNT) to 0 enables you to suspend data transfer on all channels of the DMAC. Clearing the DMA transfer enable bit (DEN) of DMA control register B (DMCNTBn) of the corresponding channel to 0 also enables you to suspend data transfer on the channel.

If the DMST or DEN bit is cleared to 0 when single operand transfer is in progress in the unit operand transfer condition or sequential operand transfer condition, DMA transfer is suspended after single operand transfer has ended without reference to each transfer mode (cycle-stealing mode or pipelined transfer mode).

If the DMST or DEN bit is cleared to 0 when DMA transfer is in progress in the non-stop transfer condition, DMA transfer is not suspended and continues till the DMA transfer end condition is detected (i.e., till the byte count reaches 0).

To restart the suspended channel, set the cleared DMST and DEN bits to 1 to restart DMA transfer.

11.6.2 Stopping of DMA Transfer on Any Channel

To stop DMA transfer on any channel, suspend DMA transfer on that channel and write 1 to the DMAC internal state clear bit (DSCLR) of DMA control register B (DMCNTBn) to initialize the DMAC interior. In this case, only the internal state of the DMAC internal circuit is initialized; each register is not initialized.

11.7 DMA Requests

11.7.1 DMA Request Sources

DMA request sources include software triggers and DMA request signal inputs. DMA request signal input sources are selected from the DMA request source select bits (DCTG) of DMA control register A (DMCNTAn) of each channel.

11.7.2 Synchronization Circuits for DMA Request Signal Inputs

Each DMAC channel is provided with a synchronization circuit to cope with asynchronously input DMA request signals. Therefore, a blank period of a few clock cycles appears during the period from the point when a DMA request signal input such as DREQ0 to DREQ3 goes active until the request is reflected in the DMA request bit (DREQ) in DMA control register B (DMCNTBn).

Figure 11.7 shows an example of the DMA request bit timing of DMA request signal input.

Figure 11.7 Example of DMA Request Bit Timing for DMA Request Signal Input

11.7.3 Sense Mode for DMA Requests

If the DREQ0 to DREQ3 pins (DCTG = 000001 to 000100) are selected by the DMA request source select bits (DCTG), a level sense (01 or 11) or an edge sense (00 or 10) can be selected from the input sense mode select bits (STRG) of DMA control register A (DMCNTAn).

If the software trigger ($DCTG = 000000$) is selected as the DMA request source, select the risingedge sense (00). If the DMA request from the on-chip peripheral modules (DCTG = 000101 to 111001) is selected as the DMA request source, set the sense mode show in table 11.8.

The following describes the level sense and edge sense.

(1) Level Sense

If a level sense $(STRG = 01$ or 11) is selected, whether a DMA request is present is judged from the DMA request signal level. A DMA request is not retained in the DMAC, so retain the DMA request signal level till the acceptance of the DMA request is confirmed.

Figure 11.8 is an example of DMA request acceptance processing when a level sense is selected.

Figure 11.8 Example of DMA Request Acceptance Processing When a Level Sense Is Selected

If a level sense is selected, the DMA request signal of the channel is masked from the start of the last write access in a single operand transfer until two clock cycles after the single operand transfer has ended. This ensures a sufficient length of time for negation of the DMA request signal.

Figure 11.9 shows the DMA request signal mask period when a level sense is selected.

Figure 11.9 DMA Request Signal Mask Period When a Level Sense Is Selected

Therefore, even if a channel for which a level sense is selected retains the DMA request signal level (continues to request DMA transfer) as is after the DMA request is accepted, DMA requests from other channels are accepted, if any, because it is judged that no DMA request is made during the DMA request signal mask period.

For sequential operand transfer, however, the DMA request signal mask period becomes effective only if operand transfer ends when the byte count reaches 0. If operand transfer ends when the byte count is not 0, channel arbitration is performed without the DMA request masked. For nonstop transfer, this mask period becomes effective if DMA transfer ends when the byte count is 0.

If DMA transfer is not continuously performed, the DMA request must be canceled within three cycles after single operand transfer has ended.

(2) Edge Sense

If an edge sense (STRG = 00 or 10) is selected, transition to a rising or falling edge of a DMA request signal is recognized as a DMA request.

When a valid edge is detected, the DMA request bit (DREQ) of DMA control register B (DMCNTBn) is set to 1. The value of this bit is retained even if the input level of the DMA request signal changes later. When a DMA request is accepted and the DMA acknowledge signal is effectively output, the DMA request bit (DREQ) bit is automatically cleared to 0.

In this way, retention of DMA requests in edge sense mode is determined from the value of the DMA request bit (DREQ). For this reason, if the DMA request bit (DREQ) is set to 1, the edges selected according to new DMA request signals are ignored. Figure 11.10 shows an example of DMA acceptance processing when an edge sense is selected.

Figure 11.10 Example of DMA Request Acceptance Processing When an Edge Sense Is Selected

11.8 Determination of DMA Channel Priorities

11.8.1 Channel Priorities

Channel 0 has the highest priority. The priorities of channels are fixed in the following order:

Channel $0 >$ channel $1 >$ channel $2 > ... >$ channel $12 >$ channel 13

11.8.2 Operation at Occurrence of Multiple DMA Requests

The DMAC determines DMA channel priorities per single operand transfer.

If a DMA request with a higher channel priority occurs during operand transfer on one channel, operand transfer on the higher-priority channel is started after operand transfer on the current channel has ended. Figure 11.11 shows an example of DMAC outline operation when multiple DMA requests occur. The thick lines in figure 11.11 indicate the period during which DMA request signals are at a low level (channel 0 (ch0), channel 2 (ch2), and channel 3 (ch3) are set to a level sense and channel 1 (ch1) is set to an edge sense).

- 1. Transfer on channel 3 is started because DMA requests are assumed to be non-existent in channel 2 because channel 2 is during the mask period.
- 2. Transfer on channel 0 is started because this channel has the highest channel priority.
- 3. Transfer on channel 2 is started because this channel has the highest channel priority at this point.
- 4. Transfer on channel 3 is started because there are no other requests at this point.
- 5. If DMA requests for channel 0, channel 1, and channel 3 occur at the same time, transfer on channel 0 is started because channel 0 has the highest priority.
- 6. When transfer on channel 0 ends, transfer on channel 1 is started because channel 1 has the second highest priority.
- 7. If a DMA request (low or high level request edge) occurs during DMA transfer on channel 1, DMA transfer on channel 1 is started again after DMA transfer on channel 1 ends. If an edge sense is set, no mask period exists.
- 8. When DMA transfer on channel 1 ends, DMA transfer on channel 3 is started because there is no other transfer request.
- 9. When channel 3 is during the mask period, DMA transfer is not started because there is no other transfer request. DMA transfer on channel 3 is started after the mask period ends.

Figure 11.11 Example of Outline Operation When Multiple DMA Requests Occur

11.9 DMA Acknowledge Signal Output and DMA-Active Signal Output

(1) LSI Internal Signals

The DMAC outputs the DMA acknowledge signal (DMAACK_N) and the DMA-active signals (DMAACTS_N/DMAACTD_N) for the source and destination when a DMA request is accepted or DMA transfer is performed. These signals are LSI internal signals, so they cannot be monitored from outside the LSI. An on-chip peripheral module that requested DMA transfer recognizes that a DMA transfer request has been accepted and DMA transfer is being performed by monitoring these signals.

• DMA-active signals

Outputs of DMAACTS_N and DMAACTD_N are enabled by setting the DMA-active signal output control bits (SACT/DACT) for source and destination in the DMA mode register of the corresponding channel.

If SACT is set to 1, an active DMAACTS_N signal is output when read access is made. If DACT is set to 1, an active DMAACTD N signal is output when write access is made.

• DMA acknowledge signal

DMAACK N is output during the period from start of single operand transfer to end of single operand transfer.

Figure 11.12 shows the output timing of the DMA acknowledge signal and DMA-active signals.

Figure 11.12 Output Timing of DMA Acknowledge Signal and DMA-active signals

(2) **DMA** Active Output Signal $(\overline{DACT}, i = 0 \text{ to } 3)$

If DMA transfer is in response to the DREQi request from an external pin and the read or write destination is normal space (CS0 to CS5), the DMA active signal $(DACTi)$ can be output to outside the LSI. If the access destination is LSI interior or SDRAM space, the DACTi signal is not output.

- To read from normal space by DMA, set the source DMA-active signal output control bit (SACT) of the DMA mode register (DMMODn) to 1. DACTi is output when normal space is read by DMA. For the DACTi output timing, see section 10, Bus State Controller (BSC). If SACT is cleared to 0, DACTi is not output when normal space is read by DMA.
- To write to normal space by DMA, set the destination DMA-active signal output control bit (DACT) of the DMA module register (DMMODn) to 1. DACTi is output when normal space is written by DMA. See section 10, Bus State Controller (BSC). If DACT is cleared to 0, DACTi is not output when normal space is written by DMA.
- If both SACT and DACT are set to 1 in DMA transfer from one normal space to another, DACTi is output when normal space is read or written by DMA.

(3) DMA Acknowledge Output Signal (DACKi**, i = 0 to 3)**

If DMA transfer is performed in response to the DREQi request from an external pin, the DMA acknowledge signal (DACKi) is output. DACKi is output with the same timing as DMAACK_N that is the DMA acknowledge signal in the LSI.

Note: The DACKi signal indicates the DMA operation timing in the LSI. If normal space or SDRAM space is written by DMA, DMA write access observed outside the LSI may be delayed several cycles compared with DMA write access in the LSI. In this case, DMA write access may be observed outside the LSI after DACKi is negated.

11.10 Units of Transfer and Transfer Byte Positions

The bit size for single data transfer can be set to a byte (8 bits), word (16 bits), or longword (32 bits).

Figure 11.13 Example of DMA Data-Byte Control for 32-Bit Bus Width

11.11 Reload Function

To set the reload function, set each reload function enable bit of DMA control register A (DMCNTAn) per channel and per transfer parameter (source address, destination address or byte count).

To configure the reloading function on a channel supporting two-dimensional addressing, set the two-dimensional reload function enable bit (2DRLOD) of DMA control register A (DMCNTAn) for the corresponding channel. This enables reloading of the six registers related to twodimensional addressing: column setting register, row setting register, block setting register, next row offset register, next block offset register, and next line offset register.

When the DMA transfer end condition is detected, DMA transfer parameters are automatically reloaded.

(1) Reload Register and Current Register

When not using the reload function, set data in the current register. When using the reload function, set data both in the reload register and current register.

Do not write data to the current register during single operand transfer. If data is written to the current register, operation is not guaranteed. The reload register can be set even during single operand transfer but this setting must be made before start of the last operand transfer (DMA transfer end). If the reload register is set after start of the last operand transfer, however, this setting may have not been reflected when data is reloaded after DMA transfer end.

(2) Continuous Transfer to Dispersed Areas

The reload function provides continuous transfer to dispersed areas. Writing values to the DMA reload source and destination address registers (DMRSADRn/DMRDADRn) and DMA reload byte count register (DMRBCTn) before DMA transfer ends enables you to prepare the next transfer parameters without affecting the current DMA transfer (current register). This enables you to continuously transfer several transfer blocks in different transfer areas and with a different number of bytes through the same channel.

Figure 11.14 shows an example of how the reload function transfers blocks between dispersed areas.

Figure 11.14 Example of Transferring Blocks between Dispersed Areas by Using Reload Function

11.12 Rotate Function

When rotation is selected in addressing mode, the address is incremented. After single operand transfer has ended, the contents of the working source address register or working destination address register for which rotation is set serve as the values in the DMA current source address register (DMCSADRn) or DMA current destination address register (DMCDADRn) set when DMA transfer was started.

Figure 11.15 is an example of transfer that uses the rotate function (source: rotation, destination: increment).

Figure 11.15 Example of Transfer That Uses Rotate Function (Source: Rotation, Destination: Increment)

11.13 Usage Note

11.13.1 Note on Transition to Software Standby Mode or Deep Standby Mode

If the SLEEP instruction for transition to software standby mode or deep standby mode during transfer by the DMAC, DMA transfer is not guaranteed because the DMAC stops without waiting for transfer end. Therefore, when making transition to software standby mode or deep standby mode, wait till DMA transfer ends or stop DMA transfer and then execute the SLEEP instruction. Also when changing the PLL multiplication rate, stop DMA transfer in advance.

Section 12 Multi-Function Timer Pulse Unit 2 (MTU2)

This LSI has an on-chip multi-function timer pulse unit 2 (MTU2) that comprises five 16-bit timer channels.

12.1 Features

- Maximum 16 pulse input/output lines
- Selection of eight counter input clocks for each channel
- The following operations can be set:
	- Waveform output at compare match
	- Input capture function
	- Counter clear operation
	- Multiple timer counters (TCNT) can be written to simultaneously
	- $\overline{}$ Simultaneous clearing by compare match and input capture is possible
	- Register simultaneous input/output is possible by synchronous counter operation
	- \overline{A} A maximum 12-phase PWM output is possible in combination with synchronous operation
- Buffer operation settable for channels 0, 3, and 4
- Phase counting mode settable independently for each of channels 1 and 2
- Cascade connection operation
- Fast access via internal 16-bit bus
- 28 interrupt sources
- Automatic transfer of register data
- A/D converter start trigger can be generated
- Module standby mode can be settable
- A total of six-phase waveform output, which includes complementary PWM output, and positive and negative phases of reset PWM output by interlocking operation of channels 3 and 4, is possible.
- AC synchronous motor (brushless DC motor) drive mode using complementary PWM output and reset PWM output is settable by interlocking operation of channels 0, 3, and 4, and the selection of two types of waveform outputs (chopping and level) is possible.
- In complementary PWM mode, interrupts at the crest and trough of the counter value and A/D converter start triggers can be skipped.

Table 12.1 MTU2 Functions

[Legend]

√: Available

—: Not available

Figure 12.1 shows a block diagram of the MTU2.

Figure 12.1 Block Diagram of MTU2

12.2 Input/Output Pins

Table 12.2 Pin Configuration

Note: For the pin configuration in complementary PWM mode, see table 12.54.

12.3 Register Descriptions

The MTU2 has the following registers. For details on register addresses and register states during each process, refer to section 32, List of Registers. To distinguish registers in each channel, an underscore and the channel number are added as a suffix to the register name; TCR for channel 0 is expressed as TCR_0.

12.3.1 Timer Control Register (TCR)

The TCR registers are 8-bit readable/writable registers that control the TCNT operation for each channel. The MTU2 has a total of five TCR registers, one each for channels 0 to 4. TCR register settings should be conducted only when TCNT operation is stopped.

x: Don't care

Table 12.4 CCLR0 to CCLR2 (Channels 0, 3, and 4)

Notes: 1. Synchronous operation is set by setting the SYNC bit in TSYR to 1.

 2. When TGRC or TGRD is used as a buffer register, TCNT is not cleared because the buffer register setting has priority, and compare match/input capture does not occur.

Notes: 1. Synchronous operation is selected by setting the SYNC bit in TSYR to 1.

2. Bit 7 is reserved in channels 1 and 2. It is always read as 0 and cannot be modified.

Table 12.6 TPSC0 to TPSC2 (Channel 0)

Table 12.7 TPSC0 to TPSC2 (Channel 1)

Note: This setting is ignored when channel 1 is in phase counting mode.

Table 12.8 TPSC0 to TPSC2 (Channel 2)

Note: This setting is ignored when channel 2 is in phase counting mode.

12.3.2 Timer Mode Register (TMDR)

The TMDR registers are 8-bit readable/writable registers that are used to set the operating mode of each channel. The MTU2 has five TMDR registers, one each for channels 0 to 4. TMDR register settings should be changed only when TCNT operation is stopped.

Table 12.10 Setting of Operation Mode by Bits MD0 to MD3

[Legend]

X: Don't care

Notes: 1. PWM mode 2 cannot be set for channels 3 and 4.

- 2. Phase counting mode cannot be set for channels 0, 3, and 4.
- 3. Reset synchronous PWM mode, complementary PWM mode can only be set for channel 3. When channel 3 is set to reset synchronous PWM mode or complementary PWM mode, the channel 4 settings become ineffective and automatically conform to the channel 3 settings. However, do not set channel 4 to reset synchronous PWM mode or complementary PWM mode. Reset synchronous PWM mode and complementary PWM mode cannot be set for channels 0, 1, and 2.

12.3.3 Timer I/O Control Register (TIOR)

The TIOR registers are 8-bit readable/writable registers that control the TGR registers. The MTU2 has a total of eight TIOR registers, two each for channels 0, 3, and 4, one each for channels 1 and $2²$

TIOR should be set while TMDR is set in normal operation, PWM mode, or phase counting mode.

The initial output specified by TIOR is valid when the counter is stopped (the CST bit in TSTR is cleared to 0). Note also that, in PWM mode 2, the output at the point at which the counter is cleared to 0 is specified.

When TGRC or TGRD is designated for buffer operation, this setting is invalid and the register operates as a buffer register.

• TIORH_0, TIOR_1, TIOR_2, TIORH_3, TIORH_4

• TIORL_0, TIORL_3, TIORL_4

X: Don't care

Table 12.12 TIORL_0 (Channel 0)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

 2. When the BFB bit in TMDR_0 is set to 1 and TGRD_0 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

Table 12.13 TIOR_1 (Channel 1)

 Description

[Legend]

X: Don't care

 Description

Table 12.14 TIOR_2 (Channel 2)

[Legend]

X: Don't care

X: Don't care

Table 12.16 TIORL_3 (Channel 3)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

 2. When the BFB bit in TMDR_3 is set to 1 and TGRD_3 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

X: Don't care

Table 12.18 TIORL_4 (Channel 4)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

2. When the BFB bit in TMDR_4 is set to 1 and TGRD_4 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

X: Don't care

Table 12.20 TIORL_0 (Channel 0)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

 2. When the BFA bit in TMDR_0 is set to 1 and TGRC_0 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

X: Don't care

Table 12.22 TIOR_2 (Channel 2)

[Legend]

X: Don't care

X: Don't care

Table 12.24 TIORL_3 (Channel 3)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

2. When the BFA bit in TMDR_3 is set to 1 and TGRC_3 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

X: Don't care

Table 12.26 TIORL_4 (Channel 4)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

 2. When the BFA bit in TMDR_4 is set to 1 and TGRC_4 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

12.3.4 Timer Interrupt Enable Register (TIER)

The TIER registers are 8-bit readable/writable registers that control enabling or disabling of interrupt requests for each channel. The MTU2 has six TIER registers, two for channel 0 and one each for channels 1 to 4.

• TIER_0, TIER_1, TIER_2, TIER_3, TIER_4

• TIER2_0

12.3.5 Timer Status Register (TSR)

The TSR registers are 8-bit readable/writable registers that indicate the status of each channel. The MTU2 has six TSR registers, two for channel 0 and one each for channels 1 to 4.

• TSR_0, TSR_1, TSR_2, TSR_3, TSR_4

Bit [.]					
	I TCFD	TCFU TCFV TGFD TGFC TGFB TGFA			
Initial value:					
B/M ¹		R R/(W)*1R/(W)*1R/(W)*1R/(W)*1R/(W)*1R/(W)*1			

Note: 1. Writing 0 to this bit after reading it as 1 clears the flag and is the only allowed way.

 2. If the next flag is set before cleared to 0 after reading 1, this bit remains 1 even when 0 is written to. In this case, read 1 again to clear to 0.

 \bullet TSR2_0

	-		$\overline{}$	$\omega_{\rm{max}}$	$\overline{}$	TGFF	TGFE	
Initial value:		υ						
B/W				R	R.	R/(W)* ¹ R/(W)* ¹		

Note: Writing 0 to this bit after reading it as 1 clears the flag and is the only allowed way. 1.

Notes: 1. Writing 0 to this bit after reading it as 1 clears the flag and is the only allowed way.

 2. If the next flag is set before cleared to 0 after reading 1, this bit remains 1 even when 0 is written to. In this case, read 1 again to clear to 0.

12.3.6 Timer Buffer Operation Transfer Mode Register (TBTM)

The TBTM registers are 8-bit readable/writable registers that specify the timing for transferring data from the buffer register to the timer general register in PWM mode. The MTU2 has three TBTM registers, one each for channels 0, 3, and 4.

12.3.7 Timer Input Capture Control Register (TICCR)

TICCR is an 8-bit readable/writable register that specifies input capture conditions when TCNT_1 and TCNT_2 are cascaded. The MTU2 has one TICCR in channel 1.

12.3.8 Timer A/D Converter Start Request Control Register (TADCR)

TADCR is a 16-bit readable/writable register that enables or disables A/D converter start requests and specifies whether to link A/D converter start requests with interrupt skipping operation. The MTU2 has one TADCR in channel 4.

Note: $*$ Do not set to 1 when complementary PWM mode is not selected.

Notes: 1. TADCR must not be accessed in eight bits; it should always be accessed in 16 bits.

- 2. When interrupt skipping is disabled (the T3AEN and T4VEN bits in the timer interrupt skipping set register (TITCR) are cleared to 0 or the skipping count set bits (3ACOR and 4VCOR) in TITCR are cleared to 0), do not link A/D converter start requests with interrupt skipping operation (clear the ITA3AE, ITA4VE, ITB3AE, and ITB4VE bits in the timer A/D converter start request control register (TADCR) to 0).
- 3. If link with interrupt skipping is enabled while interrupt skipping is disabled, A/D converter start requests will not be issued.
- * Do not set to 1 when complementary PWM mode is not selected.

normal operation mode.

2. These settings are prohibited when complementary PWM mode is not selected.

12.3.9 Timer A/D Converter Start Request Cycle Set Registers (TADCORA_4 and TADCORB_4)

TADCORA_4 and TADCORB_4 are 16-bit readable/writable registers. When the TCNT_4 count reaches the value in TADCORA_4 or TADCORB_4, a corresponding A/D converter start request will be issued.

TADCORA_4 and TADCORB_4 are initialized to H'FFFF.

Note: TADCORA_4 and TADCORB_4 must not be accessed in eight bits; they should always be accessed in 16 bits.

12.3.10 Timer A/D Converter Start Request Cycle Set Buffer Registers (TADCOBRA_4 and TADCOBRB_4)

TADCOBRA_4 and TADCOBRB_4 are 16-bit readable/writable registers. When the crest or trough of the TCNT_4 count is reached, these register values are transferred to TADCORA_4 and TADCORB_4, respectively.

TADCOBRA_4 and TADCOBRB_4 are initialized to H'FFFF.

Note: TADCOBRA_4 and TADCOBRB_4 must not be accessed in eight bits; they should always be accessed in 16 bits.

12.3.11 Timer Counter (TCNT)

The TCNT counters are 16-bit readable/writable counters. The MTU2 has five TCNT counters, one each for channels 0 to 4.

The TCNT counters must not be accessed in eight bits; they should always be accessed in 16 bits.

Note: The TCNT counters must not be accessed in eight bits; they should always be accessed in 16 bits.

12.3.12 Timer General Register (TGR)

The TGR registers are 16-bit readable/writable registers. The MTU2 has eighteen TGR registers, six for channel 0, two each for channels 1 and 2, four each for channels 3 and 4.

TGRA, TGRB, TGRC, and TGRD function as either output compare or input capture registers. TGRC and TGRD for channels 0, 3, and 4 can also be designated for operation as buffer registers. TGR buffer register combinations are TGRA and TGRC, and TGRB and TGRD.

TGRE_0 and TGRF_0 function as compare registers. When the TCNT_0 count matches the TGRE_0 value, an A/D converter start request can be issued. TGRF can also be designated for operation as a buffer register. TGR buffer register combination is TGRE and TGRF.

Note: The TGR registers must not be accessed in eight bits; they should always be accessed in 16 bits. TGR registers are initialized to H'FFFF.

12.3.13 Timer Start Register (TSTR)

TSTR is an 8-bit readable/writable register that selects operation/stoppage of TCNT for channels 0 to 4.

When setting the operating mode in TMDR or setting the count clock in TCR, first stop the TCNT counter.

12.3.14 Timer Synchronous Register (TSYR)

TSYR is an 8-bit readable/writable register that selects independent operation or synchronous operation for the channel 0 to 4 TCNT counters. A channel performs synchronous operation when the corresponding bit in TSYR is set to 1.

12.3.15 Timer Read/Write Enable Register (TRWER)

TRWER is an 8-bit readable/writable register that enables or disables access to the registers and counters which have write-protection capability against accidental modification in channels 3 and 4.

• Registers and counters having write-protection capability against accidental modification 22 registers: TCR_3, TCR_4, TMDR_3, TMDR_4, TIORH_3, TIORH_4, TIORL_3, TIORL_4, TIER_3, TIER_4, TGRA_3, TGRA_4, TGRB_3, TGRB_4, TOER, TOCR1, TOCR2, TGCR, TCDR, TDDR, TCNT_3, and TCNT4.

12.3.16 Timer Output Master Enable Register (TOER)

TOER is an 8-bit readable/writable register that enables/disables output settings for output pins TIOC4D, TIOC4C, TIOC3D, TIOC4B, TIOC4A, and TIOC3B. These pins do not output correctly if the TOER bits have not been set. Set TOER of CH3 and CH4 prior to setting TIOR of CH3 and CH4. Set TOER when count operation of TCNT channels 3 and 4 is halted.

Note: * The inactive level is determined by the settings in timer output control registers 1 and 2 (TOCR1 and TOCR2). For details, refer to section 12.3.17, Timer Output Control Register 1 (TOCR1), and section 12.3.18, Timer Output Control Register 2 (TOCR2). Set these bits to 1 to enable MTU2 output in other than complementary PWM or resetsynchronized PWM mode. When these bits are set to 0, low level is output.

12.3.17 Timer Output Control Register 1 (TOCR1)

TOCR1 is an 8-bit readable/writable register that enables/disables PWM synchronized toggle output in complementary PWM mode/reset synchronized PWM mode, and controls output level inversion of PWM output.

Note: $*$ This bit can be set to 1 only once after a power-on reset. After 1 is written, 0 cannot be written to the bit.

Notes: 1. Setting the TOCL bit to 1 prevents accidental modification when the CPU goes out of control.

- 2. Clearing the TOCS0 bit to 0 makes this bit setting valid.
- 3. After power-on reset, 1 can be written only once. After 1 has been written, 0 cannot be written.
- 4. If there is no dead time, the reverse phase output is the inversion of the forward phase. Set OLSP and OLSN to the same value.

Table 12.28 Output Level Select Function

Note: The reverse phase waveform initial output value changes to active level after elapse of the dead time after count start.

Table 12.29 Output Level Select Function

Figure 12.2 shows an example of complementary PWM mode output (1 phase) when $OLSN = 1$, $OLSP = 1.$

Figure 12.2 Complementary PWM Mode Output Level Example

12.3.18 Timer Output Control Register 2 (TOCR2)

TOCR2 is an 8-bit readable/writable register that controls output level inversion of PWM output in complementary PWM mode and reset-synchronized PWM mode.

Note: * Setting the TOCS bit in TOCR1 to 1 makes this bit setting valid. If there is no dead time, the reverse phase output is the inversion of the forward phase. Set OLSiP and OLSiN to the same value $(i = 1, 2, or 3)$.

Table 12.30 Setting of Bits BF1 and BF0

Table 12.31 TIOC4D Output Level Select Function

Note: The reverse phase waveform initial output value changes to the active level after elapse of the dead time after count start.

Table 12.32 TIOC4B Output Level Select Function

Table 12.33 TIOC4C Output Level Select Function

Note: The reverse phase waveform initial output value changes to the active level after elapse of the dead time after count start.

Table 12.34 TIOC4A Output Level Select Function

Table 12.35 TIOC3D Output Level Select Function

Note: The reverse phase waveform initial output value changes to the active level after elapse of the dead time after count start.

Table 12.36 TIOC4B Output Level Select Function

12.3.19 Timer Output Level Buffer Register (TOLBR)

TOLBR is an 8-bit readable/writable register that functions as a buffer for TOCR2 and specifies the PWM output level in complementary PWM mode and reset-synchronized PWM mode.

Figure 12.3 shows an example of the PWM output level setting procedure in buffer operation.

12.3.20 Timer Gate Control Register (TGCR)

TGCR is an 8-bit readable/writable register that controls the waveform output necessary for brushless DC motor control in reset-synchronized PWM mode/complementary PWM mode. These register settings are ineffective for anything other than complementary PWM mode/resetsynchronized PWM mode.

Table 12.37 Output level Select Function

12.3.21 Timer Subcounter (TCNTS)

TCNTS is a 16-bit read-only counter that is used only in complementary PWM mode.

The initial value of TCNTS is H'0000.

Note: Accessing the TCNTS in 8-bit units is prohibited. Always access in 16-bit units.

12.3.22 Timer Dead Time Data Register (TDDR)

TDDR is a 16-bit register, used only in complementary PWM mode that specifies the TCNT_3 and TCNT_4 counter offset values. In complementary PWM mode, when the TCNT_3 and TCNT_4 counters are cleared and then restarted, the TDDR register value is loaded into the TCNT_3 counter and the count operation starts.

The initial value of TDDR is H'FFFF.

Note: Accessing the TDDR in 8-bit units is prohibited. Always access in 16-bit units.

12.3.23 Timer Cycle Data Register (TCDR)

TCDR is a 16-bit register used only in complementary PWM mode. Set half the PWM carrier sync value (a value of two times TDDR + 3 or greater) as the TCDR register value. This register is constantly compared with the TCNTS counter in complementary PWM mode, and when a match occurs, the TCNTS counter switches direction (decrement to increment).

The initial value of TCDR is H'FFFF.

Note: Accessing the TCDR in 8-bit units is prohibited. Always access in 16-bit units.

12.3.24 Timer Cycle Buffer Register (TCBR)

TCBR is a 16-bit register used only in complementary PWM mode. It functions as a buffer register for the TCDR register. The TCBR register values are transferred to the TCDR register with the transfer timing set in the TMDR register.

Note: Accessing the TCBR in 8-bit units is prohibited. Always access in 16-bit units.

12.3.25 Timer Interrupt Skipping Set Register (TITCR)

TITCR is an 8-bit readable/writable register that enables or disables interrupt skipping and specifies the interrupt skipping count. The MTU2 has one TITCR.

Note: * When 0 is specified for the interrupt skipping count, no interrupt skipping will be performed. Before changing the interrupt skipping count, be sure to clear the T3AEN and T4VEN bits to 0 to clear the skipping counter (TICNT).

Table 12.38 Setting of Interrupt Skipping Count by Bits 3ACOR2 to 3ACOR0

Table 12.39 Setting of Interrupt Skipping Count by Bits 4VCOR2 to 4VCOR0

12.3.26 Timer Interrupt Skipping Counter (TITCNT)

TITCNT is an 8-bit readable/writable counter. The MTU2 has one TITCNT. TITCNT retains its value even after stopping the count operation of TCNT_3 and TCNT_4.

Note: To clear the TITCNT, clear the bits T3AEN and T4VEN in TITCR to 0.

12.3.27 Timer Buffer Transfer Set Register (TBTER)

TBTER is an 8-bit readable/writable register that enables or disables transfer from the buffer registers* used in complementary PWM mode to the temporary registers and specifies whether to link the transfer with interrupt skipping operation. The MTU2 has one TBTER.

Table 12.40 Setting of Bits BTE1 and BTE0

 2. When interrupt skipping is disabled (the T3AEN and T4VEN bits are cleared to 0 in the timer interrupt skipping set register (TITCR) or the skipping count set bits (3ACOR and 4VCOR) in TITCR are cleared to 0)), be sure to disable link of buffer transfer with interrupt skipping (clear the BTE1 bit in the timer buffer transfer set register (TBTER) to 0). If link with interrupt skipping is enabled while interrupt skipping is disabled, buffer transfer will not be performed.

12.3.28 Timer Dead Time Enable Register (TDER)

TDER is an 8-bit readable/writable register that controls dead time generation in complementary PWM mode. The MTU2 has one TDER in channel 3. TDER must be modified only while TCNT stops.

Note: $*$ TDDR must be set to 1 or a larger value.

12.3.29 Timer Waveform Control Register (TWCR)

TWCR is an 8-bit readable/writable register that controls the waveform when synchronous counter clearing occurs in TCNT_3 and TCNT_4 in complementary PWM mode and specifies whether to clear the counters at TGRA_3 compare match. The CCE bit and WRE bit in TWCR must be modified only while TCNT stops.

Note: $*$ Do not set to 1 when complementary PWM mode is not selected.

Note: $*$ Do not set to 1 when complementary PWM mode is not selected.

12.3.30 Bus Master Interface

The timer counters (TCNT), general registers (TGR), timer subcounter (TCNTS), timer cycle buffer register (TCBR), timer dead time data register (TDDR), timer cycle data register (TCDR), timer A/D converter start request control register (TADCR), timer A/D converter start request cycle set registers (TADCOR), and timer A/D converter start request cycle set buffer registers (TADCOBR) are 16-bit registers. A 16-bit data bus to the bus master enables 16-bit read/writes. 8 bit read/write is not possible. Always access in 16-bit units.

All registers other than the above registers are 8-bit registers. These are connected to the CPU by a 16-bit data bus, so 16-bit read/writes and 8-bit read/writes are both possible.

12.4 Operation

12.4.1 Basic Functions

Each channel has a TCNT and TGR register. TCNT performs up-counting, and is also capable of free-running operation, cycle counting, and external event counting.

Each TGR can be used as an input capture register or output compare register.

Always select MTU2 external pins set function using the pin function controller (PFC).

(1) Counter Operation

When one of bits CST0 to CST4 in TSTR is set to 1, the TCNT counter for the corresponding channel begins counting. TCNT can operate as a free-running counter, periodic counter, for example.

(a) Example of Count Operation Setting Procedure

Figure 12.4 shows an example of the count operation setting procedure.

Figure 12.4 Example of Counter Operation Setting Procedure

(b) Free-Running Count Operation and Periodic Count Operation:

Immediately after a reset, the MTU2's TCNT counters are all designated as free-running counters. When the relevant bit in TSTR is set to 1 the corresponding TCNT counter starts up-count operation as a free-running counter. When TCNT overflows (from H'FFFF to H'0000), the TCFV bit in TSR is set to 1. If the value of the corresponding TCIEV bit in TIER is 1 at this point, the MTU2 requests an interrupt. After overflow, TCNT starts counting up again from H'0000.

Figure 12.5 illustrates free-running counter operation.

When compare match is selected as the TCNT clearing source, the TCNT counter for the relevant channel performs periodic count operation. The TGR register for setting the period is designated as an output compare register, and counter clearing by compare match is selected by means of bits CCLR0 to CCLR2 in TCR. After the settings have been made, TCNT starts up-count operation as a periodic counter when the corresponding bit in TSTR is set to 1. When the count value matches the value in TGR, the TGF bit in TSR is set to 1 and TCNT is cleared to H'0000.

If the value of the corresponding TGIE bit in TIER is 1 at this point, the MTU2 requests an interrupt. After a compare match, TCNT starts counting up again from H'0000.

Figure 12.6 illustrates periodic counter operation.

(2) Waveform Output by Compare Match

The MTU2 can perform 0, 1, or toggle output from the corresponding output pin using compare match.

(a) Example of Setting Procedure for Waveform Output by Compare Match

Figure 12.7 shows an example of the setting procedure for waveform output by compare match

Figure 12.7 Example of Setting Procedure for Waveform Output by Compare Match

(b) Examples of Waveform Output Operation:

Figure 12.8 shows an example of 0 output/1 output.

In this example TCNT has been designated as a free-running counter, and settings have been made such that 1 is output by compare match A, and 0 is output by compare match B. When the set level and the pin level coincide, the pin level does not change.

Figure 12.8 Example of 0 Output/1 Output Operation

Figure 12.9 shows an example of toggle output.

In this example, TCNT has been designated as a periodic counter (with counter clearing on compare match B), and settings have been made such that the output is toggled by both compare match A and compare match B.

Figure 12.9 Example of Toggle Output Operation

(3) Input Capture Function

The TCNT value can be transferred to TGR on detection of the TIOC pin input edge.

Rising edge, falling edge, or both edges can be selected as the detected edge. For channels 0 and 1, it is also possible to specify another channel's counter input clock or compare match signal as the input capture source.

Note: When another channel's counter input clock is used as the input capture input for channels 0 and 1, Pφ/1 should not be selected as the counter input clock used for input capture input. Input capture will not be generated if P $\phi/1$ is selected.

(a) Example of Input Capture Operation Setting Procedure

Figure 12.10 shows an example of the input capture operation setting procedure.

Figure 12.10 Example of Input Capture Operation Setting Procedure

(b) Example of Input Capture Operation

Figure 12.11 shows an example of input capture operation.

In this example both rising and falling edges have been selected as the TIOCA pin input capture input edge, the falling edge has been selected as the TIOCB pin input capture input edge, and counter clearing by TGRB input capture has been designated for TCNT.

Figure 12.11 Example of Input Capture Operation

12.4.2 Synchronous Operation

In synchronous operation, the values in a number of TCNT counters can be rewritten simultaneously (synchronous presetting). Also, a number of TCNT counters can be cleared simultaneously by making the appropriate setting in TCR (synchronous clearing).

Synchronous operation enables TGR to be incremented with respect to a single time base.

Channels 0 to 4 can all be designated for synchronous operation.

(1) Example of Synchronous Operation Setting Procedure

Figure 12.12 shows an example of the synchronous operation setting procedure.

[5] Set to 1 the CST bits in TSTR for the relevant channels, to start the count operation.

Figure 12.12 Example of Synchronous Operation Setting Procedure

(2) Example of Synchronous Operation

Figure 12.13 shows an example of synchronous operation.

In this example, synchronous operation and PWM mode 1 have been designated for channels 0 to 2, TGRB_0 compare match has been set as the channel 0 counter clearing source, and synchronous clearing has been set for the channel 1 and 2 counter clearing source.

Three-phase PWM waveforms are output from pins TIOC0A, TIOC1A, and TIOC2A. At this time, synchronous presetting, and synchronous clearing by TGRB_0 compare match, are performed for channel 0 to 2 TCNT counters, and the data set in TGRB_0 is used as the PWM cycle.

For details of PWM modes, see section 12.4.5, PWM Modes.

Figure 12.13 Example of Synchronous Operation

12.4.3 Buffer Operation

Buffer operation, provided for channels 0, 3, and 4, enables TGRC and TGRD to be used as buffer registers. In channel 0, TGRF can also be used as a buffer register.

Buffer operation differs depending on whether TGR has been designated as an input capture register or as a compare match register.

Note: TGRE 0 cannot be designated as an input capture register and can only operate as a compare match register.

Table 12.41 shows the register combinations used in buffer operation.

Table 12.41 Register Combinations in Buffer Operation

Channel	Timer General Register	Buffer Register
$\mathbf 0$	TGRA 0	TGRC_0
	TGRB_0	TGRD_0
	TGRE 0	TGRF_0
3	TGRA ₃	TGRC_3
	TGRB_3	TGRD 3
4	TGRA_4	TGRC_4
	TGRB_4	TGRD_4

When TGR is an output compare register

When a compare match occurs, the value in the buffer register for the corresponding channel is transferred to the timer general register.

This operation is illustrated in figure 12.14.

Figure 12.14 Compare Match Buffer Operation

• When TGR is an input capture register

When input capture occurs, the value in TCNT is transferred to TGR and the value previously held in the timer general register is transferred to the buffer register.

This operation is illustrated in figure 12.15.

Figure 12.15 Input Capture Buffer Operation

(1) Example of Buffer Operation Setting Procedure

Figure 12.16 shows an example of the buffer operation setting procedure.

Figure 12.16 Example of Buffer Operation Setting Procedure

(2) Examples of Buffer Operation

(a) When TGR is an output compare register

Figure 12.17 shows an operation example in which PWM mode 1 has been designated for channel 0, and buffer operation has been designated for TGRA and TGRC. The settings used in this example are TCNT clearing by compare match B, 1 output at compare match A, and 0 output at compare match B. In this example, the TTSA bit in TBTM is cleared to 0.

As buffer operation has been set, when compare match A occurs the output changes and the value in buffer register TGRC is simultaneously transferred to timer general register TGRA. This operation is repeated each time that compare match A occurs.

For details of PWM modes, see section 12.4.5, PWM Modes.

Figure 12.17 Example of Buffer Operation (1)

(b) When TGR is an input capture register

Figure 12.18 shows an operation example in which TGRA has been designated as an input capture register, and buffer operation has been designated for TGRA and TGRC.

Counter clearing by TGRA input capture has been set for TCNT, and both rising and falling edges have been selected as the TIOCA pin input capture input edge.

As buffer operation has been set, when the TCNT value is stored in TGRA upon the occurrence of input capture A, the value previously stored in TGRA is simultaneously transferred to TGRC.

Figure 12.18 Example of Buffer Operation (2)

(3) Selecting Timing for Transfer from Buffer Registers to Timer General Registers in Buffer Operation

The timing for transfer from buffer registers to timer general registers can be selected in PWM mode 1 or 2 for channel 0 or in PWM mode 1 for channels 3 and 4 by setting the buffer operation transfer mode registers (TBTM_0, TBTM_3, and TBTM_4). Either compare match (initial setting) or TCNT clearing can be selected for the transfer timing. TCNT clearing as transfer timing is one of the following cases.

- When TCNT overflows (H'FFFF to H'0000)
- When H'0000 is written to TCNT during counting
- When TCNT is cleared to H'0000 under the condition specified in the CCLR2 to CCLR0 bits in TCR

Note: TBTM must be modified only while TCNT stops.

Figure 12.19 shows an operation example in which PWM mode 1 is designated for channel 0 and buffer operation is designated for TGRA_0 and TGRC_0. The settings used in this example are TCNT_0 clearing by compare match B, 1 output at compare match A, and 0 output at compare match B. The TTSA bit in TBTM 0 is set to 1.

Figure 12.19 Example of Buffer Operation When TCNT_0 Clearing is Selected for TGRC_0 to TGRA_0 Transfer Timing

12.4.4 Cascaded Operation

In cascaded operation, two 16-bit counters for different channels are used together as a 32-bit counter.

This function works by counting the channel 1 counter clock upon overflow/underflow of TCNT_2 as set in bits TPSC0 to TPSC2 in TCR.

Underflow occurs only when the lower 16-bit TCNT is in phase-counting mode.

Table 12.42 shows the register combinations used in cascaded operation.

Note: When phase counting mode is set for channel 1, the counter clock setting is invalid and the counters operates independently in phase counting mode.

Table 12.42 Cascaded Combinations

For simultaneous input capture of TCNT 1 and TCNT 2 during cascaded operation, additional input capture input pins can be specified by the input capture control register (TICCR). The edge detection that is the condition for input capture uses a signal representing the logical OR of the original input pin and the added input pins. For details, see (4) Cascaded Operation Example (c).

For input capture in cascade connection, refer to section 12.7.22, Simultaneous Capture of TCNT 1 and TCNT 2 in Cascade Connection.

Table 12.43 shows the TICCR setting and input capture input pins.

(1) Example of Cascaded Operation Setting Procedure

Figure 12.20 shows an example of the setting procedure for cascaded operation.

(2) Cascaded Operation Example (a)

Figure 12.21 illustrates the operation when TCNT_2 overflow/underflow counting has been set for TCNT 1 and phase counting mode has been designated for channel 2.

TCNT_1 is incremented by TCNT_2 overflow and decremented by TCNT_2 underflow.

Figure 12.21 Cascaded Operation Example (a)

(3) Cascaded Operation Example (b)

Figure 12.22 illustrates the operation when TCNT 1 and TCNT 2 have been cascaded and the I2AE bit in TICCR has been set to 1 to include the TIOC2A pin in the TGRA_1 input capture conditions. In this example, the IOA0 to IOA3 bits in TIOR_1 have selected the TIOC1A rising edge for the input capture timing while the IOA0 to IOA3 bits in TIOR_2 have selected the TIOC2A rising edge for the input capture timing.

Under these conditions, the rising edge of both TIOC1A and TIOC2A is used for the TGRA_1 input capture condition. For the TGRA_2 input capture condition, the TIOC2A rising edge is used.

Figure 12.22 Cascaded Operation Example (b)

(4) Cascaded Operation Example (c)

Figure 12.23 illustrates the operation when TCNT_1 and TCNT_2 have been cascaded and the I2AE and I1AE bits in TICCR have been set to 1 to include the TIOC2A and TIOC1A pins in the TGRA_1 and TGRA_2 input capture conditions, respectively. In this example, the IOA0 to IOA3 bits in both TIOR_1 and TIOR_2 have selected both the rising and falling edges for the input capture timing. Under these conditions, the ORed result of TIOC1A and TIOC2A input is used for the TGRA_1 and TGRA_2 input capture conditions.

Figure 12.23 Cascaded Operation Example (c)

(5) Cascaded Operation Example (d)

Figure 12.24 illustrates the operation when TCNT 1 and TCNT 2 have been cascaded and the I2AE bit in TICCR has been set to 1 to include the TIOC2A pin in the TGRA_1 input capture conditions. In this example, the IOA0 to IOA3 bits in TIOR_1 have selected TGRA_0 compare match or input capture occurrence for the input capture timing while the IOA0 to IOA3 bits in TIOR_2 have selected the TIOC2A rising edge for the input capture timing.

Under these conditions, as TIOR 1 has selected TGRA 0 compare match or input capture occurrence for the input capture timing, the TIOC2A edge is not used for TGRA_1 input capture condition although the I2AE bit in TICCR has been set to 1.

Figure 12.24 Cascaded Operation Example (d)

12.4.5 PWM Modes

In PWM mode, PWM waveforms are output from the output pins. The output level can be selected as 0, 1, or toggle output in response to a compare match of each TGR.

TGR registers settings can be used to output a PWM waveform in the range of 0% to 100% duty.

Designating TGR compare match as the counter clearing source enables the period to be set in that register. All channels can be designated for PWM mode independently. Synchronous operation is also possible.

There are two PWM modes, as described below.

PWM mode 1

PWM output is generated from the TIOCA and TIOCC pins by pairing TGRA with TGRB and TGRC with TGRD. The output specified by bits IOA0 to IOA3 and IOC0 to IOC3 in TIOR is output from the TIOCA and TIOCC pins at compare matches A and C, and the output specified by bits IOB0 to IOB3 and IOD0 to IOD3 in TIOR is output at compare matches B and D. The initial output value is the value set in TGRA or TGRC. If the set values of paired TGRs are identical, the output value does not change when a compare match occurs. In PWM mode 1, a maximum 8-phase PWM output is possible.

• PWM mode 2

PWM output is generated using one TGR as the cycle register and the others as duty registers. The output specified in TIOR is performed by means of compare matches. Upon counter clearing by a cycle register compare match, the output value of each pin is the initial value set in TIOR. If the set values of the cycle and duty registers are identical, the output value does not change when a compare match occurs.

In PWM mode 2, a maximum 8-phase PWM output is possible in combination use with synchronous operation.

The correspondence between PWM output pins and registers is shown in table 12.44.

Table 12.44 PWM Output Registers and Output Pins

Note: In PWM mode 2, PWM output is not possible for the TGR register in which the period is set.

(1) Example of PWM Mode Setting Procedure

Figure 12.25 shows an example of the PWM mode setting procedure.

Figure 12.25 Example of PWM Mode Setting Procedure

(2) Examples of PWM Mode Operation

Figure 12.26 shows an example of PWM mode 1 operation.

In this example, TGRA compare match is set as the TCNT clearing source, 0 is set for the TGRA initial output value and output value, and 1 is set as the TGRB output value.

In this case, the value set in TGRA is used as the period, and the values set in the TGRB registers are used as the duty levels.

Figure 12.26 Example of PWM Mode Operation (1)

Figure 12.27 shows an example of PWM mode 2 operation.

In this example, synchronous operation is designated for channels 0 and 1, TGRB_1 compare match is set as the TCNT clearing source, and 0 is set for the initial output value and 1 for the output value of the other TGR registers (TGRA_0 to TGRD_0, TGRA_1), outputting a 5-phase PWM waveform.

In this case, the value set in TGRB 1 is used as the cycle, and the values set in the other TGRs are used as the duty levels.

Figure 12.27 Example of PWM Mode Operation (2)

Figure 12.28 shows examples of PWM waveform output with 0% duty and 100% duty in PWM mode.

Figure 12.28 Example of PWM Mode Operation (3)

12.4.6 Phase Counting Mode

In phase counting mode, the phase difference between two external clock inputs is detected and TCNT is incremented/decremented accordingly. This mode can be set for channels 1 and 2.

When phase counting mode is set, an external clock is selected as the counter input clock and TCNT operates as an up/down-counter regardless of the setting of bits TPSC0 to TPSC2 and bits CKEG0 and CKEG1 in TCR. However, the functions of bits CCLR0 and CCLR1 in TCR, and of TIOR, TIER, and TGR, are valid, and input capture/compare match and interrupt functions can be used.

This can be used for two-phase encoder pulse input.

If overflow occurs when TCNT is counting up, the TCFV flag in TSR is set; if underflow occurs when TCNT is counting down, the TCFU flag is set.

The TCFD bit in TSR is the count direction flag. Reading the TCFD flag reveals whether TCNT is counting up or down.

Table 12.45 shows the correspondence between external clock pins and channels.

Table 12.45 Phase Counting Mode Clock Input Pins

	External Clock Pins	
Channels	A-Phase	B-Phase
When channel 1 is set to phase counting mode	TCLKA	TCLKB
When channel 2 is set to phase counting mode	TCLKC	TCLKD

(1) Example of Phase Counting Mode Setting Procedure

Figure 12.29 shows an example of the phase counting mode setting procedure.

(2) Examples of Phase Counting Mode Operation

In phase counting mode, TCNT counts up or down according to the phase difference between two external clocks. There are four modes, according to the count conditions.

(a) Phase counting mode 1

Figure 12.30 shows an example of phase counting mode 1 operation, and table 12.46 summarizes the TCNT up/down-count conditions.

Figure 12.30 Example of Phase Counting Mode 1 Operation

Table 12.46 Up/Down-Count Conditions in Phase Counting Mode 1

 $\overline{\mathbf{r}}$: Falling edge

(b) Phase counting mode 2

Figure 12.31 shows an example of phase counting mode 2 operation, and table 12.47 summarizes the TCNT up/down-count conditions.

Figure 12.31 Example of Phase Counting Mode 2 Operation

[Legend]

Rising edge

: Falling edge

(c) Phase counting mode 3

Figure 12.32 shows an example of phase counting mode 3 operation, and table 12.48 summarizes the TCNT up/down-count conditions.

[Legend]

 \overline{F} : Rising edge

: Falling edge

(d) Phase counting mode 4

Figure 12.33 shows an example of phase counting mode 4 operation, and table 12.49 summarizes the TCNT up/down-count conditions.

Figure 12.33 Example of Phase Counting Mode 4 Operation

[Legend]

 \overline{f} : Rising edge

 \mathbf{L} : Falling edge

(3) Phase Counting Mode Application Example

Figure 12.34 shows an example in which channel 1 is in phase counting mode, and channel 1 is coupled with channel 0 to input servo motor 2-phase encoder pulses in order to detect position or speed.

Channel 1 is set to phase counting mode 1, and the encoder pulse A-phase and B-phase are input to TCLKA and TCLKB.

Channel 0 operates with TCNT counter clearing by TGRC_0 compare match; TGRA_0 and TGRC_0 are used for the compare match function and are set with the speed control period and position control period. TGRB_0 is used for input capture, with TGRB_0 and TGRD_0 operating in buffer mode. The channel 1 counter input clock is designated as the TGRB_0 input capture source, and the pulse widths of 2-phase encoder 4-multiplication pulses are detected.

TGRA_1 and TGRB_1 for channel 1 are designated for input capture, and channel 0 TGRA_0 and TGRC_0 compare matches are selected as the input capture source and store the up/down-counter values for the control periods.

This procedure enables the accurate detection of position and speed.

Figure 12.34 Phase Counting Mode Application Example

12.4.7 Reset-Synchronized PWM Mode

In the reset-synchronized PWM mode, three-phase output of positive and negative PWM waveforms that share a common wave transition point can be obtained by combining channels 3 and 4.

When set for reset-synchronized PWM mode, the TIOC3B, TIOC3D, TIOC4A, TIOC4C, TIOC4B, and TIOC4D pins function as PWM output pins and TCNT3 functions as an upcounter.

Table 12.50 shows the PWM output pins used. Table 12.51 shows the settings of the registers.

Channel	Output Pin	Description
3	TIOC3B	PWM output pin 1
	TIOC3D	PWM output pin 1' (negative-phase waveform of PWM output 1)
4	TIOC4A	PWM output pin 2
	TIOC4C	PWM output pin 2' (negative-phase waveform of PWM output 2)
	TIOC4B	PWM output pin 3
	TIOC4D	PWM output pin 3' (negative-phase waveform of PWM output 3)

Table 12.50 Output Pins for Reset-Synchronized PWM Mode

(1) Procedure for Selecting the Reset-Synchronized PWM Mode

Figure 12.35 shows an example of procedure for selecting the reset synchronized PWM mode.

Figure 12.35 Procedure for Selecting Reset-Synchronized PWM Mode

(2) Reset-Synchronized PWM Mode Operation

Figure 12.36 shows an example of operation in the reset-synchronized PWM mode. TCNT_3 and TCNT_4 operate as upcounters. The counter is cleared when a TCNT_3 and TGRA_3 comparematch occurs, and then begins incrementing from H'0000. The PWM output pin output toggles with each occurrence of a TGRB_3, TGRA_4, TGRB_4 compare-match, and upon counter clears.

Figure 12.36 Reset-Synchronized PWM Mode Operation Example (When TOCR's $OLSN = 1$ and $OLSP = 1$)

12.4.8 Complementary PWM Mode

In the complementary PWM mode, three-phase output of non-overlapping positive and negative PWM waveforms can be obtained by combining channels 3 and 4. PWM waveforms without nonoverlapping interval are also available.

In complementary PWM mode, TIOC3B, TIOC3D, TIOC4A, TIOC4B, TIOC4C, and TIOC4D pins function as PWM output pins, the TIOC3A pin can be set for toggle output synchronized with the PWM period. TCNT_3 and TCNT_4 function as up/down counters.

Table 12.52 shows the PWM output pins used. Table 12.53 shows the settings of the registers used.

Table 12.52 Output Pins for Complementary PWM Mode

Note: * Avoid setting the TIOC3C pin as a timer I/O pin in the complementary PWM mode.

(timer read/write enable register).

Figure 12.37 Block Diagram of Channels 3 and 4 in Complementary PWM Mode

(1) Example of Complementary PWM Mode Setting Procedure

An example of the complementary PWM mode setting procedure is shown in figure 12.38.

Figure 12.38 Example of Complementary PWM Mode Setting Procedure

(2) Outline of Complementary PWM Mode Operation

In complementary PWM mode, 6-phase PWM output is possible. Figure 12.39 illustrates counter operation in complementary PWM mode, and figure 12.40 shows an example of complementary PWM mode operation.

(a) Counter Operation

In complementary PWM mode, three counters—TCNT_3, TCNT_4, and TCNTS—perform up/down-count operations.

TCNT_3 is automatically initialized to the value set in TDDR when complementary PWM mode is selected and the CST bit in TSTR is 0.

When the CST bit is set to 1, TCNT 3 counts up to the value set in TGRA 3, then switches to down-counting when it matches TGRA_3. When the TCNT3 value matches TDDR, the counter switches to up-counting, and the operation is repeated in this way.

TCNT_4 is initialized to H'0000.

When the CST bit is set to 1, TCNT4 counts up in synchronization with TCNT 3, and switches to down-counting when it matches TCDR. On reaching H'0000, TCNT4 switches to up-counting, and the operation is repeated in this way.

TCNTS is a read-only counter. It need not be initialized.

When TCNT_3 matches TCDR during TCNT_3 and TCNT_4 up/down-counting, down-counting is started, and when TCNTS matches TCDR, the operation switches to up-counting. When TCNTS matches TGRA_3, it is cleared to H'0000.

When TCNT₋₄ matches TDDR during TCNT₋₃ and TCNT₋₄ down-counting, up-counting is started, and when TCNTS matches TDDR, the operation switches to down-counting. When TCNTS reaches H'0000, it is set with the value in TGRA_3.

TCNTS is compared with the compare register and temporary register in which the PWM duty is set during the count operation only.

Figure 12.39 Complementary PWM Mode Counter Operation

(b) Register Operation

In complementary PWM mode, nine registers are used, comprising compare registers, buffer registers, and temporary registers. Figure 12.40 shows an example of complementary PWM mode operation.

The registers which are constantly compared with the counters to perform PWM output are TGRB_3, TGRA_4, and TGRB_4. When these registers match the counter, the value set in bits OLSN and OLSP in the timer output control register (TOCR) is output.

The buffer registers for these compare registers are TGRD_3, TGRC_4, and TGRD_4.

Between a buffer register and compare register there is a temporary register. The temporary registers cannot be accessed by the CPU.

Data in a compare register is changed by writing the new data to the corresponding buffer register. The buffer registers can be read or written at any time.

The data written to a buffer register is constantly transferred to the temporary register in the Ta interval. Data is not transferred to the temporary register in the Tb interval. Data written to a buffer register in this interval is transferred to the temporary register at the end of the Tb interval.

The value transferred to a temporary register is transferred to the compare register when TCNTS for which the Tb interval ends matches TGRA_3 when counting up, or H'0000 when counting down. The timing for transfer from the temporary register to the compare register can be selected with bits MD3 to MD0 in the timer mode register (TMDR). Figure 12.40 shows an example in which the mode is selected in which the change is made in the trough.

In the tb interval (tb1 in figure 12.40) in which data transfer to the temporary register is not performed, the temporary register has the same function as the compare register, and is compared

with the counter. In this interval, therefore, there are two compare match registers for one-phase output, with the compare register containing the pre-change data, and the temporary register containing the new data. In this interval, the three counters—TCNT_3, TCNT_4, and TCNTS and two registers—compare register and temporary register—are compared, and PWM output controlled accordingly.

Figure 12.40 Example of Complementary PWM Mode Operation

(c) Initialization

In complementary PWM mode, there are six registers that must be initialized. In addition, there is a register that specifies whether to generate dead time (it should be used only when dead time generation should be disabled).

Before setting complementary PWM mode with bits MD3 to MD0 in the timer mode register (TMDR), the following initial register values must be set.

TGRC_3 operates as the buffer register for TGRA_3, and should be set with 1/2 the PWM carrier cycle + dead time Td. The timer cycle buffer register (TCBR) operates as the buffer register for the timer cycle data register (TCDR), and should be set with 1/2 the PWM carrier cycle. Set dead time Td in the timer dead time data register (TDDR).

When dead time is not needed, the TDER bit in the timer dead time enable register (TDER) should be cleared to 0, TGRC_3 and TGRA_3 should be set to 1/2 the PWM carrier cycle + 1, and TDDR should be set to 1.

Set the respective initial PWM duty values in buffer registers TGRD_3, TGRC_4, and TGRD_4.

The values set in the five buffer registers excluding TDDR are transferred simultaneously to the corresponding compare registers when complementary PWM mode is set.

Set TCNT_4 to H'0000 before setting complementary PWM mode.

Note: The TGRC_3 set value must be the sum of 1/2 the PWM carrier cycle set in TCBR and dead time Td set in TDDR. When dead time generation is disabled by TDER, TGRC_3 must be set to $1/2$ the PWM carrier cycle $+1$.
(d) PWM Output Level Setting

In complementary PWM mode, the PWM pulse output level is set with bits OLSN and OLSP in timer output control register 1 (TOCR1) or bits OLS1P to OLS3P and OLS1N to OLS3N in timer output control register 2 (TOCR2).

The output level can be set for each of the three positive phases and three negative phases of 6 phase output.

Complementary PWM mode should be cleared before setting or changing output levels.

(e) Dead Time Setting

In complementary PWM mode, PWM pulses are output with a non-overlapping relationship between the positive and negative phases. This non-overlap time is called the dead time.

The non-overlap time is set in the timer dead time data register (TDDR). The value set in TDDR is used as the TCNT_3 counter start value, and creates non-overlap between TCNT_3 and TCNT_4. Complementary PWM mode should be cleared before changing the contents of TDDR.

(f) Dead Time Suppressing

Dead time generation is suppressed by clearing the TDER bit in the timer dead time enable register (TDER) to 0. TDER can be cleared to 0 only when 0 is written to it after reading TDER $=$ 1.

TGRA_3 and TGRC_3 should be set to 1/2 PWM carrier cycle + 1 and the timer dead time data register (TDDR) should be set to 1.

By the above settings, PWM waveforms without dead time can be obtained. Figure 12.41 shows an example of operation without dead time.

Figure 12.41 Example of Operation without Dead Time

(g) PWM Cycle Setting

In complementary PWM mode, the PWM pulse cycle is set in two registers—TGRA_3, in which the TCNT_3 upper limit value is set, and TCDR, in which the TCNT_4 upper limit value is set. The settings should be made so as to achieve the following relationship between these two registers:

```
With dead time: TGRA_3 set value = TCDR set value + TDDR set value
                  TCDR set value > two times TDDR + 2 
Without dead time: TGRA 3 set value = TCDR set value +1TCDR set value > 4
```
The TGRA_3 and TCDR settings are made by setting the values in buffer registers TGRC_3 and TCBR. The values set in TGRC_3 and TCBR are transferred simultaneously to TGRA_3 and TCDR in accordance with the transfer timing selected with bits MD3 to MD0 in the timer mode register (TMDR).

The updated PWM cycle is reflected from the next cycle when the data update is performed at the crest, and from the current cycle when performed in the trough. Figure 12.42 illustrates the operation when the PWM cycle is updated at the crest.

See the following section, Register Data Updating, for the method of updating the data in each buffer register.

Figure 12.42 Example of PWM Cycle Updating

(h) Register Data Updating

In complementary PWM mode, the buffer register is used to update the data in a compare register. The update data can be written to the buffer register at any time. There are five PWM duty and carrier cycle registers that have buffer registers and can be updated during operation.

There is a temporary register between each of these registers and its buffer register. When subcounter TCNTS is not counting, if buffer register data is updated, the temporary register value is also rewritten. Transfer is not performed from buffer registers to temporary registers when TCNTS is counting; in this case, the value written to a buffer register is transferred after TCNTS halts.

The temporary register value is transferred to the compare register at the data update timing set with bits MD3 to MD0 in the timer mode register (TMDR). Figure 12.43 shows an example of data updating in complementary PWM mode. This example shows the mode in which data updating is performed at both the counter crest and trough.

When rewriting buffer register data, a write to TGRD 4 must be performed at the end of the update. Data transfer from the buffer registers to the temporary registers is performed simultaneously for all five registers after the write to TGRD 4.

A write to TGRD_4 must be performed after writing data to the registers to be updated, even when not updating all five registers, or when updating the TGRD_4 data. In this case, the data written to TGRD_4 should be the same as the data prior to the write operation.

Figure 12.43 Example of Data Update in Complementary PWM Mode

(i) Initial Output in Complementary PWM Mode

In complementary PWM mode, the initial output is determined by the setting of bits OLSN and OLSP in timer output control register 1 (TOCR1) or bits OLS1N to OLS3N and OLS1P to OLS3P in timer output control register 2 (TOCR2).

This initial output is the PWM pulse non-active level, and is output from when complementary PWM mode is set with the timer mode register (TMDR) until TCNT_4 exceeds the value set in the dead time register (TDDR). Figure 12.44 shows an example of the initial output in complementary PWM mode.

An example of the waveform when the initial PWM duty value is smaller than the TDDR value is shown in figure 12.45.

Figure 12.44 Example of Initial Output in Complementary PWM Mode (1)

Figure 12.45 Example of Initial Output in Complementary PWM Mode (2)

(j) Complementary PWM Mode PWM Output Generation Method

In complementary PWM mode, 3-phase output is performed of PWM waveforms with a nonoverlap time between the positive and negative phases. This non-overlap time is called the dead time.

A PWM waveform is generated by output of the output level selected in the timer output control register in the event of a compare-match between a counter and compare register. While TCNTS is counting, compare register and temporary register values are simultaneously compared to create consecutive PWM pulses from 0 to 100%. The relative timing of on and off compare-match occurrence may vary, but the compare-match that turns off each phase takes precedence to secure the dead time and ensure that the positive phase and negative phase on times do not overlap. Figures 12.46 to 12.48 show examples of waveform generation in complementary PWM mode.

The positive phase/negative phase off timing is generated by a compare-match with the solid-line counter, and the on timing by a compare-match with the dotted-line counter operating with a delay of the dead time behind the solid-line counter. In the T1 period, compare-match **a** that turns off the negative phase has the highest priority, and compare-matches occurring prior to **a** are ignored. In the T2 period, compare-match **c** that turns off the positive phase has the highest priority, and compare-matches occurring prior to **c** are ignored.

In normal cases, compare-matches occur in the order $\mathbf{a} \to \mathbf{b} \to \mathbf{c} \to \mathbf{d}$ (or $\mathbf{c} \to \mathbf{d} \to \mathbf{a}' \to \mathbf{b}'$), as shown in figure 12.46.

If compare-matches deviate from the $\mathbf{a} \to \mathbf{b} \to \mathbf{c} \to \mathbf{d}$ order, since the time for which the negative phase is off is less than twice the dead time, the figure shows the positive phase is not being turned on. If compare-matches deviate from the $c \rightarrow d \rightarrow a' \rightarrow b'$ order, since the time for which the positive phase is off is less than twice the dead time, the figure shows the negative phase is not being turned on.

If compare-match **c** occurs first following compare-match **a**, as shown in figure 12.47, comparematch **b** is ignored, and the negative phase is turned off by compare-match **d**. This is because turning off of the positive phase has priority due to the occurrence of compare-match **c** (positive phase off timing) before compare-match **b** (positive phase on timing) (consequently, the waveform does not change since the positive phase goes from off to off).

Similarly, in the example in figure 12.48, compare-match **a'** with the new data in the temporary register occurs before compare-match **c**, but other compare-matches occurring up to **c**, which turns off the positive phase, are ignored. As a result, the negative phase is not turned on.

Thus, in complementary PWM mode, compare-matches at turn-off timings take precedence, and turn-on timing compare-matches that occur before a turn-off timing compare-match are ignored.

Figure 12.46 Example of Complementary PWM Mode Waveform Output (1)

Figure 12.47 Example of Complementary PWM Mode Waveform Output (2)

Figure 12.48 Example of Complementary PWM Mode Waveform Output (3)

Figure 12.49 Example of Complementary PWM Mode 0% and 100% Waveform Output (1)

Figure 12.50 Example of Complementary PWM Mode 0% and 100% Waveform Output (2)

Figure 12.51 Example of Complementary PWM Mode 0% and 100% Waveform Output (3)

Figure 12.52 Example of Complementary PWM Mode 0% and 100% Waveform Output (4)

Figure 12.53 Example of Complementary PWM Mode 0% and 100% Waveform Output (5)

(k) Complementary PWM Mode 0% and 100% Duty Output

In complementary PWM mode, 0% and 100% duty cycles can be output as required. Figures 12.49 to 12.53 show output examples.

100% duty output is performed when the compare register value is set to H'0000. The waveform in this case has a positive phase with a 100% on-state. 0% duty output is performed when the compare register value is set to the same value as TGRA_3. The waveform in this case has a positive phase with a 100% off-state.

On and off compare-matches occur simultaneously, but if a turn-on compare-match and turn-off compare-match for the same phase occur simultaneously, both compare-matches are ignored and the waveform does not change.

(l) Toggle Output Synchronized with PWM Cycle

In complementary PWM mode, toggle output can be performed in synchronization with the PWM carrier cycle by setting the PSYE bit to 1 in the timer output control register (TOCR). An example of a toggle output waveform is shown in figure 12.54.

This output is toggled by a compare-match between TCNT_3 and TGRA_3 and a compare-match between TCNT4 and H'0000.

The output pin for this toggle output is the TIOC3A pin. The initial output is 1.

(m) Counter Clearing by Another Channel

In complementary PWM mode, by setting a mode for synchronization with another channel by means of the timer synchronous register (TSYR), and selecting synchronous clearing with bits CCLR2 to CCLR0 in the timer control register (TCR), it is possible to have TCNT_3, TCNT_4, and TCNTS cleared by another channel.

Figure 12.55 illustrates the operation.

Use of this function enables counter clearing and restarting to be performed by means of an external signal.

Figure 12.55 Counter Clearing Synchronized with Another Channel

(n) Output Waveform Control at Synchronous Counter Clearing in Complementary PWM Mode

Setting the WRE bit in TWCR to 1 suppresses initial output when synchronous counter clearing occurs in the Tb interval at the trough in complementary PWM mode and controls abrupt change in duty cycle at synchronous counter clearing.

Initial output suppression is applicable only when synchronous clearing occurs in the Tb interval at the trough as indicated by (10) or (11) in figure 12.56. When synchronous clearing occurs outside that interval, the initial value specified by the OLS bits in TOCR is output. Even in the Tb interval at the trough, if synchronous clearing occurs in the initial value output period (indicated by (1) in figure 12.56) immediately after the counters start operation, initial value output is not suppressed.

When using the initial output suppression function, make sure to set compare registers TGRB 3, TGRA_4, and TGRB_4 to a value twice or more the setting of dead time data register TDDR. If synchronous clearing occurs with the compare registers set to a value less than twice the setting of TDDR, the PWM output dead time may be too short (or nonexistent) or illegal active-level PWM negative-phase output may occur during the initial output suppression interval. For details, see section 12.7.23, Notes on Output Waveform Control During Synchronous Counter Clearing in Complementary PWM Mode.

Figure 12.56 Timing for Synchronous Counter Clearing

• Example of Procedure for Setting Output Waveform Control at Synchronous Counter Clearing in Complementary PWM Mode

An example of the procedure for setting output waveform control at synchronous counter clearing in complementary PWM mode is shown in figure 12.57.

Figure 12.57 Example of Procedure for Setting Output Waveform Control at Synchronous Counter Clearing in Complementary PWM Mode

• Examples of Output Waveform Control at Synchronous Counter Clearing in Complementary PWM Mode

Figures 12.58 to 12.61 show examples of output waveform control in which the MTU2 operates in complementary PWM mode and synchronous counter clearing is generated while the WRE bit in TWCR is set to 1. In the examples shown in figures 12.58 to 12.61, synchronous counter clearing occurs at timing (3), (6), (8), and (11) shown in figure 12.56, respectively.

Figure 12.58 Example of Synchronous Clearing in Dead Time during Up-Counting (Timing (3) in Figure 12.56; Bit WRE of TWCR in MTU2 is 1)

Figure 12.59 Example of Synchronous Clearing in Interval Tb at Crest (Timing (6) in Figure 12.56; Bit WRE of TWCR in MTU2 is 1)

Figure 12.60 Example of Synchronous Clearing in Dead Time during Down-Counting (Timing (8) in Figure 12.56; Bit WRE of TWCR is 1)

Figure 12.61 Example of Synchronous Clearing in Interval Tb at Trough (Timing (11) in Figure 12.56; Bit WRE of TWCR is 1)

(o) Counter Clearing by TGRA_3 Compare Match

In complementary PWM mode, by setting the CCE bit in the timer waveform control register (TWCR), it is possible to have TCNT_3, TCNT_4, and TCNTS cleared by TGRA_3 compare match.

Figure 12.62 illustrates an operation example.

Notes: 1. Use this function only in complementary PWM mode 1 (transfer at crest)

- 2. Do not specify synchronous clearing by another channel (do not set the SYNC0 to SYNC4 bits in the timer synchronous register (TSYR) to 1.
- 3. Do not set the PWM duty value to H'0000.
- 4. Do not set the PSYE bit in timer output control register 1 (TOCR1) to 1.

Figure 12.62 Example of Counter Clearing Operation by TGRA_3 Compare Match

(p) Example of AC Synchronous Motor (Brushless DC Motor) Drive Waveform Output

In complementary PWM mode, a brushless DC motor can easily be controlled using the timer gate control register (TGCR). Figures 12.63 to 12.66 show examples of brushless DC motor drive waveforms created using TGCR.

When output phase switching for a 3-phase brushless DC motor is performed by means of external signals detected with a Hall element, etc., clear the FB bit in TGCR to 0. In this case, the external signals indicating the polarity position are input to channel 0 timer input pins TIOC0A, TIOC0B, and TIOC0C (set with PFC). When an edge is detected at pin TIOC0A, TIOC0B, or TIOC0C, the output on/off state is switched automatically.

When the FB bit is 1, the output on/off state is switched when the UF, VF, or WF bit in TGCR is cleared to 0 or set to 1.

The drive waveforms are output from the complementary PWM mode 6-phase output pins. With this 6-phase output, in the case of on output, it is possible to use complementary PWM mode output and perform chopping output by setting the N bit or P bit to 1. When the N bit or P bit is 0, level output is selected.

The 6-phase output active level (on output level) can be set with the OLSN and OLSP bits in the timer output control register (TOCR) regardless of the setting of the N and P bits.

Figure 12.63 Example of Output Phase Switching by External Input (1)

Figure 12.64 Example of Output Phase Switching by External Input (2)

Figure 12.65 Example of Output Phase Switching by Means of UF, VF, WF Bit Settings (1)

Figure 12.66 Example of Output Phase Switching by Means of UF, VF, WF Bit Settings (2)

(q) A/D Converter Start Request Setting

In complementary PWM mode, an A/D converter start request can be issued using a TGRA_3 compare-match, TCNT_4 underflow (trough), or compare-match on a channel other than channels 3 and 4.

When start requests using a TGRA_{_3} compare-match are specified, A/D conversion can be started at the crest of the TCNT_3 count.

A/D converter start requests can be set by setting the TTGE bit to 1 in the timer interrupt enable register (TIER). To issue an A/D converter start request at a TCNT_4 underflow (trough), set the TTGE2 bit in TIER_4 to 1.

(3) Interrupt Skipping in Complementary PWM Mode

Interrupts TGIA 3 (at the crest) and TCIV 4 (at the trough) in channels 3 and 4 can be skipped up to seven times by making settings in the timer interrupt skipping set register (TITCR).

Transfers from a buffer register to a temporary register or a compare register can be skipped in coordination with interrupt skipping by making settings in the timer buffer transfer register (TBTER). For the linkage with buffer registers, refer to description (c), Buffer Transfer Control Linked with Interrupt Skipping, below.

A/D converter start requests generated by the A/D converter start request delaying function can also be skipped in coordination with interrupt skipping by making settings in the timer A/D converter request control register (TADCR). For the linkage with the A/D converter start request delaying function, refer to section 12.4.9, A/D Converter Start Request Delaying Function.

The setting of the timer interrupt skipping setting register (TITCR) must be done while the TGIA_3 and TCIV_4 interrupt requests are disabled by the settings of TIER_3 and TIER_4 along with under the conditions in which TGFA 3 and TCFV 4 flag settings by compare match never occur. Before changing the skipping count, be sure to clear the T3AEN and T4VEN bits to 0 to clear the skipping counter.

(a) Example of Interrupt Skipping Operation Setting Procedure

Figure 12.67 shows an example of the interrupt skipping operation setting procedure. Figure 12.68 shows the periods during which interrupt skipping count can be changed.

Figure 12.67 Example of Interrupt Skipping Operation Setting Procedure

Figure 12.68 Periods during which Interrupt Skipping Count can be Changed

(b) Example of Interrupt Skipping Operation

Figure 12.69 shows an example of TGIA_3 interrupt skipping in which the interrupt skipping count is set to three by the 3ACOR bit and the T3AEN bit is set to 1 in the timer interrupt skipping set register (TITCR).

Figure 12.69 Example of Interrupt Skipping Operation

(c) Buffer Transfer Control Linked with Interrupt Skipping

In complementary PWM mode, whether to transfer data from a buffer register to a temporary register and whether to link the transfer with interrupt skipping can be specified with the BTE1 and BTE0 bits in the timer buffer transfer set register (TBTER).

Figure 12.70 shows an example of operation when buffer transfer is suppressed (BTE1 $= 0$ and $BTE0 = 1$). While this setting is valid, data is not transferred from the buffer register to the temporary register.

Figure 12.71 shows an example of operation when buffer transfer is linked with interrupt skipping $(BTE1 = 1$ and $BET0 = 0)$. While this setting is valid, data is not transferred from the buffer register outside the buffer transfer-enabled period.

Due to the buffer register rewrite timing after an interrupt, the timing of transfers from a buffer register to a temporary register differs from the timing of transfers from a temporary register to a general register.

Note that the buffer transfer-enabled period depends on the T3AEN and T4VEN bit settings in the timer interrupt skipping set register (TITCR). Figure 12.72 shows the relationship between the T3AEN and T4VEN bit settings in TITCR and buffer transfer-enabled period.

Note: This function must always be used in combination with interrupt skipping. When interrupt skipping is disabled (the T3AEN and T4VEN bits in the timer interrupt skipping set register (TITCR) are cleared to 0 or the skipping count set bits (3ACOR and 4VCOR) in TITCR are cleared to 0), make sure that buffer transfer is not linked with interrupt skipping (clear the BTE1 bit in the timer buffer transfer set register (TBTER) to 0). If buffer transfer is linked with interrupt skipping while interrupt skipping is disabled, buffer transfer is never performed.

Figure 12.70 Example of Operation when Buffer Transfer is Suppressed $(BTE1 = 0 \text{ and } BTE0 = 1)$

(2)When rewriting the buffer register after passing 1 carrier cycle from TGIA_3 interrupt

Figure 12.71 Example of Operation when Buffer Transfer is Linked with Interrupt Skipping (BTE1 = 1 and BTE0 = 0)

Figure 12.72 Relationship between Bits T3AEN and T4VEN in TITCR and Buffer Transfer-Enabled Period

(4) Complementary PWM Mode Output Protection Function

Complementary PWM mode output has the following protection function.

(a) Register and counter miswrite prevention function

With the exception of the buffer registers, which can be rewritten at any time, access by the CPU can be enabled or disabled for the mode registers, control registers, compare registers, and counters used in complementary PWM mode by means of the RWE bit in the timer read/write enable register (TRWER). The applicable registers are some (21 in total) of the registers in channels 3 and 4 shown in the following:

• TCR 3 and TCR 4, TMDR 3 and TMDR 4, TIORH 3 and TIORH 4, TIORL 3 and TIORL_4, TIER_3 and TIER_4, TCNT_3 and TCNT_4, TGRA_3 and TGRA_4, TGRB_3 and TGRB_4, TOER, TOCR, TGCR, TCDR, and TDDR.

This function enables miswriting due to CPU runaway to be prevented by disabling CPU access to the mode registers, control registers, and counters. When the applicable registers are read in the access-disabled state, undefined values are returned. Writing to these registers is ignored.

12.4.9 A/D Converter Start Request Delaying Function

A/D converter start requests can be issued in channel 4 by making settings in the timer A/D converter start request control register (TADCR), timer A/D converter start request cycle set registers (TADCORA_4 and TADCORB_4), and timer A/D converter start request cycle set buffer registers (TADCOBRA 4 and TADCOBRB 4).

The A/D converter start request delaying function compares TCNT_4 with TADCORA_4 or TADCORB_4, and when their values match, the function issues a respective A/D converter start request (TRG4AN or TRG4BN).

A/D converter start requests (TRG4AN and TRG4BN) can be skipped in coordination with interrupt skipping by making settings in the ITA3AE, ITA4VE, ITB3AE, and ITB4VE bits in **TADCR**

• Example of Procedure for Specifying A/D Converter Start Request Delaying Function Figure 12.73 shows an example of procedure for specifying the A/D converter start request delaying function.

Figure 12.73 Example of Procedure for Specifying A/D Converter Start Request Delaying Function

• Basic Operation Example of A/D Converter Start Request Delaying Function

Figure 12.74 shows a basic example of A/D converter request signal (TRG4AN) operation when the trough of TCNT 4 is specified for the buffer transfer timing and an A/D converter start request signal is output during TCNT_4 down-counting.

Figure 12.74 Basic Example of A/D Converter Start Request Signal (TRG4AN) Operation

• Buffer Transfer

The data in the timer A/D converter start request cycle set registers (TADCORA_4 and TADCORB_4) is updated by writing data to the timer A/D converter start request cycle set buffer registers (TADCOBRA_4 and TADCOBRB_4). Data is transferred from the buffer registers to the respective cycle set registers at the timing selected with the BF1 and BF0 bits in the timer A/D converter start request control register (TADCR_4).

• A/D Converter Start Request Delaying Function Linked with Interrupt Skipping A/D converter start requests (TRG4AN and TRG4BN) can be issued in coordination with interrupt skipping by making settings in the ITA3AE, ITA4VE, ITB3AE, and ITB4VE bits in the timer A/D converter start request control register (TADCR).

Figure 12.75 shows an example of A/D converter start request signal (TRG4AN) operation when TRG4AN output is enabled during TCNT_4 up-counting and down-counting and A/D converter start requests are linked with interrupt skipping.

Figure 12.76 shows another example of A/D converter start request signal (TRG4AN) operation when TRG4AN output is enabled during TCNT_4 up-counting and A/D converter start requests are linked with interrupt skipping.

Note: This function must be used in combination with interrupt skipping. When interrupt skipping is disabled (the T3AEN and T4VEN bits in the timer interrupt skipping set register (TITCR) are cleared to 0 or the skipping count set bits (3ACOR and 4VCOR) in TITCR are cleared to 0), make sure that A/D converter start requests are not linked with interrupt skipping (clear the ITA3AE, ITA4VE, ITB3AE, and ITB4VE bits in the timer A/D converter start request control register (TADCR) to 0).

Figure 12.75 Example of A/D Converter Start Request Signal (TRG4AN) Operation Linked with Interrupt Skipping

Figure 12.76 Example of A/D Converter Start Request Signal (TRG4AN) Operation Linked with Interrupt Skipping

12.4.10 TCNT Capture at Crest and/or Trough in Complementary PWM Operation

The TCNT value is captured in TGR at either the crest or trough or at both the crest and trough during complementary PWM operation. The timing for capturing in TGR can be selected by TIOR.

Figure 12.77 shows an example in which TCNT is used as a free-running counter without being cleared, and the TCNT value is captured in TGR at the specified timing (either crest or trough, or both crest and trough).

Figure 12.77 TCNT Capturing at Crest and/or Trough in Complementary PWM Operation

12.5 Interrupt Sources

12.5.1 Interrupt Sources and Priorities

There are three kinds of MTU2 interrupt source; TGR input capture/compare match, TCNT overflow, and TCNT underflow. Each interrupt source has its own status flag and enable/disabled bit, allowing the generation of interrupt request signals to be enabled or disabled individually.

When an interrupt request is generated, the corresponding status flag in TSR is set to 1. If the corresponding enable/disable bit in TIER is set to 1 at this time, an interrupt is requested. The interrupt request is cleared by clearing the status flag to 0.

Relative channel priorities can be changed by the interrupt controller, however the priority order within a channel is fixed. For details, see section 7, Interrupt Controller (INTC).

Table 12.55 lists the MTU2 interrupt sources.

Table 12.55 MTU2 Interrupts

Note: This table shows the initial state immediately after a reset. The relative channel priorities can be changed by the interrupt controller.

(1) Input Capture/Compare Match Interrupt

An interrupt is requested if the TGIE bit in TIER is set to 1 when the TGF flag in TSR is set to 1 by the occurrence of a TGR input capture/compare match on a particular channel. The interrupt request is cleared by clearing the TGF flag to 0. The MTU2 has eighteen input capture/compare match interrupts, six for channel 0, four each for channels 3 and 4, and two each for channels 1 and 2. The TGFE_0 and TGFF_0 flags in channel 0 are not set by the occurrence of an input capture.

(2) Overflow Interrupt

An interrupt is requested if the TCIEV bit in TIER is set to 1 when the TCFV flag in TSR is set to 1 by the occurrence of TCNT overflow on a channel. The interrupt request is cleared by clearing the TCFV flag to 0. The MTU2 has five overflow interrupts, one for each channel.

(3) Underflow Interrupt

An interrupt is requested if the TCIEU bit in TIER is set to 1 when the TCFU flag in TSR is set to 1 by the occurrence of TCNT underflow on a channel. The interrupt request is cleared by clearing the TCFU flag to 0. The MTU2 has two underflow interrupts, one each for channels 1 and 2.

12.5.2 DMAC Activation

The DMAC can be activated by the TGRA input capture/compare match interrupt in each channel. For details, see section 11, Direct Memory Access Controller (DMAC).

In the MTU2, a total of five TGRA input capture/compare match interrupts can be used as DMAC activation sources, one each for channels 0 to 4.

12.5.3 A/D Converter Activation

The A/D converter can be activated by one of the following three methods in the MTU2. Table 12.56 shows the relationship between interrupt sources and A/D converter start request signals.

(1) A/D Converter Activation by TGRA Input Capture/Compare Match or at TCNT_4 Trough in Complementary PWM Mode

The A/D converter can be activated by the occurrence of a TGRA input capture/compare match in each channel. In addition, if complementary PWM operation is performed while the TTGE2 bit in TIER_4 is set to 1, the A/D converter can be activated at the trough of $TCNT_4$ count (TCNT_4 = H'0000).

A/D converter start request signal TRGAN is issued to the A/D converter under either one of the following conditions.

- When the TGFA flag in TSR is set to 1 by the occurrence of a TGRA input capture/compare match on a particular channel while the TTGE bit in TIER is set to 1
- When the TCNT 4 count reaches the trough (TCNT $4 = H'0000$) during complementary PWM operation while the TTGE2 bit in TIER_4 is set to 1

When either condition is satisfied, if A/D converter start signal TRGAN from the MTU2 is selected as the trigger in the A/D converter, A/D conversion will start.

(2) A/D Converter Activation by Compare Match between TCNT_0 and TGRE_0

The A/D converter can be activated by generating A/D converter start request signal TRG0N when a compare match occurs between TCNT_0 and TGRE_0 in channel 0.

When the TGFE flag in TSR2_0 is set to 1 by the occurrence of a compare match between TCNT_0 and TGRE_0 in channel 0 while the TTGE2 bit in TIER2_0 is set to 1, A/D converter start request TGR0N is issued to the A/D converter. If A/D converter start signal TGR0N from the MTU2 is selected as the trigger in the A/D converter, A/D conversion will start.

(3) A/D Converter Activation by A/D Converter Start Request Delaying Function

The A/D converter can be activated by generating A/D converter start request signal TRG4AN or TRG4BN when the TCNT_4 count matches the TADCORA or TADCORB value if the UT4AE, DT4AE, UT4BE, and DT4BE bit in the A/D converter start request control register (TADCR) is set to 1. For details, refer to section 12.4.9, A/D Converter Start Request Delaying Function.

A/D conversion will start if A/D converter start signal TRG4AN from the MTU2 is selected as the trigger in the A/D converter when TRG4AN is generated or if TRG4BN from the MTU2 is selected as the trigger in the A/D converter when TRG4BN is generated.

Table 12.56 Interrupt Sources and A/D Converter Start Request Signals

12.6 Operation Timing

12.6.1 Input/Output Timing

(1) TCNT Count Timing

Figure 12.78 shows TCNT count timing in internal clock operation, and figure 12.79 shows TCNT count timing in external clock operation (normal mode), and figure 12.80 shows TCNT count timing in external clock operation (phase counting mode).

Figure 12.80 Count Timing in External Clock Operation (Phase Counting Mode)

(2) Output Compare Output Timing

A compare match signal is generated in the final state in which TCNT and TGR match (the point at which the count value matched by TCNT is updated). When a compare match signal is generated, the output value set in TIOR is output at the output compare output pin (TIOC pin). After a match between TCNT and TGR, the compare match signal is not generated until the TCNT input clock is generated.

Figure 12.81 shows output compare output timing (normal mode and PWM mode) and figure 12.82 shows output compare output timing (complementary PWM mode and reset synchronous PWM mode).

Figure 12.81 Output Compare Output Timing (Normal Mode/PWM Mode)

Figure 12.82 Output Compare Output Timing (Complementary PWM Mode/Reset Synchronous PWM Mode)

(3) Input Capture Signal Timing

Figure 12.83 shows input capture signal timing.

Figure 12.83 Input Capture Input Signal Timing

(4) Timing for Counter Clearing by Compare Match/Input Capture

Figure 12.84 shows the timing when counter clearing on compare match is specified, and figure 12.85 shows the timing when counter clearing on input capture is specified.

Figure 12.84 Counter Clear Timing (Compare Match)

Figure 12.85 Counter Clear Timing (Input Capture)

(5) Buffer Operation Timing

Figures 12.86 to 12.88 show the timing in buffer operation.

Figure 12.86 Buffer Operation Timing (Compare Match)

Figure 12.87 Buffer Operation Timing (Input Capture)

(6) Buffer Transfer Timing (Complementary PWM Mode)

Figures 12.89 to 12.91 show the buffer transfer timing in complementary PWM mode.

Figure 12.89 Transfer Timing from Buffer Register to Temporary Register (TCNTS Stop)

Figure 12.90 Transfer Timing from Buffer Register to Temporary Register (TCNTS Operating)

Figure 12.91 Transfer Timing from Temporary Register to Compare Register

12.6.2 Interrupt Signal Timing

(1) TGF Flag Setting Timing in Case of Compare Match

Figure 12.92 shows the timing for setting of the TGF flag in TSR on compare match, and TGI interrupt request signal timing.

Figure 12.92 TGI Interrupt Timing (Compare Match)

(2) TGF Flag Setting Timing in Case of Input Capture

Figure 12.93 shows the timing for setting of the TGF flag in TSR on input capture, and TGI interrupt request signal timing.

Figure 12.93 TGI Interrupt Timing (Input Capture)

(3) TCFV Flag/TCFU Flag Setting Timing

Figure 12.94 shows the timing for setting of the TCFV flag in TSR on overflow, and TCIV interrupt request signal timing.

Figure 12.95 shows the timing for setting of the TCFU flag in TSR on underflow, and TCIU interrupt request signal timing.

Figure 12.94 TCIV Interrupt Setting Timing

Figure 12.95 TCIU Interrupt Setting Timing

(4) Status Flag Clearing Timing

After a status flag is read as 1 by the CPU, it is cleared by writing 0 to it. When the DMAC is activated, the flag is cleared automatically. Figure 12.96 shows the timing for status flag clearing by the CPU, and figure 12.97 shows the timing for status flag clearing by the DMAC.

Figure 12.96 Timing for Status Flag Clearing by CPU

Figure 12.97 Timing for Status Flag Clearing by DMAC Activation

12.7 Usage Notes

12.7.1 Module Standby Mode Setting

MTU2 operation can be disabled or enabled using the standby control register. The initial setting is for MTU2 operation to be halted. Register access is enabled by clearing module standby mode. For details, refer to section 30, Power-Down Modes.

12.7.2 Input Clock Restrictions

The input clock pulse width must be at least 1.5 states in the case of single-edge detection, and at least 2.5 states in the case of both-edge detection. The MTU2 will not operate properly at narrower pulse widths.

In phase counting mode, the phase difference and overlap between the two input clocks must be at least 1.5 states, and the pulse width must be at least 2.5 states. Figure 12.98 shows the input clock conditions in phase counting mode.

Figure 12.98 Phase Difference, Overlap, and Pulse Width in Phase Counting Mode

12.7.3 Caution on Period Setting

When counter clearing on compare match is set, TCNT is cleared in the final state in which it matches the TGR value (the point at which the count value matched by TCNT is updated). Consequently, the actual counter frequency is given by the following formula:

$$
f = \frac{P\phi}{(N+1)}
$$

Where f: Counter frequency

Pφ: Peripheral clock operating frequency

N: TGR set value

12.7.4 Contention between TCNT Write and Clear Operations

If the counter clear signal is generated in the T2 state of a TCNT write cycle, TCNT clearing takes precedence and the TCNT write is not performed.

Figure 12.99 shows the timing in this case.

Figure 12.99 Contention between TCNT Write and Clear Operations

12.7.5 Contention between TCNT Write and Increment Operations

If incrementing occurs in the T2 state of a TCNT write cycle, the TCNT write takes precedence and TCNT is not incremented.

Figure 12.100 shows the timing in this case.

Figure 12.100 Contention between TCNT Write and Increment Operations

12.7.6 Contention between TGR Write and Compare Match

If a compare match occurs in the T2 state of a TGR write cycle, the TGR write is executed and the compare match signal is also generated.

Figure 12.101 shows the timing in this case.

Figure 12.101 Contention between TGR Write and Compare Match

12.7.7 Contention between Buffer Register Write and Compare Match

If a compare match occurs in the T2 state of a TGR write cycle, the data that is transferred to TGR by the buffer operation is the data after write.

Figure 12.102 shows the timing in this case.

Figure 12.102 Contention between Buffer Register Write and Compare Match

12.7.8 Contention between Buffer Register Write and TCNT Clear

When the buffer transfer timing is set at the TCNT clear by the buffer transfer mode register (TBTM), if TCNT clear occurs in the T2 state of a TGR write cycle, the data that is transferred to TGR by the buffer operation is the data before write.

Figure 12.103 Contention between Buffer Register Write and TCNT Clear

12.7.9 Contention between TGR Read and Input Capture

If an input capture signal is generated in the T1 state of a TGR read cycle, the data that is read will be the data in the buffer before input capture transfer.

Figure 12.104 shows the timing in this case.

Figure 12.104 Contention between TGR Read and Input Capture

12.7.10 Contention between TGR Write and Input Capture

If an input capture signal is generated in the T2 state of a TGR write cycle, the input capture operation takes precedence and the write to TGR is not performed.

Figure 12.105 shows the timing in this case.

Figure 12.105 Contention between TGR Write and Input Capture

12.7.11 Contention between Buffer Register Write and Input Capture

If an input capture signal is generated in the T2 state of a buffer register write cycle, the buffer operation takes precedence and the write to the buffer register is not performed.

Figure 12.106 Contention between Buffer Register Write and Input Capture

12.7.12 TCNT2 Write and Overflow/Underflow Contention in Cascade Connection

With timer counters TCNT1 and TCNT2 in a cascade connection, when a contention occurs during $TCNT_1$ count (during a $TCNT_2$ overflow/underflow) in the T , state of the $TCNT_2$ write cycle, the write to $TCNT$ 2 is conducted, and the $TCNT$ 1 count signal is disabled. At this point, if there is match with TGRA_1 and the TCNT_1 value, a compare signal is issued. Furthermore, when the TCNT 1 count clock is selected as the input capture source of channel 0, TGRA_0 to D_0 carry out the input capture operation. In addition, when the compare match/input capture is selected as the input capture source of TGRB_1, TGRB_1 carries out input capture operation. The timing is shown in figure 12.107.

For cascade connections, be sure to synchronize settings for channels 1 and 2 when setting TCNT clearing.

Figure 12.107 TCNT_2 Write and Overflow/Underflow Contention with Cascade Connection

12.7.13 Counter Value during Complementary PWM Mode Stop

When counting operation is suspended with TCNT 3 and TCNT 4 in complementary PWM mode, TCNT 3 has the timer dead time register (TDDR) value, and TCNT 4 is held at H'0000.

When restarting complementary PWM mode, counting begins automatically from the initialized state. This explanatory diagram is shown in figure 12.108.

When counting begins in another operating mode, be sure that TCNT_3 and TCNT_4 are set to the initial values.

Figure 12.108 Counter Value during Complementary PWM Mode Stop

12.7.14 Buffer Operation Setting in Complementary PWM Mode

In complementary PWM mode, conduct rewrites by buffer operation for the PWM cycle setting register (TGRA_3), timer cycle data register (TCDR), and duty setting registers (TGRB_3, TGRA_4, and TGRB_4).

In complementary PWM mode, channel 3 and channel 4 buffers operate in accordance with bit settings BFA and BFB of TMDR 3. When TMDR 3's BFA bit is set to 1, TGRC 3 functions as a buffer register for TGRA_3. At the same time, TGRC_4 functions as the buffer register for TGRA_4, and TCBR functions as the TCDR's buffer register.

12.7.15 Reset Sync PWM Mode Buffer Operation and Compare Match Flag

When setting buffer operation for reset sync PWM mode, set the BFA and BFB bits of TMDR 4 to 0. The TIOC4C pin will be unable to produce its waveform output if the BFA bit of TMDR_4 is set to 1.

In reset sync PWM mode, the channel 3 and channel 4 buffers operate in accordance with the BFA and BFB bit settings of TMDR_3. For example, if the BFA bit of TMDR_3 is set to 1, TGRC_3 functions as the buffer register for TGRA_3. At the same time, TGRC_4 functions as the buffer register for TGRA_4.

The TGFC bit and TGFD bit of TSR_3 and TSR_4 are not set when TGRC_3 and TGRD_3 are operating as buffer registers.

Figure 12.109 shows an example of operations for TGR_3, TGR_4, TIOC3, and TIOC4, with TMDR_3's BFA and BFB bits set to 1, and TMDR_4's BFA and BFB bits set to 0.

Figure 12.109 Buffer Operation and Compare-Match Flags in Reset Synchronous PWM Mode

12.7.16 Overflow Flags in Reset Synchronous PWM Mode

When set to reset synchronous PWM mode, TCNT_3 and TCNT_4 start counting when the CST3 bit of TSTR is set to 1. At this point, TCNT_4's count clock source and count edge obey the TCR_3 setting.

In reset synchronous PWM mode, with cycle register TGRA_3's set value at H'FFFF, when specifying TGR3A compare-match for the counter clear source, TCNT_3 and TCNT_4 count up to H'FFFF, then a compare-match occurs with TGRA_3, and TCNT_3 and TCNT_4 are both cleared. At this point, TSR's overflow flag TCFV bit is not set.

Figure 12.110 shows a TCFV bit operation example in reset synchronous PWM mode with a set value for cycle register TGRA_3 of H'FFFF, when a TGRA_3 compare-match has been specified without synchronous setting for the counter clear source.

Figure 12.110 Reset Synchronous PWM Mode Overflow Flag

12.7.17 Contention between Overflow/Underflow and Counter Clearing

If overflow/underflow and counter clearing occur simultaneously, the TCFV/TCFU flag in TSR is not set and TCNT clearing takes precedence.

Figure 12.111 shows the operation timing when a TGR compare match is specified as the clearing source, and when H'FFFF is set in TGR.

Figure 12.111 Contention between Overflow and Counter Clearing

12.7.18 Contention between TCNT Write and Overflow/Underflow

If there is an up-count or down-count in the T2 state of a TCNT write cycle, and overflow/underflow occurs, the TCNT write takes precedence and the TCFV/TCFU flag in TSR is not set.

Figure 12.112 shows the operation timing when there is contention between TCNT write and overflow.

Figure 12.112 Contention between TCNT Write and Overflow

12.7.19 Cautions on Transition from Normal Operation or PWM Mode 1 to Reset-Synchronized PWM Mode

When making a transition from channel 3 or 4 normal operation or PWM mode 1 to resetsynchronized PWM mode, if the counter is halted with the output pins (TIOC3B, TIOC3D, TIOC4A, TIOC4C, TIOC4B, TIOC4D) in the high-level state, followed by the transition to resetsynchronized PWM mode and operation in that mode, the initial pin output will not be correct.

When making a transition from normal operation to reset-synchronized PWM mode, write H'11 to registers TIORH_3, TIORL_3, TIORH_4, and TIORL_4 to initialize the output pins to low level output, then set an initial register value of H'00 before making the mode transition.

When making a transition from PWM mode 1 to reset-synchronized PWM mode, first switch to normal operation, then initialize the output pins to low level output and set an initial register value of H'00 before making the transition to reset-synchronized PWM mode.

12.7.20 Output Level in Complementary PWM Mode and Reset-Synchronized PWM Mode

When channels 3 and 4 are in complementary PWM mode or reset-synchronized PWM mode, the PWM waveform output level is set with the OLSP and OLSN bits in the timer output control register (TOCR). In the case of complementary PWM mode or reset-synchronized PWM mode, TIOR should be set to H'00.

12.7.21 Interrupts in Module Standby Mode

If module standby mode is entered when an interrupt has been requested, it will not be possible to clear the CPU interrupt source or the DMAC activation source. Interrupts should therefore be disabled before entering module standby mode.

12.7.22 Simultaneous Capture of TCNT_1 and TCNT_2 in Cascade Connection

When timer counters 1 and 2 (TCNT 1 and TCNT 2) are operated as a 32-bit counter in cascade connection, the cascade counter value cannot be captured successfully even if input-capture input is simultaneously done to TIOC1A and TIOC2A or to TIOC1B and TIOC2B. This is because the input timing of TIOC1A and TIOC2A or of TIOC1B and TIOC2B may not be the same when external input-capture signals to be input into TCNT_1 and TCNT_2 are taken in synchronization with the internal clock. For example, TCNT 1 (the counter for upper 16 bits) does not capture the count-up value by overflow from TCNT_2 (the counter for lower 16 bits) but captures the count value before the count-up. In this case, the values of $TCNT_1 = H'FFT1$ and $TCNT_2 = H'0000$ should be transferred to TGRA_1 and TGRA_2 or to TGRB_1 and TGRB_2, but the values of TCNT_1 = H'FFF0 and TCNT_2 = H'0000 are erroneously transferred.

12.7.23 Notes on Output Waveform Control During Synchronous Counter Clearing in Complementary PWM Mode

In complementary PWM mode, when output waveform control during synchronous counter clearing is enabled (WRE in the TWCR register set to 1), the following problems may occur when condition (1) or condition (2), below, is satisfied.

- Dead time for the PWM output pins may be too short (or nonexistent).
- Active-level output from the PWM negative-phase pins may occur outside the correct activelevel output interval

Condition (1): When synchronous clearing occurs in the PWM output dead time interval within initial output suppression interval (10) (figure 12.113).

Condition (2): When synchronous clearing occurs within initial output suppression interval (10) or (11) and TGRB $3 \leq$ TDDR, TGRA $4 \leq$ TDDR, or TGRB $4 \leq$ TDDR is true (figure 12.114)

Figure 12.113 Condition (1) Synchronous Clearing Example

Figure 12.114 Condition (2) Synchronous Clearing Example

The following workaround can be used to avoid these problems.

When using synchronous clearing, make sure to set compare registers TGRB_3, TGRA_4, and TGRB_4 to a value twice or more the setting of dead time data register TDDR.

12.8 MTU2 Output Pin Initialization

12.8.1 Operating Modes

The MTU2 has the following six operating modes. Waveform output is possible in all of these modes.

- Normal mode (channels 0 to 4)
- PWM mode 1 (channels 0 to 4)
- PWM mode 2 (channels 0 to 2)
- Phase counting modes 1 to 4 (channels 1 and 2)
- Complementary PWM mode (channels 3 and 4)
- Reset-synchronized PWM mode (channels 3 and 4)

The MTU2 output pin initialization method for each of these modes is described in this section.

12.8.2 Reset Start Operation

The MTU2 output pins (TIOC*) are initialized low by a power-on reset and in deep standby mode. Since MTU2 pin function selection is performed by the pin function controller (PFC), when the PFC is set, the MTU2 pin states at that point are output to the ports. When MTU2 output is selected by the PFC immediately after a reset, the MTU2 output initial level, low, is output directly at the port. When the active level is low, the system will operate at this point, and therefore the PFC setting should be made after initialization of the MTU2 output pins is completed.

Note: Channel number and port notation are substituted for *.

12.8.3 Operation in Case of Re-Setting Due to Error During Operation, etc.

If an error occurs during MTU2 operation, MTU2 output should be cut by the system. Cutoff is performed by switching the pin output to port output with the PFC and outputting the inverse of the active level. The pin initialization procedures for re-setting due to an error during operation, etc., and the procedures for restarting in a different mode after re-setting, are shown below.

The MTU2 has six operating modes, as stated above. There are thus 36 mode transition combinations, but some transitions are not available with certain channel and mode combinations. Possible mode transition combinations are shown in table 12.57.

Table 12.57 Mode Transition Combinations

[Legend]

Normal: Normal mode

PWM1: PWM mode 1

PWM2: PWM mode 2

PCM: Phase counting modes 1 to 4

CPWM: Complementary PWM mode

RPWM: Reset-synchronized PWM mode

12.8.4 Overview of Initialization Procedures and Mode Transitions in Case of Error during Operation, etc.

- When making a transition to a mode (Normal, PWM1, PWM2, PCM) in which the pin output level is selected by the timer I/O control register (TIOR) setting, initialize the pins by means of a TIOR setting.
- In PWM mode 1, since a waveform is not output to the TIOC*B (TIOC *D) pin, setting TIOR will not initialize the pins. If initialization is required, carry it out in normal mode, then switch to PWM mode 1.
- In PWM mode 2, since a waveform is not output to the cycle register pin, setting TIOR will not initialize the pins. If initialization is required, carry it out in normal mode, then switch to PWM mode 2.
- In normal mode or PWM mode 2, if TGRC and TGRD operate as buffer registers, setting TIOR will not initialize the buffer register pins. If initialization is required, clear buffer mode, carry out initialization, then set buffer mode again.
- In PWM mode 1, if either TGRC or TGRD operates as a buffer register, setting TIOR will not initialize the TGRC pin. To initialize the TGRC pin, clear buffer mode, carry out initialization, then set buffer mode again.
- When making a transition to a mode (CPWM, RPWM) in which the pin output level is selected by the timer output control register (TOCR) setting, switch to normal mode and perform initialization with TIOR, then restore TIOR to its initial value, and temporarily disable channel 3 and 4 output with the timer output master enable register (TOER). Then operate the unit in accordance with the mode setting procedure (TOCR setting, TMDR setting, TOER setting).

Note: Channel number is substituted for * indicated in this article.

Pin initialization procedures are described below for the numbered combinations in table 12.57. The active level is assumed to be low.

(1) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in Normal Mode

Figure 12.115 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in normal mode after re-setting.

Figure 12.115 Error Occurrence in Normal Mode, Recovery in Normal Mode

- 1. After a reset, MTU2 output is low and ports are in the high-impedance state.
- 2. After a reset, the TMDR setting is for normal mode.
- 3. For channels 3 and 4, enable output with TOER before initializing the pins with TIOR.
- 4. Initialize the pins with TIOR. (The example shows initial high output, with low output on compare-match occurrence.)
- 5. Set MTU2 output with the PFC.
- 6. The count operation is started by TSTR.
- 7. Output goes low on compare-match occurrence.
- 8. An error occurs.
- 9. Set port output with the PFC and output the inverse of the active level.
- 10. The count operation is stopped by TSTR.
- 11. Not necessary when restarting in normal mode.
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.
(2) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in PWM Mode 1

Figure 12.116 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in PWM mode 1 after re-setting.

Figure 12.116 Error Occurrence in Normal Mode, Recovery in PWM Mode 1

- 11. Set PWM mode 1.
- 12. Initialize the pins with TIOR. (In PWM mode 1, the TIOC*B side is not initialized. If initialization is required, initialize in normal mode, then switch to PWM mode 1.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(3) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in PWM Mode 2

Figure 12.117 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in PWM mode 2 after re-setting.

Figure 12.117 Error Occurrence in Normal Mode, Recovery in PWM Mode 2

- 11. Set PWM mode 2.
- 12. Initialize the pins with TIOR. (In PWM mode 2, the cycle register pins are not initialized. If initialization is required, initialize in normal mode, then switch to PWM mode 2.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.
- Note: PWM mode 2 can only be set for channels 0 to 2, and therefore TOER setting is not necessary.

(4) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in Phase Counting Mode

Figure 12.118 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in phase counting mode after re-setting.

Figure 12.118 Error Occurrence in Normal Mode, Recovery in Phase Counting Mode

- 11. Set phase counting mode.
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.
- Note: Phase counting mode can only be set for channels 1 and 2, and therefore TOER setting is not necessary.

(5) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in Complementary PWM Mode

Figure 12.119 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in complementary PWM mode after re-setting.

- 11. Initialize the normal mode waveform generation section with TIOR.
- 12. Disable operation of the normal mode waveform generation section with TIOR.
- 13. Disable channel 3 and 4 output with TOER.
- 14. Select the complementary PWM output level and cyclic output enabling/disabling with TOCR.
- 15. Set complementary PWM.
- 16. Enable channel 3 and 4 output with TOER.
- 17. Set MTU2 output with the PFC.
- 18. Operation is restarted by TSTR.

(6) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in Reset-Synchronized PWM Mode

Figure 12.120 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in reset-synchronized PWM mode after re-setting.

- 14. Select the reset-synchronized PWM output level and cyclic output enabling/disabling with TOCR.
- 15. Set reset-synchronized PWM.
- 16. Enable channel 3 and 4 output with TOER.
- 17. Set MTU2 output with the PFC.
- 18. Operation is restarted by TSTR.

(7) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in Normal Mode

Figure 12.121 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in normal mode after re-setting.

Figure 12.121 Error Occurrence in PWM Mode 1, Recovery in Normal Mode

- 1. After a reset, MTU2 output is low and ports are in the high-impedance state.
- 2. Set PWM mode 1.
- 3. For channels 3 and 4, enable output with TOER before initializing the pins with TIOR.
- 4. Initialize the pins with TIOR. (The example shows initial high output, with low output on compare-match occurrence. In PWM mode 1, the TIOC*B side is not initialized.)
- 5. Set MTU2 output with the PFC.
- 6. The count operation is started by TSTR.
- 7. Output goes low on compare-match occurrence.
- 8. An error occurs.
- 9. Set port output with the PFC and output the inverse of the active level.
- 10. The count operation is stopped by TSTR.
- 11. Set normal mode.
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(8) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in PWM Mode 1

Figure 12.122 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in PWM mode 1 after re-setting.

- 11. Not necessary when restarting in PWM mode 1.
- 12. Initialize the pins with TIOR. (In PWM mode 1, the TIOC*B side is not initialized.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(9) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in PWM Mode 2

Figure 12.123 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in PWM mode 2 after re-setting.

Figure 12.123 Error Occurrence in PWM Mode 1, Recovery in PWM Mode 2

- 11. Set PWM mode 2.
- 12. Initialize the pins with TIOR. (In PWM mode 2, the cycle register pins are not initialized.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.
- Note: PWM mode 2 can only be set for channels 0 to 2, and therefore TOER setting is not necessary.

(10) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in Phase Counting Mode

Figure 12.124 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in phase counting mode after re-setting.

- 11. Set phase counting mode.
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.
- Note: Phase counting mode can only be set for channels 1 and 2, and therefore TOER setting is not necessary.

(11) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in Complementary PWM Mode

Figure 12.125 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in complementary PWM mode after re-setting.

- 11. Set normal mode for initialization of the normal mode waveform generation section.
- 12. Initialize the PWM mode 1 waveform generation section with TIOR.
- 13. Disable operation of the PWM mode 1 waveform generation section with TIOR.
- 14. Disable channel 3 and 4 output with TOER.
- 15. Select the complementary PWM output level and cyclic output enabling/disabling with TOCR.
- 16. Set complementary PWM.
- 17. Enable channel 3 and 4 output with TOER.
- 18. Set MTU2 output with the PFC.
- 19. Operation is restarted by TSTR.

(12) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in Reset-Synchronized PWM Mode

Figure 12.126 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in reset-synchronized PWM mode after re-setting.

- 15. Select the reset-synchronized PWM output level and cyclic output enabling/disabling with TOCR.
- 16. Set reset-synchronized PWM.
- 17. Enable channel 3 and 4 output with TOER.
- 18. Set MTU2 output with the PFC.
- 19. Operation is restarted by TSTR.

(13) Operation when Error Occurs during PWM Mode 2 Operation, and Operation is Restarted in Normal Mode

Figure 12.127 shows an explanatory diagram of the case where an error occurs in PWM mode 2 and operation is restarted in normal mode after re-setting.

Figure 12.127 Error Occurrence in PWM Mode 2, Recovery in Normal Mode

- 1. After a reset, MTU2 output is low and ports are in the high-impedance state.
- 2. Set PWM mode 2.
- 3. Initialize the pins with TIOR. (The example shows initial high output, with low output on compare-match occurrence. In PWM mode 2, the cycle register pins are not initialized. In the example, TIOC *A is the cycle register.)
- 4. Set MTU2 output with the PFC.
- 5. The count operation is started by TSTR.
- 6. Output goes low on compare-match occurrence.
- 7. An error occurs.
- 8. Set port output with the PFC and output the inverse of the active level.
- 9. The count operation is stopped by TSTR.
- 10. Set normal mode.
- 11. Initialize the pins with TIOR.
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(14) Operation when Error Occurs during PWM Mode 2 Operation, and Operation is Restarted in PWM Mode 1

Figure 12.128 shows an explanatory diagram of the case where an error occurs in PWM mode 2 and operation is restarted in PWM mode 1 after re-setting.

- 10. Set PWM mode 1.
- 11. Initialize the pins with TIOR. (In PWM mode 1, the TIOC*B side is not initialized.)
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(15) Operation when Error Occurs during PWM Mode 2 Operation, and Operation is Restarted in PWM Mode 2

Figure 12.129 shows an explanatory diagram of the case where an error occurs in PWM mode 2 and operation is restarted in PWM mode 2 after re-setting.

Figure 12.129 Error Occurrence in PWM Mode 2, Recovery in PWM Mode 2

- 10. Not necessary when restarting in PWM mode 2.
- 11. Initialize the pins with TIOR. (In PWM mode 2, the cycle register pins are not initialized.)
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(16) Operation when Error Occurs during PWM Mode 2 Operation, and Operation is Restarted in Phase Counting Mode

Figure 12.130 shows an explanatory diagram of the case where an error occurs in PWM mode 2 and operation is restarted in phase counting mode after re-setting.

Figure 12.130 Error Occurrence in PWM Mode 2, Recovery in Phase Counting Mode

- 10. Set phase counting mode.
- 11. Initialize the pins with TIOR.
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(17) Operation when Error Occurs during Phase Counting Mode Operation, and Operation is Restarted in Normal Mode

Figure 12.131 shows an explanatory diagram of the case where an error occurs in phase counting mode and operation is restarted in normal mode after re-setting.

Figure 12.131 Error Occurrence in Phase Counting Mode, Recovery in Normal Mode

- 1. After a reset, MTU2 output is low and ports are in the high-impedance state.
- 2. Set phase counting mode.
- 3. Initialize the pins with TIOR. (The example shows initial high output, with low output on compare-match occurrence.)
- 4. Set MTU2 output with the PFC.
- 5. The count operation is started by TSTR.
- 6. Output goes low on compare-match occurrence.
- 7. An error occurs.
- 8. Set port output with the PFC and output the inverse of the active level.
- 9. The count operation is stopped by TSTR.
- 10. Set in normal mode.
- 11. Initialize the pins with TIOR.
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(18) Operation when Error Occurs during Phase Counting Mode Operation, and Operation is Restarted in PWM Mode 1

Figure 12.132 shows an explanatory diagram of the case where an error occurs in phase counting mode and operation is restarted in PWM mode 1 after re-setting.

- 10. Set PWM mode 1.
- 11. Initialize the pins with TIOR. (In PWM mode 1, the TIOC *B side is not initialized.)
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(19) Operation when Error Occurs during Phase Counting Mode Operation, and Operation is Restarted in PWM Mode 2

Figure 12.133 shows an explanatory diagram of the case where an error occurs in phase counting mode and operation is restarted in PWM mode 2 after re-setting.

Figure 12.133 Error Occurrence in Phase Counting Mode, Recovery in PWM Mode 2

- 10. Set PWM mode 2.
- 11. Initialize the pins with TIOR. (In PWM mode 2, the cycle register pins are not initialized.)
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(20) Operation when Error Occurs during Phase Counting Mode Operation, and Operation is Restarted in Phase Counting Mode

Figure 12.134 shows an explanatory diagram of the case where an error occurs in phase counting mode and operation is restarted in phase counting mode after re-setting.

- 10. Not necessary when restarting in phase counting mode.
- 11. Initialize the pins with TIOR.
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(21) Operation when Error Occurs during Complementary PWM Mode Operation, and Operation is Restarted in Normal Mode

Figure 12.135 shows an explanatory diagram of the case where an error occurs in complementary PWM mode and operation is restarted in normal mode after re-setting.

- 1. After a reset, MTU2 output is low and ports are in the high-impedance state.
- 2. Select the complementary PWM output level and cyclic output enabling/disabling with TOCR.
- 3. Set complementary PWM.
- 4. Enable channel 3 and 4 output with TOER.
- 5. Set MTU2 output with the PFC.
- 6. The count operation is started by TSTR.
- 7. The complementary PWM waveform is output on compare-match occurrence.
- 8. An error occurs.
- 9. Set port output with the PFC and output the inverse of the active level.
- 10. The count operation is stopped by TSTR. (MTU2 output becomes the complementary PWM output initial value.)
- 11. Set normal mode. (MTU2 output goes low.)
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(22) Operation when Error Occurs during Complementary PWM Mode Operation, and Operation is Restarted in PWM Mode 1

Figure 12.136 shows an explanatory diagram of the case where an error occurs in complementary PWM mode and operation is restarted in PWM mode 1 after re-setting.

- 11. Set PWM mode 1. (MTU2 output goes low.)
- 12. Initialize the pins with TIOR. (In PWM mode 1, the TIOC *B side is not initialized.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(23) Operation when Error Occurs during Complementary PWM Mode Operation, and Operation is Restarted in Complementary PWM Mode

Figure 12.137 shows an explanatory diagram of the case where an error occurs in complementary PWM mode and operation is restarted in complementary PWM mode after re-setting (when operation is restarted using the cycle and duty settings at the time the counter was stopped).

Figure 12.137 Error Occurrence in Complementary PWM Mode, Recovery in Complementary PWM Mode

- 11. Set MTU2 output with the PFC.
- 12. Operation is restarted by TSTR.
- 13. The complementary PWM waveform is output on compare-match occurrence.

(24) Operation when Error Occurs during Complementary PWM Mode Operation, and Operation is Restarted in Complementary PWM Mode

Figure 12.138 shows an explanatory diagram of the case where an error occurs in complementary PWM mode and operation is restarted in complementary PWM mode after re-setting (when operation is restarted using completely new cycle and duty settings).

- 11. Set normal mode and make new settings. (MTU2 output goes low.)
- 12. Disable channel 3 and 4 output with TOER.
- 13. Select the complementary PWM mode output level and cyclic output enabling/disabling with TOCR.
- 14. Set complementary PWM.
- 15. Enable channel 3 and 4 output with TOER.
- 16. Set MTU2 output with the PFC.
- 17. Operation is restarted by TSTR.

(25) Operation when Error Occurs during Complementary PWM Mode Operation, and Operation is Restarted in Reset-Synchronized PWM Mode

Figure 12.139 shows an explanatory diagram of the case where an error occurs in complementary PWM mode and operation is restarted in reset-synchronized PWM mode.

- 11. Set normal mode. (MTU2 output goes low.)
- 12. Disable channel 3 and 4 output with TOER.
- 13. Select the reset-synchronized PWM mode output level and cyclic output enabling/disabling with TOCR
- 14. Set reset-synchronized PWM.
- 15. Enable channel 3 and 4 output with TOER.
- 16. Set MTU2 output with the PFC.
- 17. Operation is restarted by TSTR.

(26) Operation when Error Occurs during Reset-Synchronized PWM Mode Operation, and Operation is Restarted in Normal Mode

Figure 12.140 shows an explanatory diagram of the case where an error occurs in resetsynchronized PWM mode and operation is restarted in normal mode after re-setting.

- 1. After a reset, MTU2 output is low and ports are in the high-impedance state.
- 2. Select the reset-synchronized PWM output level and cyclic output enabling/disabling with TOCR.
- 3. Set reset-synchronized PWM.
- 4. Enable channel 3 and 4 output with TOER.
- 5. Set MTU2 output with the PFC.
- 6. The count operation is started by TSTR.
- 7. The reset-synchronized PWM waveform is output on compare-match occurrence.
- 8. An error occurs.
- 9. Set port output with the PFC and output the inverse of the active level.
- 10. The count operation is stopped by TSTR. (MTU2 output becomes the reset-synchronized PWM output initial value.)
- 11. Set normal mode. (MTU2 positive phase output is low, and negative phase output is high.)
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(27) Operation when Error Occurs during Reset-Synchronized PWM Mode Operation, and Operation is Restarted in PWM Mode 1

Figure 12.141 shows an explanatory diagram of the case where an error occurs in resetsynchronized PWM mode and operation is restarted in PWM mode 1 after re-setting.

- 11. Set PWM mode 1. (MTU2 positive phase output is low, and negative phase output is high.)
- 12. Initialize the pins with TIOR. (In PWM mode 1, the TIOC *B side is not initialized.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(28) Operation when Error Occurs during Reset-Synchronized PWM Mode Operation, and Operation is Restarted in Complementary PWM Mode

Figure 12.142 shows an explanatory diagram of the case where an error occurs in resetsynchronized PWM mode and operation is restarted in complementary PWM mode after resetting.

- 11. Disable channel 3 and 4 output with TOER.
- 12. Select the complementary PWM output level and cyclic output enabling/disabling with TOCR.
- 13. Set complementary PWM. (The MTU2 cyclic output pin goes low.)
- 14. Enable channel 3 and 4 output with TOER.
- 15. Set MTU2 output with the PFC.
- 16. Operation is restarted by TSTR.

(29) Operation when Error Occurs during Reset-Synchronized PWM Mode Operation, and Operation is Restarted in Reset-Synchronized PWM Mode

Figure 12.143 shows an explanatory diagram of the case where an error occurs in resetsynchronized PWM mode and operation is restarted in reset-synchronized PWM mode after resetting.

- 11. Set MTU2 output with the PFC.
- 12. Operation is restarted by TSTR.
- 13. The reset-synchronized PWM waveform is output on compare-match occurrence.

Section 13 Compare Match Timer (CMT)

This LSI has an on-chip compare match timer module (CMT) consisting of two units of twochannel 16-bit timers, which makes a total of four channels. The CMT has a 16-bit counter, and can generate interrupts at set intervals.

13.1 Features

- Independent selection of four counter input clocks at two channels Any of four internal clocks (Pφ/8, Pφ/32, Pφ/128, and Pφ/512) can be selected.
- Selection of DMA transfer request or interrupt request generation on compare match by DMAC setting
- When not in use, the CMT can be stopped by halting its clock supply to reduce power consumption.

Figure 13.1 shows a block diagram of CMT.

Figure 13.1 Block Diagram of CMT

13.2 Register Descriptions

The CMT has the following registers.

Table 13.1 Register Configuration

13.2.1 Compare Match Timer Start Register (CMSTR)

CMSTR is a 16-bit register that selects whether compare match counter (CMCNT) operates or is stopped.

13.2.2 Compare Match Timer Control/Status Register (CMCSR)

CMCSR is a 16-bit register that indicates compare match generation, enables or disables interrupts, and selects the counter input clock.

13.2.3 Compare Match Counter (CMCNT)

CMCNT is a 16-bit register used as an up-counter. When the counter input clock is selected with bits CKS[1:0] in CMCSR, and the STR bit in CMSTR is set to 1, CMCNT starts counting using the selected clock. When the value in CMCNT and the value in compare match constant register (CMCOR) match, CMCNT is cleared to H'0000 and the CMF flag in CMCSR is set to 1.

CMCNT is initialized to H'0000 by clearing any channels of the counter start bit from 1 to 0 in the compare match timer start register (CMSTR).

13.2.4 Compare Match Constant Register (CMCOR)

CMCOR is a 16-bit register that sets the interval up to a compare match with CMCNT.

13.3 Operation

13.3.1 Interval Count Operation

When an internal clock is selected with the CKS[1:0] bits in CMCSR and the STR bit in CMSTR is set to 1, CMCNT starts incrementing using the selected clock. When the values in CMCNT and CMCOR match, CMCNT is cleared to H'0000 and the CMF flag in CMCSR is set to 1. When the CMIE bit in CMCSR is set to 1 at this time, a compare match interrupt (CMI) is requested. CMCNT then starts counting up again from H'0000.

Figure 13.2 shows the operation of the compare match counter.

Figure 13.2 Counter Operation

13.3.2 CMCNT Count Timing

One of four clocks ($P\phi/8$, $P\phi/32$, $P\phi/128$, and $P\phi/512$) obtained by dividing the peripheral clock (Pφ) can be selected with the CKS1 and CKS0 bits in CMCSR. Figure 13.3 shows the timing.

Figure 13.3 Count Timing

13.4 Interrupts

13.4.1 Interrupt Sources and DMA Transfer Requests

The CMT has channels and each of them to which a different vector address is allocated has a compare match interrupt. When both the compare match flag (CMF) and the interrupt enable bit (CMIE) are set to 1, the corresponding interrupt request is output. When the interrupt is used to activate a CPU interrupt, the priority of channels can be changed by the interrupt controller settings. For details, see section 7, Interrupt Controller (INTC).

Clear the CMF bit to 0 by the user exception handling routine. If this operation is not carried out, another interrupt will be generated. By configuring the interrupt controller (INTC), the direct memory access controller (DMAC) can be activated when a compare match interrupt is requested. In this case, an interrupt is not issued to the CPU. If the setting to activate the DMAC has not been made, an interrupt request is sent to the CPU. The CMF bit is automatically cleared to 0 when data is transferred by the DMAC.

13.4.2 Timing of Compare Match Flag Setting

When CMCOR and CMCNT match, a compare match signal is generated at the last state in which the values match (the timing when the CMCNT value is updated to H'0000) and the CMF bit in CMCSR is set to 1. That is, after a match between CMCOR and CMCNT, the compare match signal is not generated until the next CMCNT counter clock input. Figure 13.4 shows the timing of CMF bit setting.

Figure 13.4 Timing of CMF Setting

13.4.3 Timing of Compare Match Flag Clearing

The CMF bit in CMCSR is cleared by first, reading as 1 then writing to 0. However, in the case of the DMAC being activated, the CMF bit is automatically cleared to 0 when data is transferred by the DMAC.
13.5 Usage Notes

13.5.1 Conflict between Write and Compare-Match Processes of CMCNT

When the compare match signal is generated in the T2 cycle while writing to CMCNT, clearing CMCNT has priority over writing to it. In this case, CMCNT is not written to. Figure 13.5 shows the timing to clear the CMCNT counter.

Figure 13.5 Conflict between Write and Compare Match Processes of CMCNT

13.5.2 Conflict between Word-Write and Count-Up Processes of CMCNT

Even when the count-up occurs in the T2 cycle while writing to CMCNT in words, the writing has priority over the count-up. In this case, the count-up is not performed. Figure 13.6 shows the timing to write to CMCNT in words.

Figure 13.6 Conflict between Word-Write and Count-Up Processes of CMCNT

13.5.3 Conflict between Byte-Write and Count-Up Processes of CMCNT

Even when the count-up occurs in the T2 cycle while writing to CMCNT in bytes, the writing has priority over the count-up. In this case, the count-up is not performed. The byte data on the other side, which is not written to, is also not counted and the previous contents are retained.

Figure 13.7 shows the timing when the count-up occurs in the T2 cycle while writing to CMCNTH in bytes.

Figure 13.7 Conflict between Byte-Write and Count-Up Processes of CMCNT

13.5.4 Compare Match between CMCNT and CMCOR

Do not set the same value in CMCNT and CMCOR while CMCNT is not counting.

Section 14 Watchdog Timer (WDT)

This LSI includes the watchdog timer (WDT), which externally outputs an overflow signal (WDTOVF) on overflow of the counter when the value of the counter has not been updated because of a system malfunction. The WDT can simultaneously generate an internal reset signal for the entire LSI.

The WDT is a single channel timer for each of CPU0 and CPU1, and WDT0 for CPU0 counts up the clock oscillation settling period when the system leaves software standby mode or the temporary standby periods that occur when the clock frequency is changed. Both WDT0 and WDT1 can be used as a general watchdog timer or interval timer.

14.1 Features

- Can be used to ensure the clock oscillation settling time (WDT0) The WDT is used in leaving software standby mode or the temporary standby periods that occur when the clock frequency is changed.
- Can switch between watchdog timer mode and interval timer mode (WDT0, WDT1)
- Outputs WDTOVF signal in watchdog timer mode (WDT0, WDT1) When the counter overflows in watchdog timer mode, the WDTOVF signal is output externally. It is possible to select whether to reset the LSI internally when this happens. Either the power-on reset or manual reset signal can be selected as the internal reset type.
- Issues an interrupt in interval timer mode (WDT0, WDT1) An interval timer interrupt is issued when the counter overflows.
- Can select one of eight counter input clocks (WDT0, WDT1) Eight clocks ($P\phi \times 1$ to $P\phi \times 1/16384$) that are obtained by dividing the peripheral clock can be selected.

Figure 14.1 shows a block diagram of the WDT.

As shown in the figure, the reset output signals from WDT0 and WDT1 upon occurrence of an overflow are ORed and then output to both CPUs.

Figure 14.1 Block Diagram of WDT

14.2 Input/Output Pin

Table 14.1 shows the pin configuration of the WDT.

Table 14.1 Pin Configuration

14.3 Register Descriptions

The WDT has the following registers.

Table 14.2 Register Configuration

Note: * For the access size, see section 14.3.4, Notes on Register Access.

14.3.1 Watchdog Timer Counter (WTCNT0, WTCNT1)

WTCNT is an 8-bit readable/writable register that is incremented by cycles of the selected clock signal. When an overflow occurs, it generates a watchdog timer overflow signal (WDTOVF) in watchdog timer mode and an interrupt in interval timer mode.

Use word access to write to WTCNT with H'5A set in the upper byte. Use byte access to read from WTCNT.

Note: The method for writing to WTCNT differs from that for other registers to prevent erroneous writes. See section 14.3.4, Notes on Register Access, for details.

14.3.2 Watchdog Timer Control/Status Register (WTCSR0, WTCSR1)

WTCSR is an 8-bit readable/writable register composed of bits to select the clock used for the count, overflow flags, and timer enable bit.

Use word access to write to WTCSR with H'A5 set in the upper byte. Use byte access to read from WTCSR.

Note: The method for writing to WTCSR differs from that for other registers to prevent erroneous writes. See section 14.3.4, Notes on Register Access, for details.

14.3.3 Watchdog Reset Control/Status Register (WRCSR0, WRCSR1)

WRCSR is an 8-bit readable/writable register that controls output of the internal reset signal generated by watchdog timer counter (WTCNT) overflow.

Note: The method for writing to WRCSR differs from that for other registers to prevent erroneous writes. See section 14.3.4, Notes on Register Access, for details.

Notes: 1. 0 for WRCSR0 and 1 for WRCSR1

2. The LSI is not reset, but WTCNT and WTCSR in WDT are reset.

14.3.4 Notes on Register Access

The watchdog timer counter (WTCNT), watchdog timer control/status register (WTCSR), and watchdog reset control/status register (WRCSR) are more difficult to write to than other registers. The procedures for reading or writing to these registers are given below.

(1) Writing to WTCNT and WTCSR

These registers must be written by a word transfer instruction. They cannot be written by a byte or longword transfer instruction.

When writing to the WTCNT, set the upper byte to H^{'5}A and transfer the lower byte as the write data, as shown in figure 14.2. When writing to the WTCSR, set the upper byte to H'A5 and transfer the lower byte as the write data. This transfer procedure writes the lower byte data to the WTCNT or WTCSR.

Figure 14.2 Writing to WTCNT and WTCSR

(2) Writing to WRCSR

WRCSR must be written by a word access to its address. It cannot be written by byte transfer or longword transfer instructions.

Procedures for writing 0 to WOVF (bit 7) and for writing to RSTE (bit 6) and RSTS (bit 5) are different, as shown in figure 14.3.

To write 0 to the WOVF bit, the write data must be H'A5 in the upper byte and H'00 in the lower byte. This clears the WOVF bit to 0. The RSTE and RSTS bits are not affected. To write to the RSTE and RSTS bits, the upper byte must be H'5A and the lower byte must be the write data. The values of bits 6 and 5 of the lower byte are transferred to the RSTE and RSTS bits, respectively. The WOVF bit is not affected.

(3) Reading from WTCNT, WTCSR, and WRCSR

The registers of WDT0 are read in a method similar to other registers. WTCSR0 is allocated to address H'FFFE0000, WTCNT0 to address H'FFFE0002, and WRCSR0 to address H'FFFE0004. Byte transfer instructions must be used for reading from these registers.

The registers of WDT1 are read in a method similar to other registers. WTCSR1 is allocated to address H'FFFE0008, WTCNT1 to address H'FFFE000A, and WRCSR1 to address H'FFFE000C. Byte transfer instructions must be used for reading from these registers.

14.4 WDT Usage

14.4.1 Canceling Software Standby Mode

The WDT0 can be used to cancel software standby mode with an interrupt such as an NMI interrupt.

For details on the procedure, see section 30, Power-Down Modes.

14.4.2 Changing the PLL Multiplication Ratio

When changing the clock frequency by the PLL, use the WDT0. When changing the frequency only by switching the divider, do not use the WDT. For details on the procedure, see section 5, Clock Pulse Generator (CPG).

14.4.3 Using Watchdog Timer Mode

WDT0 should be used for the watchdog of CPU0 and WDT1 for the watchdog of CPU1.

- 1. Set the WT/ \overline{IT} bit in WTCSR to 1, the type of count clock in the CKS2 to CKS0 bits in WTCSR, whether this LSI is to be reset internally or not in the RSTE bit in WRCSR, the reset type if it is generated in the RSTS bit in WRCSR0, and the initial value of the counter in WTCNT.
- 2. Set the TME bit in WTCSR to 1 to start the count in watchdog timer mode.
- 3. While operating in watchdog timer mode, rewrite the counter periodically to H'00 to prevent the counter from overflowing. When the counter overflows, the WDT sets the WOVF flag in WRCSR to 1, and the WDTOVF signal is output outside the LSI (figure 14.4). The WDTOVF signal can be used to reset the system. The \overline{WDTOVF} signal is output for 64 \times P ϕ clock cycles.
- 4. If the RSTE bit in WRCSR is set to 1, a signal to reset the inside of this LSI can be generated simultaneously with the \overline{WDTOVF} signal. Either power-on reset or manual reset can be selected for this interrupt by the RSTS bit in WRCSR0. The internal reset signal is output for $128 \times$ P ϕ clock cycles.
- 5. When a WDT overflow reset is generated simultaneously with a reset input on the RES pin, the RES pin reset takes priority, and the WOVF bit in WRCSR is cleared to 0.

Figure 14.4 Operation in Watchdog Timer Mode

14.4.4 Using Interval Timer Mode

When operating in interval timer mode, interval timer interrupts are generated at every overflow of the counter. This enables interrupts to be generated at set periods.

- 1. Clear the WT/ \overline{IT} bit in WTCSR to 0, set the type of count clock in the CKS[2:0] bits in WTCSR, and set the initial value of the counter in WTCNT.
- 2. Set the TME bit in WTCSR to 1 to start the count in interval timer mode.
- 3. When the counter overflows, the WDT sets the IOVF bit in WTCSR to 1 and an interval timer interrupt request is sent to the INTC. The counter then resumes counting.

Figure 14.5 Operation in Interval Timer Mode

14.5 Usage Notes

Pay attention to the following points when using the WDT in either the interval timer or watchdog timer mode.

14.5.1 Timer Variation

After timer operation has started, the period from the power-on reset point to the first count up timing of WTCNT varies depending on the time period that is set by the TME bit of WTCSR. The shortest such time period is thus one cycle of the peripheral clock, Pφ, while the longest is the result of frequency division according to the value in the CKS[2:0] bits. The timing of subsequent increment is in accordance with the selected frequency division ratio. Accordingly, this time difference is referred to as timer variation.

14.5.2 Prohibition against Setting H'FF to WTCNT

When the value in WTCNT reaches H'FF, the WDT assumes that an overflow has occurred. Accordingly, when H'FF is set in WTCNT, an interval timer interrupt or WDT reset will occur immediately, regardless of the current clock selection by the CKS[2:0] bits.

14.5.3 Interval Timer Overflow Flag

The IOVF bit in WTCSR cannot be cleared when the value in WTCNT is H'FF. Clear the IOVF bit when the value in WTCNT is set to H'00 or the value other than H'FF is rewritten.

14.5.4 System Reset by WDTOVF **Signal**

If the WDTOVF signal is input to the RES pin of this LSI, this LSI cannot be initialized correctly.

Avoid input of the WDTOVF signal to the \overline{RES} pin of this LSI through glue logic circuits. To reset the entire system with the WDTOVF signal, use the circuit shown in figure 14.6.

Figure 14.6 Example of System Reset Circuit Using WDTOVF **Signal**

14.5.5 Manual Reset in Watchdog Timer Mode

When a manual reset occurs in watchdog timer mode, the bus cycle is continued. If a manual reset occurs during burst transfer by the DMAC, manual reset exception handling will be pended until the CPU acquires the bus mastership.

14.5.6 Transition to Deep Standby Mode

The WDT does not directly perform transition to or release of deep standby mode. However, since the WDT may generate a watchdog timer reset or interval timer interrupt during transition to deep standby mode by CPU0 issuing the SLEEP instruction, clear the WTCSR0.TME and WTCSR1.TME bits to 0 to stop the WDT before issuance of the SLEEP instruction.

14.5.7 Internal Reset in Watchdog Timer Mode

When an internal reset occurs in watchdog timer mode due to an overflow of the watchdog timer counter (WTCNT0), the watchdog reset control/status register (WRCSR0) is not initialized and the WOVF bit of WDT0 stays set to 1. While the WOVF bit of WDT0 is set to 1, WDT0 cannot generate an internal reset even when an WTCNT0 overflow occurs.

In like manner, when an internal reset occurs in watchdog timer mode due to an overflow of the watchdog timer counter (WTCNT1), the watchdog reset control/status register (WRCSR1) is not initialized and the WOVF bit of WDT1 stays set to 1. While the WOVF bit of WDT1 is set to 1, WDT1 cannot generate an internal reset even when an WTCNT1 overflow occurs.

Section 15 Realtime Clock (RTC)

This LSI has a realtime clock (RTC) with its own 32.768-kHz crystal oscillator.

15.1 Features

- Clock and calendar functions (BCD format): Seconds, minutes, hours, date, day of the week, month, and year
- 1-Hz to 64-Hz timer (binary format) 64-Hz counter indicates the state of the RTC divider circuit between 64 Hz and 1 Hz
- Start/stop function
- 30-second adjust function
- Alarm interrupt: Frame comparison of seconds, minutes, hours, date, day of the week, month, and year can be used as conditions for the alarm interrupt
- Periodic interrupts: the interrupt cycle may be $1/256$ second, $1/64$ second, $1/16$ second, $1/4$ second, 1/2 second, 1 second, or 2 seconds
- Carry interrupt: a carry interrupt indicates when a carry occurs during a counter read
- Automatic leap year adjustment

Figure 15.1 shows the block diagram of RTC.

Figure 15.1 RTC Block Diagram

15.2 Input/Output Pin

Table 15.1 shows the RTC pin configuration.

Table 15.1 Pin Configuration

15.3 Register Descriptions

The RTC has the following registers.

Table 15.2 Register Configuration

15.3.1 64-Hz Counter (R64CNT)

R64CNT indicates the state of the divider circuit between 64 Hz and 1 Hz.

Reading this register, when carry from 128-Hz divider stage is generated, sets the CF bit in the RTC control register 1 (RCR1) to 1 so that the carrying and reading 64 Hz counter are performed at the same time is indicated. In this case, the R64CNT should be read again after writing 0 to the CF bit in RCR1 since the read value is not valid.

After the RESET bit or ADJ bit in the RTC control register 2 (RCR2) is set to 1, the RTC divider circuit is initialized and R64CNT is initialized.

15.3.2 Second Counter (RSECCNT)

RSECCNT is used for setting/counting in the BCD-coded second section. The count operation is performed by a carry for each second of the 64-Hz counter.

The assignable range is from 00 through 59 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

15.3.3 Minute Counter (RMINCNT)

RMINCNT is used for setting/counting in the BCD-coded minute section. The count operation is performed by a carry for each minute of the second counter.

The assignable range is from 00 through 59 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

15.3.4 Hour Counter (RHRCNT)

RHRCNT is used for setting/counting in the BCD-coded hour section. The count operation is performed by a carry for each 1 hour of the minute counter.

The assignable range is from 00 through 23 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

15.3.5 Day of Week Counter (RWKCNT)

RWKCNT is used for setting/counting day of week section. The count operation is performed by a carry for each day of the date counter.

The assignable range is from 0 through 6 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

15.3.6 Date Counter (RDAYCNT)

RDAYCNT is used for setting/counting in the BCD-coded date section. The count operation is performed by a carry for each day of the hour counter.

The assignable range is from 01 through 31 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

The range of date changes with each month and in leap years. Confirm the correct setting. Leap years are recognized by dividing the year counter (RYRCNT) values by 400, 100, and 4 and obtaining a fractional result of 0.

15.3.7 Month Counter (RMONCNT)

RMONCNT is used for setting/counting in the BCD-coded month section. The count operation is performed by a carry for each month of the date counter.

The assignable range is from 01 through 12 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

15.3.8 Year Counter (RYRCNT)

RYRCNT is used for setting/counting in the BCD-coded year section. The count operation is performed by a carry for each year of the month counter.

The assignable range is from 0000 through 9999 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

15.3.9 Second Alarm Register (RSECAR)

RSECAR is an alarm register corresponding to the BCD coded second counter RSECCNT of the RTC. When the ENB bit is set to 1, a comparison with the RSECCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 00 through $59 + ENB$ bits (practically in BCD), otherwise operation errors occur.

15.3.10 Minute Alarm Register (RMINAR)

RMINAR is an alarm register corresponding to the minute counter RMINCNT. When the ENB bit is set to 1, a comparison with the RMINCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 00 through $59 + ENB$ bits (practically in BCD), otherwise operation errors occur.

15.3.11 Hour Alarm Register (RHRAR)

RHRAR is an alarm register corresponding to the BCD coded hour counter RHRCNT of the RTC. When the ENB bit is set to 1, a comparison with the RHRCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 00 through $23 + ENB$ bits (practically in BCD), otherwise operation errors occur.

15.3.12 Day of Week Alarm Register (RWKAR)

Initial

RWKAR is an alarm register corresponding to the BCD coded day of week counter RWKCNT. When the ENB bit is set to 1, a comparison with the RWKCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 0 through $6 + ENB$ bits (practically in BCD), otherwise operation errors occur.

15.3.13 Date Alarm Register (RDAYAR)

RDAYAR is an alarm register corresponding to the BCD coded date counter RDAYCNT. When the ENB bit is set to 1, a comparison with the RDAYCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 01 through $31 + ENB$ bits (practically in BCD), otherwise operation errors occur.

15.3.14 Month Alarm Register (RMONAR)

RMONAR is an alarm register corresponding to the BCD coded month counter RMONCNT. When the ENB bit is set to 1, a comparison with the RMONCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 01 through $12 + ENB$ bits (practically in BCD), otherwise operation errors occur.

15.3.15 Year Alarm Register (RYRAR)

RYRAR is an alarm register corresponding to the year counter RYRCNT. The assignable range is from 0000 through 9999 (practically in BCD), otherwise operation errors occur.

15.3.16 RTC Control Register 1 (RCR1)

RCR1 is a register that affects carry flags and alarm flags. It also selects whether to generate interrupts for each flag.

The CF flag remains undefined until the divider circuit is reset (the RESET and ADJ bits in RCR2 are set to 1). When using the CF flag, make sure to reset the divider circuit beforehand.

15.3.17 RTC Control Register 2 (RCR2)

RCR2 is a register for periodic interrupt control, 30-second adjustment ADJ, divider circuit RESET, and RTC count control.

RCR2 is initialized by a power-on reset or in deep standby mode. Bits other than the RTCEN and START bits are initialized by a manual reset.

15.3.18 RTC Control Register 3 (RCR3)

When the ENB bit is set to 1, RCR3 performs a comparison with the RYRCNT. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

15.4 Operation

RTC usage is shown below.

15.4.1 Initial Settings of Registers after Power-On

All the registers should be set after the power is turned on.

15.4.2 Setting Time

Figure 15.2 shows how to set the time when the clock is stopped.

Figure 15.2 Setting Time

15.4.3 Reading Time

Figure 15.3 shows how to read the time.

Figure 15.3 Reading Time

If a carry occurs while reading the time, the correct time will not be obtained, so it must be read again. Part (a) in figure 15.3 shows the method of reading the time without using interrupts; part (b) in figure 15.3 shows the method using carry interrupts. To keep programming simple, method (a) should normally be used.

15.4.4 Alarm Function

Figure 15.4 shows how to use the alarm function.

Figure 15.4 Using Alarm Function

Alarms can be generated using seconds, minutes, hours, day of the week, date, month, year, or any combination of these. Set the ENB bit in the register on which the alarm is placed to 1, and then set the alarm time in the lower bits. Clear the ENB bit in the register on which the alarm is not placed to 0.

When the clock and alarm times match, 1 is set in the AF bit in RCR1. Alarm detection can be checked by reading this bit, but normally it is done by interrupt. If 1 is set in the AIE bit in RCR1, an interrupt is generated when an alarm occurs.

The alarm flag is set when the clock and alarm times match. However, the alarm flag can be cleared by writing 0.

15.5 Usage Notes

15.5.1 Register Writing during RTC Count

The following RTC registers cannot be written to during an RTC count (while bit $0 = 1$ in RCR2).

RSECCNT, RMINCNT, RHRCNT, RDAYCNT, RWKCNT, RMONCNT, RYRCNT

The RTC count must be stopped before writing to any of the above registers.

15.5.2 Use of Real-time Clock (RTC) Periodic Interrupts

The method of using the periodic interrupt function is shown in figure 15.5.

A periodic interrupt can be generated periodically at the interval set by bits PES2 to PES0 in RCR2. When the time set by bits PES2 to PES0 has elapsed, the PEF is set to 1.

The PEF is cleared to 0 upon periodic interrupt generation or when bits PES2 to PES0 are set. Periodic interrupt generation can be confirmed by reading this bit, but normally the interrupt function is used.

15.5.3 Transition to Standby Mode after Setting Register

When a transition to standby mode is made after registers in the RTC are set, sometimes counting is not performed correctly. In case the registers are set, be sure to make a transition to standby mode after performing one dummy read the register.

15.5.4 Usage Notes when Writing to and Reading the Register

- When reading a counter register such as the second counter after having written to the register, follow the procedure in section 15.4.2, Setting Time. In this case, it is necessary to write to all the counters, from second to year, in succession. Do not read the counter registers during the write processing shown as (2) in figure 15.2.
- When reading the RCR2 register after having written to it, read the register after having dummy-read it twice. The value read in both dummy-read operations will be the value before writing. The written value will be reflected in the value read the third time.
- For other registers, written values are immediately reflected in read values.

Section 16 Serial Communication Interface with FIFO (SCIF)

This LSI has a six-channel serial communication interface with FIFO (SCIF) that supports both asynchronous and clock synchronous serial communication. It also has 16-stage FIFO registers for both transmission and reception independently for each channel that enable this LSI to perform efficient high-speed continuous communication.

16.1 Features

- Asynchronous serial communication:
	- ⎯ Serial data communication is performed by start-stop in character units. The SCIF can communicate with a universal asynchronous receiver/transmitter (UART), an asynchronous communication interface adapter (ACIA), or any other communications chip that employs a standard asynchronous serial system. There are eight selectable serial data communication formats.
	- Data length: 7 or 8 bits
	- Stop bit length: 1 or 2 bits
	- Parity: Even, odd, or none
	- Receive error detection: Parity, framing, and overrun errors
	- ⎯ Break detection: Break is detected when a framing error is followed by at least one frame at the space 0 level (low level). It is also detected by reading the RxD level directly from the serial port register when a framing error occurs.
- Clock synchronous serial communication (only channel 0, 1, 2, and 5):
	- ⎯ Serial data communication is synchronized with a clock signal. The SCIF can communicate with other chips having a clock synchronous communication function. There is one serial data communication format.
	- Data length: 8 bits
	- ⎯ Receive error detection: Overrun errors
- Full duplex communication: The transmitting and receiving sections are independent, so the SCIF can transmit and receive simultaneously. Both sections use 16-stage FIFO buffering, so high-speed continuous data transfer is possible in both the transmit and receive directions.
- On-chip baud rate generator with selectable bit rates
- Internal or external transmit/receive clock source: From either baud rate generator (internal) or SCK pin (external)
- Four types of interrupts: Transmit-FIFO-data-empty interrupt, break interrupt, receive-FIFOdata-full interrupt, and receive-error interrupts are requested independently.
- When the SCIF is not in use, it can be stopped by halting the clock supplied to it, saying power.
- In asynchronous mode, on-chip modem control functions (RTS) and \overline{CTS} (only channel 0).
- The quantity of data in the transmit and receive FIFO data registers and the number of receive errors of the receive data in the receive FIFO data register can be ascertained.
- A time-out error (DR) can be detected when receiving in asynchronous mode.
- In asynchronous mode, the base clock frequency can be either 16 or 8 times the bit rate.
- When an internal clock is selected as a clock source and the SCK pin is used as an input pin in asynchronous mode, either normal mode or double-speed mode can be selected for the baud rate generator.

Figure 16.1 shows a block diagram of the SCIF (for a single channel). Some channels have no external pins.

Figure 16.1 Block Diagram of SCIF

16.2 Input/Output Pins

Table 16.1 shows the pin configuration of the SCIF.

Table 16.1 Pin Configuration

16.3 Register Descriptions

The SCIF has the following registers.

Table 16.2 Register Configuration

Notes: 1. Only 0 can be written to clear the flag. Bits 15 to 8, 3, and 2 are read-only bits that cannot be modified.

 2. Only 0 can be written to clear the flag. Bits 15 to 1 are read-only bits that cannot be modified.

16.3.1 Receive Shift Register (SCRSR)

SCRSR receives serial data. Data input at the RxD pin is loaded into SCRSR in the order received, LSB (bit 0) first, converting the data to parallel form. When one byte has been received, it is automatically transferred to the receive FIFO data register (SCFRDR).

The CPU cannot read or write to SCRSR directly.

16.3.2 Receive FIFO Data Register (SCFRDR)

SCFRDR is a 16-byte FIFO register that stores serial receive data. The SCIF completes the reception of one byte of serial data by moving the received data from the receive shift register (SCRSR) into SCFRDR for storage. Continuous reception is possible until 16 bytes are stored. The CPU can read but not write to SCFRDR. If data is read when there is no receive data in the SCFRDR, the value is undefined.

When SCFRDR is full of receive data, subsequent serial data is lost.

16.3.3 Transmit Shift Register (SCTSR)

SCTSR transmits serial data. The SCIF loads transmit data from the transmit FIFO data register (SCFTDR) into SCTSR, then transmits the data serially from the TxD pin, LSB (bit 0) first. After transmitting one data byte, the SCIF automatically loads the next transmit data from SCFTDR into SCTSR and starts transmitting again.

The CPU cannot read from or write to SCTSR directly.

16.3.4 Transmit FIFO Data Register (SCFTDR)

SCFTDR is a 16-byte FIFO register that stores data for serial transmission. When the SCIF detects that the transmit shift register (SCTSR) is empty, it moves transmit data written in the SCFTDR into SCTSR and starts serial transmission. Continuous serial transmission is performed until there is no transmit data left in SCFTDR. The CPU can write to SCFTDR at all times.

When SCFTDR is full of transmit data (16 bytes), no more data can be written. If writing of new data is attempted, the data is ignored.

16.3.5 Serial Mode Register (SCSMR)

SCSMR specifies the SCIF serial communication format and selects the clock source for the baud rate generator.

The CPU can always read from and write to SCSMR.

16.3.6 Serial Control Register (SCSCR)

SCSCR operates the SCIF transmitter/receiver, enables/disables interrupt requests, and selects the transmit/receive clock source. The CPU can always read and write to SCSCR.

16.3.7 Serial Status Register (SCFSR)

SCFSR is a 16-bit register. The upper 8 bits indicate the number of receive errors in the receive FIFO data register, and the lower 8 bits indicate the status flag indicating SCIF operating state.

The CPU can always read and write to SCFSR, but cannot write 1 to the status flags (ER, TEND, TDFE, BRK, RDF, and DR). These flags can be cleared to 0 only if they have first been read (after being set to 1). The PER flag (bits 15 to 12 and bit 2) and the FER flag (bits 11 to 8 and bit 3) are read-only bits that cannot be written.

Note: $*$ Only 0 can be written to clear the flag after 1 is read.

Note: * Only 0 can be written to clear the flag after 1 is read.

16.3.8 Bit Rate Register (SCBRR)

SCBRR is an 8-bit register that is used with the CKS1 and CKS0 bits in the serial mode register (SCSMR) and the BGDM and ABCS bits in the serial extension mode register (SCEMR) to determine the serial transmit/receive bit rate.

The CPU can always read and write to SCBRR. SCBRR is initialized to H'FF by a power-on reset. Each channel has independent baud rate generator control, so different values can be set in three channels.

The SCBRR setting is calculated as follows:

• Asynchronous mode:

When baud rate generator operates in normal mode (when the BGDM bit of SCEMR is 0):

N =
$$
\frac{P\phi}{64 \times 2^{2n-1} \times B}
$$
 × 10⁶ – 1 (Operation on a base clock with a frequency of 16 times the bit rate)
\nN = $\frac{P\phi}{10^{6} - 1}$ (Operation on a base clock with a frequency of 8 times

 $N = \frac{F\psi}{20 \times 2^{2n+1} \times P} \times 10^6 - 1$ (Operation on a base clock with a frequency of 8 times $32 \times 2^{2n-1} \times B$ the bit rate)

When baud rate generator operates in double speed mode (when the BGDM bit of SCEMR is 1):

 $N = \frac{N}{200 \times 10^{6} \text{ m}} \times 10^{6} - 1$ (Operation on a base clock with a frequency of 16 times $32 \times 2^{2n-1} \times B$ the bit rate) Pφ

 $N = \frac{N}{\sqrt{2}} \times 10^6 - 1$ (Operation on a base clock with a frequency of 8 times $16 \times 2^{2n-1} \times B$ the bit rate) Pφ

• Clock synchronous mode:

$$
N = \frac{P\varphi}{8 \times 2^{2n-1} \times B} \times 10^6 - 1
$$

- B: Bit rate (bits/s)
- N: SCBRR setting for baud rate generator $(0 \le N \le 255)$ (The setting must satisfy the electrical characteristics.)
- Pφ: Operating frequency for peripheral modules (MHz)
- n: Baud rate generator clock source $(n = 0, 1, 2, 3)$ (for the clock sources and values of n, see table 16.3.)

Table 16.3 SCSMR Settings

The bit rate error in asynchronous mode is given by the following formula:

When baud rate generator operates in normal mode (the BGDM bit of SCEMR is 0):

Error
$$
(\%) = \left\{ \frac{P\phi \times 10^6}{(N+1) \times B \times 64 \times 2^{2n-1}} - 1 \right\} \times 100
$$
 (Operation on a base clock with a frequency of 16 times the bit rate)

\nError $(\%) = \left\{ \frac{P\phi \times 10^6}{(N+1) \times B \times 32 \times 2^{2n-1}} - 1 \right\} \times 100$ (Operation on a base clock with a frequency of 8 times the bit rate)

When baud rate generator operates in double speed mode (the BGDM bit of SCEMR is 1):

Error (%) =
$$
\left\{\frac{P\phi \times 10^6}{(N+1) \times B \times 32 \times 2^{2n-1}} - 1\right\} \times 100
$$
 (Operation on a base clock with a frequency of 16 times the bit rate)

Error (%) =
$$
\left\{\frac{P\phi \times 10^6}{(N+1) \times B \times 16 \times 2^{2n-1}} - 1\right\} \times 100
$$
 (Operation on a base clock with
a frequency of 8 times the bit rate)

Table 16.4 lists the sample SCBRR settings in asynchronous mode in which a base clock frequency is 16 times the bit rate (the ABCS bit in SCEMR is 0) and the baud rate generator operates in normal mode (the BGDM bit in SCEMR is 1), and table 16.5 lists the sample SCBRR settings in clock synchronous mode.

Table 16.4 Bit Rates and SCBRR Settings (Asynchronous Mode, BGDM = 0, ABCS = 0) (1)

	$P\phi$ (MHz)											
		12.288			14.7456			16			20	
Bit Rate (bit/s)	n	N	Error (%)	n	N	Error (%)	n	N	Error (%)	n	N	Error (%)
110	2	217	0.08	3	64	0.70	3	70	0.03	3	88	-0.25
150	2	159	0.00	2	191	0.00	2	207	0.16	3	64	0.16
300	\overline{c}	79	0.00	\overline{c}	95	0.00	2	103	0.16	2	129	0.16
600	1	159	0.00	1	191	0.00	1	207	0.16	2	64	0.16
1200	1	79	0.00	1	95	0.00	1	103	0.16	1	129	0.16
2400	0	159	0.00	0	191	0.00	0	207	0.16	1	64	0.16
4800	0	79	0.00	0	95	0.00	0	103	0.16	0	129	0.16
9600	0	39	0.00	0	47	0.00	Ω	51	0.16	0	64	0.16
19200	0	19	0.00	0	23	0.00	0	25	0.16	0	32	-1.36
31250	0	11	2.40	0	14	-1.70	Ω	15	0.00	0	19	0.00
38400	0	9	0.00	0	11	0.00	0	12	0.16	0	15	1.73

Table 16.4 Bit Rates and SCBRR Settings (Asynchronous Mode, BGDM = 0, ABCS = 0) (2)

Table 16.4 Bit Rates and SCBRR Settings (Asynchronous Mode, BGDM = 0, ABCS = 0) (4)

Note: Settings with an error of 1% or less are recommended.

Table 16.5 Bit Rates and SCBRR Settings (Clock Synchronous Mode)

[Legend]

Blank: No setting possible, or the electrical characteristics of the SH7205 cannot be satisfied regardless of the device being communicated with.

—: Setting possible, but error occurs

Table 16.6 indicates the maximum bit rates in asynchronous mode when the baud rate generator is used. Table 16.7 lists the maximum bit rates in asynchronous mode when the external clock input is used. Table 16.8 lists the maximum bit rates in clock synchronous mode when the external clock input is used (when $t_{S_{\text{cyc}}} = 12t_{\text{ncyc}}$ ^{*}).

Note: * Make sure that the electrical characteristics of this LSI and that of a connected LSI are satisfied.

Table 16.6 Maximum Bit Rates for Various Frequencies with Baud Rate Generator (Asynchronous Mode)

Table 16.7 Maximum Bit Rates with External Clock Input (Asynchronous Mode)

Table 16.8 Maximum Bit Rates with External Clock Input (Clock Synchronous Mode, $t_{\text{scyc}} = 12t_{\text{sec}}$)

16.3.9 FIFO Control Register (SCFCR)

SCFCR resets the quantity of data in the transmit and receive FIFO data registers, sets the trigger data quantity, and contains an enable bit for loop-back testing. SCFCR can always be read and written to by the CPU.

16.3.10 FIFO Data Count Set Register (SCFDR)

SCFDR is a 16-bit register which indicates the quantity of data stored in the transmit FIFO data register (SCFTDR) and the receive FIFO data register (SCFRDR).

It indicates the quantity of transmit data in SCFTDR with the upper 8 bits, and the quantity of receive data in SCFRDR with the lower 8 bits. SCFDR can always be read by the CPU.

16.3.11 Serial Port Register (SCSPTR)

The CPU can always read and write to SCSPTR.

SCSPTR controls input/output and data of pins multiplexed to SCIF function. Bits 7 and 6 can control input/output data of RTS pin. Bits 5 and 4 can control input/output data of CTS pin. Bits 3 and 2 can control input/output data of SCK pin. Bits 1 and 0 can input data from RxD pin and output data to TxD pin, so they control break of serial transmitting/receiving.

Bit Bit Name Initial Value R/W Description 15 to 8 — All 0 R Reserved These bits are always read as 0. The write value should always be 0. 7 RTSIO 0 R/W RTS Port Input/Output Indicates input or output of the serial port RTS pin. When the RTS pin is actually used as a port outputting the RTSDT bit value, the MCE bit in SCFCR should be cleared to 0. 0: RTSDT bit value not output to RTS pin 1: RTSDT bit value output to RTS pin 6 RTSDT 1 R/W RTS Port Data Indicates the input/output data of the serial port RTS pin. Input/output is specified by the RTSIO bit. For output, the RTSDT bit value is output to the RTS pin. The RTS pin status is read from the RTSDT bit regardless of the RTSIO bit setting. However, RTS input/output must be set in the PFC. 0: Input/output data is low level 1: Input/output data is high level

16.3.12 Line Status Register (SCLSR)

The CPU can always read or write to SCLSR, but cannot write 1 to the ORER flag. This flag can be cleared to 0 only if it has first been read (after being set to 1).

Note: * Only 0 can be written to clear the flag after 1 is read.

16.3.13 Serial Extension Mode Register (SCEMR)

The CPU can always read from or write to SCEMR. Setting the BGDM bit in this register to 1 allows the baud rate generator in the SCIF operates in double-speed mode when asynchronous mode is selected (by setting the C/\overline{A} bit in SCSMR to 0) and an internal clock is selected as a clock source and the SCK pin is set as an input pin (by setting the CKE[1:0] bits in SCSCR to 00).

The base clock frequency in asynchronous mode can be selected by modifying the ABCS bit setting.

16.4 Operation

16.4.1 Overview

For serial communication, the SCIF has an asynchronous mode in which characters are synchronized individually, and a clock synchronous mode in which communication is synchronized with clock pulses. However, clock synchronous mode is not available on channels 3 and 4.

The SCIF has a 16-stage FIFO buffer for both transmission and receptions, reducing the overhead of the CPU, and enabling continuous high-speed communication. Furthermore, channel 0 has RTS and CTS signals to be used as modem control signals.

The transmission format is selected in the serial mode register (SCSMR), as shown in table 16.9. The SCIF clock source is selected by the combination of the CKE1 and CKE0 bits in the serial control register (SCSCR), as shown in table 16.10.

(1) Asynchronous Mode

- Data length is selectable: 7 or 8 bits
- Parity bit is selectable. So is the stop bit length (1 or 2 bits). The combination of the preceding selections constitutes the communication format and character length.
- In receiving, it is possible to detect framing errors, parity errors, receive FIFO data full, overrun errors, receive data ready, and breaks.
- The number of stored data bytes is indicated for both the transmit and receive FIFO registers.
- An internal or external clock can be selected as the SCIF clock source.
	- ⎯ When an internal clock is selected, the SCIF operates using the clock of on-chip baud rate generator.
	- When an external clock is selected, the external clock input must have a frequency 16 or 8 times the bit rate. (The on-chip baud rate generator is not used.)

(2) Clock Synchronous Mode (Channels 0, 1, 2 and 5 only)

- The transmission/reception format has a fixed 8-bit data length.
- In receiving, it is possible to detect overrun errors (ORER).
- An internal or external clock can be selected as the SCIF clock source.
	- When an internal clock is selected, the SCIF operates using the clock of the on-chip baud rate generator, and outputs this clock to external devices as the synchronous clock.
	- When an external clock is selected, the SCIF operates on the input external synchronous clock not using the on-chip baud rate generator.

Table 16.9 SCSMR Settings and SCIF Communication Formats

[Legend]

x: Don't care

Table 16.10 SCSMR and SCSCR Settings and SCIF Clock Source Selection

[Legend]

x: Don't care

Note: When using the baud rate generator in double-speed mode (BGMD = 1), select asynchronous mode by setting the C/A bit to 0, and select an internal clock as a clock source and the SCK pin is not used (the CKE[1:0] bits set to 00).

16.4.2 Operation in Asynchronous Mode

In asynchronous mode, each transmitted or received character begins with a start bit and ends with a stop bit. Serial communication is synchronized one character at a time.

The transmitting and receiving sections of the SCIF are independent, so full duplex communication is possible. The transmitter and receiver are 16-byte FIFO buffered, so data can be written and read while transmitting and receiving are in progress, enabling continuous transmitting and receiving.

Figure 16.2 shows the general format of asynchronous serial communication.

In asynchronous serial communication, the communication line is normally held in the mark (high) state. The SCIF monitors the line and starts serial communication when the line goes to the space (low) state, indicating a start bit. One serial character consists of a start bit (low), data (LSB first), parity bit (high or low), and stop bit (high), in that order.

When receiving in asynchronous mode, the SCIF synchronizes at the falling edge of the start bit. The SCIF samples each data bit on the eighth or fourth pulse of a clock with a frequency 16 or 8 times the bit rate. Receive data is latched at the center of each bit.

Figure 16.2 Example of Data Format in Asynchronous Communication (8-Bit Data with Parity and Two Stop Bits)

(1) Transmit/Receive Formats

Table 16.11 lists the eight communication formats that can be selected in asynchronous mode. The format is selected by settings in the serial mode register (SCSMR).

[Legend]

START: Start bit

STOP: Stop bit

P: Parity bit

(2) Clock

An internal clock generated by the on-chip baud rate generator or an external clock input from the SCK pin can be selected as the SCIF transmit/receive clock. The clock source is selected by the C/\overline{A} bit in the serial mode register (SCSMR) and the CKE1 and CKE0 bits in the serial control register (SCSCR). For clock source selection, refer to table 16.10, SCSMR and SCSCR Settings and SCIF Clock Source Selection.

When an external clock is input at the SCK pin, it must have a frequency equal to 16 or 8 times the desired bit rate.

When the SCIF operates on an internal clock, it can output a clock signal on the SCK pin. The frequency of this output clock is 16 or 8 times the desired bit rate.

(3) Transmitting and Receiving Data

• SCIF Initialization (Asynchronous Mode)

Before transmitting or receiving, clear the TE and RE bits to 0 in the serial control register (SCSCR), then initialize the SCIF as follows.

When changing the operation mode or the communication format, always clear the TE and RE bits to 0 before following the procedure given below. Clearing TE to 0 initializes the transmit shift register (SCTSR). Clearing TE and RE to 0, however, does not initialize the serial status register (SCFSR), transmit FIFO data register (SCFTDR), or receive FIFO data register (SCFRDR), which retain their previous contents. Clear TE to 0 after all transmit data has been transmitted and the TEND flag in the SCFSR is set. The TE bit can be cleared to 0 during transmission, but the transmit data goes to the Mark state after the bit is cleared to 0. Set the TFRST bit in SCFCR to 1 and reset SCFTDR before TE is set again to start transmission.

When an external clock is used, the clock should not be stopped during initialization or subsequent operation. SCIF operation becomes unreliable if the clock is stopped.

Figure 16.3 shows a sample flowchart for initializing the SCIF.

Figure 16.3 Sample Flowchart for SCIF Initialization

• Transmitting Serial Data (Asynchronous Mode)

Figure 16.4 shows a sample flowchart for serial transmission.

Use the following procedure for serial data transmission after enabling the SCIF for transmission.

In serial transmission, the SCIF operates as described below.

- 1. When data is written into the transmit FIFO data register (SCFTDR), the SCIF transfers the data from SCFTDR to the transmit shift register (SCTSR) and starts transmitting. Confirm that the TDFE flag in the serial status register (SCFSR) is set to 1 before writing transmit data to SCFTDR. The number of data bytes that can be written is (16 – transmit trigger setting).
- 2. When data is transferred from SCFTDR to SCTSR and transmission is started, consecutive transmit operations are performed until there is no transmit data left in SCFTDR. When the number of transmit data bytes in SCFTDR falls below the transmit trigger number set in the FIFO control register (SCFCR), the TDFE flag is set. If the TIE bit in the serial control register (SCSR) is set to 1 at this time, a transmit-FIFO-data-empty interrupt (TXI) request is generated.

The serial transmit data is sent from the TxD pin in the following order.

- A. Start bit: One-bit 0 is output.
- B. Transmit data: 8-bit or 7-bit data is output in LSB-first order.
- C. Parity bit: One parity bit (even or odd parity) is output. (A format in which a parity bit is not output can also be selected.)
- D. Stop bit(s): One or two 1 bits (stop bits) are output.
- E. Mark state: 1 is output continuously until the start bit that starts the next transmission is sent.
- 3. The SCIF checks the SCFTDR transmit data at the timing for sending the stop bit. If data is present, the data is transferred from SCFTDR to SCTSR, the stop bit is sent, and then serial transmission of the next frame is started.

Figure 16.5 shows an example of the operation for transmission.

4. When modem control is enabled in channel 0, transmission can be stopped and restarted in accordance with the \overline{CTS} input value. When \overline{CTS} is set to 1, if transmission is in progress, the line goes to the mark state after transmission of one frame. When \overline{CTS} is set to 0, the next transmit data is output starting from the start bit.

Figure 16.6 shows an example of the operation when modem control is used.

Figure 16.6 Example of Operation Using Modem Control (CTS**)**

• Receiving Serial Data (Asynchronous Mode)

Figures 16.7 and 16.8 show sample flowcharts for serial reception.

Use the following procedure for serial data reception after enabling the SCIF for reception.

Figure 16.7 Sample Flowchart for Receiving Serial Data

Figure 16.8 Sample Flowchart for Receiving Serial Data (cont)

In serial reception, the SCIF operates as described below.

- 1. The SCIF monitors the transmission line, and if a 0 start bit is detected, performs internal synchronization and starts reception.
- 2. The received data is stored in SCRSR in LSB-to-MSB order.
- 3. The parity bit and stop bit are received.

After receiving these bits, the SCIF carries out the following checks.

- A. Stop bit check: The SCIF checks whether the stop bit is 1. If there are two stop bits, only the first is checked.
- B. The SCIF checks whether receive data can be transferred from the receive shift register (SCRSR) to SCFRDR.
- C. Overrun check: The SCIF checks that the ORER flag is 0, indicating that the overrun error has not occurred.
- D. Break check: The SCIF checks that the BRK flag is 0, indicating that the break state is not set.

If all the above checks are passed, the receive data is stored in SCFRDR.

Note: When a parity error or a framing error occurs, reception is not suspended.

4. If the RIE bit in SCSCR is set to 1 when the RDF or DR flag changes to 1, a receive-FIFOdata-full interrupt (RXI) request is generated. If the RIE bit or the REIE bit in SCSCR is set to 1 when the ER flag changes to 1, a receive-error interrupt (ERI) request is generated. If the RIE bit or the REIE bit in SCSCR is set to 1 when the BRK or ORER flag changes to 1, a break reception interrupt (BRI) request is generated.

Figure 16.9 shows an example of the operation for reception.

5. When modem control is enabled in channel 0, the RTS signal is output when SCFRDR is empty. When \overline{RTS} is 0, reception is possible. When \overline{RTS} is 1, this indicates that SCFRDR exceeds the number set for the RTS output active trigger.

Figure 16.10 shows an example of the operation when modem control is used.

Figure 16.10 Example of Operation Using Modem Control (RTS**)**

16.4.3 Operation in Clock Synchronous Mode

In clock synchronous mode, the SCIF transmits and receives data in synchronization with clock pulses. This mode is suitable for high-speed serial communication.

The SCIF transmitter and receiver are independent, so full-duplex communication is possible while sharing the same clock. The transmitter and receiver are also 16-byte FIFO buffered, so continuous transmitting or receiving is possible by reading or writing data while transmitting or receiving is in progress.

Figure 16.11 shows the general format in clock synchronous serial communication.

In clock synchronous serial communication, each data bit is output on the communication line from one falling edge of the serial clock to the next. Data is guaranteed valid at the rising edge of the serial clock.

In each character, the serial data bits are transmitted in order from the LSB (first) to the MSB (last). After output of the MSB, the communication line remains in the state of the MSB.

In clock synchronous mode, the SCIF receives data by synchronizing with the rising edge of the serial clock.

(1) Transmit/Receive Formats

The data length is fixed at eight bits. No parity bit can be added.

(2) Clock

An internal clock generated by the on-chip baud rate generator by the setting of the C/A bit in SCSMR and CKE[1:0] in SCSCR, or an external clock input from the SCK pin can be selected as the SCIF transmit/receive clock.

When the SCIF operates on an internal clock, it outputs the clock signal at the SCK pin. Eight clock pulses are output per transmitted or received character. When the SCIF is not transmitting or receiving, the clock signal remains in the high state. When only receiving, the clock signal outputs while the RE bit of SCSCR is 1 and the number of data in receive FIFO is more than the receive FIFO data trigger number.

(3) Transmitting and Receiving Data

• SCIF Initialization (Clock Synchronous Mode)

Before transmitting, receiving, or changing the mode or communication format, the software must clear the TE and RE bits to 0 in the serial control register (SCSCR), then initialize the SCIF. Clearing TE to 0 initializes the transmit shift register (SCTSR). Clearing RE to 0, however, does not initialize the RDF, PER, FER, and ORER flags and receive data register (SCRDR), which retain their previous contents.

Figure 16.12 shows a sample flowchart for initializing the SCIF.

- [1] Leave the TE and RE bits cleared to 0 until the initialization almost ends. Be sure to clear the TIE, RIE, TE, and RE bits to 0.
- [2] Set the data transfer format in SCSMR.
- Set CKE[1:0]. [3]
- Write a value corresponding to [4] the bit rate into SCBRR. This is not necessary if an external clock is used.
- [5] Sets PFC for external pins used. Set as RxD input at receiving and TxD at transmission.
- [6] Set the TE or RE bit in SCSCR to 1. Also set the TIE, RIE, and REIE bits to enable the TxD, RxD, and SCK pins to be used. When transmitting, the TxD pin will go to the mark state. When receiving in clocked synchronous mode with the synchronization clock output (clock master) selected, a clock starts to be output from the SCK pin at this point.

Figure 16.12 Sample Flowchart for SCIF Initialization

• Transmitting Serial Data (Clock Synchronous Mode)

Figure 16.13 shows a sample flowchart for transmitting serial data.

Use the following procedure for serial data transmission after enabling the SCIF for transmission.

Figure 16.13 Sample Flowchart for Transmitting Serial Data

In serial transmission, the SCIF operates as described below.

- 1. When data is written into the transmit FIFO data register (SCFTDR), the SCIF transfers the data from SCFTDR to the transmit shift register (SCTSR) and starts transmitting. Confirm that the TDFE flag in the serial status register (SCFSR) is set to 1 before writing transmit data to SCFTDR. The number of data bytes that can be written is (16 – transmit trigger setting).
- 2. When data is transferred from SCFTDR to SCTSR and transmission is started, consecutive transmit operations are performed until there is no transmit data left in SCFTDR. When the number of transmit data bytes in SCFTDR falls below the transmit trigger number set in the FIFO control register (SCFCR), the TDFE flag is set. If the TIE bit in the serial control register (SCSR) is set to 1 at this time, a transmit-FIFO-data-empty interrupt (TXI) request is generated.

If clock output mode is selected, the SCIF outputs eight synchronous clock pulses. If an external clock source is selected, the SCIF outputs data in synchronization with the input clock. Data is output from the TxD pin in order from the LSB (bit 0) to the MSB (bit 7).

- 3. The SCIF checks the SCFTDR transmit data at the timing for sending the MSB (bit 7). If data is present, the data is transferred from SCFTDR to SCTSR, and then serial transmission of the next frame is started. If there is no data, the TxD pin holds the state after the TEND flag in SCFSR is set to 1 and the MSB (bit 7) is sent.
- 4. After the end of serial transmission, the SCK pin is held in the high state.

Figure 16.14 shows an example of SCIF transmit operation.

Figure 16.14 Example of SCIF Transmit Operation

• Receiving Serial Data (Clock Synchronous Mode)

Figures 16.15 and 16.16 show sample flowcharts for receiving serial data. When switching from asynchronous mode to clock synchronous mode without SCIF initialization, make sure that ORER, PER, and FER are cleared to 0.

Figure 16.15 Sample Flowchart for Receiving Serial Data (1)

Figure 16.16 Sample Flowchart for Receiving Serial Data (2)

In serial reception, the SCIF operates as described below.

- 1. The SCIF synchronizes with serial clock input or output and starts the reception.
- 2. Receive data is shifted into SCRSR in order from the LSB to the MSB. After receiving the data, the SCIF checks the receive data can be loaded from SCRSR into SCFRDR or not. If this check is passed, the RDF flag is set to 1 and the SCIF stores the received data in SCFRDR. If the check is not passed (overrun error is detected), further reception is prevented.
- 3. After setting RDF to 1, if the receive FIFO data full interrupt enable bit (RIE) is set to 1 in SCSCR, the SCIF requests a receive-data-full interrupt (RXI). If the ORER bit is set to 1 and the receive-data-full interrupt enable bit (RIE) or the receive error interrupt enable bit (REIE) in SCSCR is also set to 1, the SCIF requests a break interrupt (BRI).

Figure 16.17 shows an example of SCIF receive operation.

Figure 16.17 Example of SCIF Receive Operation

• Transmitting and Receiving Serial Data Simultaneously (Clock Synchronous Mode) Figure 16.18 shows a sample flowchart for transmitting and receiving serial data simultaneously.

Use the following procedure for the simultaneous transmission/reception of serial data, after enabling the SCIF for transmission/reception.

16.5 SCIF Interrupts

The SCIF has four interrupt sources: transmit-FIFO-data-empty (TXI), receive-error (ERI), receive FIFO data full (RXI), and break (BRI).

Table 16.12 shows the interrupt sources and their order of priority. The interrupt sources are enabled or disabled by means of the TIE, RIE, and REIE bits in SCSCR. A separate interrupt request is sent to the interrupt controller for each of these interrupt sources.

When a TXI request is enabled by the TIE bit and the TDFE flag in the serial status register (SCFSR) is set to 1, a TXI interrupt request is generated. This TXI interrupt request activates the DMA transfer or the CPU interrupt to perform data transfer. The DMA transfer or the CPU interrupt is selectable by the DMA transfer request enable register (DREQER) of the INTC.

When an RXI request is enabled by the RIE bit and the RDF flag or the DR flag in SCFSR is set to 1, an RXI interrupt request is generated. This RXI interrupt request activates the DMA transfer or the CPU interrupt to perform data transfer. The DMA transfer or the CPU interrupt is selectable by the DMA transfer request enable register (DREQER) of the INTC. The RXI interrupt request caused by the DR flag is generated only in asynchronous mode.

When the RIE bit is set to 0 and the REIE bit is set to 1, the SCIF requests an ERI or BRI interrupt without requesting an RXI interrupt.

The TXI indicates that transmit data can be written, and the RXI indicates that there is receive data in SCFRDR.

Table 16.12 SCIF Interrupt Sources

16.6 Usage Notes

Note the following when using the SCIF.

16.6.1 SCFTDR Writing and TDFE Flag

The TDFE flag in the serial status register (SCFSR) is set when the number of transmit data bytes written in the transmit FIFO data register (SCFTDR) has fallen below the transmit trigger number set by bits TTRG[1:0] in the FIFO control register (SCFCR). After the TDFE flag is set, transmit data up to the number of empty bytes in SCFTDR can be written, allowing efficient continuous transmission.

However, if the number of data bytes written in SCFTDR is equal to or less than the transmit trigger number, the TDFE flag will be set to 1 again after being read as 1 and cleared to 0. TDFE flag clearing should therefore be carried out when SCFTDR contains more than the transmit trigger number of transmit data bytes.

The number of transmit data bytes in SCFTDR can be found from the upper 8 bits of the FIFO data count register (SCFDR).

16.6.2 SCFRDR Reading and RDF Flag

The RDF flag in the serial status register (SCFSR) is set when the number of receive data bytes in the receive FIFO data register (SCFRDR) has become equal to or greater than the receive trigger number set by bits RTRG[1:0] in the FIFO control register (SCFCR). After RDF flag is set, receive data equivalent to the trigger number can be read from SCFRDR, allowing efficient continuous reception.

However, if the number of data bytes in SCFRDR exceeds the trigger number, the RDF flag will be set to 1 again if it is cleared to 0. The RDF flag should therefore be cleared to 0 after being read as 1 after reading the number of the received data in the receive FIFO data register (SCFRDR) which is less than the trigger number.

The number of receive data bytes in SCFRDR can be found from the lower 8 bits of the FIFO data count register (SCFDR).

16.6.3 Restriction on DMAC Usage

When the DMAC writes data to SCFTDR due to a TXI interrupt request, the state of the TEND flag becomes undefined. Therefore, the TEND flag should not be used as the transfer end flag in such a case.

16.6.4 Break Detection and Processing

Break signals can be detected by reading the RxD pin directly when a framing error (FER) is detected. In the break state the input from the RxD pin consists of all 0s, so the FER flag is set and the parity error flag (PER) may also be set.

Note that, although transfer of receive data to SCFRDR is halted in the break state, the SCIF receiver continues to operate.

16.6.5 Sending a Break Signal

The I/O condition and level of the TxD pin are determined by the SPB2IO and SPB2DT bits in the serial port register (SCSPTR). This feature can be used to send a break signal.

Until TE bit is set to 1 (enabling transmission) after initializing, the TxD pin does not work. During the period, mark status is performed by the SPB2DT bit. Therefore, the SPB2IO and SPB2DT bits should be set to 1 (high level output).

To send a break signal during serial transmission, clear the SPB2DT bit to 0 (designating low level), then clear the TE bit to 0 (halting transmission). When the TE bit is cleared to 0, the transmitter is initialized regardless of the current transmission state, and 0 is output from the TxD pin.

16.6.6 Receive Data Sampling Timing and Receive Margin (Asynchronous Mode)

The SCIF operates on a base clock with a frequency 16 or 8 times the bit rate. In reception, the SCIF synchronizes internally with the fall of the start bit, which it samples on the base clock. Receive data is latched at the rising edge of the eighth or fourth base clock pulse. When the SCIF operates on a base clock with a frequency 16 times the bit rate, the receive data is sampled at the timing shown in figure 16.19.

Figure 16.19 Receive Data Sampling Timing in Asynchronous Mode (Operation on a Base Clock with a Frequency 16 Times the Bit Rate)

The receive margin in asynchronous mode can therefore be expressed as shown in equation 1.

Equation 1:

$$
M = \left| (0.5 - \frac{1}{2N}) - (L - 0.5) F - \frac{|D - 0.5|}{N} (1 + F) \right| \times 100\%
$$

Where: M: Receive margin $(\%)$

N: Ratio of clock frequency to bit rate $(N = 16$ or 8)

D: Clock duty $(D = 0$ to 1.0)

L: Frame length $(L = 9$ to 12)

F: Absolute deviation of clock frequency

From equation 1, if $F = 0$, $D = 0.5$ and $N = 16$, the receive margin is 46.875%, as given by equation 2.

Equation 2:

When $D = 0.5$ and $F = 0$: $M = (0.5 - 1/(2 \times 16)) \times 100\%$ $= 46.875\%$

This is a theoretical value. A reasonable margin to allow in system designs is 20% to 30%.

16.6.7 Selection of Base Clock in Asynchronous Mode

In this LSI, when asynchronous mode is selected, the base clock frequency within a bit period can be set to the frequency 16 or 8 times the bit rate by setting the ABCS bit in SCEMR.

Note that, however, if the base clock frequency 8 times the bit rate is used, receive margin is decreased as calculated using equation 1 in section 16.6.6, Receive Data Sampling Timing and Receive Margin (Asynchronous Mode).

If the desired bit rate can be set simply by setting SCBRR and the CKS1and CKS0 bits in SCSMR, it is recommended to use the base clock frequency within a bit period 16 times the bit rate (by setting the ABCS bit in SCEMR to 0). If an internal clock is selected as a clock source and the SCK pin is not used, the bit rate can be increased without decreasing receive margin by selecting double-speed mode for the baud rate generator (setting the BGDM bit in SCEMR to 1).

Section 17 Synchronous Serial Communication Unit (SSU)

This LSI has two synchronous serial communication unit (SSU) channels. The SSU has master mode in which this LSI outputs clocks as a master device for synchronous serial communication and slave mode in which clocks are input from an external device for synchronous serial communication. Synchronous serial communication can be performed with devices having different clock polarity and clock phase.

17.1 Features

- Choice of SSU mode and clock synchronous mode
- Choice of master mode and slave mode
- Choice of standard mode and bidirectional mode
- Synchronous serial communication with devices with different clock polarity and clock phase
- Choice of 8/16/32-bit width of transmit/receive data
- Full-duplex communication capability The shift register is incorporated, enabling transmission and reception to be executed simultaneously.
- Consecutive serial communication
- Choice of LSB-first or MSB-first transfer
- Choice of a clock source

Pφ/4, Pφ/8, Pφ/16, Pφ/32, Pφ/64, Pφ/128, Pφ/256, or an external clock

• Five interrupt sources

Transmit end, transmit data register empty, receive data full, overrun error, and conflict error. The direct memory access controller (DMAC) can be activated by a transmit data register empty request or a receive data full request to transfer data.

• Module standby mode can be set

To reduce power consumption, the operation of the SSU can be suspended by stopping the clock supply to the SSU.

Figure 17.1 shows a block diagram of the SSU.

Figure 17.1 Block Diagram of SSU

17.2 Input/Output Pins

Table 17.1 shows the SSU pin configuration.

Table 17.1 Pin Configuration

17.3 Register Descriptions

The SSU has the following registers. For details on the addresses of these registers and the states of these registers in each processing state, see section 32, List of Registers.

Table 17.2 Register Configuration

17.3.1 SS Control Register H (SSCRH)

SSCRH specifies master/slave device selection, bidirectional mode enable, SSO pin output value selection, SSCK pin selection, and \overline{SCS} pin selection.

17.3.2 SS Control Register L (SSCRL)

SSCRL selects operating mode, software reset, and transmit/receive data length.

17.3.3 SS Mode Register (SSMR)

SSMR selects the MSB first/LSB first, clock polarity, clock phase, and clock rate of synchronous serial communication.

Mar. 27, 2014

17.3.4 SS Enable Register (SSER)

SSER enables or disables transmission, reception, and interrupt requests.

17.3.5 SS Status Register (SSSR)

SSSR is a status flag register for interrupts.

17.3.6 SS Control Register 2 (SSCR2)

SSCR2 is a register that selects the assert timing of the SCS pin, data output timing of the SSO pin, and set timing of the TEND bit.

17.3.7 SS Transmit Data Registers 0 to 3 (SSTDR0 to SSTDR3)

SSTDR is an 8-bit register that stores transmit data. When 8-bit data length is selected by bits DATS1 and DATS0 in SSCRL, SSTDR0 is valid. When 16-bit data length is selected, SSTDR0 and SSTDR1 are valid. When 32-bit data length is selected, SSTDR0 to SSTDR3 are valid. The SSTDR that has not been enabled must not be accessed.

When the SSU detects that SSTRSR is empty, it transfers the transmit data written in SSTDR to SSTRSR and starts serial transmission. If the next transmit data has already been written to SSTDR during serial transmission, the SSU performs consecutive serial transmission.

Although SSTDR can always be read from or written to by the CPU and DMAC, to achieve reliable serial transmission, write transmit data to SSTDR after confirming that the TDRE bit in SSSR is set to 1.

DATS[1:0] (SSCRL[1:0])

17.3.8 SS Receive Data Registers 0 to 3 (SSRDR0 to SSRDR3)

SSRDR is an 8-bit register that stores receive data. When 8-bit data length is selected by bits DATS1 and DATS0 in SSCRL, SSRDR0 is valid. When 16-bit data length is selected, SSRDR0 and SSRDR1 are valid. When 32-bit data length is selected, SSRDR0 to SSRDR3 are valid. The SSRDR that has not been enabled must not be accessed.

When the SSU has received 1-byte data, it transfers the received serial data from SSTRSR to SSRDR where it is stored. After this, SSTRSR is ready for reception. Since SSTRSR and SSRDR function as a double buffer in this way, consecutive receive operations can be performed.

Read SSRDR after confirming that the RDRF bit in SSSR is set to 1.

SSRDR is a read-only register, therefore, cannot be written to by the CPU.

Table 17.4 Correspondence between DATS Bit Setting and SSRDR

17.3.9 SS Shift Register (SSTRSR)

SSTRSR is a shift register that transmits and receives serial data.

When data is transferred from SSTDR to SSTRSR, bit 0 of transmit data is bit 0 in the SSTDR contents ($MLS = 0$: LSB first communication) and is bit 7 in the SSTDR contents ($MLS = 1$: MSB first communication). The SSU transfers data from the LSB (bit 0) in SSTRSR to the SSO pin to perform serial data transmission.

In reception, the SSU sets serial data that has been input via the SSI pin in SSTRSR from the LSB (bit 0). When 1-byte data has been received, the SSTRSR contents are automatically transferred to SSRDR. SSTRSR cannot be directly accessed by the CPU.

17.4 Operation

17.4.1 Transfer Clock

A transfer clock can be selected from among seven internal clocks and an external clock. Before using this module, enable the SSCK pin function in the PFC. When the MSS bit in SSCRH is 1, an internal clock is selected and the SSCK pin is used as an output pin. When transfer is started, the clock with the transfer rate set by bits CKS2 to CKS0 in SSMR is output from the SSCK pin. When $MSS = 0$, an external clock is selected and the SSCK pin is used as an input pin.

17.4.2 Relationship of Clock Phase, Polarity, and Data

The relationship of clock phase, polarity, and transfer data depends on the combination of the CPOS and CPHS bits in SSMR when the value of the SSUMS bit in SSCRL is 0. Figure 17.2 shows the relationship. When SSUMS = 1, the CPHS setting is invalid although the CPOS setting is valid. Transmit data change timing and receive data fetch timing in $SSUMS = 1$ are the same timings shown in figure 17.2, (1) When CPHS = 0 .

Setting the MLS bit in SSMR selects either MSB first or LSB first communication. When MLS = 0, data is transferred from the LSB to the MSB. When $MLS = 1$, data is transferred from the MSB to the LSB.

Figure 17.2 Relationship of Clock Phase, Polarity, and Data

17.4.3 Relationship between Data Input/Output Pins and Shift Register

The connection between data input/output pins and the SS shift register (SSTRSR) depends on the combination of the MSS and BIDE bits in SSCRH and the SSUMS bit in SSCRL. Figure 17.3 shows the relationship.

The SSU transmits serial data from the SSO pin and receives serial data from the SSI pin when operating with $BIDE = 0$ and $MSS = 1$ (standard, master mode) (see figure 17.3 (1)). The SSU transmits serial data from the SSI pin and receives serial data from the SSO pin when operating with $BIDE = 0$ and $MSS = 0$ (standard, slave mode) (see figure 17.3 (2)).

The SSU transmits and receives serial data from the SSO pin regardless of master or slave mode when operating with $BIDE = 1$ (bidirectional mode) (see figures 17.3 (3) and (4)).

However, even if both the TE and RE bits are set to 1, transmission and reception are not performed simultaneously. Either the TE or RE bit must be selected.

The SSU transmits serial data from the SSO pin and receives serial data from the SSI pin when operating with $SSUMS = 1$. The SSCK pin outputs the internal clock when $MSS = 1$ and function as an input pin when $MSS = 0$ (see figures 17.3 (5) and (6)).

Figure 17.3 Relationship between Data Input/Output Pins and the Shift Register

17.4.4 Communication Modes and Pin Functions

The SSU switches the input/output pin (SSI, SSO, SSCK, and \overline{SCS}) functions according to the communication modes and register settings. The relationship of communication modes and input/output pin functions are shown in tables 17.5 to 17.7.

[Legend]

 $\overline{-}$: Not used as SSU pin (but can be used as an I/O port)

Table 17.6 Communication Modes and Pin States of SSCK Pin

Table 17.7 Communication Modes and Pin States of SCS **Pin**

[Legend]

x: Don't care

 $\overline{}$: Not used as SSU pin (but can be used as an I/O port)

17.4.5 SSU Mode

In SSU mode, data communications are performed via four lines: clock line (SSCK), data input line (SSI or SSO), data output line (SSI or SSO), and chip select line (SCS).

In addition, the SSU supports bidirectional mode in which a single pin functions as data input and data output lines.

(1) Initial Settings in SSU Mode

Figure 17.4 shows an example of the initial settings in SSU mode. Before data transfer, clear both the TE and RE bits in SSER to 0 to set the initial values.

Note: Before changing operating modes and communications formats, clear both the TE and RE bits to 0. Although clearing the TE bit to 0 sets the TDRE bit to 1, clearing the RE bit to 0 does not change the values of the RDRF and ORER bits and SSRDR. Those bits retain the previous values.

Figure 17.4 Example of Initial Settings in SSU Mode

(2) Data Transmission

Figure 17.5 shows an example of transmission operation, and figure 17.6 shows a flowchart example of data transmission.

When transmitting data, the SSU operates as shown below.

In master mode, the SSU outputs a transfer clock and data. In slave mode, when a low level signal is input to the SCS pin and a transfer clock is input to the SSCK pin, the SSU outputs data in synchronization with the transfer clock.

Writing transmit data to SSTDR after the TE bit is set to 1 clears the TDRE bit in SSSR to 0, and the SSTDR contents are transferred to SSTRSR. After that, the SSU sets the TDRE bit to 1 and starts transmission. At this time, if the TIE bit in SSER is set to 1, an SSTXI interrupt in the transmit data empty state is generated.

When 1-frame data has been transferred with $TDRE = 0$, the SSTDR contents are transferred to SSTRSR to start the next frame transmission. When the 8th bit of transmit data has been transferred with $TDRE = 1$, the TEND bit in SSSR is set to 1 and the state is retained. At this time, if the TEIE bit is set to 1, an SSTXI interrupt at the end of transmission is generated. After transmission, the output level of the SSCK pin is fixed high when $CPOS = 0$ and low when CPOS $= 1$

While the ORER bit in SSSR is set to 1, transmission is not performed. Check that the ORER bit is cleared to 0 before transmission.

Figure 17.5 Example of Transmission Operation (SSU Mode)

Figure 17.6 Flowchart Example of Data Transmission (SSU Mode)

(3) Data Reception

Figure 17.7 shows an example of reception operation, and figure 17.8 shows a flowchart example of data reception. When receiving data, the SSU operates as shown below.

After setting the RE bit to 1 and dummy-reading SSRDR, the SSU starts data reception.

In master mode, the SSU outputs a transfer clock and receives data. In slave mode, when a low level signal is input to the SCS pin and a transfer clock is input to the SSCK pin, the SSU receives data in synchronization with the transfer clock.

When 1-frame data has been received, the RDRF bit in SSSR is set to 1 and the receive data is stored in SSRDR. At this time, if the RIE bit in SSER is set to 1, an SSRXI interrupt is generated. The RDRF bit is automatically cleared to 0 by reading SSRDR.

In continuous reception as the slave device in SSU mode, be sure to read the SS receive data register (SSRDR) before reception of the next frame starts (before the externally connected master device starts the next transmission). If reception of a next frame starts while the receive data full (RDRF) bit in the SS status register (SSSR) is set (to 1) because SSRDR has not yet been read, and SSRDR is then read before reception of the next frame is complete, the conflict/incomplete error (CE) bit in SSRDR will be set at the end of reception of the next frame. Furthermore, if reception of a next frame starts after RDRF has been set (to 1) but before SSRDR has been read, and SSRDR still has not been read by the end of reception of the next frame, neither the CE nor the overrun error (ORER) bit in SSSR will be set, but the received data will be discarded.

Figure 17.7 Example of Reception Operation (SSU Mode)

Figure 17.8 Flowchart Example of Data Reception (SSU Mode)

(4) Data Transmission/Reception

Figure 17.9 shows a flowchart example of simultaneous transmission/reception. The data transmission/reception is performed combining the data transmission and data reception as mentioned above. The data transmission/reception is started by writing transmit data to SSTDR with $TE = RE = 1$. When the RDRF bit has been set to 1 at the 8th rising edge of the transfer clock, the ORER bit in SSSR is set to 1. This indicates that an overrun error (SSERI) has occurred, at this time, the data transmission/reception is stopped. While the ORER bit in SSSR is set to 1, transmission/reception is not performed. To resume the transmission/reception, clear the ORER bit to 0 .

Before switching transmission mode (TE = 1) or reception mode (RE = 1) to transmission/reception mode (TE = $RE = 1$), clear the TE and RE bits to 0. When starting the transfer, confirm that the TEND, RDRF, and ORER bits are cleared to 0 before setting the TE or RE bit to 1.

Figure 17.9 Flowchart Example of Simultaneous Transmission/Reception (SSU Mode)

17.4.6 SCS **Pin Control and Conflict Error**

When bits CSS1 and CSS0 in SSCRH are specified to B'10 and the SSUMS bit in SSCRL is cleared to 0, the SCS pin functions as an input (Hi-Z) to detect a conflict error. A conflict error detection period is from setting the MSS bit in SSCRH to 1 to starting serial transfer and after transfer ends. When a low level signal is input to the SCS pin within the period, a conflict error occurs. At this time, the CE bit in SSSR is set to 1 and the MSS bit is cleared to 0.

Note: While the CE bit is set to 1, transmission or reception is not resumed. Clear the CE bit to 0 before resuming the transmission or reception.

Figure 17.10 Conflict Error Detection Timing (Before Transfer)

Figure 17.11 Conflict Error Detection Timing (After Transfer End)

17.4.7 Clock Synchronous Communication Mode

In clock synchronous communication mode, data communications are performed via three lines: clock line (SSCK), data input line (SSI), and data output line (SSO).

(1) Initial Settings in Clock Synchronous Communication Mode

Figure 17.12 shows an example of the initial settings in clock synchronous communication mode. Before data transfer, clear both the TE and RE bits in SSER to 0 to set the initial values.

Note: Before changing operating modes and communications formats, clear both the TE and RE bits to 0. Although clearing the TE bit to 0 sets the TDRE bit to 1, clearing the RE bit to 0 does not change the values of the RDRF and ORER bits and SSRDR. Those bits retain the previous values.

Figure 17.12 Example of Initial Settings in Clock Synchronous Communication Mode

(2) Data Transmission

Figure 17.13 shows an example of transmission operation, and figure 17.14 shows a flowchart example of data transmission. When transmitting data in clock synchronous communication mode, the SSU operates as shown below.

In master mode, the SSU outputs a transfer clock and data. In slave mode, when a transfer clock is input to the SSCK pin, the SSU outputs data in synchronization with the transfer clock.

Writing transmit data to SSTDR after the TE bit is set to 1 clears the TDRE bit in SSSR to 0, and the SSTDR contents are transferred to SSTRSR. After that, the SSU sets the TDRE bit to 1 and starts transmission. At this time, if the TIE bit in SSER is set to 1, an SSTXI interrupt in the transmit data empty state is generated.

When 1-frame data has been transferred with $TDRE = 0$, the SSTDR contents are transferred to SSTRSR to start the next frame transmission. When the 8th bit of transmit data has been transferred with $TDRE = 1$, the TEND bit in SSSR is set to 1 and the state is retained. At this time, if the TEIE bit is set to 1, an SSTXI interrupt at the end of transmission is generated.

While the ORER bit in SSSR is set to 1, transmission is not performed. Check that the ORER bit is cleared to 0 before transmission.

Figure 17.14 Flowchart Example of Transmission Operation (Clock Synchronous Communication Mode)

(3) Data Reception

Figure 17.15 shows an example of reception operation, and figure 17.16 shows a flowchart example of data reception. When receiving data, the SSU operates as shown below.

After setting the RE bit in SSER to 1, the SSU starts data reception.

In master mode, the SSU outputs a transfer clock and receives data. In slave mode, when a transfer clock is input to the SSCK pin, the SSU receives data in synchronization with the transfer clock.

When 1-frame data has been received, the RDRF bit in SSSR is set to 1 and the receive data is stored in SSRDR. At this time, if the RIE bit is set to 1, an SSRXI interrupt is generated. The RDRF bit is automatically cleared to 0 by reading SSRDR.

When the RDRF bit has been set to 1 at the 8th rising edge of the transfer clock, the ORER bit in SSSR is set to 1. This indicates that an overrun error (SSERI) has occurred. At this time, data reception is stopped. While the ORER bit in SSSR is set to 1, reception is not performed. To resume the reception, clear the ORER bit to 0.

Figure 17.16 Flowchart Example of Data Reception (Clock Synchronous Communication Mode)

(4) Data Transmission/Reception

Figure 17.17 shows a flowchart example of simultaneous transmission/reception. The data transmission/reception is performed combining the data transmission and data reception as mentioned above. The data transmission/reception is started by writing transmit data to SSTDR with $TE = RE = 1$. When the RDRF bit has been set to 1 at the 8th rising edge of the transfer clock, the ORER bit in SSSR is set to 1. This indicates that an overrun error (SSERI) has occurred. At this time, data transmission/reception is stopped. While the ORER bit in SSSR is set to 1, transmission/reception is not performed. To resume the transmission/reception, clear the ORER bit to 0 .

Before switching transmission mode (TE = 1) or reception mode (RE = 1) to transmission/reception mode (TE = $RE = 1$), clear the TE and RE bits to 0. When starting the transfer, confirm that the TEND, RDRF, and ORER bits are cleared to 0 before setting the TE or RE bits to 1.

Figure 17.17 Flowchart Example of Simultaneous Transmission/Reception (Clock Synchronous Communication Mode)

17.5 SSU Interrupt Sources and DMAC

The SSU interrupt requests are an overrun error, a conflict error, a receive data register full, transmit data register empty, and a transmit end interrupts. Of these interrupt sources, a receive data register full, and a transmit data register empty can activate the DMAC for data transfer.

Since both an overrun error and a conflict error interrupts are allocated to the SSERI vector address, and both a transmit data register empty and a transmit end interrupts are allocated to the SSTXI vector address, the interrupt source should be decided by their flags. Table 17.8 lists the interrupt sources.

When an interrupt condition shown in table 17.8 is satisfied, an interrupt is requested. Clear the interrupt source by CPU or DMAC data transfer.

Table 17.8 SSU Interrupt Sources

17.6 Usage Note

17.6.1 Module Standby Mode Setting

The SSU operation can be disabled or enabled using the standby control register. The initial setting is for SSU operation to be halted. Access to registers is enabled by clearing module standby mode. For details, refer to section 30, Power-Down Modes.

17.6.2 Consecutive Data Transmission/Reception in SSU Slave Mode

In the continuous reception or transmission of data in SSU slave mode, the SCS pin is negated (placed at the high level) after every frame. If the SCS pin remains asserted (low level) over a period longer than that for one frame, transmission or reception was not correct.

17.6.3 Note in the Master Transmission Operation or the Master Transmission/Reception Operation of SSU Mode

In the master transmission operation or the master transmission/reception operation of SSU mode, please operate one of the following three ways.

- (1) Write the next transmission data to the SSTDR after the TDRE bit of the SSSR is set to 1 and before the transmission of one bit before the last bit starts.
- (2) Write the next transmission data to the SSTDR after the TEND bit of the SSSR is set to 1.
- (3) Set the SSCR2 as "TENDSTS = 0 " or "TENDSTS = 1 and SCSATS = 1 ".

Section 18 I^2C Bus Interface 3 (IIC3)

The I^2C bus interface 3 conforms to and provides a subset of the Philips I^2C (Inter-IC) bus interface functions. However, the configuration of the registers that control the $I²C$ bus differs partly from the Philips register configuration.

The I^2C bus interface 3 has four channels.

18.1 Features

- Selection of I^2C format or clocked synchronous serial format
- Continuous transmission/reception

Since the shift register, transmit data register, and receive data register are independent from each other, the continuous transmission/reception can be performed.

I 2 C bus format:

- Start and stop conditions generated automatically in master mode
- Selection of acknowledge output levels when receiving
- Automatic loading of acknowledge bit when transmitting
- Bit synchronization function

In master mode, the state of SCL is monitored per bit, and the timing is synchronized automatically. If transmission/reception is not yet possible, set the SCL to low until preparations are completed.

• Six interrupt sources

Transmit data empty (including slave-address match), transmit end, receive data full (including slave-address match), arbitration lost, NACK detection, and stop condition detection

- The direct memory access controller (DMAC) can be activated by a transmit-data-empty request or receive-data-full request to transfer data.
- Direct bus drive

Two pins, SCL and SDA pins, function as NMOS open-drain outputs when the bus drive function is selected.

Clocked synchronous serial format:

• Four interrupt sources

Transmit-data-empty, transmit-end, receive-data-full, and overrun error

• The direct memory access controller (DMAC) can be activated by a transmit-data-empty request or receive-data-full request to transfer data.

Figure 18.1 shows a block diagram of the $I²C$ bus interface 3.

Figure 18.1 Block Diagram of I²C Bus Interface 3

18.2 Input/Output Pins

Table 18.1 shows the pin configuration of the $I²C$ bus interface 3.

Table 18.1 Pin Configuration

Pin Name	Symbol	I/O	Function
Serial clock	SCL0 to SCL3	I/O	I ² C serial clock input/output
Serial data	SDA0 to SDA3	I/O	I ² C serial data input/output

Figure 18.2 shows an example of I/O pin connections to external circuits.

Figure 18.2 External Circuit Connections of I/O Pins

18.3 Register Descriptions

The I^2C bus interface 3 has the following registers.

Table 18.2 Register Configuration

j.

18.3.1 I 2 C Bus Control Register 1 (ICCR1)

ICCR1 is an 8-bit readable/writable register that enables or disables the I^2C bus interface 3, controls transmission or reception, and selects master or slave mode, transmission or reception, and transfer clock frequency in master mode.

ICCR1 is initialized to H'00 by a power-on reset.

Table 18.3 Transfer Rate

Note: The settings should satisfy external specifications.

18.3.2 I 2 C Bus Control Register 2 (ICCR2)

ICCR2 is an 8-bit readable/writable register that issues start/stop conditions, manipulates the SDA pin, monitors the SCL pin, and controls reset in the control part of the $I²C$ bus.

18.3.3 I 2 C Bus Mode Register (ICMR)

ICMR is an 8-bit readable/writable register that selects whether the MSB or LSB is transferred first, performs master mode wait control, and selects the transfer bit count.

Bits BC[2:0] are initialized to H'0 by the IICRST bit in ICCR2.

18.3.4 I 2 C Bus Interrupt Enable Register (ICIER)

ICIER is an 8-bit readable/writable register that enables or disables interrupt sources and acknowledge bits, sets acknowledge bits to be transferred, and confirms acknowledge bits received.

18.3.5 I 2 C Bus Status Register (ICSR)

ICSR is an 8-bit readable/writable register that confirms interrupt request flags and their status.

18.3.6 Slave Address Register (SAR)

SAR is an 8-bit readable/writable register that selects the communications format and sets the slave address. In slave mode with the I^2C bus format, if the upper seven bits of SAR match the upper seven bits of the first frame received after a start condition, this module operates as the slave device.

18.3.7 I 2 C Bus Transmit Data Register (ICDRT)

ICDRT is an 8-bit readable/writable register that stores the transmit data. When ICDRT detects the space in the shift register (ICDRS), it transfers the transmit data which is written in ICDRT to ICDRS and starts transferring data. If the next transfer data is written to ICDRT while transferring data of ICDRS, continuous transfer is possible.

18.3.8 I 2 C Bus Receive Data Register (ICDRR)

ICDRR is an 8-bit register that stores the receive data. When data of one byte is received, ICDRR transfers the receive data from ICDRS to ICDRR and the next data can be received. ICDRR is a receive-only register, therefore the CPU cannot write to this register.

18.3.9 I 2 C Bus Shift Register (ICDRS)

ICDRS is a register that is used to transfer/receive data. In transmission, data is transferred from ICDRT to ICDRS and the data is sent from the SDA pin. In reception, data is transferred from ICDRS to ICDRR after data of one byte is received. This register cannot be read directly from the CPU.

18.3.10 NF2CYC Register (NF2CYC)

NF2CYC is an 8-bit readable/writable register that selects the range of the noise filtering for the SCL and SDA pins. For details of the noise filter, see section 18.4.7, Noise Filter.

18.4 Operation

The I^2C bus interface 3 can communicate either in I^2C bus mode or clocked synchronous serial mode by setting FS in SAR.

18.4.1 I 2 C Bus Format

Figure 18.3 shows the I^2C bus formats. Figure 18.4 shows the I^2C bus timing. The first frame following a start condition always consists of eight bits.

Figure 18.3 I2 C Bus Formats

Figure 18.4 I²C Bus Timing

[Legend]

- S: Start condition. The master device drives SDA from high to low while SCL is high.
- SLA: Slave address
- R/\overline{W} : Indicates the direction of data transfer: from the slave device to the master device when R/W is 1, or from the master device to the slave device when R/W is 0.
- A: Acknowledge. The receive device drives SDA to low.

DATA: Transfer data

P: Stop condition. The master device drives SDA from low to high while SCL is high.
18.4.2 Master Transmit Operation

In master transmit mode, the master device outputs the transmit clock and transmit data, and the slave device returns an acknowledge signal. For master transmit mode operation timing, refer to figures 18.5 and 18.6. The transmission procedure and operations in master transmit mode are described below.

- 1. Set the ICE bit in ICCR1 to 1. Also, set bits CKS[3:0] in ICCR1. (Initial setting)
- 2. Read the BBSY flag in ICCR2 to confirm that the bus is released. Set the MST and TRS bits in ICCR1 to select master transmit mode. Then, write 1 to BBSY and 0 to SCP. (Start condition issued) This generates the start condition.
- 3. After confirming that TDRE in ICSR has been set, write the transmit data (the first byte data show the slave address and R/\overline{W}) to ICDRT. At this time, TDRE is automatically cleared to 0, and data is transferred from ICDRT to ICDRS. TDRE is set again.
- 4. When transmission of one byte data is completed while TDRE is 1, TEND in ICSR is set to 1 at the rise of the 9th transmit clock pulse. Read the ACKBR bit in ICIER, and confirm that the slave device has been selected. Then, write second byte data to ICDRT. When ACKBR is 1, the slave device has not been acknowledged, so issue the stop condition. To issue the stop condition, write 0 to BBSY and SCP. SCL is fixed low until the transmit data is prepared or the stop condition is issued.
- 5. The transmit data after the second byte is written to ICDRT every time TDRE is set.
- 6. Write the number of bytes to be transmitted to ICDRT. Wait until TEND is set (the end of last byte data transmission) while TDRE is 1, or wait for NACK (NACKF in ICSR = 1) from the receive device while ACKE in ICIER is 1. Then, issue the stop condition to clear TEND or NACKF.
- 7. When the STOP bit in ICSR is set to 1, the operation returns to the slave receive mode.

Figure 18.5 Master Transmit Mode Operation Timing (1)

Figure 18.6 Master Transmit Mode Operation Timing (2)

18.4.3 Master Receive Operation

In master receive mode, the master device outputs the receive clock, receives data from the slave device, and returns an acknowledge signal. For master receive mode operation timing, refer to figures 18.7 and 18.8. The reception procedure and operations in master receive mode are shown below.

- 1. Clear the TEND bit in ICSR to 0, then clear the TRS bit in ICCR1 to 0 to switch from master transmit mode to master receive mode. Then, clear the TDRE bit to 0.
- 2. When ICDRR is read (dummy data read), reception is started, and the receive clock is output, and data received, in synchronization with the internal clock. The master device outputs the level specified by ACKBT in ICIER to SDA, at the 9th receive clock pulse.
- 3. After the reception of first frame data is completed, the RDRF bit in ICSR is set to 1 at the rise of 9th receive clock pulse. At this time, the receive data is read by reading ICDRR, and RDRF is cleared to 0.
- 4. The continuous reception is performed by reading ICDRR every time RDRF is set. If 8th receive clock pulse falls after reading ICDRR by the other processing while RDRF is 1, SCL is fixed low until ICDRR is read.
- 5. If next frame is the last receive data, set the RCVD bit in ICCR1 to 1 before reading ICDRR. This enables the issuance of the stop condition after the next reception.
- 6. When the RDRF bit is set to 1 at rise of the 9th receive clock pulse, issue the stage condition.
- 7. When the STOP bit in ICSR is set to 1, read ICDRR. Then clear the RCVD bit to 0.
- 8. The operation returns to the slave receive mode.
- Note: If only one byte is received, read ICDRR (dummy-read) after the RCVD bit in ICCR1 is set.

Figure 18.7 Master Receive Mode Operation Timing (1)

18.4.4 Slave Transmit Operation

In slave transmit mode, the slave device outputs the transmit data, while the master device outputs the receive clock and returns an acknowledge signal. For slave transmit mode operation timing, refer to figures 18.9 and 18.10.

The transmission procedure and operations in slave transmit mode are described below.

- 1. Set the ICE bit in ICCR1 to 1. Set bits CKS[3:0] in ICCR1. (Initial setting) Set the MST and TRS bits in ICCR1 to select slave receive mode, and wait until the slave address matches.
- 2. When the slave address matches in the first frame following detection of the start condition, the slave device outputs the level specified by ACKBT in ICIER to SDA, at the rise of the 9th clock pulse. At this time, if the 8th bit data (R/W) is 1, the TRS bit in ICCR1 and the TDRE bit in ICSR are set to 1, and the mode changes to slave transmit mode automatically. The continuous transmission is performed by writing transmit data to ICDRT every time TDRE is set.
- 3. If TDRE is set after writing last transmit data to ICDRT, wait until TEND in ICSR is set to 1, with TDRE = 1. When TEND is set, clear TEND.
- 4. Clear TRS for the end processing, and read ICDRR (dummy read). SCL is opened.
- 5. Clear TDRE.

Figure 18.9 Slave Transmit Mode Operation Timing (1)

Figure 18.10 Slave Transmit Mode Operation Timing (2)

18.4.5 Slave Receive Operation

In slave receive mode, the master device outputs the transmit clock and transmit data, and the slave device returns an acknowledge signal. For slave receive mode operation timing, refer to figures 18.11 and 18.12. The reception procedure and operations in slave receive mode are described below.

- 1. Set the ICE bit in ICCR1 to 1. Set bits CKS[3:0] in ICCR1. (Initial setting) Set the MST and TRS bits in ICCR1 to select slave receive mode, and wait until the slave address matches.
- 2. When the slave address matches in the first frame following detection of the start condition, the slave device outputs the level specified by ACKBT in ICIER to SDA, at the rise of the 9th clock pulse. At the same time, RDRF in ICSR is set to read ICDRR (dummy read). (Since the read data show the slave address and R/W, it is not used.)
- 3. Read ICDRR every time RDRF is set. If 8th receive clock pulse falls while RDRF is 1, SCL is fixed low until ICDRR is read. The change of the acknowledge before reading ICDRR, to be returned to the master device, is reflected to the next transmit frame.
- 4. The last byte data is read by reading ICDRR.

Figure 18.11 Slave Receive Mode Operation Timing (1)

Figure 18.12 Slave Receive Mode Operation Timing (2)

18.4.6 Clocked Synchronous Serial Format

This module can be operated with the clocked synchronous serial format, by setting the FS bit in SAR to 1. When the MST bit in ICCR1 is 1, the transfer clock output from SCL is selected. When MST is 0, the external clock input is selected.

(1) Data Transfer Format

Figure 18.13 shows the clocked synchronous serial transfer format.

The transfer data is output from the fall to the fall of the SCL clock, and the data at the rising edge of the SCL clock is guaranteed. The MLS bit in ICMR sets the order of data transfer, in either the MSB first or LSB first. The output level of SDA can be changed during the transfer wait, by the SDAO bit in ICCR2.

Figure 18.13 Clocked Synchronous Serial Transfer Format

(2) Transmit Operation

In transmit mode, transmit data is output from SDA, in synchronization with the fall of the transfer clock. The transfer clock is output when MST in ICCR1 is 1, and is input when MST is 0. For transmit mode operation timing, refer to figure 18.14. The transmission procedure and operations in transmit mode are described below.

- 1. Set the ICE bit in ICCR1 to 1. Set the MST and CKS[3:0] bits in ICCR1. (Initial setting)
- 2. Set the TRS bit in ICCR1 to select the transmit mode. Then, TDRE in ICSR is set.
- 3. Confirm that TDRE has been set. Then, write the transmit data to ICDRT. The data is transferred from ICDRT to ICDRS, and TDRE is set automatically. The continuous transmission is performed by writing data to ICDRT every time TDRE is set. When changing from transmit mode to receive mode, clear TRS while TDRE is 1.

Figure 18.14 Transmit Mode Operation Timing

(3) Receive Operation

In receive mode, data is latched at the rise of the transfer clock. The transfer clock is output when MST in ICCR1 is 1, and is input when MST is 0. For receive mode operation timing, refer to figure 18.15. The reception procedure and operations in receive mode are described below.

- 1. Set the ICE bit in ICCR1 to 1. Set bits CKS[3:0] in ICCR1. (Initial setting)
- 2. When the transfer clock is output, set MST to 1 to start outputting the receive clock.
- 3. When the receive operation is completed, data is transferred from ICDRS to ICDRR and RDRF in ICSR is set. When $MST = 1$, the next byte can be received, so the clock is continually output. The continuous reception is performed by reading ICDRR every time RDRF is set. When the 8th clock is risen while RDRF is 1, the overrun is detected and AL/OVE in ICSR is set. At this time, the previous reception data is retained in ICDRR.
- 4. To stop receiving when MST = 1, set RCVD in ICCR1 to 1, then read ICDRR. Then, SCL is fixed high after receiving the next byte data.
- Notes: Follow the steps below to receive only one byte with MST = 1 specified. See figure 18.16 for the operation timing.
	- 1. Set the ICE bit in ICCR1 to 1. Set bits CKS[3:0] in ICCR1. (Initial setting)
	- 2. Set MST = 1 while the RCVD bit in ICCR1 is 0. This causes the receive clock to be output.
	- 3. Check if the BC2 bit in ICMR is set to 1 and then set the RCVD bit in ICCR1 to 1. This causes the SCL to be fixed to the high level after outputting one byte of the receive clock.

Figure 18.15 Receive Mode Operation Timing

Figure 18.16 Operation Timing For Receiving One Byte (MST = 1)

18.4.7 Noise Filter

The logic levels at the SCL and SDA pins are routed through noise filters before being latched internally. Figure 18.17 shows a block diagram of the noise filter circuit.

The noise filter consists of three cascaded latches and a match detector. The SCL (or SDA) input signal is sampled on the peripheral clock. When NF2CYC is set to 0, this signal is not passed forward to the next circuit unless the outputs of both latches agree. When NF2CYC is set to 1, this signal is not passed forward to the next circuit unless the outputs of three latches agree. If they do not agree, the previous value is held.

Figure 18.17 Block Diagram of Noise Filter

18.4.8 Example of Use

Flowcharts in respective modes that use the I^2C bus interface 3 are shown in figures 18.18 to 18.21.

Figure 18.19 Sample Flowchart for Master Receive Mode

Figure 18.20 Sample Flowchart for Slave Transmit Mode

Figure 18.21 Sample Flowchart for Slave Receive Mode

18.5 Interrupt Requests

There are six interrupt requests in this module; transmit data empty, transmit end, receive data full, NACK detection, STOP recognition, and arbitration lost/overrun error. Table 18.4 shows the contents of each interrupt request.

Table 18.4 Interrupt Requests

When the interrupt condition described in table 18.4 is 1, the CPU executes an interrupt exception handling. Note that a TXI or RXI interrupt can activate the DMAC if the setting for DMAC activation has been made. In such a case, an interrupt request is not sent to the CPU. Interrupt sources should be cleared in the exception handling. The TDRE and TEND bits are automatically cleared to 0 by writing the transmit data to ICDRT. The RDRF bit is automatically cleared to 0 by reading ICDRR. The TDRE bit is set to 1 again at the same time when the transmit data is written to ICDRT. Therefore, when the TDRE bit is cleared to 0, then an excessive data of one byte may be transmitted.

18.6 Bit Synchronous Circuit

In master mode, this module has a possibility that high level period may be short in the two states described below.

- When SCL is driven to low by the slave device
- When the rising speed of SCL is lowered by the load of the SCL line (load capacitance or pullup resistance)

Therefore, it monitors SCL and communicates by bit with synchronization.

Figure 18.22 shows the timing of the bit synchronous circuit and table 18.5 shows the time when the SCL output changes from low to Hi-Z then SCL is monitored.

Figure 18.22 Bit Synchronous Circuit Timing

Table 18.5 Time for Monitoring SCL

Note: $*$ tpcyc indicates the frequency of the peripheral clock ($P\phi$).

18.7 Usage Notes

18.7.1 Note on Setting for Multi-Master Operation

In multi-master operation, when the setting for IIC transfer rate (ICCR1.CKS[3:0]) makes this LSI slower than the other masters, pulse cycles with an unexpected length will infrequently be output on SCL.

Be sure to specify a transfer rate that is at least 1/1.8 of the fastest transfer rate among the other masters.

18.7.2 Note on Master Receive Mode

Reading ICDRR around the falling edge of the 8th clock might fail to fetch the receive data.

In addition, when RCVD is set to 1 around the falling edge of the 8th clock and the receive buffer full, a stop condition may not be issued.

Use either 1 or 2 below as a measure against the situations above.

- 1. In master receive mode, read ICDRR before the rising edge of the 8th clock.
- 2. In master receive mode, set the RCVD bit to 1 so that transfer proceeds in byte units.

18.7.3 Note on Setting ACKBT in Master Receive Mode

In master receive mode operation, set ACKBT before the falling edge of the 8th SCL cycle of the last data being continuously transferred. Not doing so can lead to an overrun for the slave transmission device.

18.7.4 Note on the States of Bits MST and TRN when Arbitration Is Lost

When sequential bit-manipulation instructions are used to set the MST and TRS bits to select master transmission in multi-master operation, a conflicting situation where AL in ICSR = 1 but the mode is master transmit mode ($MST = 1$ and $TRS = 1$) may arise; this depends on the timing of the loss of arbitration when the bit manipulation instruction for TRS is executed.

This can be avoided in either of the following ways.

- In multi-master operation, use the MOV instruction to set the MST and TRS bits.
- When arbitration is lost, check whether the MST and TRS bits are 0. If the MST and TRS bits have been set to a value other than 0, clear the bits to 0.

18.7.5 Note on I²C-bus Interface Master Receive Mode

After a master receive operation is completed, confirm the falling edge of the ninth clock cycle of the SCL signal and generate a stop condition or regenerate a start condition.

18.7.6 Note on IICRST and BBSY bits

When 1 is written to IICRST in ICCR2, this LSI release SCL and SDA pins. Then, if the SDA level changes from low to high under the condition of $SCL = high$, BBSY in ICCR2 is cleared to 0 assuming that the stop condition has been issued.

18.7.7 Note on Issuance of Stop Conditions in Master Transmit Mode while ACKE = 1

When a stop condition is issued in master transmit mode while the ACKE bit in the $\mathbf{I}^2\mathbf{C}$ bus interrupt enable register (ICIER) is 1, the stop condition may not be normally output depending on the issued timing. To avoid this, recognize the falling edge of the ninth clock before issuance of the stop condition.

The falling edge of the ninth clock can be recognized by checking the SCLO bit in the $I²C$ control register 2 (ICCR2).

Section 19 Serial Sound Interface with FIFO (SSIF)

The serial sound interface with FIFO (SSIF) is a module designed to send or receive audio data interface with various devices offering I^2S bus compatibility. It also provides additional modes for other common formats, as well as support for multi-channel mode.

19.1 Features

- Number of channels: Six channels
- Operating mode: Non-compressed mode The non-compressed mode supports serial audio streams divided by channels.
- Serves as both a transmitter and a receiver
- Capable of using serial bus format
- Asynchronous transfer takes place between the data buffer and the shift register.
- It is possible to select a value as the dividing ratio for the clock used by the serial bus interface.
- It is possible to control data transmission or reception with DMAC and interrupt requests.
- Selects the oversampling clock input from among the following pins:

EXTAL, XTAL (Clock operation modes 0 and 1: 10 to 33.33 MHz)

CKIO (Clock operation mode 2: 40 to 50 MHz*)

AUDIO_CLK (1 to 40 MHz)

AUDIO_X1, AUDIO_X2 (when connecting a crystal oscillator: 10 to 40 MHz, when used to input external clock: 1 to 40 MHz)

Note: * Do not select CKIO as the supply source for the oversampling clock if the frequency of CKIO is over 50 MHz and the mode is clock operating mode 2.

• Includes an 8-stage FIFO buffer for transmission and reception

Figure 19.1 shows a schematic diagram of the four channels in the SSIF module.

Figure 19.1 Schematic Diagram of SSIF Module

Figure 19.2 shows a block diagram of the SSIF module.

19.2 Input/Output Pins

Table 19.1 shows the pin assignments relating to the SSIF module.

Table 19.1 Pin Assignments

19.3 Register Description

The SSIF has the following registers. Note that explanation in the text does not refer to the channels.

Table 19.2 Register Description

Notes: 1. Although bits 26 and 27 in these registers can be read from or written to, bits other than these are read-only. For details, refer to section 19.3.2, Status Register (SSISR).

 2. To bits 1 and 0 in these registers, only 0 can be written to clear the flags. Other bits are read-only. For details, refer to section 19.3.6, FIFO Status Register (SSIFSR).

 ^{3.} These registers cannot be written to during reception. For details, refer to section 19.3.7, FIFO Data Register (SSIFDR).

19.3.1 Control Register (SSICR)

SSICR is a readable/writable 32-bit register that controls the IRQ, selects the polarity status, and sets operating mode.

19.3.2 Status Register (SSISR)

SSISR consists of status flags indicating the operational status of the SSIF module and bits indicating the current channel numbers and word numbers.

Note: $*$ The bit can be read or written to. Writing 0 initializes the bit, but writing 1 is ignored.

Note: * The bit can be read or written to. Writing 0 initializes the bit, but writing 1 is ignored.

19.3.3 Transmit Data Register (SSITDR)

SSITDR is a 32-bit register that stores data to be transmitted. The data for transmission to be stored to SSITDR is automatically transferred from the FIFO data register.

Data written to this register is transferred to the shift register upon transmission request. If the data word length is less than 32 bits, the alignment is determined by the setting of the PDTA control bit in SSICR.

The CPU cannot read or write data from/to SSITDR.

19.3.4 Receive Data Register (SSIRDR)

SSIRDR is a 32-bit register that stores received data. The received data stored in SSIRDR is automatically transferred to the FIFO data register.

Data in this register is transferred from the shift register each time data word is received. If the data word length is less than 32 bits, the alignment is determined by the setting of the PDTA control bit in SSICR.

The CPU cannot read or write data from/to SSIRDR.

19.3.5 FIFO Control Register (SSIFCR)

SSIFCR is a readable/writable 32-bit register that specifies the data trigger numbers and selects between transmission and reception for the FIFO data register, and enables or disables FIFO data reset and interrupt requests.

SSIFCR can always be read or written by the CPU.

19.3.6 FIFO Status Register (SSIFSR)

SSIFSR consists of status flags indicating the operating status of the FIFO data register.

Note: $*$ The bit can be read or written to. Writing 0 initializes the bit, but writing 1 is ignored.

Note: * The bit can be read or written to. Writing 0 initializes the bit, but writing 1 is ignored.

19.3.7 FIFO Data Register (SSIFDR)

In transmission, SSIFDR operates as a FIFO register consisting of eight stages of 32-bit registers for storing data to be serially transmitted. On detecting that the transmit data register (SSITDR) is empty, the SSIF transfers the data for transmission written to SSIFDR to SSITDR to start serial transmission, which can continue until SSIFDR becomes empty. SSIFDR can be written to by the CPU at any time. Note that when SSIFDR is full of data (32 bytes), the next data cannot be written to it and will be ignored if writing is attempted.

In reception, SSIFDR operates as a FIFO register consisting of eight stages of 32-bit registers for storing serially received data. When four bytes of data have been received, the SSIF transfers the received data in the receive data register (SSIRDR) to SSIFDR to complete reception operation. Reception can continue until 32 bytes of data have been stored to SSIFDR. SSIFDR can be read by the CPU but cannot be written to. Note that when SSIFDR is read when it stores no received data, undefined values will be read. After SSIFDR becomes full of received data, the data received thereafter will be lost.

Note: $*$ Not writable during reception.

19.4 Operation Description

19.4.1 Bus Format

The SSIF module can operate as a transmitter or a receiver and can be configured into many serial bus formats in either mode.

The bus format can be selected from one of the four major modes shown in table 19.3.

Table 19.3 Bus Format for SSIF Module

19.4.2 Non-Compressed Modes

The non-compressed modes support all serial audio streams split into channels. It supports the I^2S compatible format as well as many more variants on these modes.

(1) Slave Receiver

This mode allows the module to receive serial data from another device. The clock and word select signal used for the serial data stream is also supplied from an external device. If these signals do not conform to the format specified in the configuration fields of the SSIF module, operation is not guaranteed.

(2) Slave Transmitter

This mode allows the module to transmit serial data to another device. The clock and word select signal used for the serial data stream is also supplied from an external device. If these signals do not conform to the format specified in the configuration fields of the SSIF module, operation is not guaranteed.

(3) Master Receiver

This mode allows the module to receive serial data from another device. The clock and word select signals are internally derived from the oversampling clock. The format of these signals is defined in the configuration fields of the SSIF module. If the incoming data does not follow the configured format, operation is not guaranteed.

(4) Master Transmitter

This mode allows the module to transmit serial data to another device. The clock and word select signals are internally derived from the oversampling clock. The format of these signals is defined in the configuration fields of the SSIF module.

(5) Operating Setting Related to Word Length

All bits related to the SSICR's word length are valid in non-compressed modes. There are many configurations the SSIF module supports, but some of the combinations are shown below for the I 2 S compatible format, MSB-first and left-aligned format, and MSB-first and right-aligned format.

Figures 19.3 and 19.4 demonstrate the supported I^2S compatible format both without and with padding. Padding occurs when the data word length is smaller than the system word length.

 \bullet I²S Compatible Format

Figure 19.3 I²S Compatible Format (without Padding)

Figure 19.4 I'S Compatible Format (with Padding)

Figure 19.5 shows the MSB-first and left-aligned format and figure 19.6 shows the MSB-first and right-aligned format. Padding is assumed in both cases, but may not be present in a final implementation if the system word length equals the data word length.

• MSB-first and Left-Aligned Format

Figure 19.5 MSB-first and Left-Aligned Format (Transmitted and received in the order of serial data and padding bits)

• MSB-first and Right-Aligned Format

Figure 19.6 MSB-first and Right-Aligned Format (Transmitted and received in the order of padding bits and serial data)

(6) Multi-channel Formats

Some devices extend the definition of the specification by I^2S bus and allow more than 2 channels to be transferred within two system words.

The SSIF module supports the transfer of 4, 6 and 8 channels by using the CHNL, SWL and DWL bits only when the system word length (SWL) is greater than or equal to the data word length (DWL) multiplied by channels (CHNL).

Table 19.4 shows the number of padding bits for each of the valid setting. If setting is not valid, "
—" is indicated instead of a number.

Padding Bits Per System Word			DWL[2:0]	000	001	010	011	100	101	110
CHNL [1:0]	Decoded Channels per System Word	SWL [2:0]	Decoded Word Length	8	16	18	20	22	24	32
00	1	000	8	0						
		001	16	8	0					
		010	24	16	8	6	$\overline{4}$	\overline{c}	0	
		011	32	24	16	14	12	10	8	0
		100	48	40	32	30	28	26	24	16
		101	64	56	48	46	44	42	40	32
		110	128	120	112	110	108	106	104	96
		111	256	248	240	238	236	234	232	224
01	\overline{c}	000	8							
		001	16	0						
		010	24	8						
		011	32	16	0					
		100	48	32	16	12	8	$\overline{4}$	0	
		101	64	48	32	28	24	20	16	0
		110	128	112	96	92	88	84	80	64
		111	256	240	224	220	216	212	208	192

Table 19.4 The Number of Padding Bits for Each Valid Setting

When the SSIF module acts as a transmitter, each word written to SSITDR is transmitted to the serial audio bus in the order they are written. When the SSIF module acts as a receiver, each word received by the serial audio bus is read in the order received from the SSIRDR register.

Figures 19.7 to 19.9 show how the data on 4, 6 and 8 channels are transferred to the serial audio bus. Note that there are no padding bits in the first example, the second example is left-aligned and the third is right-aligned. The other conditions in these examples have been selected arbitrarily.

Figure 19.7 Multi-Channel Format (4 Channels Without Padding)

Figure 19.8 Multi-Channel Format (6 Channels with High Padding)

Figure 19.9 Multi-Channel Format (8 Channels; Transmitting and Receiving in the order of serial data and Padding Bits; with padding)

(7) Bit Setting Configuration Format

Several more configuration bits in non-compressed mode are shown below. These bits are not mutually exclusive, but some combinations may not be useful for any other device.

These configuration bits are described below with reference to figure 19.10.

Figure 19.10 Basic Sample Format (Transmit Mode with Example System/Data Word Length)

Figure 19.10 uses a system word length of 6 bits and a data word length of 4 bits. These settings are not possible with the SSIF module but are used only for clarification of the other configuration bits.

• Inverted Clock

Figure 19.11 Inverted Clock

• Inverted Word Select

Figure 19.12 Inverted Word Select

• Inverted Padding Polarity

Figure 19.13 Inverted Padding Polarity

• Transmitting and Receiving in the Order of Padding Bits and Serial Data; with Delay

Figure 19.14 Transmitting and Receiving in the Order of Padding Bits and Serial Data; with Delay

• Transmitting and Receiving in the Order of Padding Bits and Serial Data; without Delay

Figure 19.15 Transmitting and Receiving in the Order of Padding Bits and Serial Data; without Delay

• Transmitting and Receiving in the Order of Serial Data and Padding Bits; without Delay

Figure 19.16 Transmitting and Receiving in the Order of Serial Data and Padding Bits; without Delay

• Parallel Right-Aligned with Delay

Figure 19.17 Parallel Right-Aligned with Delay

• Mute Enabled

Figure 19.18 Mute Enabled

19.4.3 Operation Modes

There are three modes of operation: configuration, enabled and disabled. Figure 19.19 shows how the module enters each of these modes.

Figure 19.19 Operation Modes

(1) Configuration Mode

This mode is entered after the module is released from reset. All required configuration fields in the control register should be defined in this mode, before the SSIF module is enabled by setting the EN bit.

Setting the EN bit causes the module to enter the module enabled mode.

(2) Module Enabled Mode

Operation of the module in this mode is dependent on the operation mode selected. For details, refer to section 19.4.4, Transmit Operation and section 19.4.5, Receive Operation, below.

19.4.4 Transmit Operation

Transmission can be controlled either by DMA or interrupt.

DMA control is preferred to reduce the processor load. In DMA control mode the processor will only receive interrupts if there is an underflow or overflow of data or the DMAC has finished its transfer.

The alternative method is using the interrupts that the SSIF module generates to supply data as required.

When disabling the SSIF module, the clock* must be kept supplied to the SSIF until the IIRO bit indicates that the module is in the idle state.

Figure 19.20 shows the transmit operation in DMA control mode, and figure 19.21 shows the transmit operation in interrupt control mode.

Note: $*$ Input clock from the SSISCK pin when SCKD = 0. Oversampling clock when $SCKD = 1$.

(1) Transmission Using DMA Controller

Figure 19.20 Transmission Using DMA Controller

(2) Transmission Using Interrupt Data Flow Control

Figure 19.21 Transmission Using Interrupt Data Flow Control

19.4.5 Receive Operation

Like transmission, reception can be controlled either by DMA or interrupt.

Figures 19.22 and 19.23 show the flow of operation.

When disabling the SSIF module, the clock* must be kept supplied to the SSIF until the IIRQ bit indicates that the module is in the idle state.

Note: $*$ Input clock from the SSISCK pin when SCKD = 0. Oversampling clock when $SCKD = 1$.

(1) Reception Using DMA Controller

Figure 19.22 Reception Using DMA Controller

(2) Reception Using Interrupt Data Flow Control

Figure 19.23 Reception Using Interrupt Data Flow Control

When an underflow or overflow error condition has matched, the CHNO [1:0] bit and the SWNO bit can be used to recover the SSIF module to a known status. When an underflow or overflow occurs, the host can read the channel number and system word number to determine what point the serial audio stream has reached. In the transmitter case, the host can skip forward through the data it wants to transmit until it finds the sample data that matches what the SSIF module is expecting to transmit next, and so resynchronize with the audio data stream. In the receiver case the host CPU can store null data to make the number of receive data items consistent until it is ready to store the sample data that the SSIF module is indicating will be received next, and so resynchronize with the audio data stream.

19.4.6 Serial Bit Clock Control

This function is used to control and select which clock is used for the serial bus interface.

If the serial clock direction is set to input $(SCKD = 0)$, the SSIF module is in clock slave mode and the shift register uses the bit clock that was input to the SSISCK pin.

If the serial clock direction is set to output $(SCKD = 1)$, the SSIF module is in clock master mode, and the shift register uses the oversampling clock or a divided oversampling clock as the bit clock. The oversampling clock is divided by the ratio specified by the serial oversampling clock division ratio bits (CKDV) in SSICR for use as the bit clock by the shift register.

In either case the module pin, SSISCK, is the same as the bit clock.

19.5 Usage Notes

19.5.1 Limitations from Underflow or Overflow during DMA Operation

If an underflow or overflow occurs while the DMA is in operation, the module should be restarted. The transmit and receive buffers in the SSI consists of 32-bit registers that share the L and R channels. Therefore, data to be transmitted and received at the L channel may sometimes be transmitted and received at the R channel if an underflow or overflow occurs, for example, under the following condition: the control register (SSICR) has a 32-bit setting for both data word length (DWL2 to DWL0) and system word length (SWL2 to SWL0).

If an error occurrence is confirmed with two types of error interrupts (underflow and overflow) or the corresponding error status flag (the bits UIRQ and OIRQ in SSISR), write 0 to the EN bit in SSICR to disable DMA transfer requests in this module, thus stopping the operation. (In this case, the direct memory access controller setting should also be stopped.) After this, write 0 to the error status flag bit to clear the error status, set the direct memory access controller again and restart the transfer.

Section 20 Controller Area Network (RCAN-TL1)

20.1 Summary

20.1.1 Overview

This document primarily describes the programming interface for the RCAN-TL1 (Renesas CAN Time Trigger Level 1) module. It serves to facilitate the hardware/software interface so that engineers involved in the RCAN-TL1 implementation can ensure the design is successful.

20.1.2 Scope

The CAN Data Link Controller function is not described in this document. It is the responsibility of the reader to investigate the CAN Specification Document (see references). The interfaces from the CAN Controller are described, in so far as they pertain to the connection with the User Interface.

The programming model is described in some detail. It is not the intention of this document to describe the implementation of the programming interface, but to simply present the interface to the underlying CAN functionality.

The document places no constraints upon the implementation of the RCAN-TL1 module in terms of process, packaging or power supply criteria. These issues are resolved where appropriate in implementation specifications.

20.1.3 Audience

In particular this document provides the design reference for software authors who are responsible for creating a CAN application using this module.

In the creation of the RCAN-TL1 user interface LSI engineers must use this document to understand the hardware requirements.

20.1.4 References

- 1. CAN Specification Version 2.0 part A, Robert Bosch GmbH, 1991
- 2. CAN Specification Version 2.0 part B, Robert Bosch GmbH, 1991
- 3. Implementation Guide for the CAN Protocol, CAN Specification 2.0 Addendum, CAN In Automation, Erlangen, Germany, 1997
- 4. Road vehicles Controller area network (CAN): Part 1: Data link layer and physical signalling (ISO-11898-1, 2003)

5. Road vehicles - Controller area network (CAN): Part 4: Time triggered communication (ISO-11898-4, 2004)

20.1.5 Features

- Supports CAN specification 2.0B
- Bit timing compliant with ISO-11898-1
- 32 Mailbox version
- Clock 16 to 33 MHz
- 31 programmable Mailboxes for transmit / receive $+1$ receive-only mailbox
- Sleep mode for low power consumption and automatic recovery from sleep mode by detecting CAN bus activity
- Programmable receive filter mask (standard and extended identifier) supported by all Mailboxes
- Programmable CAN data rate up to 1MBit/s
- Transmit message queuing with internal priority sorting mechanism against the problem of priority inversion for real-time applications
- Data buffer access without SW handshake requirement in reception
- Flexible micro-controller interface
- Flexible interrupt structure
- 16-bit free running timer with flexible clock sources and pre-scaler, 3 Timer Compare Match Registers
- 6-bit Basic Cycle Counter for Time Trigger Transmission
- Timer Compare Match Registers with interrupt generation
- Timer counter clear / set capability
- Registers for Time-Trigger: Local_Time, Cycle_time, Ref_Mark, Tx_Enable Window, Ref_Trigger_Offset
- Flexible TimeStamp at SOF for both transmission and reception supported
- Time-Trigger Transmission, Periodic Transmission supported (on top of Event Trigger Transmission)
- Basic Cycle value can be embedded into a CAN frame and transmitted

20.2 Architecture

The RCAN-TL1 device offers a flexible and sophisticated way to organise and control CAN frames, providing the compliance to CAN2.0B Active and ISO-11898-1. The module is formed from 5 different functional entities. These are the Micro Processor Interface (MPI), Mailbox, Mailbox Control, Timer, and CAN Interface. The figure below shows the block diagram of the RCAN-TL1 Module. The bus interface timing is designed according to the peripheral bus I/F required for each product.

Figure 20.1 RCAN-TL1 architecture

Important: Although core of RCAN-TL1 is designed based on a 32-bit bus system, the whole RCAN-TL1 including MPI for the CPU has 16-bit bus interface to CPU. In that case, LongWord (32-bit) access must be implemented as 2 consecutive word (16-bit) accesses. In this manual, LongWord access means the two consecutive accesses.

• Micro Processor Interface (MPI)

The MPI allows communication between the Renesas CPU and the RCAN-TL1's registers/mailboxes to control the memory interface. It also contains the Wakeup Control logic that detects the CAN bus activities and notifies the MPI and the other parts of RCAN-TL1 so that the RCAN-TL1 can automatically exit the Sleep mode. It contains registers such as MCR, IRR, GSR and IMR.

• Mailbox

The Mailboxes consists of RAM configured as message buffers and registers. There are 32 Mailboxes, and each mailbox has the following information.

 $<$ RAM $>$

- CAN message control (identifier, rtr, ide, etc)
- CAN message data (for CAN Data frames)
- Local Acceptance Filter Mask for reception

<Registers>

- CAN message control (dlc)
- Time Stamp for message reception/transmission
- ⎯ 3-bit wide Mailbox Configuration, Disable Automatic Re-Transmission bit, Auto-Transmission for Remote Request bit, New Message Control bit
- Tx-Trigger Time
- Mailbox Control

The Mailbox Control handles the following functions.

- ⎯ For received messages, compare the IDs and generate appropriate RAM addresses/data to store messages from the CAN Interface into the Mailbox and set/clear appropriate registers accordingly.
- ⎯ To transmit event-triggered messages, run the internal arbitration to pick the correct priority message, and load the message from the Mailbox into the Tx-buffer of the CAN Interface and set/clear appropriate registers accordingly. In the case of time-triggered transmission, compare match of Tx-Trigger time invoke loading the messages.
- ⎯ Arbitrates Mailbox accesses between the CPU and the Mailbox Control.
- ⎯ Contains registers such as TXPR, TXCR, TXACK, ABACK, RXPR, RFPR, UMSR and MBIMR.

• Timer

The Timer function is the functional entity, which provides RCAN-TL1 with support for transmitting messages at a specific time frame and recording the result.

The Timer is a 16-bit free running up counter which can be controlled by the CPU. It provides one 16-bit Compare Match Register to compare with Local Time and two 16-bit ones to compare with Cycle Time. The Compare Match Registers can generate interrupt signals and clear the Counter.

The clock period of this Timer offers a wide selection derived from the system clock or can be programmed to be incremented with one nominal bit timing of CAN Bus.

Contains registers such as TCNTR, TTCR0, CMAX_TEW, RFTROFF, TSR, CCR, CYCTR, RFMK, TCMR0, TCMR1, TCMR2 and TTTSEL.

• CAN Interface

This block conforms to the requirements for a CAN Bus Data Link Controller which is specified in Ref. [2, 4]. It fulfils all the functions of a standard DLC as specified by the OSI 7 Layer Reference model. This functional entity also provides the registers and the logic which are specific to a given CAN bus, which includes the Receive Error Counter, Transmit Error Counter, the Bit Configuration Registers and various useful Test Modes. This block also contains functional entities to hold the data received and the data to be transmitted for the CAN Data Link Controller.

20.3 Programming Model—Overview

The purpose of this programming interface is to allow convenient, effective access to the CAN bus for efficient message transfer. Please bear in mind that the user manual reports all settings allowed by the RCAN-TL1 IP. Different use of RCAN-TL1 is not allowed.

20.3.1 Memory Map

The diagram of the memory map is shown below.

Base address RCAN0: H'FFFE 5000 RCAN1: H'FFFE 5800

Figure 20.2 RCAN-TL1 Memory Map

The locations not used (between H'000 and H'4F3) are reserved and cannot be accessed.

20.3.2 Mailbox Structure

Mailboxes play a role as message buffers to transmit/receive CAN frames. Each Mailbox is comprised of 3 identical storage fields that are 1): Message Control, 2): Local Acceptance Filter Mask, 3): Message Data. In addition some Mailboxes contain the following extra Fields: 4): Time Stamp, 5): Time Trigger configuration and 6): Time Trigger Control. The following table shows the address map for the control, LAFM, data, timestamp, Transmission Trigger Time and Time Trigger Control addresses for each mailbox.

Mailbox-0 is a receive-only box, and all the other Mailboxes can operate as both receive and transmit boxes, dependant upon the MBC (Mailbox Configuration) bits in the Message Control. The following diagram shows the structure of a Mailbox in detail.

Table 20.1 Roles of Mailboxes

(ET) shows that it works during merged arbitrating window, after completion of time-triggered transmission.

Figure 20.3 Mailbox-N Structure

MB23 to 16 (MB without timestamp)																						
Address		Data Bus															Access Size	Field Name				
	15	14	13	12	11	10	9	8	$\overline{7}$	6		5	$\overline{4}$		3	\overline{c}	$\overline{1}$	$\mathbf 0$				
$H'100 + N*32$	IDE	RTR Ω EXTID[17:16] STDID[10:0]														Word/LW	Control 0					
$H'102 + N*32$	EXTID[15:0]															Word						
$H'104 + N*32$	IDE LAFM	EXTID LAFM[17:16] STDID LAFM[10:0]										Word/LW	LAFM									
$H'106 + N*32$	EXTID LAFM[15:0]														Word							
$H'108 + N*32$	MSG DATA 0 (first Rx/Tx Byte)														MSG DATA 1	Byte/Word/LW						
$H'10A + N*32$	MSG DATA 2									MSG DATA 3									Byte/Word	Data		
$H'10C + N*32$	MSG DATA 4									MSG DATA 5									Byte/Word/LW			
$H'10E + N*32$	MSG DATA 6									MSG DATA 7									Byte/Word			
H'110 + N*32	Ω	ATX DART NMC MBC[2:0] $\mathbf 0$ \circ $\mathbf 0$ 0 0										DLC[3:0]			Byte/Word	Control 1						
Address	15	14	13	12	11	10	9	8	$\overline{7}$	6		5	4		3	\overline{c}		$\mathbf 0$	Access Size Field Name			
								Data Bus														
$H'100 + N*32$	IDE	RTR	$\overline{0}$						STDID[10:0]									EXTID[17:16]	Word/LW	Control 0		
$H'102 + N*32$	EXTID[15:0] EXTID														Word Word/LW							
$H'104 + N*32$	IDE LAFM STDID LAFM[10:0] $\mathbf{0}$ $\mathbf{0}$													LAFM[17:16]	Word	LAFM						
$H'106 + N*32$	EXTID LAFM[15:0]																					
$H'108 + N*32$ $H'10A + N*32$	MSG DATA 0 (first Rx/Tx Byte)									MSG DATA 1									Byte/Word/LW	Data		
	MSG DATA 2 MSG DATA 4									MSG DATA 3									Byte/Word Byte/Word/LW			
$H'10C + N*32$								MSG DATA 5														
$H'10E + N*32$	MSG DATA 6									MSG DATA 7									Byte/Word			
$H'110 + N*32$	NMC ATX DART \circ $\mathbf 0$ MBC[2:0] $\mathbf 0$ $\mathbf 0$ DLC[3:0] $\mathbf 0$ $\mathbf 0$												Byte/Word	Control 1								
$H'112 + N*32$									reserved											÷.		
$H'114 + N*32$		Tx-Triggered Time (TTT)												Word	Trigger Time							
$H'116 + N*32$	TTW[1:0] offset							Ω	0		0	0		0		Rep Factor		Word	TT control			

Figure 20.3 Mailbox-N Structure (continued)

- Notes: 1. All bits shadowed in grey are reserved and must be written LOW. The value returned by a read may not always be '0' and should not be relied upon.
	- 2. ATX and DART are not supported by Mailbox-0, and the MBC setting of Mailbox-0 is limited.
	- 3. ID Reorder (MCR15) can change the order of STDID, RTR, IDE and EXTID of both message control and LAFM.

(1) Message Control Field

STDID[10:0]: These bits set the identifier (standard identifier) of data frames and remote frames.

EXTID[17:0]: These bits set the identifier (extended identifier) of data frames and remote frames.

RTR (Remote Transmission Request bit): Used to distinguish between data frames and remote frames. This bit is overwritten by received CAN Frames depending on Data Frames or Remote Frames.

Important: Please note that, when ATX bit is set with the setting MBC = 001(bin), the RTR bit will never be set. When a Remote Frame is received, the CPU can be notified by the corresponding RFPR set or IRR[2] (Remote Frame Receive Interrupt), however, as RCAN-TL1 needs to transmit the current message as a Data Frame, the RTR bit remains unchanged.

Important: In order to support automatic answer to remote frame when MBC = 001 (bin) is used and ATX = 1 the RTR flag must be programmed to zero to allow data frame to be transmitted.

Note: when a Mailbox is configured to send a remote frame request the DLC used for transmission is the one stored into the Mailbox.

IDE (Identifier Extension bit): Used to distinguish between the standard format and extended format of CAN data frames and remote frames.

Mailbox-0

Note: MBC[1] of MB0 is always "1".

Mailbox-31 to 1

NMC (New Message Control): When this bit is set to '0', the Mailbox of which the RXPR or RFPR bit is already set does not store the new message but maintains the old one and sets the UMSR correspondent bit. When this bit is set to '1', the Mailbox of which the RXPR or RFPR bit is already set overwrites with the new message and sets the UMSR correspondent bit.

Important: Please note that if a remote frame is overwritten with a data frame or vice versa could be that both RXPR and RFPR flags (together with UMSR) are set for the same Mailbox. In this case the RTR bit within the Mailbox Control Field should be relied upon.

Important: Please note that when the Time Triggered mode is used NMC needs to be set to '1' for Mailbox 31 to allow synchronization with all incoming reference messages even when RXPR[31] is not cleared.

ATX (Automatic Transmission of Data Frame): When this bit is set to '1' and a Remote Frame is received into the Mailbox DLC is stored. Then, a Data Frame is transmitted from the same Mailbox using the current contents of the message data and updated DLC by setting the corresponding TXPR automatically. The scheduling of transmission is still governed by ID priority or Mailbox priority as configured with the Message Transmission Priority control bit (MCR.2). In order to use this function, MBC[2:0] needs to be programmed to be '001' (Bin). When a transmission is performed by this function, the DLC (Data Length Code) to be used is the one that has been received. Application needs to guarantee that the DLC of the remote frame correspond to the DLC of the data frame requested.

Important: When ATX is used and MBC = 001 (Bin) the filter for the IDE bit cannot be used since ID of remote frame has to be exactly the same as that of data frame as the reply message.

Important: Please note that, when this function is used, the RTR bit will never be set despite receiving a Remote Frame. When a Remote Frame is received, the CPU will be notified by the corresponding RFPR set, however, as RCAN-TL1 needs to transmit the current message as a Data Frame, the RTR bit remains unchanged.

Important: Please note that in case of overrun condition (UMSR flag set when the Mailbox has its $NMC = 0$) the message received is discarded. In case a remote frame is causing overrun into a Mailbox configured with $ATX = 1$, the transmission of the corresponding data frame may be triggered only if the related PFPR flag is cleared by the CPU when the UMSR flag is set. In such case PFPR flag would get set again.

DART (Disable Automatic Re-Transmission): When this bit is set, it disables the automatic retransmission of a message in the event of an error on the CAN bus or an arbitration lost on the CAN bus. In effect, when this function is used, the corresponding TXCR bit is automatically set at the start of transmission. When this bit is set to '0', RCAN-TL1 tries to transmit the message as many times as required until it is successfully transmitted or it is cancelled by the TXCR.

MBC[2:0] (Mailbox Configuration): These bits configure the nature of each Mailbox as follows. When $MBC = 111$ (Bin), the Mailbox is inactive, i.e., it does not receive or transmit a message regardless of TXPR or other settings. The MBC = '110', '101' and '100' settings are prohibited. When the MBC is set to any other value, the LAFM field becomes available. Please don't set TXPR when MBC is set as reception as there is no hardware protection, and TXPR will remain set. MBC[1] of Mailbox-0 is fixed to "1" by hardware. This is to ensure that MB0 cannot be configured to transmit Messages.

Notes: $*$ In order to support automatic retransmission, RTR shall be "0" when MBC = 001(bin) and $ATX = 1$.

When $ATX = 1$ is used the filter for IDE must not be used.

DLC[3:0] (Data Length Code): These bits encode the number of data bytes from 0,1, 2, … 8 that will be transmitted in a data frame. Please note that when a remote frame request is transmitted the DLC value to be used must be the same as the DLC of the data frame that is requested.

(2) Local Acceptance Filter Mask (LAFM)

This area is used as Local Acceptance Filter Mask (LAFM) for receive boxes**.**

LAFM: When MBC is set to 001, 010, 011(Bin), this field is used as LAFM Field. It allows a Mailbox to accept more than one identifier. The LAFM is comprised of two 16-bit read/write areas as follows.

Figure 20.4 Acceptance filter

If a bit is set in the LAFM, then the corresponding bit of a received CAN identifier is ignored when the RCAN-TL1 searches a Mailbox with the matching CAN identifier. If the bit is cleared, then the corresponding bit of a received CAN identifier must match to the STDID/IDE/EXTID set in the mailbox to be stored. The structure of the LAFM is same as the message control in a Mailbox. If this function is not required, it must be filled with '0'.

Important: RCAN-TL1 starts to find a matching identifier from Mailbox-31 down to Mailbox-0. As soon as RCAN-TL1 finds one matching, it stops the search. The message will be stored or not depending on the NMC and RXPR/RFPR flags. This means that, even using LAFM, a received message can only be stored into 1 Mailbox.

Important: When a message is received and a matching Mailbox is found, the whole message is stored into the Mailbox. This means that, if the LAFM is used, the STDID, RTR, IDE and EXTID may differ to the ones originally set as they are updated with the STDID, RTR, IDE and EXTID of the received message.

STD_LAFM[10:0] — Filter mask bits for the CAN base identifier [10:0] bits.

EXT_LAFM[17:0] — Filter mask bits for the CAN Extended identifier [17:0] bits.

IDE_LAFM — Filter mask bit for the CAN IDE bit.

(3) Message Data Fields

Storage for the CAN message data that is transmitted or received. MSG_DATA[0] corresponds to the first data byte that is transmitted or received. The bit order on the CAN bus is bit 7 through to \overline{b} it $\overline{0}$.

When $CMAX!=3' b111/MBC[30] = 3' b000$ and $TXPR[30]$ is set, Mailbox-30 is configured as transmission of time reference. Its DLC must be greater than 0 and its RTR must be zero (as specified for TTCAN Level 1) so that the Cycle count (CCR register) is embedded in the first byte of the data field instead of MSG_DATA_0[5:0] when this Mailbox starts transmission. This function shall be used when RCAN-TL1 is enabled to work in TTCAN mode to perform a Potential Time Master role to send the Time reference message. MSG_DATA_0[7:6] is still transmitted as stored in the Mailbox. User can set MSG_DATA_0[7] when a Next_is_Gap needs to be transmitted.

Please note that the CCR value is only embedded on the frame transmitted but not stored back into Mailbox 30.

When $CMAX!=3'b111$, $MBC[31]=3'b011$ and $TXPR[31]$ is cleared, Mailbox-31 is configured as reception of time reference. When a valid reference message is received (DLC > 0) RCAN-TL1 performs internal synchronisation (modifying its RFMK and basic cycle CCR).

Figure 20.5 Message Data Field

(4) Timestamp

Storage for the Timestamp recorded on messages for transmit/receive. The Timestamp will be a useful function to monitor if messages are received/transmitted within expected schedule.

• Timestamp

Message Receive: For received messages of Mailbox-15 to 0, Timestamp always captures the CYCTR (Cycle Time Register) value or Cycle Counter CCR $[5:0]$ + CYCTR $[15:6]$ value, depending on the programmed value in the bit 14 of TTCR0 (Timer Trigger Control Register 0) at SOF.

For messages received into Mailboxes 30 and 31, Timestamp captures the TCNTR (Timer Counter Register) value at SOF.

Message Transmit: For transmitted messages of Mailbox-15 to 1, Timestamp always captures the CYCTR (Cycle Time Register) value or Cycle_Counter CCR[5:0] + CYCTR[15:6] value, depending on the programmed value in the bit 14 of TTCR0 (Timer Trigger Control Register 0), at SOF.

For messages transmitted from Mailboxes30 and 31, Timestamp captures the TCNTR (Timer Counter Register) value at SOF.

Important: Please note that the TimeStamp is stored in a temporary register. Only after a successful transmission or reception the value is then copied into the related Mailbox field. The TimeStamp may also be updated if the CPU clears RXPR[N]/RFPR[N] at the same time that UMSR[N] is set in overrun, however it can be read properly before clearing RXPR[N]/RFPR[N].

(5) Tx-Trigger Time (TTT) and Time Trigger control

For Mailbox-29 to 24, when MBC is set to 000 (Bin) in time trigger mode (CMAX!= 3'b111), Tx-Trigger Time works as Time_Mark to determine the boundary between time windows. The TTT and TT control are comprised of two 16-bit read/write areas as follows. Mailbox-30 doesn't have TT control and works as Time_Ref.

Mailbox 30 to 24 can be used for reception if not used for transmission in TT mode. However they cannot join the event trigger transmission queue when the TT mode is used.

• Tx-Trigger Time

• Time Trigger control

The following figure shows the differences between all Mailboxes supporting Time Triggered mode.

• **TTW[1:0] (Time Trigger Window):** These bits show the attribute of time windows. Please note that once a merged arbitrating window is opened by $TTW = 2'b10$, the window must be closed by $TTW = 2'b11$. Several messages with $TTW = 2'b10$ may be used within the start and the end of a merged arbitrating window.

The first 16-bit area specifies the time that triggers the transmission of the message in cycle time. The second 16-bit area specifies the basic cycle in the system matrix where the transmission must start (Offset) and the frequency for periodic transmission. When the internal TTT register matches to the CYCTR value, and the internal Offset matches to CCR value transmission is attempted from the corresponding Mailbox. In order to enable this function, the CMAX (Cycle Maximum Register) must be set to a value different from 3'b111, the Timer (TCNTR) must be running (TTCR0 bit15 = 1), the corresponding MBC must be set to 3'b000 and the corresponding TXPR bit must be set. Once TXPR is set by S/W, RCAN-TL1 does not clear the corresponding TXPR bit (among Mailbox-30 to 24) to carry on performing the periodic transmission. In order to stop the periodic transmission, TXPR must be cleared by TXCR. Please note that in this case it is possible that both TXACK and ABACK are set for the same Mailbox if TXACK is not cleared right after completion of transmission. Please refer to figure 20.7.

Please note that for Mailbox 30 TTW is fixed to '01', Offset to '00' and rep_factor to '0'.The following tables report the combinations for the rep_factor and the offset.

The Offset Field determines the first cycle in which a Time Triggered Mailbox may start transmitting its Message.

The following relation must be maintained:

Cycle_Count_Maximum + 1 >= Repeat_Factor > Offset

Cycle_Count_Maximum = 2^{CMAX} - 1

Repeat Factor = $2^{\text{rep_factor}}$

CMAX, Repeat_Factor, and Offset are register values

Figure 20.8 System Matrix

Tx-Trigger Times must be set in ascending order such that the difference between them satisfies the following condition.

TTT(mailbox i) –1 TTT(mailbox i-1) > TEW + Maximum frame length + 9

20.3.3 RCAN-TL1 Control Registers

The following sections describe RCAN-TL1 control registers. The address is mapped as follow.

Important: These registers can only be accessed in Word size (16-bit).

Figure 20.9 RCAN-TL1 control registers

(1) Master Control Register (MCR)

The Master Control Register (MCR) is a 16-bit read/write register that controls RCAN-TL1.

• MCR (Address = $H'000$)

Bit 15 — ID Reorder (MCR15): This bit changes the order of STDID, RTR, IDE and EXTID of both message control and LAFM.

Figure 20.10 ID Reorder

This bit can be modified only in reset mode.

Bit 14 — Auto Halt Bus Off (MCR14): If both this bit and MCR6 are set, MCR1 is automatically set as soon as RCAN-TL1 enters BusOff.

This bit can be modified only in reset mode.

Bit 13 — Reserved. The written value should always be '0' and the returned value is '0'.

Bit 12 — Reserved. The written value should always be '0' and the returned value is '0'.

Bit 11 — Reserved. The written value should always be '0' and the returned value is '0'.

Bit 10 - 8 — Test Mode (TST[2:0]): This bit enables/disables the test modes. Please note that before activating the Test Mode it is requested to move RCAN-TL1 into Halt mode or Reset mode. This is to avoid that the transition to Test Mode could affect a transmission/reception in progress. For details, please refer to section 20.4.1, Test Mode Settings.

Please note that the test modes are allowed only for diagnosis and tests and not when RCAN-TL1 is used in normal operation.

Bit 7 — Auto-wake Mode (MCR7): MCR7 enables or disables the Auto-wake mode. If this bit is set, the RCAN-TL1 automatically cancels the sleep mode (MCR5) by detecting CAN bus activity (dominant bit). If MCR7 is cleared the RCAN-TL1 does not automatically cancel the sleep mode.

RCAN-TL1 cannot store the message that wakes it up.

Note: This bit can be modified only Reset or Halt mode.

Bit 6 — Halt during Bus Off (MCR6): MCR6 enables or disables entering Halt mode immediately when MCR1 is set during Bus Off. This bit can be modified only in Reset or Halt mode. Please note that when Halt is entered in Bus Off the CAN engine is also recovering immediately to Error Active mode.

Bit 5 — Sleep Mode (MCR5): Enables or disables Sleep mode transition. If this bit is set, while RCAN-TL1 is in halt mode, the transition to sleep mode is enabled. Setting MCR5 is allowed after entering Halt mode. The two Error Counters (REC, TEC) will remain the same during Sleep mode. This mode will be exited in two ways:

- 1. by writing a '0' to this bit position,
- 2. or, if MCR[7] is enabled, after detecting a dominant bit on the CAN bus.

If Auto wake up mode is disabled, RCAN-TL1 will ignore all CAN bus activities until the sleep mode is terminated. When leaving this mode the RCAN-TL1 will synchronise to the CAN bus (by checking for 11 recessive bits) before joining CAN Bus activity. This means that, when the No.2 method is used, RCAN-TL1 will miss the first message to receive. CAN transceivers stand-by mode will also be unable to cope with the first message when exiting stand by mode, and the S/W needs to be designed in this manner.

In sleep mode only the following registers can be accessed: MCR, GSR, IRR and IMR.

Important: RCAN-TL1 is required to be in Halt mode before requesting to enter in Sleep mode. That allows the CPU to clear all pending interrupts before entering sleep mode. Once all interrupts are cleared RCAN-TL1 must leave the Halt mode and enter Sleep mode simultaneously (by writing $MCR[5] = 1$ and $MCR[1] = 0$ at the same time).

Bit 4 — Reserved. The written value should always be '0' and the returned value is '0'.

Bit 3 — Reserved. The written value should always be '0' and the returned value is '0'.

Bit 2 — Message Transmission Priority (MCR2): MCR2 selects the order of transmission for pending transmit data. If this bit is set, pending transmit data are sent in order of the bit position in the Transmission Pending Register (TXPR). The order of transmission starts from Mailbox-31 as the highest priority, and then down to Mailbox-1 (if those mailboxes are configured for transmission). Please note that this feature cannot be used for time trigger transmission of the Mailboxes 24 to 30.

If MCR2 is cleared, all messages for transmission are queued with respect to their priority (by running internal arbitration). The highest priority message has the Arbitration Field (STDID + IDE bit + EXTID (if IDE = 1) + RTR bit) with the lowest digital value and is transmitted first. The internal arbitration includes the RTR bit and the IDE bit (internal arbitration works in the same

way as the arbitration on the CAN Bus between two CAN nodes starting transmission at the same time).

This bit can be modified only in Reset or Halt mode.

Bit 1—Halt Request (MCR1): Setting the MCR1 bit causes the CAN controller to complete its current operation and then enter Halt mode (where it is cut off from the CAN bus). The RCAN-TL1 remains in Halt Mode until the MCR1 is cleared. During the Halt mode, the CAN Interface does not join the CAN bus activity and does not store messages or transmit messages. All the user registers (including Mailbox contents and TEC/REC) remain unchanged with the exception of IRR0 and GSR4 which are used to notify the halt status itself. If the CAN bus is in idle or intermission state regardless of MCR6, RCAN-TL1 will enter Halt Mode within one Bit Time. If MCR6 is set, a halt request during Bus Off will be also processed within one Bit Time. Otherwise the full Bus Off recovery sequence will be performed beforehand. Entering the Halt Mode can be notified by IRR0 and GSR4.

If both MCR14 and MCR6 are set, MCR1 is automatically set as soon as RCAN-TL1 enters BusOff.

In the Halt mode, the RCAN-TL1 configuration can be modified with the exception of the Bit Timing setting, as it does not join the bus activity. MCR[1] has to be cleared by writing a '0' in order to re-join the CAN bus. After this bit has been cleared, RCAN-TL1 waits until it detects 11 recessive bits, and then joins the CAN bus.

- Notes: 1. After issuing a Halt request the CPU is not allowed to set TXPR or TXCR or clear MCR1 until the transition to Halt mode is completed (notified by IRR0 and GSR4). After MCR1 is set this can be cleared only after entering Halt mode or through a reset operation (SW or HW).
	- 2. Transition into or recovery from HALT mode, is only possible if the BCR1 and BCR0 registers are configured to a proper Baud Rate.

Bit 0 — Reset Request (MCR0): Controls resetting of the RCAN-TL1 module. When this bit is changed from '0' to '1' the RCAN-TL1 controller enters its reset routine, re-initialising the internal logic, which then sets GSR3 and IRR0 to notify the reset mode. During a re-initialisation, all user registers are initialised.

RCAN-TL1 can be re-configured while this bit is set. This bit has to be cleared by writing a '0' to join the CAN bus. After this bit is cleared, the RCAN-TL1 module waits until it detects 11 recessive bits, and then joins the CAN bus. The Baud Rate needs to be set up to a proper value in order to sample the value on the CAN Bus.

After Power On Reset, this bit and GSR3 are always set. This means that a reset request has been made and RCAN-TL1 needs to be configured.

The Reset Request is equivalent to a Power On Reset but controlled by Software.

(2) General Status Register (GSR)

The General Status Register (GSR) is a 16-bit read-only register that indicates the status of RCAN-TL1.

 GSR (Address = $H'002$)

Bits 15 to 6: Reserved. The written value should always be '0' and the returned value is '0'.

Bit 5 — Error Passive Status Bit (GSR5): Indicates whether the CAN Interface is in Error Passive or not. This bit will be set high as soon as the RCAN-TL1 enters the Error Passive state and is cleared when the module enters again the Error Active state (this means the GSR5 will stay high during Error Passive and during Bus Off). Consequently to find out the correct state both GSR5 and GSR0 must be considered.

Bit 4 — Halt/Sleep Status Bit (GSR4): Indicates whether the CAN engine is in the halt/sleep state or not. Please note that the clearing time of this flag is not the same as the setting time of IRR12.

Please note that this flag reflects the status of the CAN engine and not of the full RCAN-TL1 IP. RCAN-TL1 exits sleep mode and can be accessed once MCR5 is cleared. The CAN engine exits sleep mode only after two additional transmission clocks on the CAN Bus.

Bit 3 — Reset Status Bit (GSR3): Indicates whether the RCAN-TL1 is in the reset state or not.

Bit 2 — Message Transmission in progress Flag (GSR2): Flag that indicates to the CPU if the RCAN-TL1 is in Bus Off or transmitting a message or an error/overload flag due to error detected during transmission. The timing to set TXACK is different from the time to clear GSR2. TXACK is set at the $7th$ bit of End Of Frame. GSR2 is set at the $3rd$ bit of intermission if there are no more messages ready to be transmitted. It is also set by arbitration lost, bus idle, reception, reset or halt transition.

Bit 1—Transmit/Receive Warning Flag (GSR1): Flag that indicates an error warning.

requested by the Bus Off recovery sequence. However the flag GSR1 is not set in Bus Off.

Bit 0—Bus Off Flag (GSR0): Flag that indicates that RCAN-TL1 is in the bus off state.

Note: Only the lower 8 bits of TEC are accessible from the user interface. The $9th$ bit is equivalent to GSR0.

(3) Bit Configuration Register (BCR0, BCR1)

The bit configuration registers (BCR0 and BCR1) are 2 X 16-bit read/write register that are used to set CAN bit timing parameters and the baud rate pre-scaler for the CAN Interface.

The Time quanta is defined as:

$$
Timequanta = \frac{2 * BRP}{f_{clk}}
$$

Where: BRP (Baud Rate Pre-scaler) is the value stored in BCR0 incremented by 1 and fclk is the used peripheral clock frequency.

 $BCR1$ (Address = H 004)

Bits 15 to 12 — Time Segment 1 (TSG1[3:0] = $BCR1[15:12]$ **):** These bits are used to set the segment TSEG1 (= PRSEG + PHSEG1) to compensate for edges on the CAN Bus with a positive phase error. A value from 4 to 16 time quanta can be set.

Bit 15: TSG1[3] TSG1[2] TSG1[1] TSG1[0] Description Bit 14: Bit 13: Bit 12:

Bit 11: Reserved. The written value should always be '0' and the returned value is '0'.

Bit 8:

Bit 10: Bit 9:

Bits 7 and 6: Reserved. The written value should always be '0' and the returned value is '0'.

Bits 5 and 4 - ReSynchronisation Jump Width (SJW[1:0] = BCR0[5:4]): These bits set the synchronisation jump width.

Bits 3 to 1: Reserved. The written value should always be '0' and the returned value is '0'.

Bit 0 — Bit Sample Point (BSP = BCR1[0]): Sets the point at which data is sampled.

$BCR0$ (Address = $H'006$)

Bits 8 to 15: Reserved. The written value should always be '0' and the returned value is '0'.

Bits 7 to 0—Baud Rate Pre-scale (BRP[7:0] = BCR0 [7:0]): These bits are used to define the peripheral clock periods contained in a Time Quantum.

• Requirements of Bit Configuration Register

- SYNC_SEG: Segment for establishing synchronisation of nodes on the CAN bus. (Normal bit edge transitions occur in this segment.)
- PRSEG: Segment for compensating for physical delay between networks.
- PHSEG1: Buffer segment for correcting phase drift (positive). (This segment is extended when synchronisation (resynchronisation) is established.)
- PHSEG2: Buffer segment for correcting phase drift (negative). (This segment is shortened when synchronisation (resynchronisation) is established)

TSEG1: TSG1 + 1

 $TSEG2: TSG2 + 1$

The RCAN-TL1 Bit Rate Calculation is:

Bit Rate = $\frac{\text{fclk}}{2 \times (\text{BRP} + 1) \times (\text{TSEG1} + \text{TSEG2} + 1)}$

Where BRP is given by the register value and TSEG1 and TSEG2 are derived values from TSG1 and TSG2 register values. The '+1' in the above formula is for the Sync-Seg which duration is 1 time quanta.

 f_{C1K} = Peripheral Clock

BCR Setting Constraints

TSEG1min > TSEG2 \geq SJWmax (SJW = 1 to 4) $8 \leq$ TSEG1 + TSEG2 + 1 \leq 25 time quanta (TSEG1 + TSEG2 + 1 = 7 is not allowed) $TSEG2 > 2$

These constraints allow the setting range shown in the table below for TSEG1 and TSEG2 in the Bit Configuration Register. The number in the table shows possible setting of SJW. "No" shows that there is no allowed combination of TSEG1 and TSEG2.

Example 1: To have a Bit rate of 500Kbps with a frequency of fclk = 32MHz it is possible to set: $BRP = 1$, $TSEG1 = 11$, $TSEG2 = 4$.

Then the configuration to write is $BCR1 = H'A300$ and $BCR0 = H'0001$.

Example 2: To have a Bit rate of 500Kbps with a frequency of fclk = 20MHz it is possible to set: $BPR = 1$, $TSEG1 = 6$, $TSEG2 = 3$.

Then the configuration to write is $BCR1 = H'5200$ and $BCR0 = 0001$.

(4) Interrupt Request Register (IRR)

The interrupt register (IRR) is a 16-bit read/write-clearable register containing status flags for the various interrupt sources.

 $IRR (Address = H'008)$

Bit 15 — Timer Compare Match Interrupt 1 (IRR15): Indicates that a Compare-Match condition occurred to the Timer Compare Match Register 1 (TCMR1). When the value set in the TCMR1 matches to Cycle Time (TCMR1 = CYCTR), this bit is set.

Bit 14 — Timer Compare Match Interrupt 0 (IRR14): Indicates that a Compare-Match condition occurred to the Timer Compare Match Register 0 (TCMR0). When the value set in the TCMR0 matches to Local Time (TCMR0 = TCNTR), this bit is set.

Bit 13 - Timer Overrun Interrupt/Next_is_Gap Reception Interrupt/Message Error Interrupt (IRR13): This interrupt assumes a different meaning depending on the RCAN-TL1 mode. It indicates that:

• The Timer (TCNTR) has overrun when RCAN-TL1 is working in event-trigger mode (including test modes)

- Time reference message with Next is Gap set has been received when working in time-trigger mode. Please note that when a Next_is_Gap is received the application is responsible to stop all transmission at the end of the current basic cycle (including test modes)
- Message error has occurred when in test mode. Note: If a Message Overload condition occurs when in Test Mode, then this bit will not be set.

Bit 12 – Bus activity while in sleep mode (IRR12): IRR12 indicates that a CAN bus activity is present. While the RCAN-TL1 is in sleep mode and a dominant bit is detected on the CAN bus, this bit is set. This interrupt is cleared by writing a '1' to this bit position. Writing a '0' has no effect. If auto wakeup is not used and this interrupt is not requested it needs to be disabled by the related interrupt mask register. If auto wake up is not used and this interrupt is requested it should be cleared only after recovering from sleep mode. This is to avoid that a new falling edge of the reception line causes the interrupt to get set again.

Please note that the setting time of this interrupt is different from the clearing time of GSR4.

Bit 11 — Timer Compare Match Interrupt 2 (IRR11): Indicates that a Compare-Match condition occurred to the Timer Compare Match Register 2 (TCMR2). When the value set in the TCMR2 matches to Cycle Time (TCMR2 = CYCTR), this bit is set.

Bit 10 — Start of new system matrix Interrupt (IRR10): Indicates that a new system matrix is starting.

When $CCR = 0$, this bit is set at the successful completion of reception/transmission of time reference message. Please note that when $CMAX = 0$ this interrupt is set at every basic cycle.

Bit 9 – Message Overrun/Overwrite Interrupt Flag (IRR9): Flag indicating that a message has been received but the existing message in the matching Mailbox has not been read as the corresponding RXPR or RFPR is already set to '1' and not yet cleared by the CPU. The received message is either abandoned (overrun) or overwritten dependant upon the NMC (New Message Control) bit. This bit is cleared when all bit in UMSR (Unread Message Status Register) are cleared (by writing '1') or by setting MBIMR (MailBox interrupt Mast Register) for all UMSR flag set. It is also cleared by writing a '1' to all the correspondent bit position in MBIMR. Writing to this bit position has no effect.

Bit 8 - Mailbox Empty Interrupt Flag (IRR8): This bit is set when one of the messages set for transmission has been successfully sent (corresponding TXACK flag is set) or has been successfully aborted (corresponding ABACK flag is set). In Event Triggered mode the related TXPR is also cleared and this mailbox is now ready to accept a new message data for the next transmission. In Time Trigger mode TXPR for the Mailboxes from 30 to 24 is not cleared after a successful transmission in order to keep transmitting at each programmed basic cycle. In effect, this bit is set by an OR'ed signal of the TXACK and ABACK bits not masked by the corresponding MBIMR flag. Therefore, this bit is automatically cleared when all the TXACK and ABACK bits are cleared. It is also cleared by writing a '1' to all the correspondent bit position in MBIMR. Writing to this bit position has no effect.

Bit 7 - Overload Frame (IRR7): Flag indicating that the RCAN-TL1 has detected a condition that should initiate the transmission of an overload frame. Note that in the condition of transmission being prevented, such as listen only mode, an Overload Frame will NOT be transmitted, but IRR7 will still be set. IRR7 remains asserted until reset by writing a '1' to this bit position - writing a '0' has no effect.

Bit 6 - Bus Off Interrupt Flag (IRR6): This bit is set when RCAN-TL1 enters the Bus-off state or when RCAN-TL1 leaves Bus-off and returns to Error-Active. The cause therefore is the existing condition TEC \geq 256 at the node or the end of the Bus-off recovery sequence (128X11) consecutive recessive bits) or the transition from Bus Off to Halt (automatic or manual). This bit remains set even if the RCAN-TL1 node leaves the bus-off condition, and needs to be explicitly cleared by S/W. The S/W is expected to read the GSR0 to judge whether RCAN-TL1 is in the busoff or error active status. It is cleared by writing a '1' to this bit position even if the node is still bus-off. Writing a '0' has no effect.

Bit 5 - Error Passive Interrupt Flag (IRR5): Interrupt flag indicating the error passive state caused by the transmit or receive error counter or by Error Passive forced by test mode. This bit is reset by writing a '1' to this bit position, writing a '0' has no effect. If this bit is cleared the node may still be error passive. Please note that the SW needs to check GSR0 and GSR5 to judge whether RCAN-TL1 is in Error Passive or Bus Off status.

Bit 4 - Receive Error Counter Warning Interrupt Flag (IRR4): This bit becomes set if the receive error counter (REC) reaches a value greater than 95 when RCAN-TL1 is not in the Bus Off status. The interrupt is reset by writing a '1' to this bit position, writing '0' has no effect.

Bit 3 - Transmit Error Counter Warning Interrupt Flag (IRR3): This bit becomes set if the transmit error counter (TEC) reaches a value greater than 95. The interrupt is reset by writing a '1' to this bit position, writing '0' has no effect.

Bit 2 - Remote Frame Receive Interrupt Flag (IRR2): Flag indicating that a remote frame has been received in a mailbox. This bit is set if at least one receive mailbox, with related MBIMR not set, contains a remote frame transmission request. This bit is automatically cleared when all bits in the Remote Frame Receive Pending Register (RFPR), are cleared. It is also cleared by writing a '1' to all the correspondent bit position in MBIMR. Writing to this bit has no effect.

Bit 1 – Data Frame Received Interrupt Flag (IRR1): IRR1 indicates that there are pending Data Frames received. If this bit is set at least one receive mailbox contains a pending message. This bit is cleared when all bits in the Data Frame Receive Pending Register (RXPR) are cleared, i.e. there is no pending message in any receiving mailbox. It is in effect a logical OR of the RXPR flags from each configured receive mailbox with related MBIMR not set. It is also cleared by writing a '1' to all the correspondent bit position in MBIMR. Writing to this bit has no effect.

Bit 0 – Reset/Halt/Sleep Interrupt Flag (IRR0): This flag can get set for three different reasons. It can indicate that:

- 1. Reset mode has been entered after a SW (MCR0) or HW reset
- 2. Halt mode has been entered after a Halt request (MCR1)
- 3. Sleep mode has been entered after a sleep request (MCR5) has been made while in Halt mode.

The GSR may be read after this bit is set to determine which state RCAN-TL1 is in.

Important: When a Sleep mode request needs to be made, the Halt mode must be used beforehand. Please refer to the MCR5 description and Figure 20.15 Halt Mode/Sleep Mode.

IRR0 is set by the transition from "0" to "1" of GSR3 or GSR4 or by transition from Halt mode to Sleep mode. So, IRR0 is not set if RCAN-TL1 enters Halt mode again right after exiting from Halt mode, without GSR4 being cleared. Similarly, IRR0 is not set by direct transition from Sleep mode to Halt Request. At the transition from Halt/Sleep mode to Transition/Reception, clearing GSR4 needs (one-bit time - TSEG2) to (one-bit time \times 2 - TSEG2).

In the case of Reset mode, IRR0 is set, however, the interrupt to the CPU is not asserted since IMR0 is automatically set by initialisation.

(5) Interrupt Mask Register (IMR)

The interrupt mask register is a 16 bit register that protects all corresponding interrupts in the Interrupt Request Register (IRR) from generating an output signal on the IRQ. An interrupt request is masked if the corresponding bit position is set to '1'. This register can be read or written at any time. The IMR directly controls the generation of IRQ, but does not prevent the setting of the corresponding bit in the IRR.

 IMR (Address = $H'00A$)

Bit 15 to 0: Maskable interrupt sources corresponding to IRR[15:0] respectively. When a bit is set, the interrupt signal is not generated, although setting the corresponding IRR bit is still performed.

(6) Transmit Error Counter (TEC) and Receive Error Counter (REC)

The Transmit Error Counter (TEC) and Receive Error Counter (REC) is a 16-bit read/(write) register that functions as a counter indicating the number of transmit/receive message errors on the CAN Interface. The count value is stipulated in the CAN protocol specification Refs. [1], [2], [3] and [4]. When not in (Write Error Counter) test mode this register is read only, and can only be modified by the CAN Interface. This register can be cleared by a Reset request (MCR0) or entering to bus off.

In Write Error Counter test mode (i.e. $TST[2:0] = 3'b100$), it is possible to write to this register. The same value can only be written to TEC/REC, and the value written into TEC is set to TEC and REC. When writing to this register, RCAN-TL1 needs to be put into Halt Mode. This feature is only intended for test purposes.

• TEC/REC (Address = H'00C)

		- 13	- 12 - 11		10 9 8 7 6 5 4 3 2 1					
										TEC7 TEC6 TEC5 TEC4 TEC3 TEC2 TEC1 TEC0 REC7 REC6 REC5 REC4 REC3 REC2 REC1 REC0
Initial value:	\sim 0	$\mathbf{0}$	$\overline{0}$	$\overline{0}$	0 0 0 0 0 0 0 0 0 0					

Note: * It is only possible to write the value in test mode when TST[2:0] in MCR is 3'b100. REC is incremented during Bus Off to count the recurrences of 11 recessive bits as requested by the Bus Off recovery sequence.

20.3.4 RCAN-TL1 Mailbox Registers

The following sections describe RCAN-TL1 Mailbox registers that control/flag individual Mailboxes. The address is mapped as follows.

Important: LongWord access is carried out as two consecutive Word accesses.

Figure 20.11 RCAN-TL1 Mailbox registers

(1) Transmit Pending Register (TXPR1, TXPR0)

The concatenation of TXPR1 and TXPR0 is a 32-bit register that contains any transmit pending flags for the CAN module. In the case of 16-bit bus interface, Long Word access is carried out as two consecutive word accesses.

<Longword Write Operation>

The TXPR1 controls Mailbox-31 to Mailbox-16, and the TXPR0 controls Mailbox-15 to Mailbox-1. The CPU may set the TXPR bits to affect any message being considered for transmission by writing a '1' to the corresponding bit location. Writing a '0' has no effect, and TXPR cannot be cleared by writing a '0' and must be cleared by setting the corresponding TXCR bits. TXPR may be read by the CPU to determine which, if any, transmissions are pending or in progress. In effect there is a transmit pending bit for all Mailboxes except for the Mailbox-0. Writing a '1' to a bit location when the mailbox is not configured to transmit is not allowed.

In Event Triggered Mode RCAN-TL1 will clear a transmit pending flag after successful transmission of its corresponding message or when a transmission abort is requested successfully from the TXCR. In Time Trigger Mode, TXPR for the Mailboxes from 30 to 24 is NOT cleared after a successful transmission, in order to keep transmitting at each programmed basic cycle. The TXPR flag is not cleared if the message is not transmitted due to the CAN node losing the arbitration process or due to errors on the CAN bus, and RCAN-TL1 automatically tries to transmit it again unless its DART bit (Disable Automatic Re-Transmission) is set in the Message-Control of the corresponding Mailbox. In such case (DART set), the transmission is cleared and notified through Mailbox Empty Interrupt Flag (IRR8) and the correspondent bit within the Abort Acknowledgement Register (ABACK).

If the status of the TXPR changes, the RCAN-TL1 shall ensure that in the identifier priority scheme (MCR2 = 0), the highest priority message is always presented for transmission in an intelligent way even under circumstances such as bus arbitration losses or errors on the CAN bus. Please refer to the Application Note for details.

When the RCAN-TL1 changes the state of any TXPR bit position to a '0', an empty slot interrupt (IRR8) may be generated. This indicates that either a successful or an aborted mailbox transmission has just been made. If a message transmission is successful it is signalled in the TXACK register, and if a message transmission abortion is successful it is signalled in the ABACK register. By checking these registers, the contents of the Message of the corresponding Mailbox may be modified to prepare for the next transmission.

• TXPR1

Note: * It is possible only to write a '1' for a Mailbox configured as transmitter.

Bit 15 to 0 — Requests the corresponding Mailbox to transmit a CAN Frame. The bit 15 to 0 corresponds to Mailbox-31 to 16 respectively. When multiple bits are set, the order of the transmissions is governed by the MCR2 – CAN-ID or Mailbox number.

• TXPR0

Note: * It is possible only to write a '1' for a Mailbox configured as transmitter.

Bit 15 to 1 — Indicates that the corresponding Mailbox is requested to transmit a CAN Frame. The bit 15 to 1 corresponds to Mailbox-15 to 1 respectively. When multiple bits are set, the order of the transmissions is governed by the MCR2 – CAN-ID or Mailbox number.

Bit 0— Reserved: This bit is always '0' as this is a receive-only Mailbox. Writing a '1' to this bit position has no effect. The returned value is '0'.

(2) Transmit Cancel Register (TXCR1, TXCR0)

The TXCR1 and TXCR0 are 16-bit read/conditionally-write registers. The TXCR1 controls Mailbox-31 to Mailbox-16, and the TXCR0 controls Mailbox-15 to Mailbox-1.This register is used by the CPU to request the pending transmission requests in the TXPR to be cancelled. To clear the corresponding bit in the TXPR the CPU must write a '1' to the bit position in the TXCR. Writing a '0' has no effect.

When an abort has succeeded the CAN controller clears the corresponding TXPR + TXCR bits, and sets the corresponding ABACK bit. However, once a Mailbox has started a transmission, it cannot be cancelled by this bit. In such a case, if the transmission finishes in success, the CAN controller clears the corresponding TXPR + TXCR bit, and sets the corresponding TXACK bit, however, if the transmission fails due to a bus arbitration loss or an error on the bus, the CAN controller clears the corresponding TXPR + TXCR bit, and sets the corresponding ABACK bit. If an attempt is made by the CPU to clear a mailbox transmission that is not transmit-pending it has no effect. In this case the CPU will be not able at all to set the TXCR flag.

• TXCR1

Bit:	15	14	13	12	11	10	9		8 7	6	- 5	4	- 3			
	TXCR1[15:0]															
Initial value:			Ω	Ω	Ω	$\mathbf{0}$	Ω	Ω	Ω	$\mathbf{0}$	Ω	Ω				
R/W :	B/W^*	B/W^*	B/W^*	B/W^*	B/W^*	R/W^*	R/W* R/W* R/W*			R/W^*	R/W^*	B/W^*	B/W^*	B/W^*	B/W^*	B/W^*

Note: * Only writing a '1' to a Mailbox that is requested for transmission and is configured as transmit.

Bit 15 to 0 — Requests the corresponding Mailbox, that is in the queue for transmission, to cancel its transmission. The bit 15 to 0 corresponds to Mailbox-31 to 16 (and TXPR1[15:0]) respectively.

• TXCR0

Note: * Only writing a '1' to a Mailbox that is requested for transmission and is configured as transmit.

Bit 15 to 1 — Requests the corresponding Mailbox, that is in the queue for transmission, to cancel its transmission. The bit 15 to 1 corresponds to Mailbox-15 to 1 (and TXPR0[15:1]) respectively.

Bit 0 — This bit is always '0' as this is a receive-only mailbox. Writing a '1' to this bit position has no effect and always read back as a '0'.

(3) Transmit Acknowledge Register (TXACK1, TXACK0)

The TXACK1 and TXACK0 are 16-bit read/conditionally-write registers. These registers are used to signal to the CPU that a mailbox transmission has been successfully made. When a transmission has succeeded the RCAN-TL1 sets the corresponding bit in the TXACK register. The CPU may clear a TXACK bit by writing a '1' to the corresponding bit location. Writing a '0' has no effect.

• TXACK1

Bit 15 to 0 — Notifies that the requested transmission of the corresponding Mailbox has been finished successfully. The bit 15 to 0 corresponds to Mailbox-31 to 16 respectively.

• TXACK0

Note: * Only when writing a '1' to clear.

Bit 15 to 1 — Notifies that the requested transmission of the corresponding Mailbox has been finished successfully. The bit 15 to 1 corresponds to Mailbox-15 to 1 respectively.

Bit[15:1]:TXACK0 Description

Bit 0 — This bit is always '0' as this is a receive-only mailbox. Writing a '1' to this bit position has no effect and always read back as a '0'.

(4) Abort Acknowledge Register (ABACK1, ABACK0)

The ABACK1 and ABACK0 are 16-bit read/conditionally-write registers. These registers are used to signal to the CPU that a mailbox transmission has been aborted as per its request. When an abort has succeeded the RCAN-TL1 sets the corresponding bit in the ABACK register. The CPU may clear the Abort Acknowledge bit by writing a '1' to the corresponding bit location. Writing a '0' has no effect. An ABACK bit position is set by the RCAN-TL1 to acknowledge that a TXPR bit has been cleared by the corresponding TXCR bit.

• ABACK1

Note: $*$ Only when writing a '1' to clear.

Bit[15:0]:ABACK1 Description

Bit 15 to 0 — Notifies that the requested transmission cancellation of the corresponding Mailbox has been performed successfully. The bit 15 to 0 corresponds to Mailbox-31 to 16 respectively.

• ABACK0

Note: $*$ Only when writing a '1' to clear.

Bit 15 to 1 — Notifies that the requested transmission cancellation of the corresponding Mailbox has been performed successfully. The bit 15 to 1 corresponds to Mailbox-15 to 1 respectively.

Bit[15:1]:ABACK0 Description

Bit 0 — This bit is always '0' as this is a receive-only mailbox. Writing a '1' to this bit position has no effect and always read back as a '0'.

(5) Data Frame Receive Pending Register (RXPR1, RXPR0)

The RXPR1 and RXPR0 are 16-bit read/conditionally-write registers. The RXPR is a register that contains the received Data Frames pending flags associated with the configured Receive Mailboxes. When a CAN Data Frame is successfully stored in a receive mailbox the corresponding bit is set in the RXPR. The bit may be cleared by writing a '1' to the corresponding bit position. Writing a '0' has no effect. However, the bit may only be set if the mailbox is configured by its MBC (Mailbox Configuration) to receive Data Frames. When a RXPR bit is set, it also sets IRR1 (Data Frame Received Interrupt Flag) if its MBIMR (Mailbox Interrupt Mask Register) is not set, and the interrupt signal is generated if IMR1 is not set. Please note that these bits are only set by receiving Data Frames and not by receiving Remote frames.

• RXPR1

Note : * Only when writing a '1' to clear.

Bit 15 to 0 — Configurable receive mailbox locations corresponding to each mailbox position from 31 to 16 respectively.

• RXPR0

Note: * Only when writing a '1' to clear.

Bit 15 to 0 — Configurable receive mailbox locations corresponding to each mailbox position from 15 to 0 respectively.

(6) Remote Frame Receive Pending Register (RFPR1, RFPR0)

The RFPR1 and RFPR0 are 16-bit read/conditionally-write registers. The RFPR is a register that contains the received Remote Frame pending flags associated with the configured Receive Mailboxes. When a CAN Remote Frame is successfully stored in a receive mailbox the corresponding bit is set in the RFPR. The bit may be cleared by writing a '1' to the corresponding bit position. Writing a '0' has no effect. In effect there is a bit position for all mailboxes. However, the bit may only be set if the mailbox is configured by its MBC (Mailbox Configuration) to receive Remote Frames. When a RFPR bit is set, it also sets IRR2 (Remote Frame Receive Interrupt Flag) if its MBIMR (Mailbox Interrupt Mask Register) is not set, and the interrupt signal is generated if IMR2 is not set. Please note that these bits are only set by receiving Remote Frames and not by receiving Data frames.

• RFPR1

Note: $*$ Only when writing a '1' to clear.

• RFPR0

Bit 15 to 0 — Remote Request pending flags for mailboxes 15 to 0 respectively.

(7) Mailbox Interrupt Mask Register (MBIMR)

The MBIMR1 and MBIMR0 are 16-bit read/write registers. The MBIMR only prevents the setting of IRR related to the Mailbox activities, that are IRR[1] – Data Frame Received Interrupt, IRR[2] – Remote Frame Receive Interrupt, IRR[8] – Mailbox Empty Interrupt, and IRR[9] – Message OverRun/OverWrite Interrupt. If a mailbox is configured as receive, a mask at the corresponding bit position prevents the generation of a receive interrupt (IRR[1] and IRR[2] and IRR[9]) but does not prevent the setting of the corresponding bit in the RXPR or RFPR or UMSR. Similarly when a mailbox has been configured for transmission, a mask prevents the generation of an Interrupt signal and setting of an Mailbox Empty Interrupt due to successful transmission or abortion of transmission (IRR[8]), however, it does not prevent the RCAN-TL1 from clearing the corresponding TXPR/TXCR bit + setting the TXACK bit for successful transmission, and it does not prevent the RCAN-TL1 from clearing the corresponding TXPR/TXCR bit + setting the ABACK bit for abortion of the transmission.

A mask is set by writing a '1' to the corresponding bit position for the mailbox activity to be masked. At reset all mailbox interrupts are masked.

• MBIMR1

Bit[15:0]: MBIMR1 Description

• MBIMR0

Bit 15 to 0 — Enable or disable interrupt requests from individual Mailbox-15 to Mailbox-0 respectively.

Bit[15:0]: MBIMR0 Description

(8) Unread Message Status Register (UMSR)

This register is a 32-bit read/conditionally write register and it records the mailboxes whose contents have not been accessed by the CPU prior to a new message being received. If the CPU has not cleared the corresponding bit in the RXPR or RFPR when a new message for that mailbox is received, the corresponding UMSR bit is set to '1'. This bit may be cleared by writing a '1' to the corresponding bit location in the UMSR. Writing a '0' has no effect.

If a mailbox is configured as transmit box, the corresponding UMSR will not be set.

• UMSR1

Bit 15 to 0 — Indicate that an unread received message has been overwritten or overrun condition has occurred for Mailboxes 31 to 16.

Bit[15:0]: UMSR1 Description

• UMSR0

Bit 15 to 0 — Indicate that an unread received message has been overwritten or overrun condition has occurred for Mailboxes 15 to 0.

20.3.5 Timer Registers

The Timer is 16 bits and supports several source clocks. A pre-scale counter can be used to reduce the speed of the clock. It also supports three Compare Match Registers (TCMR2, TCMR1, TCMR0). The address map is as follows.

Important: These registers can only be accessed in Word size (16-bit).

Figure 20.12 RCAN-TL1 Timer registers

(1) Time Trigger Control Register0 (TTCR0)

The Time Trigger Control Register0 is a 16-bit read/write register and provides functions to control the operation of the Timer. When operating in Time Trigger Mode, please refer to section 20.4.3 (1), Time Triggered Transmission.

• TTCR0 (Address $=$ H'080)

Bit 15 — Enable Timer: When this bit is set, the timer TCNTR is running. When this bit is cleared, TCNTR and CCR are cleared.

Bit 14 — TimeStamp value: Specifies if the Timestamp for transmission and reception in Mailboxes 15 to 0 must contain the Cycle Time (CYCTR) or the concatenation of CCR[5:0] + CYCTR[15:6]. This feature is very useful for time triggered transmission to monitor Rx_Trigger.

This register does not affect the TimeStamp for Mailboxes 30 and 31.

Bit 13 — Cancellation by TCMR2: The messages in the transmission queue are cancelled by setting TXCR, when both this bit and bit12 are set and compare match occurs when RCAN-TL1 is not in the Halt status, causing the setting of all TXCR bits with the corresponding TXPR bits set.

Bit 12 — TCMR2 compare match enable: When this bit is set, IRR11 is set by TCMR2 compare match.

Bit12 TTCR0 12 Description

Bit 11 — TCMR1 compare match enable: When this bit is set, IRR15 is set by TCMR1 compare match.

Bit 10 — TCMR0 compare match enable: When this bit is set, IRR14 is set by TCMR0 compare match.

Bits 9 to 7: Reserved. The written value should always be '0' and the returned value is '0'.

Bit 6 — Timer Clear-Set Control by TCMR0: Specifies if the Timer is to be cleared and set to H'0000 when the TCMR0 matches to the TCNTR. Please note that the TCMR0 is also capable to generate an interrupt signal to the CPU via IRR14.

Note: If RCAN-TL1 is working in TTCAN mode (CMAX isn't 3'b111), TTCR0 bit6 has to be '0' to avoid clearing Local Time.

Bit5 to 0 — RCAN-TL1 Timer Prescaler (TPSC[5:0]): This control field allows the timer source clock (4*[RCAN-TL1 system clock]) to be divided before it is used for the timer. This function is available only in event-trigger mode. In time trigger mode (CMAX is not 3'b111), one nominal Bit Timing (= one bit length of CAN bus) is automatically chosen as source clock of **TCNTR**

The following relationship exists between source clock period and the timer period.

(2) Cycle Maximum/Tx-Enable Window Register (CMAX_TEW)

This register is a 16-bit read/write register. CMAX specifies the maximum value for the cycle counter (CCR) for TT Transmissions to set the number of basic cycles in the matrix system. When the Cycle Counter reaches the maximum value (CCR = CMAX), after a full basic cycle, it is cleared to zero and an interrupt is generated on IRR.10.

TEW specifies the width of Tx-Enable window.

 $CMAX$ TEW (Address = H'084)

Bits 15 to 11: Reserved. The written value should always be '0' and the returned value is '0'.

Bit 10 to 8 — Cycle Count Maximum (CMAX): Indicates the maximum number of CCR. The number of basic cycles available in the matrix cycle for Timer Triggered transmission is (Cycle Count Maximum + 1).

Unless $CMAX = 3th111$, $RCAN-TL1$ is in time-trigger mode and time trigger function is available. If CMAX = 3'b111, RCAN-TL1 is in event-trigger mode.

Important: Please set CMAX = 3'b111 when event-trigger mode is used.

Bits 7 to 4: Reserved. The written value should always be '0' and the returned value is '0'.

Bit 3 to $0 - Tx$ **-Enable Window (TEW):** Indicates the width of Tx-Enable Window. TEW = H'00 shows the width is one nominal Bit Timing. All values from 0 to 15 are allowed to be set.

Note: The CAN core always needs a time between 1 to 2 bit timing to initiate transmission. The above values are not considering this accuracy.

(3) Reference Trigger Offset Register (RFTROFF)

This is a 8-bit read/write register that affects Tx-Trigger Time (TTT) of Mailbox-30. The TTT of Mailbox-30 is compared with CYCTR after RFTROFF extended with sign is added to the TTT. However, the value of TTT is not modified. The offset value doesn't affect others except Mailbox-30.

• RFTROFF (Address = $H'086$)

Bit 15 to 8 — Indicate the value of Reference Trigger Offset.

(4) Timer Status Register (TSR)

This register is a 16-bit read-only register, and allows the CPU to monitor the Timer Compare Match status and the Timer Overrun Status.

• TSR $(Address = H'088)$

Bits 15 to 5: Reserved. The written value should always be '0' and the returned value is '0'.

Bit 4 to 0 — RCAN-TL1 Timer Status (TSR[4:0]): This read-only field allows the CPU to monitor the status of the Cycle Counter, the Timer and the Compare Match registers. Writing to this field has no effect.

Bit 4 — Start of New System Matrix (TSR4): Indicates that a new system matrix is starting. When $CCR = 0$, this bit is set at the successful completion of reception/transmission of time reference message.

Bit 3 — Timer Compare Match Flag 2 (TSR3): Indicates that a Compare-Match condition occurred to the Timer Compare Match Register 2 (TCMR2). When the value set in the TCMR2 matches to Cycle Time Register (TCMR2 = CYCTR), this bit is set if TTCR0 bit12 = 1. Please note that this bit is read-only and is cleared when IRR11 (Timer Compare Match Interrupt 2) is cleared.

Bit 2 — Timer Compare Match Flag 1 (TSR2): Indicates that a Compare-Match condition occurred to the Timer Compare Match Register 1 (TCMR1). When the value set in the TCMR1 matches to Cycle Time Register (TCMR1 = CYCTR), this bit is set if TTCR0 bit11 = 1. Please note that this bit is read-only and is cleared when IRR15 (Timer Compare Match Interrupt 1) is cleared.

Bit 1 — Timer Compare Match Flag 0 (TSR1): Indicates that a Compare-Match condition occurred to the Compare Match Register 0 (TCMR0). When the value set in the TCMR0 matches to the Timer value (TCMR0 = TCNTR), this bit is set if TTCR0 bit10 = 1. Please note that this bit is read-only and is cleared when IRR14 (Timer Compare Match Interrupt 0) is cleared.

Bit 0 — Timer Overrun/Next_is_Gap Reception/Message Error (TSR0): This flag is assigned to three different functions. It indicates that the Timer has overrun when working in event-trigger mode, time reference message with Next_is_Gap set has been received in time-trigger mode, and error detected on the CAN bus has occurred in test mode, respectively. Test mode has higher priority with respect to the other settings.

(5) Cycle Counter Register (CCR)

This register is a 6-bit read/write register. Its purpose is to store the number of the basic cycle for Time -Triggered Transmissions. Its value is updated in different fashions depending if RCAN-TL1 is programmed to work as a potential time master or as a time slave. If RCAN-TL1 is working as (potential) time master, CCR is:

- ⎯ Incremented by one every time the cycle time (CYCTR) matches to Tx-Trigger Time of Mailbox-30 or
- Overwritten with the value contained in MSG_DATA_0[5:0] of Mailbox 31 when a valid reference message is received.

If RCAN-TL1 is working as a time slave, CCR is only overwritten with the value of MSG_DATA_0[5:0] of Mailbox 31 when a valid reference message is received.

If $CMAX = 3'111$, CCR is always H'0000.

 CCR (Address = H'08A)

Bits 15 to 6: Reserved. The written value should always be '0' and the returned value is '0'.

Bit 5 to 0 — Cycle Counter Register (CCR): Indicates the number of the current Base Cycle of the matrix cycle for Timer Triggered transmission.

(6) Timer Counter Register (TCNTR)

This is a 16-bit read/write register that allows the CPU to monitor and modify the value of the Free Running Timer Counter. When the Timer meets TCMR0 (Timer Compare Match Register 0) + TTCR0 [6] is set to '1', the TCNTR is cleared to H'0000 and starts running again. In Time-Trigger mode, this timer can be used as Local Time and TTCR0[6] has to be cleared to work as a free running timer.

- Notes: 1. It is possible to write into this register only when it is enabled by the bit 15 in TTCR0. If TTCR0 bit15 = 0, TCNTR is always $H'0000$.
	- 2. There could be a delay of a few clock cycles between the enabling of the timer and the moment where TCNTR starts incrementing. This is caused by the internal logic used for the pre-scaler.
- TCNTR (Address = H'08C)

Note: * The register can be written only when enabled in TTCR0[15]. Write operation is not allowed in Time Trigger mode (i.e. CMAX is not 3'b111).

Bit 15 to 0 — Indicate the value of the Free Running Timer.

(7) Cycle Time register (CYCTR)

This register is a 16-bit read-only register. This register shows Cycle Time = Local Time (TCNTR) - Reference_Mark (RFMK). In ET mode this register is the exact copy of TCNTR as RFMK is always fixed to zero.

 $CYCTR (Address = H'090)$

(8) Reference Mark Register (RFMK)

This register is a 16-bit read-only register. The purpose of this register is to capture Local Time (TCNTR) at SOF of the reference message when the message is received or transmitted successfully. In ET mode this register is not used and it is always cleared to zero.

 $RFMK (Address = H'094)$

Bit 15 to 0 — Reference Mark Register (RFMK): Indicates the value of TCNTR at SOF of time reference message.

(9) Timer Compare Match Registers (TCMR0, TCMR1, TCMR2)

These three registers are 16-bit read/write registers and are capable of generating interrupt signals, clearing-setting the Timer value (only supported by TCMR0) or clear the transmission messages in the queue (only supported by TCMR2). TCMR0 is compared with TCNTR, however, TCMR1 and TCMR2 are compared with CYCTR.

The value used for the compare can be configured independently for each register. In order to set flags, TTCR0 bit 12-10 needs to be set.

In Time-Trigger mode, TTCR0 bit6 has to be cleared by software to prevent TCNTR from being cleared.

TMCR0 is for Init_Watch_Trigger, and TCMR2 is for Watch_Trigger.

Interrupt:

The interrupts are flagged by the Bit11, Bit15 and 14 in the IRR accordingly when a Compare Match occurs, and setting these bits can be enabled by Bit12, Bit11, Bit10 in TTCR0. The generation of interrupt signals itself can be prevented by the Bit11, Bit15 and Bit14 in the IMR. When a Compare Match occurs and the IRR11 (or IRR15 or IRR14) is set, the Bit3 or Bit2 or Bit1 in the TSR (Timer Status Register) is also set. Clearing the IRR bit also clears the corresponding bit of TSR.

Timer Clear-Set:

The Timer value can only be cleared when a Compare Match occurs if it is enabled by the Bit6 in the TTCR0. TCMR1 and TCMR2 do not have this function.

Cancellation of the messages in the transmission queue:

The messages in the transmission queue can only be cleared by the TCMR2 through setting TXCR when a Compare Match occurs while RCAN-TL1 is not in the halt status. TCMR1 and TCMR0 do not have this function.

 $TCMR0$ (Address = $H'098$)

Bit 15 to 0 — Timer Compare Match Register (TCMR0): Indicates the value of TCNTR when compare match occurs.

 $TCMR1$ (Address = H'09C)

Bit 15 to 0 — Timer Compare Match Register (TCMR1): Indicates the value of CYCTR when compare match occurs.

• TCMR2 (Address = $H'0A0$)

Bit 15 to 0 — Timer Compare Match Register (TCMR2): Indicates the value of CYCTR when compare match occurs.

(10) Tx-Trigger Time Selection Register (TTTSEL)

This register is a 16-bit read/write register and specifies the Tx-Trigger Time waiting for compare match with Cycle Time. Only one bit is allowed to be set. Please don't set more bits than one, or clear all bits.

This register may only be modified during configuration mode. The modification algorithm is shown in figure 20.13.

Please note that this register is only indented for test and diagnosis. When not in test mode, this register must not be written to and the returned value is not guaranteed.

• TTTSEL (Address = H'0A4)

Note: Only one bit is allowed to be set.

Bit 15: Reserved. The written value should always be '0' and the returned value is '0'.

Bit 14 to 8 — Specifies the Tx-Trigger Time waiting for compare match with CYCTR The bit 14 to 8 corresponds to Mailbox-30 to 24, respectively.

Bits 7 to 0: Reserved. The written value should always be '0' and the returned value is '0'.

Figure 20.13 TTTSEL modification algorithm

20.4 Application Note

20.4.1 Test Mode Settings

The RCAN-TL1 has various test modes. The register TST[2:0] (MCR[10:8]) is used to select the RCAN-TL1 test mode. The default (initialised) settings allow RCAN-TL1 to operate in Normal mode. The following table is examples for test modes.

Test Mode can be selected only while in configuration mode. The user must then exit the configuration mode (ensuring BCR0/BCR1 is set) in order to run the selected test mode.

When message error occurs, IRR13 is set in all test modes.

20.4.2 Configuration of RCAN-TL1

RCAN-TL1 is considered in configuration mode or after a H/W (Power On Reset)/S/W (MCR[0]) reset or when in Halt mode. In both conditions RCAN-TL1 cannot join the CAN Bus activity and configuration changes have no impact on the traffic on the CAN Bus.

After a Reset request

The following sequence must be implemented to configure the RCAN-TL1 after (S/W or H/W) reset. After reset, all the registers are initialised, therefore, RCAN-TL1 needs to be configured before joining the CAN bus activity. Please read the notes carefully.

Figure 20.14 Reset Sequence

• Halt mode

When RCAN-TL1 is in Halt mode, it cannot take part to the CAN bus activity. Consequently the user can modify all the requested registers without influencing existing traffic on the CAN Bus. It is important for this that the user waits for the RCAN-TL1 to be in halt mode before to modify the requested registers - note that the transition to Halt Mode is not always immediate (transition will occurs when the CAN Bus is idle or in intermission). After RCAN-TL1 transit to Halt Mode, GSR4 is set.

Once the configuration is completed the Halt request needs to be released. RCAN-TL1 will join CAN Bus activity after the detection of 11 recessive bits on the CAN Bus.

• Sleep mode

When RCAN-TL1 is in sleep mode the clock for the main blocks of the IP is stopped in order to reduce power consumption. Only the following user registers are clocked and can be accessed: MCR, GSR, IRR and IMR. Interrupt related to transmission (TXACK and ABACK) and reception (RXPR and RFPR) cannot be cleared when in sleep mode (as TXACK, ABACK, RXPR and RFPR are not accessible) and must to be cleared beforehand.

Figure 20.15 shows allowed state transitions.

- Please don't set MCR5 (Sleep Mode) without entering Halt Mode.
- ⎯ After MCR1 is set, please don't clear it before GSR4 is set and RCAN-TL1 enters Halt Mode.

Figure 20.15 Halt Mode/Sleep Mode

- Notes: 1. MCR5 can be cleared by automatically by detecting a dominant bit on the CAN Bus if MCR7 is set or by writing '0'.
	- 2. MCR1 is cleared in SW. Clearing MCR1 and setting MCR5 have to be carried out by the same instruction.
	- 3. MCR1 must not be cleared in SW, before GSR4 is set. MCR1 can be set automatically in HW when RCAN-TL1 moves to Bus Off and MCR14 and MCR6 are both set.
	- 4. When MCR5 is cleared and MCR1 is set at the same time, RCAN-TL1 moves to Halt Request. Right after that, it moves to Halt Mode with no reception/transmission.

The following table shows conditions to access registers.

RCAN-TL1 Registers

Notes: 1. No hardware protection.

2. When TXPR is not set.

20.4.3 Message Transmission Sequence

• Message Transmission Request

The following sequence is an example to transmit a CAN frame onto the bus. As described in the previous register section, please note that IRR8 is set when one of the TXACK or ABACK bits is set, meaning one of the Mailboxes has completed its transmission or transmission abortion and is now ready to be updated for the next transmission, whereas, the GSR2 means that there is currently no transmission request made (No TXPR flags set).

Figure 20.16 Transmission request

Internal Arbitration for transmission

The following diagram explains how RCAN-TL1 manages to schedule transmission-requested messages in the correct order based on the CAN identifier. 'Internal arbitration' picks up the highest priority message amongst transmit-requested messages.

Figure 20.17 Internal Arbitration for transmission

The RCAN-TL1 has two state machines. One is for transmission, and the other is for reception.

- 1-1: When a TXPR bit(s) is set while the CAN bus is idle, the internal arbitration starts running immediately and the transmission is started.
- 1-2: Operations for both transmission and reception starts at SOF. Since there is no reception frame, RCAN-TL1 becomes transmitter.
- 2-1: At crc delimiter, internal arbitration to search next message transmitted starts.
- 2-2: Operations for both transmission and reception starts at SOF. Because of a reception frame with higher priority, RCAN-TL1 becomes receiver. Therefore, Reception is carried out instead of transmitting Frame-3.
- 3-1: At crc delimiter, internal arbitration to search next message transmitted starts.
- 3-2: Operations for both transmission and reception starts at SOF. Since a transmission frame has higher priority than reception one, RCAN-TL1 becomes transmitter.

Internal arbitration for the next transmission is also performed at the beginning of each error delimiter in case of an error is detected on the CAN Bus. It is also performed at the beginning of error delimiters following overload frame.

As the arbitration for transmission is performed at CRC delimiter, in case a remote frame request is received into a Mailbox with $ATX = 1$ the answer can join the arbitration for transmission only at the following Bus Idle, CRC delimiter or Error Delimiter.

Depending on the status of the CAN bus, following the assertion of the TXCR, the corresponding Message abortion can be handled with a delay of maximum 1 CAN Frame.

(1) Time Triggered Transmission

RCAN-TL1 offers a H/W support to perform communication in Time Trigger mode in line with the emerging ISO-11898-4 Level 1 Specification.

This section reports the basic procedures to use this mode.

• Setting Time Trigger Mode

In order to set up the time trigger mode the following settings need to be used.

- CMAX in CMAX_TEW must be programmed to a value different from 3'b111.
- ⎯ Bit 15 in TTCR0 has to be set, to start TCNTR.
- $-$ Bit 6 in TTCR0 has to be cleared to prevent TCNTR from being cleared after a match.
- DART in Mailboxes used for time-triggered transmission cannot be used, since for Time Triggered Mailboxes, TXPR is not cleared to support periodic transmission.
- Roles of Registers

The user registers of RCAN-TL1 can be used to handle the main functions requested by the TTCAN standard.

• Time Master/Time Slave

RCAN-TL1 can be programmed to work as a potential time master of the network or as a time slave. The following table shows the settings and the operation automatically performed by RCAN-TL1 in each mode.

• Setting Tx-Trigger Time

The Tx-Trigger Time(TTT) must be set in ascending order shown below, and the difference between them has to satisfy the following expressions. TEW in the following expressions is the register value.

TTT (Mailbox-24) < TTT (Mailbox-25) < TTT (Mailbox-26) < TTT (Mailbox-27) < TTT (Mailbox-28) < TTT (Mailbox-29) < TTT (Mailbox-30)

and

TTT (Mailbox-i) – TTT (Mailbox-i-1) > TEW + the maximum frame length $+9$

TTT (Mailbox-24) to TTT (Mailbox-29) correspond to Time_Marks, and TTT (Mailbox-30) corresponds to Time_Ref showing the length of a basic cycle, respectively when working as potential time master.

The above limitation is not applied to mailboxes which are not set as time-triggered transmission.

Important: Because of limitation on setting Tx-Trigger Time, only one Mailbox can be assigned to one time window.

Figure 20.18 Limitation on Tx-Trigger Time

The value of TCMR2 as Watch_Trigger has to be larger than TTT(Mailbox-30), which shows the length of a basic cycle.

Figures 20.19 and 20.20 show examples of configurations for (Potential) Time Master and Time Slave. "L" in diagrams shows the length in time of the time reference messages.

Figure 20.19 (Potential) Time Master

Figure 20.20 Time Slave

Function to be implemented by software

Some of the TTCAN functions need to be implemented in software. The main details are reported hereafter. Please refer to ISO-11898-4 for more details.

Change from Init_Watch_Trigger to Watch_Trigger

RCAN-TL1 offers the two registers TCMR0 and TCMR2 as H/W support for Init_Watch_Trigger and Watch_Trigger respectively. The SW is requested to enable TCMR0 and disable TCMR2 up to the first reference message is detected on the CAN Bus and then disable TCMR0 and enable TCMR2.- Schedule Synchronization state machine.

Only reception of Next_is_Gap interrupt is supported. The application needs to take care of stopping all transmission at the end of the current basic cycle by setting the related TXCR flags.Master-Slave Mode control.

Only automatic cycle time synchronization and CCR increment is supported.

Message status count

Software has to count scheduling errors for periodic messages in exclusive windows.

• Message Transmission Request for Time Triggered communication

When the Time Triggered mode is used communications must fulfils the ISO11898-4 requirements.

The following procedure should be used.

- ⎯ Send RCAN-TL1 to reset or halt mode
- Set TCMR0 to the Init_Watch_Trigger (0xFFFF)
- Enable TCMR0 compare match setting bit 10 of TTCR0
- Set TCMR2 to the specified Watch Trigger value
- Keep TCMR2 compare match disabled by keeping cleared the bit 12 of TTCR0
- Set CMAX to the requested value (different from 111 bin)
- Set TEW to the requested value
- Configure the necessary Mailboxes for Time Trigger transmission and reception
- Set LAFM for the 3 LSBs of Mailbox 31
- Configure MCR, BCR1 and BCR0 to the requested values
- If working as a potential time master:
	- Set RFTROFF to the requested Init_Ref_Offset value
	- Set TXPR for Mailbox 30
	- Write H'4000 into TTTSEL
- Enable the TCNTR timer through the bit 15 of TTCR0
- Move to Transmission Reception mode
- Wait for the reception or transmission of a valid reference message or for TCMR0 match
- If the local time reaches the value of TCMR0 the Init Watch Trigger is reached and the application needs to set TXCR for Mailbox 30 and start again
- If the reference message is transmitted (TXACK[30] is set) set RFTROFF to zero
- ⎯ If a valid reference message is received (RXPR[31] is set) then:
	- If 3 LSBs of ID of Mailbox 31 have high priority than the 3 LSBs of Mailbox 30 (if working as potential time master) keep RFTROFF to Init_Ref_Offset
	- If 3 LSBs of ID of Mailbox 31 have lower priority than the 3 LSBs of Mailbox 30 (if working as potential time master) decrement by 1 the value in RFTROFF
- Disable TCMR0 compare match by clearing bit 10 of TTCR0
- Enable TCMR2 compare match by setting bit 12 of TTCR0
- ⎯ Only after two reference messages have been detected on the CAN Bus (transmitted or received) can the application set TXPR for the other Time Triggered Mailboxes.

If, at any time, a reference message cannot be detected on the CAN Bus, and the cycle time CYCTR reaches TCMR2, RCAN-TL1 automatically aborts all pending transmissions (including the Reference Message).

The following is the sequence to request further transmission in Time Triggered mode.

Figure 20.21 Message transmission request

S/W has to ensure that a message is updated before a Tx trigger for transmission occurs.

When the CYCTR reaches to TTT (Tx-Trigger Time) of a Mailbox and CCR matches with the programmed cycle for transmission, RCAN-TL1 immediately transfers the message into the Tx buffer. At this point, RCAN-TL1 will attempt a transmission within the specified Time Enable Window. If RCAN-TL1 misses this time slot, it will suspend the transmission request up to the next Tx Trigger, keeping the corresponding TXPR bit set to '1' if the transmission is periodic (Mailbox-24 to 30). There are three factors that may cause RCAN-TL1 to miss the time slot –

- 1. The CAN bus currently used
- 2. An error on the CAN bus during the time triggered message transmission
- 3. Arbitration loss during the time triggered message transmission

In case of Merged Arbitrating Window the slot for transmission goes from the Tx_Trig of the Mailbox opening the Window ($TTW = 10 \text{ bin}$) to the end to the TEW of the Mailbox closing the Window ($TTW = 11$ bin). The TXPR can be modified at any time. RCAN-TL1 ensures the transmission of Time Triggered messages is always scheduled correctly. However, in order to guarantee the correct schedule, there are some important rules that are :

- ⎯ TTT (Tx Trigger Time) can be modified during configuration mode.
- ⎯ TTT cannot be set outside the range of Time_Ref, which specifies the length of basic cycle. This could cause a scheduling problem.
- ⎯ TXPR is not automatically cleared for periodic transmission. If a periodic transmission needs to be cancelled, the corresponding TXCR bit needs to be set by the application.
- Example of Time Triggered System

The following diagram shows a simple example of how time trigger system works using RCAN-TL1 in time slave mode.

Figure 20.22 Example of Time trigger system as Time Slave

The following settings were used in the above example:

During merged arbitrating window, request by time-triggered transmission is served in the way of FCFS (First Come First Served). For example, if Mailbox-25 cannot be transmitted between Tx-Trigger Time 25 (TTT25) and TTT26, Mailbox-25 has higher priority than Mailbox-26 between TTT26 and 28.

MBC needs to be set into 3'b111, in order to disable time-triggered transmission. If RCAN-TL1 is Time Master, MBC[30] has to be 3'b000 and time reference window is automatically recognized as arbitrating window.

• Timer Operation

Figure 20.23 shows the timing diagram of the timer. By setting Tx-Trigger Time = n, time trigger transmission starts between CYCTR $= n + 2$ and $CYCTR = n + 3$.

Figure 20.23 Timing Diagram of Timer

During merged arbitrating window, event-trigger transmission is served after completion of timetriggered transmission. For example, If transmission of Mailbox-25 is completed and CYCTR doesn't reach TTT26, event-trigger transmission starts based on message transmission priority specified by MCR2. TXPR of time-triggered transmission is not cleared after transmission completion, however, that of event-triggered transmission is cleared.

Note: that in the case that the TXPR is not set for the Mailbox which is assigned to close the Merged Arbitrating Window (MAW), then the MAW will still be closed (at the end of the TEW following the TTT of the assigned Mailbox.

Please refer to Table Roles of Mailboxes in section 20.3.2, Mailbox Structure.

20.4.4 Message Receive Sequence

The diagram below shows the message receive sequence.

Figure 20.24 Message receive sequence

When RCAN-TL1 recognises the end of the Arbitration field while receiving a message, it starts comparing the received identifier to the identifiers set in the Mailboxes, starting from Mailbox-31 down to Mailbox-0. It first checks the MBC if it is configured as a receive box, and reads LAFM, and reads the CAN-ID of Mailbox-31 (if configured as receive) to finally compare them to the received ID. If it does not match, the same check takes place at Mailbox-30 (if configured as receive). Once RCAN-TL1 finds a matching identifier, it stores the number of Mailbox-[N] into an internal buffer, stops the search, and goes back to idle state, waiting for the EndOfFrame (EOF) to come. When the $6th$ bit of EOF is notified by the CAN Interface logic, the received message is written or abandoned, depending on the NMC bit. No modification of configuration during communication is allowed. Entering Halt Mode is one of ways to modify configuration. If it is written into the corresponding Mailbox, including the CAN-ID, i.e., there is a possibility that the CAN-ID is overwritten by a different CAN-ID of the received message due to the LAFM used. This also implies that, if the identifier of a received message matches to ID + LAFM of 2 or more Mailboxes, the higher numbered Mailbox will always store the relevant messages and the lower numbered Mailbox will never receive messages. Therefore, the settings of the identifiers and LAFMs need to be carefully selected.

With regards to the reception of data and remote frames described in the above flow diagram the clearing of the UMSR flag after the reading of IRR is to detect situations where a message is overwritten by a new incoming message stored in the same mailbox (if its $NMC = 1$) while the interrupt service routine is running. If during the final check of UMSR a overwrite condition is detected the message needs to be discarded and read again.

In case UMSR is set and the Mailbox is configured for overrun ($NMC = 0$) the message is still valid, however it is obsolete as it is not reflecting the latest message monitored on the CAN Bus.

Please access the full Mailbox content before clearing the related RXPR/RFPR flag.

Please note that in the case a received remote frame is overwritten by a data frame, both the remote frame receive interrupt (IRR2) and data frame received interrupt (IRR1) and also the Receive Flags (RXPR and RFPR) are set. In an analogous way, the overwriting of a data frame by a remote frame, leads to setting both IRR2 and IRR1.

When a message is received and stored into a Mailbox all the fields of the data not received are stored as zero. The same applies when a standard frame is received. The extended identifier part (EXTID[17:0]) is written as zero.

20.4.5 Reconfiguration of Mailbox

When re-configuration of Mailboxes is required, the following procedures should be taken.

• Change configuration of transmit box

Two cases are possible.

- Change of ID, RTR, IDE, LAFM, Data, DLC, NMC, ATX, DART This change is possible only when $MBC = 3'6000$. Confirm that the corresponding TXPR is not set. The configuration (except MBC bit) can be changed at any time.
- Change from transmit to receive configuration (MBC)

Confirm that the corresponding TXPR is not set. The configuration can be changed only in Halt or reset state. Please note that it might take longer for RCAN-TL1 to transit to halt state if it is receiving or transmitting a message (as the transition to the halt state is delayed until the end of the reception/transmission), and also RCAN-TL1 will not be able to receive/transmit messages during the Halt state.

In case RCAN-TL1 is in the Bus Off state the transition to halt state depends on the configuration of the bit 6 of MCR and also bit and 14 of MCR.

• Change configuration (ID, RTR, IDE, LAFM, Data, DLC, NMC, ATX, DART, MBC) of receiver box or Change receiver box to transmitter box

The configuration can be changed only in Halt Mode.

RCAN-TL1 will not lose a message if the message is currently on the CAN bus and RCAN-TL1 is a receiver. RCAN-TL1 will be moving into Halt Mode after completing the current reception. Please note that it might take longer if RCAN-TL1 is receiving or transmitting a message (as the transition to the halt state is delayed until the end of the reception/transmission), and also RCAN-TL1 will not be able to receive/transmit messages during the Halt Mode.

In case RCAN-TL1 is in the Bus Off state the transition to halt mode depends on the configuration of the bit 6 and 14 of MCR.

Figure 20.25 Change ID of receive box or Change receive box to transmit box

20.5 Interrupt Sources

Table 20.2 lists the RCAN-TL1 interrupt sources. These sources can be masked. Masking is implemented using the mailbox interrupt mask registers (MBIMR) and interrupt mask register (IMR). For details on the interrupt vector of each interrupt source, see section 7, Interrupt Controller (INTC).

Table 20.2 RCAN-TL1-n*¹ **Interrupt Sources**

Notes: $1. n = 0, 1$

 2. RM0 is an interrupt generated by the remote request pending flag for mailbox 0 (RFPR0[0]) or the data frame receive flag for mailbox 0 (RXPR0[0]). RM1 is an interrupt generated by the remote request pending flag for mailbox n (RFPR0[n]) or the data frame receive flag for mailbox n (RXPR0[n]) ($n = 1$ to 31).

- 3. IRR1 is a data frame received interrupt flag for mailboxes 0 to 31, and IRR2 is a remote frame request interrupt flag for mailboxes 0 to 31.
- 4. The DMAC is activated only by an RM0n interrupt.

20.6 DMAC Interface

The DMAC can be activated by the reception of a message in RCAN-TL1 mailbox 0. When DMAC transfer ends after DMAC activation has been set, flags of RXPR0 and RFPR0 are cleared automatically. An interrupt request due to a receive interrupt from the RCAN-TL1 cannot be sent to the CPU in this case. Figure 20.26 shows a DMAC transfer flowchart.

Figure 20.26 DMAC Transfer Flowchart

20.7 CAN Bus Interface

A bus transceiver IC is necessary to connect this LSI to a CAN bus. A Renesas HA13721 transceiver IC and its compatible products are recommended. As the CRx and CTx pins use 3 V, an external level shifter is necessary. Figure 20.27 shows a sample connection diagram.

Figure 20.27 High-Speed CAN Interface Using HA13721

20.8 Setting I/O Ports for RCAN-TL1

The I/O ports for the RCAN-TL1 must be specified before or during the configuration mode. For details on the settings of I/O ports, see section 27, Pin Function Controller (PFC). Two methods are available using two channels of the RCAN-TL1 in this LSI.

- Using RCAN-TL1 as a 2-channel module (channels 0 and 1) Each channel has 32 Mailboxes.
- Using RCAN-TL1 as a 1-channel module (channels 0 and 1 functioning as a single channel)

When the second method is used, see section 20.9.1, Notes on Port Setting for Multiple Channels Used as Single Channel.

Figures 20.28 and 20.29 show connection examples for individual port settings.

Figure 20.28 Connection Example when Using RCAN-TL1 as 2-Channel Module (32 Mailboxes × **2 Channels)**

Figure 20.29 Connection Example when Using RCAN-TL1 as 1-Channel Module (64 Mailboxes × **1 Channel)**

20.9 Usage Notes

20.9.1 Notes on Port Setting for Multiple Channels Used as Single Channel

The RCAN-TL1 in this LSI has two channels and some of these channels can be used as a single channel. When using multiple channels as a single channel, keep the following in mind.

Figure 20.30 Connection Example when Using RCAN-TL1 as 1-Channel Module (64 Mailboxes × **1 Channel)**

1. No ACK error is detected even when any other nodes are not connected to the CAN bus. This occurs when channel 1 transmits an ACK in the ACK field in response to a message channel 0 has transmitted.

Channel 1 receives a message which channel 0 has transmitted on the CAN bus and then transmits an ACK in the ACK field. After that, channel 0 receives the ACK.

To avoid this, make channel 1 which is not currently used for transmission the listen-only mode $(TST[2:0] = B'001)$ or the reset state $(MCR0 = 1)$. With this setting, only a channel which transmits a message transmits an ACK.

2. Internal arbitration for channels 0 and 1 is independently controlled to determine the order of transmission.

Although the internal arbitration is performed on 31 Mailboxes at a time, it is not performed on 64 Mailboxes at a time even though multiple channels function as a single channel.

3. Do not set the same transmission message ID in both channels 0 and 1.

Two messages may be transmitted from the two channels after arbitration on the CAN bus.

Section 21 A/D Converter (ADC)

This LSI includes a 10-bit successive-approximation A/D converter allowing selection of up to eight analog input channels.

21.1 Features

- Resolution: 10 bits
- Input channels: 8
- Minimum conversion time: 3.9 μs per channel
- Absolute accuracy: ± 4 LSB
- Operating modes: 3
	- Single mode: A/D conversion on one channel
	- Multi mode: A/D conversion on one to four channels or on one to eight channels
	- Scan mode: Continuous A/D conversion on one to four channels or on one to eight channels
- Data registers: 8

Conversion results are held in a 16-bit data register for each channel

- Sample-and-hold function
- A/D conversion start methods: 3
	- Software
	- Conversion start trigger from multi-function timer pulse unit 2 (MTU2)
	- External trigger signal
- Interrupt source

An A/D conversion end interrupt (ADI) request can be generated on completion of A/D conversion.

Module standby mode can be set

Figure 21.1 Block Diagram of A/D Converter

21.2 Input/Output Pins

Table 21.1 summarizes the A/D converter's input pins.

Table 21.1 Pin Configuration

21.3 Register Descriptions

The A/D converter has the following registers.

Table 21.2 Register Configuration

21.3.1 A/D Data Registers A to H (ADDRA to ADDRH)

The sixteen A/D data registers, ADDRA to ADDRH, are 16-bit read-only registers that store the results of A/D conversion.

An A/D conversion produces 10-bit data, which is transferred for storage into the ADDR corresponding to the selected channel. The 10 bits of the result are stored in the upper bits (bits 15 to 6) of ADDR. Bits 5 to 0 of ADDR are reserved bits that are always read as 0.

Access to ADDR in 8-bit units is prohibited. ADDR must always be accessed in 16-bit units.

Table 21.3 indicates the pairings of analog input channels and ADDR.

Table 21.3 Analog Input Channels and ADDR

21.3.2 A/D Control/Status Register (ADCSR)

ADCSR is a 16-bit readable/writable register that selects the mode, controls the A/D converter, and enables or disables starting of A/D conversion by external trigger input.

Note: $*$ Only 0 can be written to clear the flag after 1 is read.

[Legend]

- x: Don't care
- Notes: 1. The flag can only be cleared by writing 0 to it after reading it as 1. However, in the following cases as well the flag is cleared by writing 0 to it:
	- (1) When the CPU reads the value of ADF as 1
	- (2) When ADF is cleared to 0 by the DMAC reading ADDR
	- (3) When the ADF flag is set to 1 at A/D conversion end
	- (4) When the CPU writes 0 to ADF
	- 2. Set the A/D conversion time to minimum or more values to meet the absolute accuracy of the A/D conversion characteristics.
- 3. t_{poyc} indicates the peripheral clock (P ϕ) cycle.

21.4 Operation

The A/D converter uses the successive-approximation method, and the resolution is 10 bits. It has three operating modes: single mode, multi mode, and scan mode. Switching the operating mode or analog input channels must be done while the ADST bit in ADCSR is 0 to prevent incorrect operation. The ADST bit can be set at the same time as the operating mode or analog input channels are changed.

21.4.1 Single Mode

Single mode should be selected when only A/D conversion on one channel is required.

In single mode, A/D conversion is performed once for the specified one analog input channel, as follows:

- 1. A/D conversion for the selected channel starts when the ADST bit in ADCSR is set to 1 by software, MTU2, or external trigger input.
- 2. When A/D conversion is completed, the A/D conversion result is transferred to the A/D data register corresponding to the channel.
- 3. After A/D conversion has completed, the ADF bit in ADCSR is set to 1. If the ADIE bit is set to 1 at this time, an ADI interrupt request is generated.
- 4. The ADST bit that remains 1 during A/D conversion is automatically cleared to 0 when A/D conversion is completed, and the A/D converter becomes idle.

When the operating mode or analog input channel selection must be changed during A/D conversion, to prevent incorrect operation, first clear the ADST bit to 0 to halt A/D conversion. After making the necessary changes, set the ADST bit to 1 to start A/D conversion again. The ADST bit can be set at the same time as the mode or channel selection is switched.

Typical operations when a single channel (AN1) is selected in single mode are described next. Figure 21.2 shows a timing diagram for this example (the bits which are set in this example belong to ADCSR).

- 1. Single mode is selected, input channel AN1 is selected $(CH[2:0] = 001)$, the A/D interrupt is enabled (ADIE = 1), and A/D conversion is started (ADST = 1).
- 2. When A/D conversion is completed, the A/D conversion result is transferred into ADDRB. At the same time the ADF flag is set to 1, the ADST bit is cleared to 0, and the A/D converter becomes idle.
- 3. Since $ADF = 1$ and $ADIE = 1$, an ADI interrupt is requested.
- 4. The A/D interrupt handling routine starts.
- 5. The routine reads $ADF = 1$, and then writes 0 to the ADF flag.
- 6. The routine reads and processes the A/D conversion result (ADDRB).
- 7. Execution of the A/D interrupts handling routine ends. Then, when the ADST bit is set to 1, A/D conversion starts and steps 2 to 7 are executed.

Figure 21.2 Example of A/D Converter Operation (Single Mode, One Channel (AN1) Selected)

21.4.2 Multi Mode

Multi mode should be selected when performing A/D conversion once on one or more channels.

In multi mode, A/D conversion is performed once for a maximum of eight specified analog input channels, as follows:

- 1. A/D conversion starts from the analog input channel with the lowest number (e.g. AN0, AN1, …, AN3) when the ADST bit in ADCSR is set to 1 by software, MTU2, or external trigger input.
- 2. When A/D conversion is completed on each channel, the A/D conversion result is sequentially transferred to the A/D data register corresponding to that channel.
- 3. After A/D conversion on all selected channels has completed, the ADF bit in ADCSR is set to 1. If the ADIE bit is set to 1 at this time, an ADI interrupt request is generated.
- 4. The ADST bit that remains 1 during A/D conversion is automatically cleared to 0 when A/D conversion is completed, and the A/D converter becomes idle. If the ADST bit is cleared to 0 during A/D conversion, A/D conversion is halted and the A/D converter becomes idle. The ADF bit is cleared by reading ADF while $ADF = 1$, then writing 0 to the ADF bit.

A/D conversion is to be performed once on all the specified channels. The conversion results are transferred for storage into the A/D data registers corresponding to the channels.

When the operating mode or analog input channel selection must be changed during A/D conversion, to prevent incorrect operation, first clear the ADST bit to 0 to halt A/D conversion. After making the necessary changes, set the ADST bit to 1. A/D conversion will start again from the first channel in the group. The ADST bit can be set at the same time as the mode or channel selection is changed.

Typical operations when three channels (AN0 to AN2) are selected in multi mode are described next. Figure 21.3 shows a timing diagram for this example.

- 1. Multi mode is selected (MDS2 = 1, MDS1 = 0), analog input channels AN0 to AN2 are selected (CH $[2:0] = 010$), and A/D conversion is started (ADST = 1).
- 2. A/D conversion of the first channel (AN0) starts. When A/D conversion is completed, the A/D conversion result is transferred into ADDRA.
- 3. Next, the second channel (AN1) is selected automatically and A/D conversion starts.
- 4. Conversion proceeds in the same way through the third channel (AN2).
- 5. When conversion of all selected channels (AN0 to AN2) is completed, the ADF flag is set to 1 and the ADST bit cleared to 0.
- 6. If the ADIE bit is set to 1 at this time, an ADI interrupt is requested.

Figure 21.3 Example of A/D Converter Operation (Multi Mode, Three Channels (AN0 to AN2) Selected)

21.4.3 Scan Mode

Scan mode is useful for monitoring analog inputs in a group of one or more channels at all times. In scan mode, A/D conversion is performed sequentially for a maximum of eight specified analog input channels, as follows:

- 1. A/D conversion starts from the analog input channel with the lowest number (e.g. AN0, AN1, …, AN3) when the ADST bit in ADCSR is set to 1 by software, MTU2, or external trigger input.
- 2. When A/D conversion is completed on each channel, the A/D conversion result is sequentially transferred to the A/D data register corresponding to that channel.
- 3. After A/D conversion on all selected channels has completed, the ADF bit in ADCSR is set to 1. If the ADIE bit is set to 1 at this time, an ADI interrupt request is generated. The A/D converter starts A/D conversion again from the channel with the lowest number.
- 4. The ADST bit is not cleared automatically, so steps 2. and 3. are repeated as long as the ADST bit remains set to 1. When the ADST bit is cleared to 0, A/D conversion halts and the A/D converter becomes idle.

The ADF bit is cleared by reading ADF while $ADE = 1$, then writing 0 to the ADF bit.

When the operating mode or analog input channel selection must be changed during A/D conversion, to prevent incorrect operation, first clear the ADST bit to 0 to halt A/D conversion. After making the necessary changes, set the ADST bit to 1. A/D conversion will start again from the first channel in the group. The ADST bit can be set at the same time as the mode or channel selection is changed.

Typical operations when three channels (AN0 to AN2) are selected in scan mode are described as follows. Figure 21.4 shows a timing diagram for this example.

- 1. Scan mode is selected (MDS2 = 1, MDS1 = 1), analog input channels AN0 to AN2 are selected (CH $[2:0] = 010$), and A/D conversion is started (ADST = 1).
- 2. A/D conversion of the first channel (AN0) starts. When A/D conversion is completed, the A/D conversion result is transferred into ADDRA.
- 3. Next, the second channel (AN1) is selected automatically and A/D conversion starts.
- 4. Conversion proceeds in the same way through the third channel (AN2).
- 5. When conversion of all the selected channels (AN0 to AN2) is completed, the ADF flag is set to 1 and conversion of the first channel (AN0) starts again. If the ADIE bit is set to 1 at this time, an ADI interrupt is requested.

6. The ADST bit is not cleared automatically, so steps 2. to 4. are repeated as long as the ADST bit remains set to 1. When steps 2. to 4. are repeated, the ADF flag is kept to 1. When the ADST bit is cleared to 0, A/D conversion stops. The ADF bit is cleared by reading ADF while $ADF = 1$, then writing 0 to the ADF bit.

If both the ADF flag and ADIE bit are set to 1 while steps 2. to 4. are repeated, an ADI interrupt is requested at all times. To generate an interrupt on completing conversion of the third channel, clear the ADF bit to 0 after an interrupt is requested.

SH7205 Group Section 21 A/D Converter (ADC)

Figure 21.4 Example of A/D Converter Operation (Scan Mode, Three Channels (AN0 to AN2) Selected)

21.4.4 A/D Converter Activation by External Trigger or MTU2

The A/D converter can be independently activated by an external trigger or an A/D conversion request from the MTU2. To activate the A/D converter by an external trigger or the MTU2, set the A/D trigger enable bits (TRGS[3:0]). When an external trigger or an A/D conversion request from the MTU2 is generated with this bit setting, the ADST bit is set to 1 to start A/D conversion. The channel combination is determined by bits CH2 to CH0 in ADCSR. The timing from setting of the ADST bit until the start of A/D conversion is the same as when 1 is written to the ADST bit by software.

21.4.5 Input Sampling and A/D Conversion Time

The A/D converter has a built-in sample-and-hold circuit. The A/D converter samples the analog input at the A/D conversion start delay time (t_D) after the ADST bit in ADCSR is set to 1, then starts conversion. Figure 21.5 shows the A/D conversion timing. Table 21.4 indicates the A/D conversion time.

As indicated in figure 21.5, the A/D conversion time (t_{conv}) includes t_{p} and the input sampling time(t_{cyl}). The length of t_p varies depending on the timing of the write access to ADCSR. The total conversion time therefore varies within the ranges indicated in table 21.4.

In multi mode and scan mode, the values given in table 21.4 apply to the first conversion. In the second and subsequent conversions, time is the values given in table 21.5.

Table 21.4 A/D Conversion Time (Single Mode)

Note: Values in the table are the numbers of t_{ocyc} . t_{ocyc} indicates the peripheral clock (P ϕ) cycle.

Table 21.5 A/D Conversion Time (Multi Mode and Scan Mode)

Note: Values in the table are the numbers of t_{new} indicates the peripheral clock (P ϕ) cycle.

21.4.6 External Trigger Input Timing

A/D conversion can also be externally triggered. When the TRGS[3:0] bits in ADCSR are set to B'1001, an external trigger is input to the ADTRG pin. The ADST bit in ADCSR is set to 1 at the falling edge of the ADTRG pin, thus starting A/D conversion. Other operations, regardless of the operating mode, are the same as when the ADST bit has been set to 1 by software. Figure 21.6 shows the timing.

Figure 21.6 External Trigger Input Timing

21.5 Interrupt Sources and DMAC Transfer Request

The A/D converter generates an A/D conversion end interrupt (ADI) at the end of A/D conversion. An ADI interrupt request is generated if the ADIE bit is set to 1 when the ADF bit in ADCSR is set to 1 on completion of A/D conversion. Note that the interrupt controller (INTC) can be activated by an ADI interrupt depending on the DMAC setting. In this case, an interrupt is not issued to the CPU. If the setting to activate the DMAC has not been made, an interrupt request is sent to the CPU. Having the converted data read by the DMAC in response to an ADI interrupt enables continuous conversion to be achieved without imposing a load on software.

In single mode, set the DMAC so that DMA transfer initiated by an ADI interrupt is performed only once. In the case of A/D conversion on multiple channels in scan mode or multi mode, setting the DMA transfer count to one causes DMA transfer to finish after transferring only one channel of data. To make the DMAC transfer all conversion data, set the ADDR where A/D conversion data is stored as the transfer source address, the number of converted channels \times 2 as the transfer byte count, and continuous operand transfer or non-stop transfer as the DMA transfer condition.

When the DMAC is activated by ADI, the ADF bit in ADCSR is automatically cleared to 0 when data is transferred by the DMAC.

21.6 Definitions of A/D Conversion Accuracy

The A/D converter compares an analog value input from an analog input channel with its analog reference value and converts it to 10-bit digital data. The absolute accuracy of this A/D conversion is the deviation between the input analog value and the output digital value. It includes the following errors:

- Offset error
- Full-scale error
- Quantization error
- Nonlinearity error

These four error quantities are explained below with reference to figure 21.7. In the figure, the 10 bit A/D converter is illustrated as the 3-bit A/D converter for explanation. Offset error is the deviation between actual and ideal A/D conversion characteristics when the digital output value changes from the minimum (zero voltage) B'0000000000 (000 in the figure) to B'000000001 (001 in the figure)(figure 21.7, item (1)). Full-scale error is the deviation between actual and ideal A/D conversion characteristics when the digital output value changes from B'1111111110 (110 in the figure) to the maximum B'11111111111 (111 in the figure)(figure 21.7, item (2)). Quantization error is the intrinsic error of the A/D converter and is expressed as 1/2 LSB (figure 21.7, item (3)). Nonlinearity error is the deviation between actual and ideal A/D conversion characteristics between zero voltage and full-scale voltage (figure 21.7, item (4)). Note that it does not include offset, full-scale, or quantization error.

Figure 21.7 Definitions of A/D Conversion Accuracy

21.7 Usage Notes

When using the A/D converter, note the following points.

21.7.1 Module Standby Mode Setting

Operation of the A/D converter can be disabled or enabled using the standby control register. The initial setting is for operation of the A/D converter to be halted. Register access is enabled by clearing module standby mode. For details, see section 30, Power-Down Modes.

21.7.2 Setting Analog Input Voltage

Permanent damage to the LSI may result if the following voltage ranges are exceeded.

1. Analog input range

During A/D conversion, voltages on the analog input pins ANn should not go beyond the following range: $AVss \leq ANn \leq AVcc$ (n = 0 to 7).

2. AVcc and AVss input voltages

Input voltages AVcc and AVss should be PVcc -0.3 V \leq AVcc \leq PVcc and AVss = PVss. Do not leave the AVcc and AVss pins open when the A/D converter or D/A converter is not in use and in software standby mode. When not in use, connect AVcc to the power supply (PVcc) and AVss to the ground (PVss).

3. Setting range of AVref input voltage Set the reference voltage range of the AVref pin as $3.0 \text{ V} \leq \text{A} \text{V} \text{ref} \leq \text{A} \text{V} \text{cc}$.

21.7.3 Notes on Board Design

In board design, digital circuitry and analog circuitry should be as mutually isolated as possible, and layout in which digital circuit signal lines and analog circuit signal lines cross or are in close proximity should be avoided as far as possible. Failure to do so may result in incorrect operation of the analog circuitry due to inductance, adversely affecting A/D conversion values.

Digital circuitry must be isolated from the analog input signals (AN0 to AN7), analog reference voltage (AVref), and analog power supply (AVcc) by the analog ground (AVss). Also, the analog ground (AVss) should be connected at one point to a stable digital ground (PVss) on the board.

21.7.4 Processing of Analog Input Pins

To prevent damage from voltage surges at the analog input pins (AN0 to AN7), connect an input protection circuit like the one shown in figure 21.8. The circuit shown also includes a CR filter to suppress noise. This circuit is shown as an example; the circuit constants should be selected according to actual application conditions.

Figure 21.9 shows an equivalent circuit diagram of the analog input ports and table 21.7 lists the analog input pin specifications.

Table 21.7 Analog Input Pin Ratings

21.7.5 Permissible Signal Source Impedance

This LSI's analog input is designed such that conversion precision is guaranteed for an input signal for which the signal source impedance is $5 \text{ k}\Omega$ or less. This specification is provided to enable the A/D converter's sample-and-hold circuit input capacitance to be charged within the sampling time; if the sensor output impedance exceeds 5 kΩ, charging may be insufficient and it may not be possible to guarantee A/D conversion precision. However, for A/D conversion in single mode with a large capacitance provided externally for A/D conversion in single mode, the input load will essentially comprise only the internal input resistance of 3 k Ω , and the signal source impedance is ignored. However, as a low-pass filter effect is obtained in this case, it may not be possible to follow an analog signal with a large differential coefficient (e.g., $5 \text{ mV/}\mu s$ or greater) (see figure 21.10). When converting a high-speed analog signal, a low-impedance buffer should be inserted.

Figure 21.10 Example of Analog Input Circuit

21.7.6 Influences on Absolute Precision

Adding capacitance results in coupling with GND, and therefore noise in GND may adversely affect absolute precision. Be sure to connect AVss, etc. to an electrically stable GND.

Care is also required to insure that filter circuits do not communicate with digital signals on the mounting board (i.e., acting as antennas).

21.7.7 Usage Notes in Scan Mode or Multi Mode

Starting conversion immediately after stopping scan mode or multi mode can cause incorrect conversion results. To continue with conversion in which cases, allow a duration equivalent to the A/D conversion time for one channel to elapse after clearing ADST to 0 before starting conversion (by setting ADST to 1). (The A/D conversion time for one channel differs depending on the peripheral register settings.)

Section 22 D/A Converter (DAC)

22.1 Features

- 8-bit resolution
- Two output channels
- Minimum conversion time of 10 μ s (with 20 pF load)
- Output voltage of 0 V to AVref
- D/A output hold function in software standby mode
- Module standby mode can be set

Figure 22.1 Block Diagram of D/A Converter

22.2 Input/Output Pins

Table 22.1 shows the pin configuration of the D/A converter.

Table 22.1 Pin Configuration

22.3 Register Descriptions

The D/A converter has the following registers.

Table 22.2 Register Configuration

22.3.1 D/A Data Registers 0 and 1 (DADR0 and DADR1)

DADR is an 8-bit readable/writable register that stores data to which D/A conversion is to be performed. Whenever analog output is enabled, the values in DADR are converted and output to the analog output pins.

22.3.2 D/A Control Register (DACR)

DACR is an 8-bit readable/writable register that controls the operation of the D/A converter.

Table 22.3 Control of D/A Conversion

22.4 Operation

The D/A converter includes D/A conversion circuits for two channels, each of which can operate independently. When the DAOE bit in DACR is set to 1, D/A conversion is enabled and the conversion result is output.

An operation example of D/A conversion on channel 0 is shown below. Figure 22.2 shows the timing of this operation.

- 1. Write the conversion data to DADR0.
- 2. Set the DAOE0 bit in DACR to 1 to start D/A conversion. The conversion result is output from the analog output pin DA0 after the conversion time t_{PCONV} has elapsed. The conversion result continues to be output until DADR0 is written to again or the DAOE0 bit is cleared to 0. The output value is expressed by the following formula:

Contents of DADR ²⁵⁶ [⋅]AVref

- 3. If DADR0 is written to again, the conversion is immediately started. The conversion result is output after the conversion time t_{DCONV} has elapsed.
- 4. If the DAOE0 bit is cleared to 0, analog output is disabled.

Figure 22.2 Example of D/A Converter Operation

22.5 Usage Notes

22.5.1 Module Standby Mode Setting

Operation of the D/A converter can be disabled or enabled using the standby control register. The initial setting is for operation of the D/A converter to be halted. Register access is enabled by canceling module standby mode. For details, see section 30, Power-Down Modes.

22.5.2 D/A Output Hold Function in Software Standby Mode

When this LSI enters software standby mode with D/A conversion enabled, the D/A outputs are retained, and the analog power supply current is equal to as during D/A conversion. If the analog power supply current needs to be reduced in software standby mode, clear the DAOE0, DAOE1, and DAE bits to 0 to disable the D/A outputs.

22.5.3 Setting Analog Input Voltage

The reliability of this LSI may be adversely affected if the following voltage ranges are exceeded.

1. AVcc and AVss input voltages

Input voltages AVcc and AVss should be PVcc -0.3 V \leq AVcc \leq PVcc and AVss = PVss. Do not leave the AVcc and AVss pins open when the A/D converter or D/A converter is not in use and in software standby mode. When not in use, connect AVcc to the power supply (PVcc) and AVss to the ground (PVss).

2. Setting range of AVref input voltage Set the reference voltage range of the AVref pin as $3.0 \text{ V} \leq \text{AVec}$.

Section 23 AND/NAND Flash Memory Controller (FLCTL)

The AND/NAND flash memory controller (FLCTL) provides interfaces for an external AND-type flash memory and NAND-type flash memory. To take measures for errors specific to flash memory, the FLCTL supports the ECC generation and error detection functions.

Up to 4-symbol ECC generator, error detector, and hardware error pattern generator have been provided in addition to the 3-symbol ECC detector of the earlier products.

23.1 Features

(1) AND/NAND-Type Flash Memory Interface

- Interface directly connectable to AND/NAND-type flash memory
- Read or write in sector units $(512 + 16 \text{ bytes})$ and ECC processing executed
- Read or write in byte units
- Supports large-block $(2048 + 64 \text{ bytes})$ flash memory
- Supports addresses for 2 Gbits and more by extension to 5-byte addresses
- Note: The FLCTL handles $512 + 16$ bytes as a sector. For products with $2048 + 64$ byte-pages, the FLCTL divide a page into 512 +16 bytes units (i.e. four sectors per page) for processing.

(2) Access Modes: The FLCTL can select one of the following two access modes.

- Command access mode: Performs an access by specifying a command to be issued from the FLCTL to flash memory, address, and data size to be input or output. Read, write, or erasure of data without ECC processing can be achieved.
- Sector access mode: Performs a read or write in sector units by specifying a sector address and controls ECC generation and check. By specifying the number of sectors, the continuous physical sectors can be read or written.

(3) Sectors and Control Codes

- A sector is the basic unit of access and comprised of 512-byte data and 16-byte control code fields. The control code field includes 8-byte ECC when the 3-symbol ECC circuit is used, and 10-byte ECC when the 4-symbol ECC circuit is used.
- The position of the ECC in the control code field can be specified in 4-byte units when the 3symbol ECC circuit is used, and in 1-byte units when the 4-symbol ECC circuit is used.
- User information can be written to the part of the control code field where ECC is not placed.

(4) 3-Symbol ECC

- 64 bits (8 bytes) of ECC is added to a sector, which consists of 512-byte data $+0/4/8$ -byte control code.
- Error correction and detection is up to three errors (30 bits at maximum) at random positions.
- In a write operation, ECC is generated for the data and control code preceding the ECC. The control code following the ECC is not considered.
- In a read operation, an ECC error is checked for data and control code preceding the ECC. The ECC on the control code in the FIFO are the results of checking replaced by the ECC circuit, not the ECC read from flash memory.
- Error correction is not performed even when an ECC error occurs. Error corrections must be performed by software.

(5) 4-Symbol ECC

- 80 bits (10 bytes) of ECC is added to a sector, which consists of 512-byte data + 1-to 6-byte control code.
- Error correction and detection is up to four errors (40 bits at maximum) at random positions.
- In a write operation, ECC is generated for the data and control code preceding the ECC. The control code following the ECC is not considered.
- In a read operation, an ECC error is checked for data and control code preceding the ECC. The ECC on the control code in the FIFO are the results of checking replaced by the ECC circuit, not the ECC read from flash memory.
- The 4-symbol ECC circuit of the FLCTL has the capability of error correction pattern generation by hardware, which is executed on a sector-by-sector basis.
- In the error correction by hardware, addresses indicating the error positions and an error pattern for correcting the errors are output. Data replacement must be performed by software.

(6) Data Error

- When a program error or erase error occurs, the error is reflected on the error source flags. Interrupts for each source can be specified.
- When a read error occurs, an ECC in the control code is other than 0. This read error is reflected on the ECC error source flag.
- When an ECC error occurs, perform an error correction, specify another sector to be replaced, and copy the contents of the block to another sector as required.

(7) Data Transfer FIFO and Data Register

- The 224-byte data FIFO register (FLDTFIFO) is incorporated for data transfer of flash memory.
- The 32-byte control code FIFO register (FLECFIFO) is incorporated for data transfer of control code.

(8) DMA Transfer

• By individually specifying the destinations of data and control code of flash memory to the DMA controller, data and control code can be sent to different areas.

(9) Access Time

- The operating clock (FCLK) on the pins for the AND-/NAND-type flash memory is generated by dividing the peripheral clock $(P\phi)$. The division ratio can be specified by the FCKSEL and QTSEL bits in the common control register (FLCMNCR).
- Before changing the CPG configuration, the FLCTL must be placed in a module stop state.
- In NAND-type flash memory, the FSC and $\overline{\text{FWE}}$ pins operate at the frequency of FCLK. In AND-type flash memory, the FSC pin operates at the frequency of FCLK and the $\overline{\rm FWE}$ pin operates at half the FCLK frequency. These operating frequencies must be specified within the maximum operating frequency of memory to be connected.

Figure 23.1 shows a block diagram of the FLCTL.

Figure 23.1 FLCTL Block Diagram

23.2 Input/Output Pins

The pin configuration of the FLCTL is listed in table 23.1.

Table 23.1 Pin Configuration

Note: * Not supported in this LSI.

23.3 Register Descriptions

Table 23.2 shows the FLCTL register configuration.

Table 23.2 Register Configuration of FLCTL

23.3.1 Common Control Register (FLCMNCR)

FLCMNCR is a 32-bit readable/writable register that specifies the type (AND/NAND) of flash memory, access mode, and other items.

23.3.2 Command Control Register (FLCMDCR)

FLCMDCR is a 32-bit readable/writable register that issues a command in command access mode, specifies address issue, and specifies source or destination of data transfer. In sector access mode, FLCMDCR specifies the number of sector transfers.

23.3.3 Command Code Register (FLCMCDR)

FLCMCDR is a 32-bit readable/writable register that specifies a command to be issued in command access or sector access.

23.3.4 Address Register (FLADR)

FLADR is a 32-bit readable/writable register that specifies the value to be output as an address.

The address of the size specified by ADRCNT[1:0] in the command control register is output sequentially from ADR1 in byte units. By the sector access address specification bit (ADRMD) of the command control register, it is possible to specify whether the sector number set in the address data bits is converted into an address to be output to the flash memory.

When $ADRMD = 1$

• When $ADRMD = 0$

23.3.5 Address Register 2 (FLADR2)

FLADR2 is a 32-bit readable/writable register, and is valid when the ADRCNT2 bit in FLCMDCR is set to 1. FLADR2 specifies an address to be output in command access mode.

23.3.6 Data Counter Register (FLDTCNTR)

FLDTCNTR is a 32-bit readable/writable register that specifies the number of bytes to be read or written in command access mode.

23.3.7 Data Register (FLDATAR)

FLDATAR is a 32-bit readable/writable register. It stores input/output data used when 0 is written to the CDSRC bit in FLCMDCR in command access mode. FLDATAR cannot be used for reading or writing of five or more bytes of contiguous data.

23.3.8 Interrupt DMA Control Register (FLINTDMACR)

FLINTDMACR is a 32-bit readable/writable register that enables or disables DMA transfer requests or interrupts. A transfer request from the FLCTL to the DMAC is issued after each access mode has been started.

Bits 9 to 5 are the flag bits that indicate various errors occurred in flash memory access and whether there is a transfer request from the FIFO. Only 0 can be written to these bits. To clear a flag, write 0 to the target flag bit and 1 to the other flag bits.

Note: * Only 0 can be written to these bits.

23.3.9 Ready Busy Timeout Setting Register (FLBSYTMR)

FLBSYTMR is a 32-bit readable/writable register that specifies the timeout time when the FRB pin is busy.

23.3.10 Ready Busy Timeout Counter (FLBSYCNT)

FLBSYCNT is a 32-bit read-only register.

The status of flash memory obtained by the status read is stored in the bits STAT[7:0].

The timeout time set in the bits RBTMOUT[19:0] in FLBSYTMR is copied to the bits RBTIMCNT[19:0] and counting down is started when the FRB pin is placed in a busy state. When values in the RBTIMCNT[19:0] become 0, 1 is set to the BTOERB bit in FLINTDMACR, thus notifying that a timeout error has occurred. In this case, an FLSTE interrupt request can be issued if an interrupt is enabled by the RBERINTE bit in FLINTDMACR.

23.3.11 Data FIFO Register (FLDTFIFO)

FLDTFIFO is used to read or write the data FIFO area.

In DMA transfer, this register must be specified as the destination or source.

Note that the direction of read or write specified by the SELRW bit in FLCMDCR must match that specified in this register. When changing the read/write direction, FLDTFIFO should be cleared by setting the AC0CLR bit in FLINTDMACR before use.

23.3.12 Control Code FIFO Register (FLECFIFO)

FLECFIFO is used to read or write the control code FIFO area.

In DMA transfer, data in this register must be specified as the destination (source).

Note that the direction of read or write specified by the SELRW bit in FLCMDCR must match that specified in this register. When changing the read/write direction, FLECFIFO should be cleared by setting the AC1CLR bit in FLINTDMACR before use.

23.3.13 Transfer Control Register (FLTRCR)

Setting the TRSTRT bit to 1 initiates access to flash memory. Access completion can be checked by the TREND bit. During the transfer (from when the TRSTRT bit is set to 1 until the TREND bit is set to 1), the processing should not be forcibly ended (by setting the TRSTRT bit to 0).

When reading from flash memory, TREND is set when reading from flash memory have been finished. However, if there is any read data remaining in the FIFO, the processing should not be forcibly ended until all data has been read from the FIFO.

23.3.14 4-Symbol ECC Processing Result Register n (FL4ECCRESn) (n = 1 to 4)

FL4ECCRESn is a 32-bit read-only register that stores the error correction pattern for the nth error generated by the 4-symbol ECC circuits and the address for the nth error. The contents of this register become valid when bits 23 (4EECEN) and 22 (4ECCCORRECT) are set to 1 and a correction pattern has been generated by the setting of the 4-symbol ECC control register (FL4ECCCR).

23.3.15 4-Symbol ECC Control Register (FL4ECCCR)

FL4ECCCR is a 32-bit readable register that indicates the processing states of the 4-symbol ECC circuit. This register consists of flag bits to which only 0 can be written. To clear a flag, write 0 to the target flag bit and 1 to the other flag bits.

Note: * Only 0 can be written to these bits.

23.3.16 4-Symbol ECC Error Count Register (FL4ECCCNT)

FL4ECCCNT is a 32-bit readable register that indicates the number of errors detected by the 4 symbol ECC circuit. Only 0 can be written to this register. To clear this register, write 0 to all bits.

23.4 Operation

23.4.1 Access Sequence

The FLCTL performs accesses in several independent stages.

For example, AND-type flash memory programming consists of the following five stages.

- First command issue stage (program setup command)
- Address issue stage (program address)
- Data stage (output)
- Second command issue stage (program start command)
- Status read stage

AND-type flash memory programming access is achieved by executing these five stages sequentially. An access to flash memory is completed at the end of the final stage (status read stage).

Figure 23.2 Programming Operation for AND-Type Flash Memory and Stages

For details on AND-type flash memory read and NAND-type flash memory read/program operation, see section 23.4.4, Command Access Mode.

23.4.2 Operating Modes

Two operating modes are supported.

- Command access mode
- Sector access mode

The ECC generation and error check are performed in sector access mode.

23.4.3 Register Setting Procedure

Figure 23.3 shows the register setting flow required for accessing the flash memory.

Figure 23.3 Register Setting Flow

23.4.4 Command Access Mode

Command access mode accesses flash memory by specifying a command to be issued to flash memory, address, data, read/write direction, and number of times to the registers. In this mode, I/O data can be transferred by the DMA via FLDTFIFO.

(1) AND-Type Flash Memory Access

Figures 23.4 and 23.5 show examples of read operation for AND-type flash memory. In these examples, the first command is specified as H'00 and address data length is specified as 2 bytes (SA1 and SA2). (Only SA1 and SA2 are specified, while CA1 and CA2 are not specified.). In addition, the number of read bytes is specified as 4 bytes in the data counter and H'FF is specified as the second command.

Figure 23.4 Read Operation Timing for AND-Type Flash Memory (1)

Figure 23.5 Read Operation Timing for AND-Type Flash Memory (2)

Figures 23.6 and 23.7 show examples of programming operation for AND-type flash memory.

Figure 23.6 Programming Operation Timing for AND-Type Flash Memory (1)

Figure 23.7 Programming Operation Timing for AND-Type Flash Memory (2)

(2) NAND-Type Flash Memory Access

Figure 23.8 shows an example of read operation for NAND-type flash memory. In this example, the first command is specified as H'00, address data length is specified as 3 bytes, and the number of read bytes is specified as 8 bytes in the data counter.

Figures 23.9 and 23.10 show examples of programming operation for NAND-type flash memory.

Figure 23.9 Programming Operation Timing for NAND-Type Flash Memory (1)

(3) NAND-Type Flash Memory (2048 + 64 Bytes) Access

Figure 23.11 shows an example of read operation for NAND-type flash memory (2048 + 64 bytes). In this example, the first command is specified as H'00, the second command is specified as H'30, and address data length is specified as 4 bytes. The number of read bytes is specified as 4 bytes in the data counter.

Figure 23.11 Read Operation Timing for NAND-Type Flash Memory

Figures 23.12 and 23.13 show examples of programming operation for NAND-type flash memory $(2048 + 64 \text{ bytes}).$

Figure 23.12 Programming Operation Timing for NAND-Type Flash Memory (1)

Figure 23.13 Programming Operation Timing for NAND-Type Flash Memory (2)

23.4.5 Sector Access Mode

In sector access mode, flash memory can be read or programmed in sector units by specifying the sector number of the sector to be accessed. In programming, an ECC is added. In read, an ECC error check (detection) is performed.

Since 512-byte data is stored in FLDTFIFO and 16-byte control code is stored in FLECFIFO, the DREQ1EN and DREQ0EN bits in FLINTDMACR can be set to transfer by the DMA.

Figure 23.14 shows the relationship of DMA transfer between sectors in flash memory (data and control code) and memory on the address space.

(1) Sector Address

Figure 23.15 shows the relationship between the physical sector address of AND/NAND-type flash memory and the address of flash memory.

Figure 23.15 Relationship between Sector Number and Address Expansion of AND-/NAND-Type Flash Memory

(2) Continuous Sector Access

A series of sectors can be read or written by specifying the start sector address of NAND-type flash memory and the number of sectors to be transferred. Figure 23.16 shows an example of physical sector specification register and transfer count specification register settings when transferring logical sectors 0 to 40, which are not contiguous because of an unusable sector in NAND-type flash memory.

Figure 23.16 Sector Access when Unusable Sector Exists in Continuous Sectors

23.4.6 ECC Error Correction

The FLCTL generates and adds ECC during write operation in sector access mode and performs ECC error check during read operation in sector access mode. ECC processing is selectable between 3-symbol ECC, the function provided in the earlier FLCTL, and 4-symbol ECC.

With 3-symbol ECC, only ECC generation and error detection are performed and error correction is not performed. So, errors must be corrected by software. On the other hand, 4-symbol ECC is capable of ECC generation, error detection, and error correction pattern generation by hardware.

(1) Overview of 4-Symbol ECC Circuit

The 4-symbol ECC circuit in the FLCTL is capable of correcting up to10 bits per symbol, which makes a maximum of 40 bits for four symbols. However, the circuit corrects up to 32 bits because the data in the flash memory data area is counted as eight bits per symbol.

Error correction pattern generation means generation of information necessary for correcting errors, not execution of error correction. For details, see (3) 4-Symbol ECC Error Correction Pattern Generation.

The 4-symbol ECC circuit is roughly divided into three stages (figure 23.17).

- 1. ECC generator
- 2. Error count detector
- 3. Error correction pattern generator

ECC generation and error count detection can be executed continuously while error correction pattern generation is executed on a sector-by-sector basis.

Figure 23.17 4-Symbol ECC Circuit

(2) 4-Symbol ECC Operation

Figure 23.18 shows a flowchart of the operation when the 4-symbol ECC circuit is used. Setting the 4ECCEN bit in FLCMNCR enables the 4-symbol ECC circuit and ECC is generated and output for each sector. If the 4ECCCORRECT bit in FLCMNCR is also set to 1, information necessary for correction pattern generation is accumulated in the 4-symbol ECC circuit.

In the case when the FLCTL is reading data from flash memory by continuous sector access, the reading operation stops when an error-containing sector has been read regardless of the number of remaining sectors. After reading of the error-containing sector has ended, generation of error correction pattern is started by setting the FL4ECCCR register. If the sector contains five or more errors, that sector is regarded as uncorrectable. Note that a sector may be uncorrectable for some error patterns even if it contains four or less errors. In such a case, invalid data are placed in the FL4ECCRES1 to FL4ECCRES4 registers.

Figure 23.18 Flow of 4-Symbol ECC Operation

(3) 4-Symbol ECC Error Correction Pattern Generation

The 4-symbol ECC circuit of the FLCTL can generate error correction patterns by hardware. The original data can be restored by using the error correction patterns. Since the hardware processing only covers generation of error correction patterns, processing for data restoration must be provided by software.

The error correction patterns are output in the following format. The bits in a correction pattern at error bit positions are set to 1, so the original data is restored by taking EOR of error data and error correction pattern.

• Example 1

Example 2

• Example 3 Original data: B'11110000 Erroneous data: B'00000000 Correction pattern: B'0011110000 (higher two bits are unnecessary data) Recovered data: B'11110000 (EOR of error pattern and correction pattern)

23.4.7 Status Read

The FLCTL can read the status register of an AND/NAND-type flash memory. The data in the status register is input through the I/O7 to I/O0 pins and stored in the bits STAT[7:0] in FLBSYCNT, which can be read by the CPU. If a program error or erase error is detected when the status register value is stored in the bits STAT[7:0] in FLBSYCNT, the STERB bit in FLINTDMACR is set to 1 and generates an interrupt to the CPU if the STERINTE bit in FLINTDMACR is enabled. If a status error is occurred during continuous sector access, the TREND bit in FLTRCR is set to 1 and the procedure stops.

(1) Status Read of AND-Type Flash Memory

The status register of AND-type flash memory can be read by asserting the output enable signal $OE (OE = 0)$. If programming is executed in command access mode or sector access mode while the DOSR bit in FLCMDCR is set to 1, the FLCTL automatically asserts the \overline{OE} signal and reads the status register of AND-type flash memory. When the status register of AND-type flash memory is read, the I/O7 to I/O0 pins indicate the following information as described in table 23.3.

Table 23.3 Status Read of AND-Type Flash Memory
(2) Status Read of NAND-Type Flash Memory

The status register of NAND-type flash memory can be read by inputting command H'70 to NAND-type flash memory. If programming is executed in command access mode or sector access mode while the DOSR bit in FLCMDCR is set to 1, the FLCTL automatically inputs command H'70 to NAND-type flash memory and reads the status register of NAND-type flash memory. When the status register of NAND-type flash memory is read, the I/O7 to I/O0 pins indicate the following information as described in table 23.4.

Table 23.4 Status Read of NAND-Type Flash Memory

23.5 Interrupt Sources

The FLCTL has seven interrupt sources: Status error, ready/busy timeout error, ECC error, 4 symbol ECC pattern generation end, transfer end, FIFO0 transfer request, and FIFO1 transfer request. Each of the interrupt sources has its corresponding interrupt flag and the interrupt can be requested independently to the CPU if the interrupt is enabled by the interrupt enable bit. Note that the status error, ready/busy timeout error, ECC error, and 4-symbol ECC pattern generation end, use the common FLSTE interrupt to the CPU.

Table 23.5 FLCTL Interrupt Requests

Note: Flags for the FIFO0 overrun error/underrun error and FIFO1 overrun error/underrun error also exist. However, no interrupt is requested to the CPU.

23.6 DMA Transfer Specifications

The FLCTL can request DMA transfers separately to the data area FLDTFIFO and control code area FLECFIFO. Table 23.6 summarizes DMA transfer enable or disable states in each access mode.

For details on DMAC settings, see section 11, Direct Memory Access Controller (DMAC).

23.7 Usage Notes

23.7.1 Writing to the Control-Code Area when 4-Symbol ECC Circuit Is in Use

Follow the procedure given below to write to the control-code area when the 4-symbol ECC circuit is in use. If this procedure is not followed, correctly writing to the control-code area of the flash memory will not be possible.

Figure 23.19 Writing Procedure to the Control-Code Area when 4-Symbol ECC is Used

23.7.2 Usage Notes for the SNAND bit

When using the SNAND bit in FLCMNCR, only the first command or the second command is corresponded in spite of the setting of the DOCMD1 or DOCMD2 bit in FLCMDCR.

When no command or only the first command is issued, 0 should be written in the SNAND bit.

Section 24 USB 2.0 Host/Function Module (USB)

The USB 2.0 host/function module (USB) is a USB controller, which provides capabilities as a USB host controller and USB function controller function.

When used as the host controller, this module supports high-speed transfer defined by USB (universal serial bus) Specification 2.0, full-speed transfer, and low-speed transfer, allowing the use of two USB ports.

When used as the function controller, this module supports high-speed transfer defined by USB Specification 2.0, and full-speed transfer, allowing the use of one USB port.

This module has a USB transceiver* and supports all of the transfer types defined by the USB specification.

This module has a 10-Kbyte buffer memory for data transfer, providing a maximum of ten pipes. Any endpoint numbers can be assigned to PIPE1 to PIPE9, based on the function devices or user system for communication.

Note: * When this module is to be used, start by making settings for the internal USB transceiver. For details, see section 24.5.1, Procedure for Setting the USB Transceiver.

24.1 Features

(1) Host Controller and Function Controller Supporting USB High-Speed Operation

- The USB host controller and USB function controller are incorporated.
- The USB host controller and USB function controller can be switched by register settings.
- Both high-speed transfer (480 Mbps) and full-speed transfer (12 Mbps) are supported.
- High-speed/full-speed/low-speed USB transceiver (shared by the USB host and USB function) is incorporated.

(2) Reduced Number of External Pins and Space-Saving Installation

- On-chip D+ pull-up resistor (during USB function operation)
- On-chip D+ and D- pull-down resistor (during USB host operation)
- On-chip D+ and D- terminal resistor (during high-speed operation)
- On-chip D+ and D- output resistor (during low-speed/full-speed operation)

(3) All Types of USB Transfers Supported

- Control transfer
- Bulk transfer
- Interrupt transfer (high bandwidth transfers not supported)
- Isochronous transfer (high bandwidth transfers not supported)

(4) Internal Bus Interfaces

• Two DMA interface channels are incorporated.

(5) Pipe Configuration

- On-chip 10-Kbyte buffer memory for USB communications
- Up to ten pipes can be selected (including the default control pipe)
- Programmable pipe configuration
- Endpoint numbers can be assigned flexibly to PIPE1 to PIPE9.
- Transfer conditions that can be set for each pipe:

(6) Features of the USB Host Controller

- High-speed transfer (480 Mbps), full-speed transfer (12 Mbps), and low-speed transfer (1.5 Mbps) are supported.
- Communications with multiple peripheral devices connected via a single hub
- Auto response for reset handshake
- Automatic scheduling for SOF and packet transmissions
- Programmable intervals for isochronous and interrupt transfers

(7) Features of the USB Function Controller

- Both high-speed transfer (480 Mbps) and full-speed transfer (12 Mbps) are supported.
- Automatic recognition of high-speed operation or full-speed operation based on automatic response to the reset handshake
- Control transfer stage control function
- Device state control function
- Auto response function for SET_ADDRESS request
- NAK response interrupt function (NRDY)
- SOF interpolation function

(8) Other Features

- Byte endian swap function that allows handling of data in both big endian and little endian formats
- Transfer ending function using transaction count
- DMA transfer ending function
- BRDY interrupt event notification timing change function (BFRE)
- Function that automatically clears the buffer memory after the data for the pipe specified at the DnFIFO $(n = 0 \text{ or } 1)$ port has been read (DCLRM)
- NAK setting function for response PID generated by end of transfer (SHTNAK)

24.2 Input/Output Pins

Table 24.1 shows the pin configuration of the USB. If the module is not in use, handle the pins as indicated below.

- Be sure to apply the power supply voltage.
- Connect DP1, DP0, DM1, DM0, and VBUS to Vss.
- Connect REFRIN to USBAPVcc via a 5.6-k Ω resistor with tolerance of $\pm 20\%$.
- Refer to section 5.3, Clock Operating Modes, for handling of ht USB_X1 and USB_X2 pins.

Table 24.1 USBH/F Pin Configuration

24.3 Register Description

Table 24.2 shows the register configuration of the USB.

Table 24.2 Register Configuration

24.3.1 System Configuration Control Register 0 (SYSCFG0)

SYSCFG0 is a register that enables supply of the USB clock to this module and high-speed operation on PORT0, selects the host controller function or function controller function, controls the DP and DM pins, and enables operation of the USB block of this module.

This register is initialized by a power-on reset.

Note: The DRPD bit should be cleared to 0 when the function controller function is selected. The DPRPU bit should be cleared to 0 when the host controller function is selected.

Table 24.3 Register Bits Initialized by Clearing the USBE Bit to 0 (when Function Controller Function is Selected)

Table 24.4 Register Bits Initialized by Clearing the USBE Bit to 0 (when Host Controller Function is Selected)

24.3.2 System Configuration Control Register 1 (SYSCFG1)

SYSCFG1 is a register that enables high-speed operation on PORT1, controls the DP and DM pins and access cycles for access to this module.

This register is initialized by a power-on reset.

Note: The DRPD bit should be cleared to 0 when the function controller function is selected.

24.3.3 System Configuration Status Register 0 (SYSSTS0)

SYSSTS0 is a register that monitors the line status (D+ and D− lines) of the USB data bus on PORT0.

This register is initialized by a power-on reset.

Note: * Depends on the states of the D+ and D- lines.

24.3.4 System Configuration Status Register 1 (SYSSTS1)

SYSSTS1 is a register that monitors the line status (D+ and D− lines) of the USB data bus on PORT1.

This register is initialized by a power-on reset.

Note: * Depends on the states of the D+ and D- lines.

Table 24.5 USB Data Bus Line Status

[Legend]

Chirp: The reset handshake protocol is being executed in high-speed operation enabled state $(HSE = 1)$.

Squelch: SE0 or idle state

Not squelch: High-speed J state or high-speed K state

Chirp J: Chirp J state

Chirp K: Chirp K state

24.3.5 Device State Control Register 0 (DVSTCTR0)

DVSTCTR0 is a register that controls and confirms the state of the USB data bus of PORT0.

This register is initialized by a power-on reset. Only the WKUP bit is initialized by a USB bus reset.

Note: When the function controller function is selected, set bits RWUPE, AUSBRST, ARESUME, and AUACT to all 0s.

When the host controller function is selected, set bit WKUP to 0.

Table 24.6 PORT0 USB Data Bus Line Statuses

24.3.6 Device State Control Register 1 (DVSTCTR1)

DVSTCTR1 is a register that controls and confirms the state of the USB data bus of PORT1.

This register is initialized by a power-on reset.

Note: Clear the value of bits RWUPE, USBRST, RESUME, and UACT when the function controller function is selected.

Table 24.7 PORT1 USB Data Bus Line Statuses

24.3.7 Test Mode Register (TESTMODE)

TESTMODE is a register that controls the USB test signal output during high-speed operation.

This register is initialized by a power-on reset.

Table 24.8 Test Mode Operation

24.3.8 DMA-FIFO Bus Configuration Registers (D0FBCFG, D1FBCFG)

D0FBCFG and D1FBCFG perform access control of the D0FIFO and D1FIFO ports.

24.3.9 FIFO Port Registers (CFIFO, D0FIFO, D1FIFO)

CFIFO, D0FIFO and D1FIFO are port registers that are used to read data from the FIFO buffer memory and writing data to the FIFO buffer memory.

The transmission/reception buffer memory of this module has a FIFO structure (FIFO buffer). Use the FIFO port registers to access the FIFO buffer. There are three FIFO ports: the CFIFO, D0FIFO and D1FIFO ports. Each FIFO port is configured of a port register (CFIFO, D0FIFO, D1FIFO) that handles reading of data from the buffer memory and writing of data to the FIFO buffer memory, a select register (CFIFOSEL, D0FIFOSEL, D1FIFOSEL) that is used to select the pipe assigned to the FIFO port, and a control register (CFIFOCTR, D0FIFOCTR, D1FIFOCTR).

 9. The valid bits in this register depend on the settings of the MBW bits (access bit width setting) and BIGEND bit (endian setting) as shown in tables 24.9 to 24.11.

Table 24.9 Endian Operation in 32-Bit Access (when MBW = 10)

Table 24.10 Endian Operation in 16-Bit Access (when MBW = 01)

Table 24.11 Endian Operation in 8-Bit Access (when MBW = 00)

24.3.10 FIFO Port Select Registers (CFIFOSEL, D0FIFOSEL, D1FIFOSEL)

CFIFOSEL, D0FIFOSEL and D1FIFOSEL are registers that select the pipes to be assigned to the FIFO ports, and control access to the corresponding ports.

The same pipe should not be specified in the CURPIPE bits of CFIFOSEL, D0FIFOSEL and D1FIFOSEL. When the CURPIPE bits in D0FIFOSEL and D1FIFOSEL are cleared to B'000, no pipe is selected.

The pipe number should not be changed while the DMA transfer is enabled.

These registers are initialized by a power-on reset.

(1) CFIFOSEL

Note: $*$ Only 0 can be read and 1 can be written to.

(2) D0FIFOSEL, D1FIFOSEL

Note: $*$ Only 0 can be read and 1 can be written to.

24.3.11 FIFO Port Control Registers (CFIFOCTR, D0FIFOCTR, D1FIFOCTR)

These registers determine whether or not writing to the buffer memory has been finished, the buffer on the CPU side has been cleared, and the FIFO port is accessible. CFIFOCTR, D0FIFOCTR, and D1FIFOCTR are provided for the corresponding FIFO ports.

Notes: 1. Only 1 can be written to.

2. Only 0 can be read and 1 can be written to.

24.3.12 Interrupt Enable Register 0 (INTENB0)

INTENB0 enables or disables various interrupts. If an interrupt for which the corresponding bit in this register is set to 1 has occurred, an interrupt request is output to the interrupt controller.

Note: * The RSME, DVSE, and CTRE bits are cleared to 0 when the host controller function is selected.

24.3.13 Interrupt Enable Register 1 (INTENB1)

INTENB1 enables or disables various interrupts. If an interrupt for which the corresponding bit in this register is set to 1 has occurred, an interrupt request is output to the interrupt controller.

Note: Clear each bit in this register to 0 when the function controller function is selected.

24.3.14 Interrupt Enable Register 2 (INTENB2)

INTENB2 enables or disables various interrupts. If an interrupt for which the corresponding bit in this register is set to 1 has occurred, an interrupt request is output to the interrupt controller.

Note: Clear each bit in this register to 0 when the function controller function is selected.

24.3.15 BRDY Interrupt Enable Register (BRDYENB)

BRDYENB enables or disables the BRDY interrupts for individual pipes. If a BRDY interrupt for which the corresponding bit in this register is set to 1 has occurred, an interrupt request is output to the interrupt controller.

24.3.16 NRDY Interrupt Enable Register (NRDYENB)

NRDYENB enables or disables the NRDY interrupts for individual pipes. If a NRDY interrupt for which the corresponding bit in this register is set to 1 has occurred, an interrupt request is output to the interrupt controller.

24.3.17 BEMP Interrupt Enable Register (BEMPENB)

BEMPENB enables or disables the BEMP interrupts for individual pipes. If a NRDY interrupt for which the corresponding bit in this register is set to 1 has occurred, an interrupt request is output to the interrupt controller.

24.3.18 SOF Output Configuration Register (SOFCFG)

SOFCFG specifies the transaction-enabled time and PIPEBRDY interrupt status clear timing.

Note: Clear the value of the TRNENSEL bit to 0 when the function controller function is selected.

24.3.19 Interrupt Status Register 0 (INTSTS0)

INTSTS0 indicates the statuses of various interrupts.

This register is initialized by a power-on reset. The DVST and DVSQ[2:0] bits are initialized by a USB bus reset.

Notes: 1. DVST is initialized to 0 and DVSQ[2:0] to 000 by a power-on reset. DVST is initialized to 1 and DVSQ[2:0] to 001 by a USB bus reset.

- 2. Only 0 can be written to.
- 3. To clear the VBINT, RESM, SOFR, DVST, or CTRT bit, write 0 only to the bits to be cleared; write 1 to the other bits. Do not write 0 to the status bits indicating 0.
- 4. This module detects the changes in the statuses indicated by the VBINT and RESM bits even while the clock supply is stopped (while SCKE is 0), and outputs the corresponding interrupt requests as long as they are enabled. Clearing the status should be done after enabling the clock supply.
- 5. Transitions in the status of the RESM, DVST, and CTRT bits only occur when the function controller function is selected; clear the corresponding interrupt enable bits to 0 (disable) when the host controller function is selected.
- 6. The DVSQ, VALID, and CTRQ bits are valid when the function controller function is selected

24.3.20 Interrupt Status Register 1 (INTSTS1)

INTSTS1 indicates the statuses of various interrupts.

Notes: 1. Only 1 can be written to.

 2. The interrupts generated by the status transitions indicated by each bit in this register should only be enabled when the host controller function is selected.

 3. This module detects the change in the status indicated by the BCHG bit even while the clock supply is stopped (while SCKE is 0), and outputs the corresponding interrupt request as long as it is enabled. Clearing the status should be done after enabling the clock supply.

24.3.21 Interrupt Status Register 2 (INTSTS2)

INTSTS2 indicates the statuses of various interrupts.

Notes: 1. Only 1 can be written to.

 2. The interrupts generated by the status transitions indicated by each bit in this register should only be enabled when the host controller function is selected.

 3. This module detects the change in the status indicated by the BCHG bit even while the clock supply is stopped (while SCKE is 0), and outputs the corresponding interrupt request as long as it is enabled. Clearing the status should be done after enabling the clock supply.

24.3.22 BRDY Interrupt Status Register (BRDYSTS)

BRDYSTS is used to confirm the BRDY interrupt status for each pipe. The conditions for generation and clearing of the BRDY interrupt status are different depending on the settings of the BRDYM bit and the BFRE bit for each pipe.

This register is initialized by a power-on reset.

(1) When BRDYM = 0 and BFRE = 0 are Set

With this setting, the BRDY interrupt indicates that the FIFO ports have become accessible.

In the event of the following conditions, an internal BRDY interrupt request trigger is generated and the bit corresponding to the pipe that caused the request trigger is set to 1.

(a) Conditions for Pipes in the Transmitting Direction

- The DIR bit is modified from 0 to 1.
- Packet transmission in the pertinent pipe has completed while the FIFO buffer assigned to the pipe is not writable by the CPU (i.e., when 0 is read from the BSTS bit). If continuous transmission/reception mode has been set, a request trigger is generated on completion of transmitting the data for one plane of FIFO buffer.
- When the FIFO buffer is configured as double buffers, writing to one buffer has completed while the other buffer is empty. Even if transmission from one FIFO buffer has completed while the other buffer is being written to, a request trigger is not generated until the writing to the buffer is completed.
- Buffer flush by this module has occurred in an isochronous transfer pipe.
- The FIFO buffer has become writable from a non-writable state by writing 1 to the ACLRM bit.

Request triggers are not generated for the DCP (that is, in data transfer by control transfer).

(b) Conditions for Pipes in the Receiving Direction

• Packet reception in the pertinent pipe has completed and the FIFO buffer has become ready for reading while the FIFO buffer assigned to the pipe is not writable by the CPU (i.e., when 0 is read from the BSTS bit).

A request trigger is not generated for a transaction in which a data PID disagreement has occurred.

In continuous transmission/reception mode, a request trigger is not generated if the size of data is MaxPacketSize and there is a free space remaining in the buffer.

When a short packet is received, a request trigger is generated even if there is a free space in the FIFO buffer.

When a transaction is used, a request trigger is generated when the set number of packets has been received. In this case, a request trigger is generated even when there is some free space in the FIFO buffer.

• When the FIFO buffer is configured as double buffers, reading from a FIFO buffer has completed while the other buffer is also ready to be read. Even if reception into one FIFO buffer has completed while the other buffer is being read, a request trigger is not generated until the reading from the buffer is completed.

The BRDY interrupt is not generated during communication in the status stage of control transfer when the function controller function is selected.

The BRDY interrupt status bit of the pertinent pipe can be cleared by writing 0 to the bit corresponding to the pipe in this register. This interrupt status clearing must be done before making access to the FIFO buffer.

(2) When BRDYM = 0 and BFRE = 1 are Set

With this setting, a BRDY interrupt is generated when all the data for one round of transfer have been read through a receiving pipe and the bit corresponding to that pipe is set.

This module determines that the last data for one round of transfer has been read in any of the following conditions.

(a) Short Packet, Including Zero-Length Packet, has been Received

(b) Packets for the Value Set in Bits TRNCNT have been Received When the Transaction Counter (Bits TRNCNT) is Used

When either of the above conditions is satisfied and reading of that data has been completed, this module determines that all the data for one round of transfer have been read out.

When a zero-length packet is received while the FIFO buffer is empty, the USB 2.0 host/function module determines that all the data for a single transfer has been read at the point at which the FRDY bit is set to 1 and the DTLN bit cleared to 0 in the FIFO port control register. In this case, to start the next transfer, write 1 to the BCLR bit in the corresponding FIFOCTR register.

With this setting, the BRDY interrupt is not detected for transmitting pipes.

The BRDY interrupt status bit of the pertinent pipe can be cleared by writing 0 to the bit corresponding to the pipe in this register.

When this mode is used, do not modify the BFRE bit setting until the processing for that round of transfer is completed. To modify the BFRE bit, use the ACLRM bit to clear all the contents of the FIFO buffer of the corresponding pipe.

(3) When BRDYM = 1 and BFRE = 0 are Set

With this setting, the values of the bits in this register reflect the value of the BSTS bit of the individual pipes. This means that the BRDY interrupt statuses are set according to the states of the FIFO buffers.

(a) Condition for Pipes in the Transmitting Direction

When data writing to the FIFO port is possible, the corresponding bit in this register is set to 1. When not possible, the bit is cleared to 0. For the transmitting DCP, however, the BRDY interrupt is not generated even if writing is possible.

(b) Condition for Pipes in the Receiving Direction

When data reading from the FIFO port is possible, the corresponding bit in this register is set to 1. After all the data have been read out (when reading has become not possible), the bit is cleared to 0. If a zero-length packet is received while the FIFO buffer is empty, the corresponding bit is set to 1 until 1 is written to BCLR, during which the BRDY interrupt is continuously generated.

With this setting, the status bits in this register cannot be cleared by writing 0.

When BRDYM $= 1$ is set, all the BFRE bits (for all pipes) must be cleared to 0.

Bit Bit Name

Initial

15 to 10 Undefined R Reserved

Undefined values are read from these bits. The write

R/W Description

value should always be 0.

24.3.23 NRDY Interrupt Status Register (NRDYSTS)

NRDYSTS is used to confirm the NRDY interrupt status for each pipe

This register is initialized by a power-on reset.

When an internal NRDY interrupt request has occurred in the pipe for which PID = BUF is set, the corresponding bit for that pipe is set to 1. An internal NRDY interrupt request is generated for a pipe in the conditions described below. Note, however, that these interrupt generating conditions do not apply to the cases when a setup transaction is being executed while the host controller function is selected. During a setup transaction with the host controller function selected, SACK interrupt or SIGN interrupt can occur. Also note that interrupt requests are not generated during execution of the status stage in control transfer when the function controller function is selected.

(1) Connection with which Split Transactions Do Not Occur when the Host Controller Function is Selected

(a) Conditions for Pipes in the Transmitting Direction

An NRDY interrupt is generated if any of the following conditions is met.

- In a pipe specified for isochronous transfer, when the time to issue an OUT token comes while the FIFO buffer contains no data for transmission: In this case, this module transmits a zero-length packet subsequent to the OUT token and sets the bit corresponding to the pipe and the OVRN bit to 1.
- In a pipe not specified for isochronous transfer and executing communication for other than setup transactions, when the function device returns no response (timeout is detected without detecting a handshake packet from the function device), or when any error has detected in the packet from the function device for consecutive three times: In this case, this module sets the bit corresponding to the pipe and modifies the PID bit setting

for the pipe to STALL.

• When a STALL handshake (not only STALL for OUT token but also STALL for PING token apply) is received from the function device during communication for other than setup transactions:

In this case, this module sets the bit corresponding to the pipe and modifies the PID bit setting for the pipe to STALL.

(b) Conditions for Pipes in the Receiving Direction

• In a pipe specified for isochronous transfer, when the time to issue an IN token comes while the FIFO buffer has no empty space:

In this case, this module discards the data received in response to the IN token and sets the bit corresponding to the pipe and the OVRN bit to 1. Further, it also sets the CFRCDE bit to 1 if a packet error has been detected in the data received in response to the IN token.

• In a pipe not specified for isochronous transfer, when the function device returns no response (timeout is detected without detecting a data packet from the function device) to the IN token sent by this module, or when any error has detected in the packet from the function device for consecutive three times:

In this case, this module sets the bit corresponding to the pipe and modifies the PID bit setting for the pipe to NAK.

- In a pipe specified for isochronous transfer, when the function device returns no response (timeout is detected without detecting a data packet from the function device) to the IN token, or when any error has detected in the packet from the function device: In this case, this module sets the bit corresponding to the pipe to 1 (it does not modify the PID bit setting for the pipe).
- In a pipe specified for isochronous transfer, when a CRC error or bit stuffing error has detected in the received data packet: In this case, this module sets the bit corresponding to the pipe and the CRCE bit to 1.
- When a STALL handshake is received: In this case, this module sets the bit corresponding to the pipe and modifies the PID bit setting for the pipe to STALL.

(2) Connection with which Split Transactions can Occur when the Host Controller Function is Selected

(a) Conditions for Pipes in the Transmitting Direction

- In a pipe specified for isochronous transfer, when the time to issue an OUT token comes while the FIFO buffer contains no data for transmission: In this case, this module sets the bit corresponding to the pipe to 1 when it issues a Start-Split transaction (S-Split) and sets the OVRN bit to 1. It also transmits a zero-length packet subsequent to the OUT token.
- In a pipe not specified for isochronous transfer, when the hub returns no response (timeout is detected without detecting a handshake packet from the hub) to the S-Split or Complete-Split transaction (C-Split), or when any error has detected in the packet from the hub for consecutive three times:

In this case, this module sets the bit corresponding to the pipe and modifies the PID bit setting

for the pipe to NAK. If an NRDY interrupt is detected when a C-Split has been issued, the module clears the CSSTS bit to 0.

- When a STALL handshake is received in response to the C-Split: In this case, this module sets the bit corresponding to the pipe, modifies the PID bit setting for the pipe to STALL and clears the CSSTS bit to 0. Note that the NRDY interrupt is not detected in setup transactions.
- In a pipe specified for interrupt transfer, when NYET is received in response to the C-Split for the micro frame number of 4:

In this case, this module sets the bit corresponding to the pipe to 1 and clears the CSSTS bit to 0 (it does not modify the PID bit setting for the pipe).

(b) Conditions for Pipes in the Receiving Direction

In a pipe specified for isochronous transfer, when the time to issue an IN token comes while the FIFO buffer has no empty space:

In this case, this module sets the bit corresponding to the pipe to 1 when it issues S-Split and sets the OVRN bit to 1. It also discards the data received in response to the IN token

- In a pipe specified for bulk transfer or during transfer for other than setup transactions in the DCP, when the hub returns no response (timeout is detected without detecting a data packet from the hub) to the IN token issued by this module at the issuance of S-Split or C-Split, or when any error has detected in the packet from the hub for consecutive three times: In this case, this module sets the bit corresponding to the pipe and modifies the PID bit setting for the pipe to NAK. If this condition arise with C-Split, the module also clears the CSSTS bit to 0 .
- For C-Split in a pipe specified for isochronous transfer or interrupt transfer, when the hub returns no response (timeout is detected without detecting a data packet from the hub) to the IN token issued by this module, or when any error has detected in the packet from the hub for consecutive three times:

If this condition arise with a pipe for interrupt transfer, this module sets the bit corresponding to the pipe, modifies the PID bit setting for the pipe to NAK, and clears the CSSTS bit to 0. If this condition arise with a pipe for isochronous transfer, this module sets the bit corresponding to the pipe and the CRCE bit to 1 and clears the CSSTS bit to 0 (it does not modify the PID bit setting for the pipe).

• For C-Split in a pipe not specified for isochronous transfer, when a STALL handshake is received:

In this case, this module sets the bit corresponding to the pipe, modifies the PID bit setting for the pipe to STALL, and clears the CSSTS bit to 0.

• For C-Split in a pipe specified for isochronous transfer or interrupt transfer, when NYET handshake is received for the micro frame number of 4:

In this case, this module sets the bit corresponding to the pipe and the CRCE bit to 1 and clears the CSSTS bit to 0 (it does not modify the PID bit setting for the pipe).

(3) When the Function Controller Function is Selected

(a) Condition for Pipes in the Transmitting Direction

When an IN token is received while the FIFO buffer contains no data for transmission: In this case, this module generates an NRDY interrupt when it has received the IN token and sets the bit corresponding to the pipe to 1. If the pipe where the interrupt has occurred is for isochronous transfer, this module transmits a zero-length packet and sets the OVRN bit to 1.

(b) Conditions for Pipes in the Receiving Direction

- When an OUT token is received while the FIFO buffer has no empty space: In the case of a pipe specified for isochronous transfer, this module generates an NRDY interrupt request when it has received the OUT token and sets the bit corresponding to the pipe and the OVRN bit to 1. In the case of a pipe not specified for isochronous transfer, this module generates an NRDY interrupt request when it issues a NAK handshake after receiving data that comes following the OUT token and sets the bit corresponding to the pipe to 1. Note that the NRDY interrupt request is not generated for re-transmission (on a mismatch of data PID). The interrupt request is also not generated when the data packet contains any error.
- When a PING token is received while the FIFO buffer has no empty space: In this case, this module generates an NRDY interrupt request when it has received the PING token and sets the bit corresponding to the pipe to 1.
- In a pipe specified for isochronous transfer, when the token is not received normally within the interval frame:

In this case, this module generates an NRDY interrupt request when it has received SOF and sets the bit corresponding to the pipe to 1.

Note: * Only 0 can be written to these bits.

24.3.24 BEMP Interrupt Status Register (BEMPSTS)

BEMPSTS is used to confirm the BEMP interrupt status for each pipe

This register is initialized by a power-on reset.

When this module detects a BEMP interrupt in the pipe for which $PID = BUF$ is set, the corresponding bit for that pipe is set to 1. This module generates an internal BEMP interrupt request in the following conditions.

(a) Pipe in the Transmitting Direction

An internal BEMP interrupt request is generated in a transmitting pipe when transmission has been completed (including zero-length packet transmission) while the FIFO buffer for the pipe is empty. With a single-buffer configuration, an internal BEMP interrupt request is generated at the same time as the BEMP interrupt for the pipes other than the DCP.

However, an internal BEMP interrupt request is not generated in the following cases.

- In a double-buffer configuration, writing to the FIFO buffer on the CPU side has been started when transmission for a single plane of data is completed.
- Buffer is cleared (has become empty) by writing 1 to the ACLRM or BCLR bit.
- When the function module function has been set, IN transfer (zero-length packet transmission) is executed for the status stage in control transfer.

(b) Pipe in the Receiving Direction

An internal BEMP interrupt request is generated when data of the amount greater than the size set by MaxPacketSize has been received normally. In this case, this module generates a BEMP interrupt request, sets the bit corresponding to the pipe to 1, discards the received data, and modifies the PID bit setting for the pipe to STALL. This module returns no response when the host controller function is selected, and carries out the STALL response when the function controller function is selected.

However, an internal BEMP interrupt request is not generated in the following cases.

- A CRC error or bit stuffing error is detected in the received data.
- A setup transaction is executed.

Note: * Only 0 can be written to these bits.

24.3.25 Frame Number Register (FRMNUM)

FRMNUM determines the source of isochronous error notification and indicates the frame number.

Note: * Only 0 can be written to.

24.3.26 μ**Frame Number Register (UFRMNUM)**

UFRMNUM indicates the μframe number.

24.3.27 USB Address Register (USBADDR)

USBADDR indicates the USB address. This register is valid only when the function controller function is selected. When the host controller function is selected, function device addresses should be set using the DEVSEL bits in PIPEMAXP.

24.3.28 USB Request Type Register (USBREQ)

USBREQ is a register for storing setup requests for control transfer.

When the function controller function is selected, the values of bRequest and bmRequestType that have been received are stored.

When the host controller function is selected, the values of bRequest and bmRequestType to be transmitted should be set. After setting SUREQ to 1, do not rewrite this register until 0 is read from SUREQ.

24.3.29 USB Request Value Register (USBVAL)

USBVAL is a register for storing setup requests for control transfer.

When the function controller function is selected, the value of wValue that has been received is stored.

When the host controller function is selected, the value of wValue to be transmitted should be set. After setting SUREQ to 1, do not rewrite this register until 0 is read from SUREQ.

24.3.30 USB Request Index Register (USBINDX)

USBINDEX is a register for storing setup requests for control transfer.

When the function controller function is selected, the value of wIndex that has been received is stored.

When the host controller function is selected, the value of wIndex to be transmitted should be set. After setting SUREQ to 1, do not rewrite this register until 0 is read from SUREQ.

24.3.31 USB Request Length Register (USBLENG)

USBLENG is a register for storing setup requests for control transfer.

When the function controller function is selected, the value of wLength that has been received is stored.

When the host controller function is selected, the value of wLength to be transmitted should be set. After setting SUREQ to 1, do not rewrite this register until 0 is read from SUREQ.

24.3.32 DCP Configuration Register (DCPCFG)

DCPCFG specifies the data transfer direction for the default control pipe (DCP).

Notes: When the function controller function is selected, the DIR bit should be cleared to 0.

When changing the setting of the DCP's PID bits after they have been changed from BUF to NAK, confirm that both CSSTS and PBUSY are cleared to 0 before changing the PID bits. However, it is not necessary to check the PBUSY bit if the PID bits were changed to NAK by the USB module.

Change the setting of the CNTMD bit when $CSSTS = 0$, PID = NAK, and the CURPIPE bits are unset. When changing the setting of the CNTMD bit after USB communication using the DCP, write 1 to BCLR and clear the FIFO buffer allocated to the DCP, in addition to ensuring that the above three registers are as described.

Change the setting of the SHTNAK bit when $CSSTS = 0$ and $PID = NAK$.

Table 24.12 Setting Value of CNTMD Bit and Method of Determining Completion of Transmit and Receive to and from FIFO Buffer

24.3.33 DCP Maximum Packet Size Register (DCPMAXP)

DCPMAXP specifies the maximum packet size for the DCP.

Note: ∗ The DEVSEL bit should be set while CSSTS is 0, PID is NAK, and SUREQ is 0. The MXPS bit should be set while CSSTS is 0 and PID is NAK. Before modifying these bits after modifying the PID bits for the DCP from BUF to NAK, check that CSSTS and PBUSY are 0. However, if the PID bits have been modified to NAK by this module, checking PBUSY is not necessary. When the function controller function is selected, the DEVSEL bit should be set to 0.

24.3.34 DCP Control Register (DCPCTR)

DCPCTR is a register that is used to confirm the buffer memory status, control setup transactions and split transactions, change and confirm the data PID sequence bit, and set the response PID for the DCP.

This register is initialized by a power-on reset. The CCPL and PID[2:0] bits are initialized by a USB bus reset.

Notes: 1. This bit is always read as 0.

2. Only 1 can be written to.

Notes: When the function controller function is selected, bits SUREQ, CSCLR, CSSTS, SUREQCLR, and PINGE should be cleared to all 0s.

When the host controller function is selected, bit CCPL should be cleared to 0.

- 1. This bit is always read as 0.
- 2. Only 1 can be written to.
- 3. These bits should be modified while CSSTS is 0 and PID is NAK. Before modifying this bit after modifying the PID bits from BUF to NAK, make sure that CSSTS and PBUSY are 0. However, if the PID bits have been modified to NAK by this module, checking PBUSY is not necessary.

24.3.35 Pipe Window Select Register (PIPESEL)

PIPESEL selects the pipe for use from among PIPE1 to PIPE9. After a pipe has been selected, configure the functions of the individual pipe using the PIPECFG, PIPEBUF, PIPEMAXP, and PIPEPERI registers. The PIPEnCTR, PIPEnTRE, and PIPEnTRN registers can be set regardless of the pipe selection in PIPESEL.

24.3.36 Pipe Configuration Register (PIPECFG)

PIPECFG is a register that specifies the transfer type, buffer memory access direction, and endpoint numbers for PIPE1 to PIPE9. It also selects continuous or non-continuous transfer mode, single or double buffer mode, and whether to continue or disable pipe operation at the end of transfer.

- Notes: Before modifying these bits after modifying the PID bits for the selected pipe from BUF to NAK, make sure that $\text{CSTS} = 0$ and $\text{PBUSY} = 0$. However, if the PID bits have been modified to NAK by this module, checking of PBUSY is not necessary.
	- ∗ Modify the TYPE bit while the PID bit using the selected pipe is NAK.

 Modify the BFRE, DBLB, CNTMD, and DIR bits while CSSTS is 0, PID is NAK, and the CURPIPE bits has not been set. When modifying the settings of bits after USB communication is done using the selected pipe, successively write 1 and then 0 to ACLRM to clear the FIFO buffer assigned to the selected pipe, in addition to satisfying the conditions of the above mentioned three registers.

Modify the SHTNAK bit while CSSTS is 0 and PID is NAK.

 Modify the EPNUM bit while CSSTS is 0, PID is NAK, and the CURPIPE bit has not been set.

Table 24.13 CNTMD Bit Setting and Method of Determining Completion of Transmission/Reception for the FIFO Buffer

24.3.37 Pipe Buffer Setting Register (PIPEBUF)

PIPEBUF is specifies the buffer size and buffer number for PIPE1 to PIPE9.

This register is initialized by a power-on reset.

Notes: Before modifying these bits after modifying the PID bits for the selected pipe from BUF to NAK, make sure that $CSSTS = 0$ and $PBUSY = 0$. However, if the PID bits have been modified to NAK by this module, checking of PBUSY is not necessary.

∗ Modify these bits while CSSTS is 0, PID is NAK, and the CURPIPE bit has not been set.

24.3.38 Pipe Maximum Packet Size Register (PIPEMAXP)

PIPEMAXP sets the function device address when the host controller function is selected and specifies the maximum packet size for PIPE1 to PIPE9.

This register is initialized by a power-on reset.

Notes: When the function controller function is selected, the DEVSEL bit is cleared to 0.

1. The DEVSEL bit should be set while CSSTS = 0 and PID = NAK.

 Modify the MXPS bit while CSSTS is 0 and PID is NAK and before The CURPIPE bit is set.

 Before modifying these bits after modifying the PID bits for the selected pipe from BUF to NAK, make sure that CSSTS = 0 and PBUSY = 0. However, if the PID bits have been modified to NAK by this module, checking of PBUSY is not necessary.

 2. The initial value of MXPS is H'000 when no pipe is selected with the PIPESEL bits in PIPESEL and H'040 when a pipe is selected with the PIPESEL bit in PIPESEL.

24.3.39 Pipe Timing Control Register (PIPEPERI)

PIPEPERI is a register that selects whether the buffer is flushed or not when an interval error occurred during isochronous IN transfer, and sets the interval error detection interval for PIPE1 to PIPE9.

This register is initialized by a power-on reset.

Notes: When the host controller function is selected, set the IFIS bit to 0.

∗ Modify these bits while CSSTS is 0 and PID is NAK and before the CURPIPE bit is selected.

 Before modifying these bits after modifying the PID bits for the selected pipe from BUF to NAK, make sure that $CSSTS = 0$ and $PBUSY = 0$. However, if the PID bits have been modified to NAK by this module, checking of PBUSY is not necessary.

The interval error detection interval bits (IITV) are explained below.

(1) When the Host Controller Function is Selected

The IITV bits can be set when the selected pipe is specified for isochronous transfer or interrupt transfer. This module controls the intervals of token issuance according to the setting of this bit. It issues a token for the selected pipe once in 2^{IITV} (micro) frames.

This module counts 1-ms frames to measure the interval for the pipe used for communication with a full-speed/low-speed function device connected to a high-speed hub.

This module starts counting for token issuance intervals from the first (micro) frame that comes after the PID bits have been set to BUF.

USB bus	S O F		S O F		S O F	O U	A \overline{A}	S O F		D A A 0
PID bit setting		NAK		BUF		BUF			BUF	
Token issuance (0: Issued -: Not issued)						Ω			Ω	
Start of interval counting										

Figure 24.1 Whether a Token is Issued or Not when IITV = 0

USB bus	S O F	S Ο F	S O U U Α Ő T E A	S O F	S D U O A T F A 0	S F	D S C A F \overline{A} 0
PID bit setting	NAK	BUF	BUF	BUF	BUF	BUF	BUF
Token issuance (0: Issued -: Not issued)	-	-	0	$\overline{}$	Ω		Ω
Start of interval counting							

Figure 24.2 Whether a Token is Issued or Not when IITV = 1

When the selected pipe is for isochronous transfer, this module performs the following operations while it controls the token issuance intervals. The module issues tokens even when the condition for NRDY interrupt generation is met.

When the selected pipe is for isochronous transfer and set in the IN direction: The module issues an IN token and, if it does not receive packets from the function device normally (i.e., no response or packet error), generates an NRDY interrupt. If the time to issue an IN token comes in a situation where the module is unable to receive data because of the full FIFO buffer, which may be caused by slow or delayed data reading from the buffer or some other reasons, the module sets the OVRN bit to 1 and generates an NRDY interrupt.

• When the selected pipe is for isochronous transfer and set in the OUT direction:

If the time to issue an OUT token comes in a situation where the FIFO buffer does not contain any data ready for transmission because, for example, data writing to the FIFO buffer is slow or delayed, the module sets the OVRN bit to 1, generates an NRDY interrupt, and transmits a zero-length packet.

The interval of token issuance is reset by a power-on reset or when the ACLRM bit is set to 1.

(2) When the Function Controller Function is Selected

When the selected pipe is for isochronous transfer and set in the OUT direction:

If the module does not receive any data packets within the (micro) frames for the interval specified by the IITV bits, it generates an NRDY interrupt.

The module also generates an NRDY interrupt when it cannot receive a data packet normally because any error, such as a CRC error, is found in the incoming data packet or when the module is unable to receive data because of the full FIFO buffer, which may be caused by slow or delayed data reading from the buffer or some other reasons.

An NRDY interrupt is generated with the timing of receiving an SOF packet. Even if the SOF packet is corrupted, the interrupt is generated with the proper timing of SOF reception by the internal interpolating function. Note that in the cases other than $IITY = 0$, an NRDY interrupt is generated upon SOF packet reception at every interval after the counting of the interval is started.

If the PID bits are set to NAK by software after the interval timer is started, this module does not generate an NRDY interrupt even when it receives an SOF packet.

The conditions for starting the interval timer differ according to the IITV bit setting.

(a) When $IITV = 0$

Counting of the interval is started from the first (micro) frame after the PID bit setting of the selected pipe is changed to BUF.

Figure 24.3 (Micro) Frames and Expectation for Token Reception when IITV = 0

(b) When Other than $\text{IITV} = 0$

Counting of the interval is started on completion of the first normal reception of data packet after the PID bit setting of the selected pipe is changed to BUF.

USB bus	S О E	S	S Ο A F \overline{A}	S	S A E A 0	S O	S A F A 0
PID bit setting	NAK	BUF	BUF	BUF	BUF	BUF	BUF
Expectation for token reception (0: Expect to receive -: Expect not to receive)			Ω		0		
Start of interval counting							

Figure 24.4 (Micro) Frames and Expectation for Token Reception when IITV = 1

• When the selected pipe is for isochronous transfer and set in the IN direction:

The IITV bits should be used in combination with IFIS = 1. With IFIS = 0, the module transmits data packets in response to the received tokens regardless of the IITV bit setting. When $IFIS = 1$ is set, the module clears the FIFO buffer if it does not receive any IN token within the (micro) frames for the interval specified by the IITV bits.

The module also clears the FIFO buffer when it cannot receive normally because a bus error, such as a CRC error, has occurred in the IN token.

The FIFO buffer is cleared with the timing of receiving an SOF packet. Even if the SOF packet is corrupted, the FIFO buffer is cleared with the proper timing of SOF reception by the internal interpolating function.

The clearing conditions for starting the interval timer differ according to the IITV bit setting (same as the case with the OUT direction).

- (a) Power-on reset
- (b) When $\text{ACLRM} = 1$ is set
- (c) When USB bus reset is detected

24.3.40 PIPEn Control Registers (PIPEnCTR) (n = 1 to 5)

The PIPEnCTR registers for PIPE1 to PIPE5 are used to confirm the buffer memory status, change and confirm the data PID sequence bit, determine whether auto response mode is set, determine whether auto buffer clear mode is set, and set a response PID for the corresponding pipe. These registers can be set regardless of the pipe selection in PIPESEL.

This register is initialized by a power-on reset. PID[1:0] are initialized by a USB bus reset.

- Notes: When the function controller function is selected, clear the CSCLR bit to 0. And when the host controller function is selected, clear the ATREPM bit to 0.
	- 1. Only 0 can be read.
	- 2. Only 1 can be written to.
	- 3. Modify ATREPM, SZCLR, and SQSET bits while CSSTS is 0 and PID is NAK. Modify the ACLRM bit while CSSTS is 0, PID is NAK, and before the CURPIPE bit is selected. Before modifying these bits after modifying the PID bits for the selected pipe from BUF to NAK, make sure that $CSSTS = 0$ and $\overline{PBUSY} = 0$. However, if the PID bits have been modified to NAK by this module, checking of PBUSY is not necessary.

Table 24.14 Meaning of BSTS Bit

Table 24.15 (1) Information Cleared by this Module by Setting ACLRM = 1

No. Information Cleared by ACLRM Bit Manipulation

- 1 All the information in the FIFO buffer assigned to the pertinent pipe (all the information in two FIFO buffer planes in double buffer mode)
- 2 The interval count value when the pertinent pipe is for isochronous transfer

Table 24.15 (2) Cases That Require Setting ACLRM to 1

No. Cases in which Clearing the Information is Necessary

Table 24.16 Operation of This Module depending on PID Setting (when Host Controller Function is Selected)

Table 24.17 Operation of This Module depending on PID Setting (when Function Controller Function is Selected)

24.3.41 PIPEn Control Registers (PIPEnCTR) (n = 6 to 9)

The PIPEnCTR registers for PIPE6 to PIPE9 are used to confirm the buffer memory status, change and confirm the data PID sequence bit, determine whether auto buffer clear mode is set, and set a response PID for the corresponding pipe. These registers can be set regardless of the pipe selection in PIPESEL.

This register is initialized by a power-on reset. PID[1:0] are initialized by a USB bus reset.

Notes: When the function controller function is selected, clear the CSCLR bit to 0.

- 1. Only 0 can be read.
- 2. Only 1 can be written to.
- 3. Modify the ACLRM bit while CSSTS is 0, PID is NAK, and before the CURPIPE bit is selected. Modify the SQCLR and SQSET bits while CSSTS is 0 and PID is NAK.

 Before modifying these bits after modifying the PID bits for the selected pipe from BUF to NAK, make sure that CSSTS = 0 and $PBUSY = 0$. However, if the PID bits have been modified to NAK by this module, checking of PBUSY is not necessary.

Table 24.18 (1) Information Cleared by this Module by Setting ACLRM = 1

No. Information Cleared by ACLRM Bit Manipulation

Table 24.18 (2) Cases That Require Setting ACLRM to 1

24.3.42 Transaction Counter Enable Registers (PIPEnTRE) (n = 1 to 5)

The PIPEnTRE registers configure transaction counter operations for PIPE1 to PIPE5. These registers can be set regardless of the pipe selection in PIPESEL.

This register is initialized by a power-on reset.

Notes: 1. Only 0 can be read.

- 2. Only 1 can be written to.
- 3. Modify each bit when $CSSTS = 0$ and $PID = NAK$.

 Before modifying these bits after modifying the PID bits for the selected pipe from BUF to NAK, make sure that CSSTS = 0 and $PBUSY = 0$. However, if the PID bits have been modified to NAK by this module, checking of PBUSY is not necessary.

24.3.43 Transaction Counter Registers (PIPEnTRN) (n = 1 to 5)

The PIPEnTRN registers are used to specify the number of transactions by DMA transfer for PIPE1 to PIPE5, and the current number of transactions can be read from them.

These registers are initialized by a power-on reset.

transferred by DMA.

 Before modifying these bits after modifying the PID bits for the corresponding pipe from BUF to NAK, make sure that $CSSTS = 0$ and $PBUSY = 0$. However, if the PID bits have been modified to NAK by this module, checking of PBUSY is not necessary.

24.3.44 Device Address Configuration Registers (DEVADDn) (n = 0 to 9, A)

When the host controller function is selected, the DEVADDn registers specify the address and port number of the hub to which the communication target function device is connected and also specifies the communication speed of the function device.

These registers are initialized by a power-on reset.

Notes: 1. When the host controller function is selected, be sure to set the bits in this register before starting the communication on each pipe.

 2. The settings of the bits in this register should be modified when there is no active pipe using the setting of these bits. The active pipe refers to a pipe that satisfies both of the following conditions:

- (1) The DEVSEL bit setting is designating this register.
- (2) The PID bits of the pipe are set to BUF, or the pipe is the DCP and SUREQ = 1.
- 3. When the function controller function is selected, all the bits in this register should be set to 0.
24.3.45 USB AC Characteristics Switching Register 0 (USBACSWR0)

The USBACSWR0 registers specify a USB transceiver that is stored in this module. This register is initialized by a power-on reset.

For details, refer to section 24.5.1, Procedure for Setting the USB Transceiver.

24.3.46 USB AC Characteristics Switching Register 1 (USBACSWR1)

The USBACSWR1 registers specify a USB transceiver that is stored in this module. This register is initialized by a power-on reset.

For details, refer to section 24.5.1, Procedure for Setting the USB Transceiver.

24.4 Operation

24.4.1 System Control

This section describes the register operations that are necessary to the initial settings of this module, and the registers necessary for power consumption control.

(1) Resets

Table 24.19 lists the types of controller resets. For the initialized states of the registers following the reset operations, see section 24.3, Register Description.

Table 24.19 Types of Reset

Name	Operation
Power-on reset	Low level input from the RES pin
USB bus reset	Automatically detected by this module from the $D+$ and $D-$ lines when the function controller function is selected

(2) Controller Function Selection

This module can select the host controller function or function controller function using the DCFM bit in SYSCFG0.

This controller selects functions for each USB port as shown in table 24.19.

Table 24.20 Functions Selected for USB Ports

When Host Mode is Selected (DCFM = 1)

(3) USB Data Bus Resistor Control

This module controls switching of the pull-up resistor for the D+ signal and a pull-down resistor for the D+ and D- signals. These signals can be pulled up or down using the DPRPU and DRPD bits in SYSCFG0 (for PORT0) and DRPD bit in SYSCFG1 (for PORT1).

When the function controller function is selected, set the DPRPU bit in the SYSCFG register to 1 after the connection to the USB host is recognized to pull up the D+ signal. When disconnection from the USB host is recognized, manipulate the DPRPU and DCFM bits as follows:

- 1. Clear the DPRPU bit to 0.
- 2. Wait for at least 1 us.
- 3. Set the DCFM bit to 1.
- 4. Wait for at least 200 ns.
- 5. Clear the DCFM bit to 0.

This module includes the terminal resistor for the D+ and D- signals during high-speed operation and the output resistor for the signals during full-speed operation. This module automatically switches the resistor after connection with the USB host or function device when reset handshake, suspended state or resume is detected.

When the function controller function is selected and the DPRPU bit in SYSCFG0 is cleared to 0 during communication with the USB host, the pull-up resistor (or the terminal resistor) of the USB data line is disabled. This allows notification of device disconnection to the USB host.

(4) Register Access Wait Control

The following restrictions apply to numbers of cycles for access to registers of this module below SYSSTS₀.

The cycle time for consecutive access to registers of this module must be at least 4 USB clock (48 MHz) cycles (83.33 ns).

To comply with this constraint, the BWIT[3:0] bits of the SYSCFG1 register must be set to apply wait control for register access. Since the initial value is the largest value (17 clock cycles for a cycle of access), select the optimal value.

Example of settings (1): Consecutive access to registers of this module Bus-clock frequency: 66 MHz Calculation: (2 cycles (access cycle for registers of this module) + 1 cycle (interval between consecutive access operations) + BWAIT) \times 1/66 MHz \geq 83.33 ns $BWAIT = 3$

Example of settings (2): Transfer of data from internal memory to the FIFO port registers Bus-clock frequency: 66 MHz Calculation: $(2 \text{ cycles} (access \text{ cycle} for \text{ registers} of \text{ this module}) + 2 \text{ cycles} (access \text{ cycle} for \text{?})$ internal memory) + BWAIT) \times 1/66 MHz \geq 83.33 ns

 $BWAIT = 2$

24.4.2 Interrupt Functions

(1) Interrupt Control Overview

Table 24.21 lists the interrupt generation conditions for this module.

When an interrupt generation condition is satisfied and the interrupt output is enabled using the corresponding interrupt enable register, this module outputs the USB interrupt request signal to the INTC.

Table 24.21 Interrupt Generation Conditions

Note: All the bits without register name indication are in INTSTS0.

Figure 24.5 shows a diagram relating to interrupts of this module.

Figure 24.5 Items Relating to Interrupts

(2) Device State Transition Interrupt (Function Controller Function)

Figure 24.6 shows a diagram of this module device state transitions. This module controls device states and generates device state transition interrupts. However, recovery from the suspended state (resume signal detection) is detected by means of the resume interrupt. The device state transition interrupts can be enabled or disabled individually using INTENB0. The device state that made a transition can be confirmed using the DVSQ bit in INTSTS0.

To make a transition to the default state, the device state transition interrupt is generated after the reset handshake protocol has been completed.

Device state can be controlled only when the function controller function is selected. Also, the device state transition interrupts can be generated only when the function controller function is selected.

Figure 24.6 Device State Transitions

(3) Control Transfer Stage Transition Interrupt (Function Controller Function)

Figure 24.7 shows a diagram of how this module handles the control transfer stage transition. This module controls the control transfer sequence and generates control transfer stage transition interrupts. Control transfer stage transition interrupts can be enabled or disabled individually using INTENB0. The transfer stage that made a transition can be confirmed using the CTSQ bit in INTSTS0.

Control transfer stage transition interrupts are only generated when the function controller function is selected.

The control transfer sequence errors are described below. If an error occurs, the PID bit in DCPCTR is set to B'1x (STALL).

- 1. During control read transfers
	- At the IN token of the data stage, an OUT or PING token is received when there have been no data transfers at all.
	- $\overline{}$ An IN token is received at the status stage
	- $\overline{}$ A packet is received at the status stage for which the data packet is DATAPID = DATA0
- 2. During control write transfers
	- At the OUT token of the data stage, an IN token is received when there have been no ACK response at all
	- $\overline{}$ A packet is received at the data stage for which the first data packet is DATAPID = DATA0
	- At the status stage, an OUT or PING token is received
- 3. During no-data control transfers
	- At the status stage, an OUT or PING token is received

At the control write transfer stage, if the number of receive data exceeds the wLength value of the USB request, it cannot be recognized as a control transfer sequence error. At the control read transfer status stage, packets other than zero-length packets are received by an ACK response and the transfer ends normally.

When a CTRT interrupt occurs in response to a sequence error (SERR $= 1$), the CTSQ $= 110$ value is retained until $CTRT = 0$ is written from the system (the interrupt status is cleared). Therefore, while $CTSQ = 110$ is being held, the CTRT interrupt that ends the setup stage will not be generated even if a new USB request is received. (This module retains the setup stage end, and after the interrupt status has been cleared by software, a CTRT interrupt is generated.)

Figure 24.7 Control Transfer Stage Transitions

24.4.3 Pipe Control

Table 24.22 lists the pipe setting items of this module. With USB data transfer, data transmission has to be carried out using the logic pipe called the endpoint. This module has ten pipes that are used for data transfer.

Settings should be entered for each of the pipes in conjunction with the specifications of the system.

Table 24.22 Pipe Setting Items

(1) Pipe Control Register Modifying Procedures

The following bits in the pipe control registers can be modified only when USB communication is disabled (PID = NAK). Figure 24.8 shows the procedure for rewriting the pipe control registers from a state in which USB communication is enabled (PID = BUF).

Registers that Should Not be Set when USB Communication is Enabled (PID = BUF):

- Bits in DCPMAXP
- The SOCLR, SOSET, and PINGE bits in DCPCTR
- Bits in PIPECFG, PIPEBUF, PIPEMAXP and PIPEPERI
- The ATREPM, ACLRM, SOCLR and SOSET bits in PIPEnCTR
- Bits in PIPEnTRE and PIPEnTRN
- Bits in DEVADDn
- Note: In addition to the above, observe the setting procedures described in the register descriptions regarding the settings of the CSCLR bit and DEVADDn register.

Figure 24.8 Procedure for Changing Pipe Information from a USB Communication Enabled State (PID = BUF)

The following bits in the pipe control registers can be modified only for the pipes that are not specified in the CURPIPE bits for any of the CPU/DMA0/DMA1-FIFO ports.

Registers that Should Not be Set When CURPIPE for FIFO Port is Set.

- Bits in DCPCFG and DCPMAXP
- Bits in PIPECFG, PIPEBUF, PIPEMAXP and PIPEPERI
- The ACLRM bit in PIPEnCTR

In order to modify pipe information, the CURPIPE bits should be set to the pipes other than the pipe to be modified. For the DCP, the buffer should be cleared using BCLR after the pipe information is modified.

(2) Maximum Packet Size Setting

The MXPS bit in DCPMAXP and PIPEMAXP is used to specify the maximum packet size for each pipe. DCP and PIPE1 to PIPE5 can be set to any of the maximum pipe sizes defined by the USB specification. For PIPE6 to PIPE9, 64 bytes are the upper limit of the maximum packet size. The maximum packet size should be set before beginning the transfer ($PID = BUF$).

- 1. DCP: 64 should be set when using high-speed operation.
- 2. DCP: Select and set 8, 16, 32, or 64 when using full-speed operation.
- 3. PIPE1 to PIPE5: 512 should be set when using high-speed bulk transfer.
- 4. PIPE1 to PIPE5: Select and set 8, 16, 32, or 64 when using full-speed bulk transfer.
- 5. PIPE1 and PIPE2: Set a value between 1 and 1024 when using high-speed isochronous transfer.
- 6. PIPE1 and PIPE2: Set a value between 1 and 1023 when using full-speed isochronous transfer.
- 7. PIPE6 to PIPE9: Set a value between 1 and 64.

The high bandwidth transfers used with interrupt transfers and isochronous transfers are not supported.

(3) Response PID

The PID bits in DCPCTR and PIPEnCTR are used to set the response PID for each pipe.

The following shows this module operation with various response PID settings:

(a) Response PID settings when the host controller function is selected

The response PID is used to specify the execution of transactions.

- NAK setting: Using pipes is disabled. No transaction is executed.
- BUF setting: Transactions are executed based on the status of the buffer memory. For OUT direction: If there are transmit data in the buffer memory, an OUT token is issued. For IN direction: If there is an area to receive data in the buffer memory, an IN token is issued.
- STALL setting: Using pipes is disabled. No transaction is executed.

Setup transactions for the DCP are set with the SUREQ bit.

(b) Response PID settings when the function controller function is selected

The response PID is used to specify the response to transactions from the host.

- NAK setting: The NAK response is always returned in response to the generated transaction.
- BUF setting: Responses are made to transactions based on the status of the buffer memory.
- STALL setting: The STALL response is always returned in response to the generated transaction.

For setup transactions, an ACK response is always returned, regardless of the PID setting, and the USB request is stored in the register.

This module may carry out writing to the PID bits, depending on the results of the transaction.

(a) When the host controller function has been selected and the response PID is set by hardware

- NAK setting: In the following cases, $PID = NAK$ is set and issuing of tokens is automatically stopped:
	- ⎯ When, during a transfer other than isochronous transfer, three receive errors such as no response, bit stuffing error, or CRC error are returned in succession after token transmission.
	- ⎯ When, during isochronous transfer, three receive errors such bit stuffing error or CRC error are returned in succession after token transmission.
	- ⎯ When a short packet is received in the data stage of control read transfer while the SHTNAK bit in the DCPCFG register is set to 1.
	- ⎯ If a short packet is received when the SHTNAK bit in PIPECFG has been set to 1 for bulk transfer.
	- ⎯ If the transaction counter ended when the SHTNAK bit has been set to 1 for bulk transfer.
- BUF setting: There is no BUF writing by this module.
- STALL setting: In the following cases, PID = STALL is set and issuing of tokens is automatically stopped:
	- ⎯ When STALL is received in response to the transmitted token.
	- ⎯ When the size of the receive data packet exceeds the maximum packet size.

(b) When the function controller function has been selected and the response PID is set by hardware

- NAK setting: In the following cases, $PID = NAK$ is set and NAK is always returned in response to transactions:
	- When the SETUP token is received normally (DCP only).
	- ⎯ If the transaction counter ended or a short packet is received when the SHTNAK bit in PIPECFG has been set to 1 for bulk transfer.
- BUF setting: There is no BUF writing by this module.
- STALL setting: In the following cases, PID = STALL is set and STALL is always returned in response to transactions:
	- ⎯ When the size of the receive data packet exceeds the maximum packet size.
	- When a control transfer sequence error has been detected (DCP only).

(4) Data PID Sequence Bit

This module automatically toggles the sequence bit in the data PID when data is transferred normally in the control transfer data stage, bulk transfer and interrupt transfer. The sequence bit of the data PID that was transmitted can be confirmed with the SQMON bit in DCPCTR and PIPEnCTR. When data is transmitted, the sequence bit switches at the timing at which the ACK handshake is received. When data is received, the sequence bit switches at the timing at which the ACK handshake is transmitted. The SQCLR bit in DCPCTR and the SQSET bit in PIPEnCTR can be used to change the data PID sequence bit.

When the function controller function has been selected and control transfer is used, this module automatically sets the sequence bit when a stage transition is made. The bit is set to DATA1 when the setup stage ends. In the status stage, DATA1 is returned without referencing the sequence bit. Therefore, settings are not required. However, when the host controller function has been selected and control transfer is used, the sequence bit should be set at the stage transition.

For the Clearfeature request transmission or reception, the data PID sequence bit should be set, regardless of whether the host controller function or function controller function is selected.

With pipes for which isochronous transfer has been set, sequence bit operation cannot be carried out using the SQSET bit.

24.4.4 FIFO Buffer

This section describes the operation related to the FIFO buffer in the USB module. Unless otherwise noted, the operation is the same regardless of whether the host controller or function controller function is selected.

(1) FIFO Buffer Allocation

Figure 24.9 shows an example of FIFO buffer mapping of this module. The FIFO buffer is an area shared by the CPU and this module. As the FIFO buffer statuses, there is a state in which the access right to the FIFO buffer is held by the system (CPU side), and a state in which the access right is held by this module (SIE side).

Areas for the FIFO buffer are set independently for each pipe. A memory area is set by specifying the first block number and the number of blocks, where one block consists of 64 bytes (the settings are made through the BUFNMB and BUFSIZE bits in PIPEBUF). When continuous transfer mode has been selected using the CNTMD bit in PIPECFG, the BUFSIZE bits should be set so that the buffer memory size should be an integral multiple of the maximum packet size. When double buffer mode has been selected using the DBLB bit in PIPECFG, two planes of the memory area specified using the BUFSIZE bits in PIPEBUF can be assigned to a single pipe.

Three FIFO ports are used for access to the FIFO buffer (reading and writing data). A pipe is assigned to the FIFO port by specifying the pipe number using the CURPIPE bit in C/DnFIFOSEL.

The FIFO buffer statuses of the various pipes can be confirmed using the BSTS bit in DCPCTR and the INBUFM bit in PIPEnCTR. Also, the access right of the FIFO port can be confirmed using the FRDY bit in C/DnFIFOCTR.

Figure 24.9 Example of FIFO Buffer Memory Mapping

(2) FIFO Buffer Clearing

Table 24.23 summarizes the clearing of the FIFO buffer by this module. The FIFO buffer can be cleared using the three bits shown below.

Table 24.23 List of Buffer Clearing Methods

(3) FIFO Port Functions

Table 24.24 shows the settings for the FIFO port functions of this module. In write access, writing data until the buffer is full (or the maximum packet size for non-continuous transfers) automatically enables sending of the data to the USB bus. To enable sending of data before the buffer is full (or before the maximum packet size for non-continuous transfers), the BVAL bit in C/DnFIFOCTR must be set to signal that the writing has ended. Also, to send a zero-length packet, the BCLR bit in the same register must be used to clear the buffer and then the BVAL bit set in order to signal the end of writing.

In read access, reception of new packets is automatically enabled if all of the data has been read. Data cannot be read when a zero-length packet is being received $(DTLN = 0)$, so the BCLR bit in the register must be used to release the buffer. The length of the data being received can be confirmed using the DTLN bit in C/DnFIFOCTR.

Table 24.24 FIFO Port Function Settings

(a) FIFO Port Selection

Table 24.25 shows the pipes that can be selected with the various FIFO ports. The pipe to be accessed is selected using the CURPIPE bit in C/DnFIFOSEL. After the pipe has been selected, check to see whether the CURPIPE value for the last-written pipe has been read correctly (if the previous pipe number is read, it indicates that the pipe changing processing is being done by this module). When a correct value has been read from CURPIPE, confirm $FRDY = 1$ and then access the FIFO port.

The bus width to be accessed should be selected using the MBW bit. The buffer memory access direction is in accord with the ISEL bit setting for the DCP, and in accord with the DIR bit in PIPECFG for other pipes.

Table 24.25 FIFO Port Access Categorized by Pipe

(b) Method of Reading Partial Data from FIFO ports

In reading data form an FIFO port when the width of the data to be read is shorter than the bit width set by the MBW bits in the FIFO port select register, read the bit width specified by the MBW bits and use software to discard the unnecessary portion of the data.

In writing data form an FIFO port when the width of the data to be written is shorter than the bit width set by the MBW bits in the FIFO port select register, proceed with access as shown in the examples below. These examples show ways of writing 24-bit data when the width for access to the FIFO port has been set to 32 bits (MBW = 10).

Example 1 of writing partial data: Writing data with 16-and 8-bit widths once each

Figure 24.10 Example 1 of Writing Partial Data to an FIFO Port

Example 2 of writing partial data: Write three times by 8-bit width

(c) Changing the MBW Bits when the Direction for the Specified Pipe is Reception

When the direction for the specified pipe is reception, write to the MBW bits of the FIFO port selection register (CFIFOSEL, D0FIFOSEL, and D1FIFOSEL) at the same time as the CURPIPE setting is made. When the CFIFO register has the DCP setting (CURPIPE $= 0$), write to the MBW bits at the same time as the settings for CURPIPE or ISEL are made.

The procedure for changing the setting of only the MBW bits for a given current pipe setting is shown below. However, once processing to read from the buffer memory has started, do not change the setting of the MBW bits until reading out of all data is complete.

This is applicable other than when the CURPIPE bits for DFIFO0, DFIFO1, or CFIFO have the DCP setting.

Figure 24.12 Example of Adjusting the MBW Bits when the CURPIPE Bits of DFIFO0, DFIFO1, or CFIFO Have a Setting Other than DCP (000)

This is applicable when the CURPIPE bits for CFIFO have the DCP setting (000).

Figure 24.13 Example of Adjusting the MBW Bits when the CURPIPE Bits of CFIFO have a DCP Setting (000)

(4) DMA Transfer (D0FIFO, D1FIFO Ports)

(a) Overview of DMA Transfer

For pipes 1 to 9, the FIFO port is accessible by the DMAC. When the buffer to the pipe set for DMA becomes accessible, a DMA transfer request is output.

In the DnFIFOSEL register, use the MBW bits to set the unit of transfer to the FIFO port for the pipe set for DMA transfer by the CURPIPE bits. Do not change the pipe number while DMA transfer is enabled.

(b) Automatic Recognition of Completion of DMA Transfer

In this module, completion of the writing of FIFO data by DMA transfer can be under the control of the input of a DMA transfer end signal. Then DMA transfer end signal causes the DMAC to proceed with DMA transfer the number of times that corresponds to the setting of the DMA current byte count register (DMCBCT) in the DMAC. Transfer to the buffer memory is enabled (the same as setting $BVAL = 1$) when the DMA transfer end signal is sampled. The setting for sampling or non-sampling of the DMA transfer end signal can be made in the TENDE bit of the DnFBCFG register. Furthermore, be sure to set the DMA transfer end signal output control bits (DTCM) in the DMA mode register (DMMOD) to 10 (output of the DMA transfer end signal on the final write cycle) whenever this function is used.

(c) Data Transferred for One Operand

In this module, the data to be transferred for each operand is selectable as one unit of data, 16 bytes, or 32 bytes by the DFACC bits in the DMA-FIFO bus-configuration registers (DnFBCFG).

- With the DFACC = 00 (access to one unit of data) setting, set the operand size for the DMAC's method of transfer to 1 and the data size to the size selected by the MBW bits.
- With the DFACC = 01 (consecutive access to 16 bytes) setting, set the operand size for the DMAC's method of transfer and the data size (the size selected by the MBW bits) such that the multiple of the two is 16 bytes.
- With the DFACC = 10 (consecutive access to 32 bytes) setting, set the operand size for the DMAC's method of transfer and the data size (the size selected by the MBW bits) such that the multiple of the two is 32 bytes.

(d) DnFIFO Auto Clear Mode (D0FIFO/D1FIFO Port Reading Direction)

If 1 is set for the DCLRM bit in DnFIFOSEL, the module automatically clears the buffer memory of the corresponding pipe when reading of the data from the buffer memory has been completed.

Table 24.26 shows the packet reception and buffer memory clearing processing for each of the various settings. As shown, the buffer clear conditions depend on the value set to the BFRE bit. Using the DCLRM bit eliminates the need for the buffer to be cleared by software even if a situation occurs that necessitates clearing of the buffer. This makes it possible to carry out DMA transfers without involving software.

This function can be set only in the buffer memory reading direction.

(e) BRDY Interrupt Timing Selection Function

By setting the BFRE bit setting in PIPECFG, it is possible to keep the BRDY interrupt from being generated when a data packet consisting of the maximum packet size is received.

When using DMA transfers, this function can be used to generate an interrupt only when the last data item has been received. The last data item refers to the reception of a short packet, or the ending of the transaction counter. When the BFRE bit is set to 1, the BRDY interrupt is generated after the received data has been read. When the DTLN bit in DnFIFOCTR is read, the length of the data received in the last data packet to have been received can be confirmed.

Table 24.27 shows the timing at which the BRDY interrupts are generated by this module.

Note: This function is valid only in the reading direction of reading from the buffer memory. In the writing direction, the BFRE bit should be fixed at 0.

24.4.5 Control Transfers (DCP)

Data transfers of the data stage of control transfers are done using the default control pipe (DCP). The DCP buffer memory is a 256-byte single buffer, and is a fixed area that is shared for both control reading and control writing. The buffer memory can be accessed through the CFIFO port.

(1) Control Transfers when the Host Controller Function is Selected

(a) Setup Stage

USQREQ, USBVAL, USBINDX, and USBLENG are the registers that are used to transmit a USB request for setup transactions. Writing setup packet data to the registers and writing 1 to the SUREQ bit in DCPCTR transmits the specified data for setup transactions. Upon completion of transactions, the SUREQ bit is cleared to 0. The above USB request registers should not be modified while $\text{SUREQ} = 1$. The device address for setup transactions is specified using the DEVSEL bits in DCPMAXP.

When the data for setup transactions has been sent, a SIGN or SACK interrupt request is generated according to the response received from the peripheral side (SIGN1 or SACK bits in INTSTS1), by means of which the result of the setup transactions can be confirmed.

A data packet of DATA0 (USB request) is transmitted as the data packet for the setup transactions regardless of the setting of the SQMON bit in DCPCTR.

(b) Data Stage

Data transfers are done using the DCP buffer memory.

The access direction of the DCP buffer memory should be specified using the ISEL bit in CFIFOSEL.

For the first data packet of the data stage, the data PID must be transferred as DATA1. Transaction is done by setting the data $PID = DATA1$ and the PID bit = BUF using the SQSET bit in DCPCTR. Completion of data transfer is detected using the BRDY or BEMP interrupts.

For control write transfers, when the number of data bytes to be sent is the integer multiple of the maximum packet size, control to send a zero-length packet at the end.

(c) Status Stage

Zero-length packet data transfers are done in the direction opposite to that in the data stage. As with the data stage, data transfers are done using the DCP buffer memory. Transactions are done in the same manner as the data stage.

For the data packets of the status stage, the data PID must be transferred as DATA1. The data PID should be set to DATA1 using the SQSET bit in DCPCFG.

For reception of a zero-length packet, the received data length must be confirmed using the DTLN bits in CFIFOCTR after the BRDY interrupt is generated, and the buffer memory must then be cleared using the BCLR bit in C/DnFIFOCTR.

(2) Control Transfers when the Function Controller Function is Selected

(a) Setup Stage

This module always sends an ACK response in response to a setup packet that is normal with respect to this module. The operation of this module operates in the setup stage is noted below.

- 1. When a new USB request is received, this module sets the following registers:
	- ⎯ Set the VALID bit in INTSTS0 to 1.
	- ⎯ Set the PID bit in DCPCTR to NAK.
	- Set the CCPL bit in DCPCTR to 0.
- 2. When a data packet is received right after the SETUP packet, the USB request parameters are stored in USBREQ, USBVAL, USBINDX, and USBLENG.

Response processing with respect to the control transfer should always be carried out after first setting VALID = 0. In the VALID = 1 state, $PID = BUF$ cannot be set, and the data stage cannot be terminated.

Using the function of the VALID bit, this module is able to interrupt the processing of a request currently being processed if a new USB request is received during a control transfer, and can send a response in response to the newest request.

Also, this module automatically judges the direction bit (bit 8 of the bmRequestType) and the request data length (wLength) of the USB request that was received, and then distinguishes between control read transfers, control write transfers, and no-data control transfers, and controls the stage transition. For a wrong sequence, the sequence error of the control transfer stage transition interrupt is generated, and the software is notified. For information on the stage control of this module, see figure 24.7.

(b) Data Stage

Data transfers corresponding to USB requests that have been received should be done using the DCP. Before accessing the DCP buffer memory, the access direction should be specified using the ISEL bit in CFIFOSEL.

A transaction is executed by setting the PID bits in the DCPCTR register to BUF. The BRDY interrupt or the BEMP interrupt can be used to detect the end of data transfer. Use the BRDY interrupt to detect the end of control write transfers and the BEMP interrupt to detect the end of control read transfers.

With control write transfers during high-speed operation, the NYET handshake response is carried out in accordance with the state of the buffer memory.

(c) Status Stage

Control transfers are terminated by setting the CCPL bit to 1 with the PID bit in DCPCTR set to $PID = RHF$.

After the above settings have been entered, this module automatically executes the status stage in accordance with the data transfer direction determined at the setup stage. The specific procedure is as follows.

1. For control read transfers:

The zero-length packet is received from the USB host, and this module sends an ACK response.

2. For control write transfers and no-data control transfers:

This module sends a zero-length packet and receives an ACK response from the USB host.

(d) Control Transfer Auto Response Function

This module automatically responds to a normal SET_ADDRESS request. If any of the following errors occur in the SET_ADDRESS request, a response from the software is necessary.

- 1. bmRequestType \neq H'00
- 2. wIndex \neq H'00
- 3. wLength \neq H'00
- 4. wValue $>$ H'7F
- 5. $DVSQ = 011$ (Configured)

For all requests other than the SET_ADDRESS request, a response is required from the corresponding software.

24.4.6 Bulk Transfers (PIPE1 to PIPE5)

The buffer memory specifications for bulk transfers (single/double buffer setting, or continuous/non-continuous transfer mode setting) can be selected. The maximum size that can be set for the buffer memory is 2 Kbytes. The buffer memory state is controlled by this module, with a response sent automatically for a PING packet/NYET handshake.

(1) PING Packet Control when the Host Controller Function is Selected

This module automatically sends a PING packet in the OUT direction.

On receiving an ACK handshake in the initial state in which PING packet sending mode is set, this module sends an OUT packet as noted below. The module returns to PING packet sending mode when a NAK or NYET is received during an OUT transaction.

- 1. Sets OUT data sending mode.
- 2. Sends a PING packet.
- 3. Receives an ACK handshake.
- 4. Sends an OUT data packet.
- 5. Receives an ACK handshake.

(Repeats steps 4 and 5.)

- 6. Sends an OUT data packet.
- 7. Receives an NAK/NYET handshake.
- 8. Sends a PING packet.

This module is returned to PING packet sending mode by a power-on reset, receiving a NYET/NAK handshake, clearing the sequence toggle bits (SQCLR), and setting the buffer clear bit (ACLRM) in PIPEnCTR.

(2) NYET Handshake Control when the Function Controller Function is Selected

Table 24.28 lists the responses to received tokens during bulk transmission and control transmission. The USB 2.0 host/function module returns a NYET when an OUT token is received while there is only one packet of empty space in the buffer memory, during both bulk transmission and control transmission. However, an ACK is returned instead of a NYET even under the above conditions when a short packet is received.

Table 24.28 List of Responses to Received Tokens

[Legend]

RCV-BRDY1: When an OUT/PING token is received, there is space in the buffer memory for two or more packets.

RCV-BRDY2: When an OUT token is received, there is only enough space in the buffer memory for one packet.

RCV-NRDY: When a PING token is received, there is no space in the buffer memory.

TRN-BRDY: When an IN token is received, there is data to be sent in the buffer memory.

TRN-NRDY: When an IN token is received, there is no data to be sent in the buffer memory.

24.4.7 Interrupt Transfers (PIPE6 to PIPE9)

When the function controller function is selected, this module carries out interrupt transfers in accordance with the timing controlled by the host controller. In interrupt transfers, PING packets are ignored (no responses are sent), and the ACK, NAK, and STALL responses are carried out without an NYET handshake response being made.

When the host controller function is selected, this module can set the timing of issuing a token using the interval timer. This module issues an OUT token without issuing a PING token even in the OUT direction.

This module does not support high bandwidth transfers of interrupt transfers.

(1) Interval Counter during Interrupt Transfers when the Host Controller Function is Selected

(a) Outline of Operation

For interrupt transfers, intervals between transactions are set in the IITV bits in PIPEPERI. This controller issues an interrupt transfer token based on the specified intervals.

(b) Counter Initialization

This controller initializes the interval counter under the following conditions.

• Power-on reset:

The IITV bits are initialized.

• Buffer memory initialization using the ACLRM bit:

The IITV bits are not initialized but the count value is. Setting the ACLRM bit to 0 starts counting from the value set in the IITV bits.

Note that the interval counter is not initialized in the following case.

• USB bus reset, USB suspended:

The IITV bits are not initialized. Setting 1 to the UACT bit starts counting from the value before entering the USB bus reset state or USB suspended state.

(c) Operation when Transmission/Reception is Impossible at Token Issuance Timing

This module cannot issue tokens even at token issuance timing in the following cases. In such a case, this module attempts transactions at the subsequent interval.

- When the PID is set to NAK or STALL.
- When the buffer memory is full at the token sending timing in the receiving (IN) direction.
- When there is no data to be sent in the buffer memory at the token sending timing in the sending (OUT) direction.

24.4.8 Isochronous Transfers (PIPE1 and PIPE2)

This module has the following functions pertaining to isochronous transfers.

- Notification of isochronous transfer error information
- Interval counter (specified by the IITV bit)
- Isochronous IN transfer data setup control (IDLY function)
- Isochronous IN transfer buffer flush function (specified by the IFIS bit)

This module does not support the High Bandwidth transfers of isochronous transfers.

When operation as a host controller is selected and two pipes are used for isochronous transfer at the same time, observe the restrictions on packets stated in the USB 2.0 Specification, section 5.6.3, Isochronous Transfer Packet Size Constraints.

(1) Error Detection with Isochronous Transfers

This module has a function for detecting the error information noted below, so that when errors occur in isochronous transfers, software can control them. Tables 24.28 and 24.29 show the priority in which errors are confirmed and the interrupts that are generated.

(a) PID errors

If the PID of the packet being received is illegal

(b) CRC errors and bit stuffing errors

If an error occurs in the CRC of the packet being received, or the bit stuffing is illegal

(c) Maximum packet size exceeded

The maximum packet size exceeded the set value.

(d) Overrun and underrun errors

- When host controller function is selected:
	- ⎯ When using isochronous IN transfers (reception), the IN token was received but the buffer memory is not empty.
	- ⎯ When using isochronous OUT transfers (transmission), the OUT token was transmitted, but the data was not in the buffer memory.
- When function controller function is selected:
	- ⎯ When using isochronous IN transfers (transmission), the IN token was received but the data was not in the buffer memory.
	- ⎯ When using isochronous OUT transfers (reception), the OUT token was received, but the buffer memory was not empty.

(e) Interval errors

Detection

When function controller function is selected, the following cases are considered as interval errors:

- During an isochronous IN transfer, the token could not be received during the interval frame.
- During an isochronous OUT transfer, the OUT token was received during frames other than the interval frame.

Table 24.29 Error Detection when a Token is Received

Detection

Table 24.30 Error Detection when a Data Packet is Received

(2) DATA-PID

This module does not support High Bandwidth transfers. When the function controller function is selected, this module operates as follows in response to the received PID.

(a) IN direction

- DATA0: Sent as data packet PID
- DATA1: Not sent
- DATA2: Not sent
- mDATA: Not sent
- **(b) OUT direction (when using full-speed operation)**
- DATA0: Received normally as data packet PID
- DATA1: Received normally as data packet PID
- DATA2: Packets are ignored
- mDATA: Packets are ignored
- **(c) OUT direction (when using high-speed operation)**
- DATA0: Received normally as data packet PID
- DATA1: Received normally as data packet PID
- DATA2: Received normally as data packet PID
- mDATA: Received normally as data packet PID

(3) Interval Counter

(a) Outline of Operation

The isochronous interval can be set using the IITV bits in PIPEPERI. The interval counter enables the functions shown in table 24.31 when the function controller function is selected. When the host controller function is selected, this module generates the token issuance timing. When the host controller function is selected, the interval counter operation is the same as the interrupt transfer operation.

Table 24.31 Functions of the Interval Counter when the Function Controller Function is Selected

The interval count is carried out when an SOF is received or for interpolated SOFs, so the isochronism can be maintained even if an SOF is damaged. The frame interval that can be set is the 2^{IITV} frame or 2^{IITV} u frames.

(b) Interval Counter Initialization when the Function Controller Function is Selected

This module initializes the interval counter under the following conditions.

• Power-on reset

The IITV bit is initialized.

- Buffer memory initialization using the ACLRM bit The IITV bits are not initialized but the count value is.
- USB bus reset

After the interval counter has been initialized, the counter is started under the following conditions 1 or 2 when a packet has been transferred normally.

- 1. An SOF is received following transmission of data in response to an IN token, in the PID = BUF state.
- 2. An SOF is received after data following an OUT token is received in the PID = BUF state.

The interval counter is not initialized under the conditions noted below.

1. When the PID bit is set to NAK or STALL

The interval timer does not stop. This module attempts the transactions at the subsequent interval.

2. The USB bus reset or the USB is suspended

The IITV bit is not initialized. When the SOF has been received, the counter is restarted from the value prior to the reception of the SOF.

(4) Setup of Data to be Transmitted using Isochronous Transfer when the Function Controller Function is Selected

With isochronous data transmission using this module in function controller function, after data has been written to the buffer memory, a data packet can be sent with the next frame in which an SOF packet is detected. This function is called the isochronous transfer transmission data setup function, and it makes it possible to designate the frame from which transmission began.

If a double buffer is used for the buffer memory, transmission will be enabled for only one of the two buffers even after the writing of data to both buffers has been completed, that buffer memory being the one to which the data writing was completed first. For this reason, even if multiple IN tokens are received, the only buffer memory that can be sent is one packet's worth of data.

When an IN token is received, if the buffer memory is in the transmission enabled state, this module transmits the data. If the buffer memory is not in the transmission enabled state, however, a zero-length packet is sent and an underrun error occurs.

Figure 24.14 shows an example of transmission using the isochronous transfer transmission data setup function with this module, when $IITV = 0$ (every frame) has been set.

Figure 24.14 Example of Data Setup Function Operation

(5) Isochronous Transfer Transmission Buffer Flush when the Function Controller Function is Selected

If an SOF packet or a μSOF packet is received without receiving an IN token in the interval frame during isochronous data transmission, this module operates as if an IN token had been corrupted, and clears the buffer for which transmission is enabled, putting that buffer in the writing enabled state.

If a double buffer is being used and writing to both buffers has been completed, the buffer memory that was cleared is seen as the data having been sent at the same interval frame, and transmission is enabled for the buffer memory that is not discarded with SOF or μSOF packets reception.

The timing at which the operation of the buffer flush function varies depending on the value set for the IITV bit.

If $IITV = 0$

The buffer flush operation starts from the next frame after the pipe becomes valid.

• In any cases other than $IITV = 0$

The buffer flush operation is carried out subsequent to the first normal transaction.

Figure 24.15 shows an example of the buffer flush function of this module. When an unanticipated token is received prior to the interval frame, this module sends the written data or a zero-length packet according to the buffer state.

Figure 24.16 shows an example of this module generating an interval error. There are five types of interval errors, as shown below. The interval error is generated at the timing indicated by (1) in the figure, and the IN buffer flush function is activated.

If an interval error occurs during an IN transfers, the buffer flush function is activated; and if it occurs during an OUT transfer, an NRDY interrupt is generated.

The OVRN bit should be used to distinguish between NRDY interrupts such as received packet errors and overrun errors.

In response to tokens that are shaded in the figure, responses occur based on the buffer memory status.

- 1. IN direction:
	- ⎯ If the buffer is in the transmission enabled state, the data is transferred as a normal response.
	- ⎯ If the buffer is in the transmission disabled state, a zero-length packet is sent and an underrun error occurs.
- 2. OUT direction:
	- ⎯ If the buffer is in the reception enabled state, the data is received as a normal response.
	- ⎯ If the buffer is in the reception disabled state, the data is discarded and an overrun error occurs.

Figure 24.16 Example of an Interval Error Being Generated when IITV = 1

24.4.9 SOF Interpolation Function

When the function controller function is selected and if data could not be received at intervals of 1 ms (when using full-speed operation) or 125 μs (when using high-speed operation) because an SOF packet was corrupted or missing, this module interpolates the SOF. The SOF interpolation operation begins when $\text{USBE} = 1$, $\text{SCKE} = 1$ and an SOF packet is received. The interpolation function is initialized under the following conditions.

- Power-on reset
- USB bus reset
- Suspended state detected

Also, the SOF interpolation operates under the following specifications.

- 125 μs/1 ms conforms to the results of the reset handshake protocol.
- The interpolation function is not activated until an SOF packet is received.
- After the first SOF packet is received, either 125 μs or 1 ms is counted with an internal clock of 48 MHz, and interpolation is carried out.
- After the second and subsequent SOF packets are received, interpolation is carried out at the previous reception interval.
- Interpolation is not carried out in the suspended state or while a USB bus reset is being received. (With suspended transitions in high-speed operation, interpolation continues for 3 ms after the last packet is received.)

This module supports the following functions based on the SOF detection. These functions also operate normally with SOF interpolation, if the SOF packet was corrupted.

- Refreshing of the frame number and the micro-frame number
- SOFR interrupt and μ SOF lock
- Isochronous transfer interval count

If an SOF packet is missing when full-speed operation is being used, the FRNM bit in FRMNUM0 is not refreshed.

If a μ SOF packet is missing during high-speed operation, the UFRNM bit in FRMNUM1 is refreshed.

However, if a μ SOF packet for which the μ FRNM = 000 is missing, the FRNM bit is not refreshed. In this case, the FRNM bit is not refreshed even if successive μSOF packets other than μ FRNM = 000 are received normally.

24.5 Usage Notes

24.5.1 Procedure for Setting the USB Transceiver

When this module is to be used, start by making settings for the internal USB transceiver.

The method for the settings is described below. Figure 24.17 also gives an example of program code to implement the procedure.

- **(1)** Write 1 to bits UACS14 and UACS5 in USBAC characteristics switching register 0 (USBACSWR0).
- **(2)** Write 1 to the UACS26 bit in USBAC characteristics switching register 1 (USBACSWR1).

Initialization routine (1) Set 1 to UACS14 and UACS5. MOVI20 #H'FFFF00C0, R0 MOV.W #H'4020, R1 MOV.W R1, @R0 (2) Set 1 to UACS26. MOVI20 #H'FFFF00C2, R0 MOV.W #H'0400, R1 MOV.W R1, @R0

Figure 24.17 Procedure for Setting the USB Transceiver

Section 25 AT Attachment Packet Interface (ATAPI)

The ATAPI interface device provides both the ATA and ATAPI physical interfaces. This device also supports both the ATA task and ATAPI packet commands.

25.1 Features

- Supporting primary channel
- Supporting master/slave
- Supporting 3.3V I/O interface
- Supporting PIO modes 0 to 4, multiword DMA modes 0 to 2, and ultra DMA modes 0 to 2

Figure 25.1 Block Diagram of ATAPI

25.2 Input/Output Pins

Table 25.1 Pin Configuration

Note: * The ATAPI interface treats the interrupt signal from the ATAPI device as a leveltriggered input.

25.3 Register Description

The following register set is allocated in the on-chip peripheral module space of this LSI device.

25.3.1 ATAPI Interface Registers

Table 25.2 ATA Task File Register Map

(These resisters are allocated to the ATAPI or ATA device, but not to this module.)

Notes: 1. The CPU must access these registers in longword (32-bit) units. Byte and word accesses are prohibited.

2. Bits 15 to 0 of the data bus are used.

3. Bits 7 to 0 of the data bus are used.

Table 25.3 ATAPI Packet Command Task File Register Map

(These resisters are allocated to the ATAPI or ATA device, but not to this module.)

Notes: 1. The CPU must access these registers in longword (32-bit) units. Byte and word accesses are prohibited.

2. Bits 15 to 0 of the data bus are used.

3. Bits 7 to 0 of the data bus are used.

Table 25.4 ATAPI Interface Control Register Map

(These resisters are allocated to this module.)

Note: * These registers must be accessed in longword (32-bit) units. Byte and word accesses are prohibited.

25.3.2 ATAPI Interface Control Register Map

All control/status registers are active high.

(1) ATAPI control register (ATAPI_CONTROL)

(2) ATAPI status register (ATAPI_STATUS)

(3) Interrupt enable (ATAPI_INT_ENABLE)

Note: Writing 1 to each bit enables the interrupt signal corresponding to each bit in the ATAPI status register.

(4) PIO timing register (ATAPI_PIO_TIMING)

Before accessing an ATAPI device, set the number of machine cycles in the following bits in this register.

A machine cycle is equal to an enhanced bus clock cycle. Its frequency is the same as of the bus clock.

Note: The prefix pS pertains to slaves, and pM, to the master.

Figure 25.2 PIO Timing Register

• PIO timing register value table (master/slave)

(5) Multiword DMA timing register (ATAPI_MULTI_TIMING)

Before accessing an ATAPI device, set the number of machine cycles in the following bits in this register.

Note: The prefix mS pertains to slaves, and mM, to the master.

Figure 25.3 Multiword DMA Timing Register

• Multiword DMA timing register value table

(6) Ultra DMA timing register (ATAPI_ULTRA_TIMING)

Before accessing an ATAPI device, set the number of machine cycles in the following bits in this register.

Note: The prefix uS pertains to slaves, and uM, to the master.

Figure 25.4 Ultra DMA Timing Register

• Ultra DMA timing register value table

(7) DMA start address register (ATAPI_DMA_START_ADR)

Notes: 1. This address will not change even after DMA becomes active; it will retain its setting.

2. The access destination is in SDRAM.

(8) DMA transfer count register (ATAPI_DMA_TRANS_CNT)

Note: This count value will not change even after DMA becomes active; it will retain its setting.

(9) ATAPI control 2 register (ATAPI_CONTROL2)

(10) ATAPI signal status register (ATAPI_SIG_ST)

(11) Byte swap register (ATAPI_BYTE_SWAP)

25.4 Operation

The ATAPI interface supports the primary channel as a host. It also supports a master/slave configuration as stipulated in the ATAPI interface specification. The FIFO read/write buffer of the ATAPI interface is designed to transfer data at up to 16 Mbytes/s in multiword DMA mode and up to 33 Mbytes/s in ultra DMA mode. The ATAPI interface supports the 3.3V I/O interface.

The ATA task file register and ATAPI packet command file register are allocated in the on-chip peripheral module space of this LSI device. Therefore, accessing these registers from the LSI device can be made by addressing the on-chip register of the DVDROM drive or the like with the DCS1, DCS0, and DSA2 to DSA0 pins.

25.4.1 Data Transfer Modes

The ATAPI interface control register supports PIO transfer modes, multiword DMA transfer modes, and ultra DMA transfer modes. The ATAPI interface control register is used to initiate each transfer mode and set the ATAPI interface timing that varies from one transfer mode to another.

The transfer modes supported with the ATAPI interface include PIO modes 0 to 4, multiword DMA modes 0 to 2, and ultra DMA modes 0 to 2.

The enhanced bus is used for multiword and ultra DMA data transfers, and the peripheral bus, for PIO transfers.

Table 25.5 Data Transfer Modes

Note: * The CPU accesses the ATA device in PIO mode. In enhanced bus DMA transfers, data is transferred between the ATAPI device and memory.

25.4.2 Initialization Procedure

(1) Setting the interface enable bit

Set the IFEN bit in ATAPI control 2 register to 1.

(2) Setting the timing register

Write an appropriate value to the following registers.

See the respective descriptions for what value is appropriate for each register.

- PIO timing register
- Multiword DMA timing register
- Ultra DMA timing register

25.4.3 PIO Transfer Mode Operation Procedure

Figure 25.5 PIO Transfer Mode Operation Procedure

25.4.4 Multiword DMA Transfer Mode Operation Procedure

Figure 25.6 Transfer to and from Memory via Enhanced Bus by Polling

(2) Transfer to and from memory via enhanced bus by interrupt

Figure 25.7 Transfer to and from Memory via Enhanced Bus by Interrupt

25.4.5 Ultra DMA Transfer Mode Operation Procedure

(1) Transfer to and from memory via enhanced bus by polling

Figure 25.8 Transfer to and from Memory via Enhanced Bus by Polling

(2) Transfer to and from memory via enhanced bus by interrupt

Figure 25.9 Transfer to and from Memory via Enhanced Bus by Interrupt

25.4.6 ATAPI Device Hardware Reset Procedure

Figure 25.10 ATAPI Device Hardware Reset Procedure

25.5 DIRECTION Pin

The DIRECTION pin outputs a low level while data is being written to an external ATA device by this LSI.

To be specific, the DIRECTION pin goes low at:

- PIO data transfer to an ATA device
- Multiword DMA transfer (data-out)
- Ultra DMA data-in CRC transmission
- Ultra DMA transfer data-out

See section 33, Electrical Characteristics for descriptions about the timing of each transfer mode.

25.6 Usage Note

When using the ATAPI module, set the frequencies of the bus clock and the peripheral clock so that their ratio is 2:1.

Section 26 2D Graphics Engine (2DG)

The 2DG engine (below referred to as the 2DG) is provided as a two-dimensional graphics accelerator that has the following functions:

alpha blending of two areas specified as rectangles followed by resizing and the output of enlarged or reduced rectangular areas: and

the resizing of externally supplied moving pictures to the size of the display panel, followed by composition with the output graphics plane, and outputting to VIDEO OUT (D/A converter) at a constant rate.

The graphics data is transferred between the SDRAM and the 2DG at high speed by the DMAC under CPU control. All of the required data, including the source and destination addresses for the SDRAM area and the source and destination addresses for the 2DG are issued by the DMAC. The 2DG blit operations are the specified processing for the data supplied to the source buffers and the output to the destination buffer.

26.1 Features

- Planes: Major examples (characters, graphics plane, and output plane, totally two planes)
- Acceleration: blitting with two inputs and one output, filling, bit-blitting, chromakey, logical operations, color gradation handling, variable blending operations
- Resizing,
	- ⎯ Blitter: Bilinear and nearest-neighbor method are independently selectable for the horizontal and vertical directions (conversion ratios from $1/2$ to 2) Selection of pre-filtering (as a measure against the Moire effect) on or off.
	- ⎯ Output block: Bi-cubic algorithm in the horizontal direction (conversion ratios: from 1/3 to 1)
- Moving picture input: BT656 format (NTSC/PAL system) input (requiring both the VIHSYNC and VIVSYNC signal inputs.)
- Superimposition on moving pictures: Alpha blending of the graphics plane and moving picture, followed by constant-rate output in RGB666 format
- Input pixel formats for blit blocks: α RGB444 (16 bits), α RGB555 (16 bits), or α (4 bits)
- Output pixel formats for blit blocks: αRGB444 (16 bits), or αRGB555 (16 bits)
- Final picture resolution: WOVGA (480 x 234) or OVGA (320 x 240)
	- Capacity of input/output buffers for graphics (each is in a double-buffer configuration)
		- Input buffer E for the output block: 16bits x 512 words x 2 planes
		- Input buffers A and B for the blitter: 16 bits x 64 words x 2 planes, each
		- Output buffer C for the blitter: 16 bits x 256 words x 2 planes

Figure 26.1 is a block diagram of the 2DG.

Figure 26.1 Block Diagram of the 2DG

26.2 Input/Output Pins

Table 26.1 Pin Configuration

26.3 Register Descriptions

The 2DG has the following registers. During operation in synchronization with the VSYNC signal, register values are applied to the 2DG when the VSYNC signals are low pulse. However, the read/write to the applicable register is irrelevant to the VSYNC synchronization.

Table 26.2 Configuration of Registers

26.3.1 Blit Function Setting Register for Graphics (GR_BLTPLY)

Register GR_BLTPLY is used to enable blitting.

Setting values of the SB_STEN and SA_STEN bits for blit operations are shown below.

- Blit operations only for source SA alone are not possible. If only a single source signal is supplied, use source SB.
- If the value 0 is written to during blitting, the operation is ended forcibly.
- The value "1" should never be written to these bits when neither buffer SA nor buffer SB is empty. Whether the buffers are empty or not is checked by the state of the GR_DOSTAT register.
- When both of the bits SB_STEN and SA_STEN are set to 1, blit operation only proceeds when the same amounts of data are to be transferred from buffers SA and SB (for details, see section 26.4.3 (2), Summary of Operations between the Blitter and External Memory.)

26.3.2 Mixing Function Setting Register for Graphics (GR_MIXPLY)

The register GR_MIXPLY specifies the display of externally input pictures and graphics. The register value is applied to the 2DG in synchronization with the VSYNC signal.

Setting values of the EXTEN and OUTEN bits for the graphics display are shown below.

26.3.3 Operation Status Register for Graphics (GR_DOSTAT)

The register GR_DOSTAT indicates the operating status of the 2DG.

• Each SEHF_STAT bit changes from 0 to 1 when the corresponding buffer SE1 or SE2 is full or when the amount of pixel data in the SE buffer coincides with the specified number of pixels.

Furthermore, the SEHF_STAT bit changes from 1 to 0 on completion of reading the data in the corresponding half of SE buffer (or the data remaining herein).

• Each DCHF_STAT bit changes from 0 to 1 when the corresponding buffer DC1 or DC2 is full or when the amount of pixel data in the DC buffer coincides with the specified number of pixels.

Furthermore, the DCHF_STAT bit changes from 1 to 0 on completion of DMA transfer of the data in the corresponding half of DC buffer (or the data remaining herein).

Each SBHF_STAT bit changes from 0 to 1 when the corresponding buffer SB1 or SB2 is full or when the amount of pixel data in the SB buffer coincides with the specified number of pixels.

Furthermore, the SBHF_STAT bit changes from 1 to 0 on completion of reading the data in the corresponding half of SB buffer (or the data remaining herein).

Each SAHF_STAT bit changes from 0 to 1 when the corresponding buffer SA1 or SA2 is full or when the amount of pixel data in the SA buffer coincides with the specified number of pixels.

Furthermore, the SAHF_STAT bit changes from 1 to 0 on completion of reading the data in the corresponding half of SA buffer (or the data remaining herein).

- "Coincide" in the above list means that the width setting in register GR_SABSET, GR_DCSET, or MGR_SESET matches the number of data transferred to or from the relevant buffer or that the corresponding half of the double-buffer is full.
- If an abnormal state arises, such as stopping of graphics operations before they are completed, use this register identify the cause of the problem.
	- If the SB REND or SA REND bit is being held at 1, check the settings of registers related to buffer DC (for example, the value of GR_DCSET).
	- If the SB_REND or SA_REND bit is being held at 0, and one buffer of the SB and SA double-buffers remains full, check the settings of registers related to buffers SB and SA (for example, the value of GR_SABSET).
	- If the blitter is reactivated, write 1 to bits SB STEN and SA STEN in register GR_BLTPLY.
	- If buffer SE is empty, write 0 to bits OUTEN and EXTEN in register GR_MIXPLY. Output can then be restarted.

26.3.4 Interrupt Status Register for Graphics (GR_IRSTAT)

The register GR_IRSTAT indicates the interrupt state of the 2DG. When an interrupt event corresponding to the IRQ_DEMPT, IRQ_ASHFUL, IRQ_DHFUL, or IRQ_SHFUL bit in this register occurs, the given bit will be set as long as the event has not been masked by the MSK_DEMPT, MSK_ASHFUL, MSK_DHFUL, or MSK_SHFUL bit in the GR_INTMSK register. For the other bits in register GR_IRSTAT, the bit will be set to "1" when the corresponding event occurs, regardless of the setting in the interrupt mask control register for graphics (GR_INTMSK). For details on interrupts, see section 26.4.5, Interrupts.

Note: The INT_UDFL and INT_VSYC bits may be set even when the output block has not been started. So, be sure to clear the INT_UDFL and INT_VSYC bits in the GR_INTDIS register before starting up the output block.

26.3.5 Interrupt Mask Control Register for Graphics (GR_INTMSK)

The register GR_INTMSK masks 2DG interrupts. When an interrupt event occurs, the interrupt status register for graphics (GR_IRSTAT) will be set even if the corresponding interrupt is not enabled (masked). For details on interrupts, see section 26.4.5, Interrupts.

• Note that the MGR_MIXHTMG, MGR_MIXHS, and MGR_MIXVTMG registers should be set according to the display panel to be used before the interrupts masked by the MSK_UDFL and MSK_FILD bits are cancelled.

26.3.6 Interrupt Reset Control Register for Graphics (GR_INTDIS)

The register GR_INTDIS cancels 2DG interrupts. Interrupt signals are deasserted by writing 1 to the corresponding bits in this register. Furthermore, the IRQ_DEMPT, IRQ_ASHFUL, IRQ_DHFUL, IRQ_SHFUL, INT_VSYC, INT_UDFL, INT_ FILD, and INT_GR bits in GR_IRSTAT are cleared by writing 1 to the corresponding bits in this register. Note, however, that the INT_DEMP, INT_ASHFUL, INT_DHFUL, and INT_SHFUL bits in GR_IRSTAT are not cleared even if 1 is written to the corresponding bits in this register (the hardware automatically handles clearing of these bits). When a 1 is written to any of these bits, the hardware automatically sets the bit to its initial value. For details on interrupts, see section 26.4.5, Interrupts.

26.3.7 DMAC-Request Control Register for Graphics (GR_DMAC)

The register GR_DMAC sets DMA transfer and CPU transfer control methods for SA, SB, DC and SE buffers. Note that the settings for this register should be the same as the corresponding settings in the DMAC.

Note: When the H1PHS_INTGR bits in GR_HSPHAS are odd (H1PS_INTGR [0] = 1), be sure to set the SZEL1 bit to 0 (16 bits).

26.3.8 Source A&B Read-In-Area Setting Register for Blitter (GR_SABSET)

The register GR_SABSET sets the SA and SB areas. In a DMA transfer, the total number of pixels to be transferred from the external memory space is obtained by $SSWIDH \times SSHIGH$.

- When both SA and SB areas are used, it is assumed that SB (foreground) > SA (background). For example, if two-plane blending is performed while SB and SA are assumed to be a character and background, respectively, blending is performed so that the character will always be foreground.
- If the pixel format for SB area is set to α (4 bits), the SSWIDH bits in GR_SABSET should be set as follows.

In 16-bit access: For the SSWIDH bits in GR_SABSET, the minimum number of pixels to be transferred and the pixel transfer unit should be set to 4 and $4 \times n$ (n: arbitrary integer), respectively. Setting 2 or 3 pixels to these bits is prohibited.

- In 32-bit access: For the SSWIDH bits in GR_SABSET, the minimum number of pixels to be transferred and the pixel transfer unit should be set to 8 and $8 \times n$ (n: arbitrary integer), respectively.
- If the pixel format is set to α RGB444 or α RGB555, only 16-bit access is enabled when the total number of pixels to be transferred (SSWIDH \times SSHIGH) is odd; both 32-bit and 16-bit accesses are enabled when the value of $SSWIDH \times SSHIGH$ is even. If duplicate-lines are set for enlargement resizing, either 32-bit access or 16-bit access should be selected according to the number of pixels on a line (SSWIDH), not according to the total number of pixels to be transferred (SSWIDH × SSHIGH).

26.3.9 Destination C Write Area Setting Register for Blitter (GR_DCSET)

The register GR_DCSET sets the DC area. In a DMA transfer, the total number of pixels to be transferred to the external memory space is obtained by D CWIDH \times DCHIGH.

• If the pixel format is set to α RGB444 or α RGB555, only 16-bit access is enabled when the total number of pixels to be transferred (DCWIDH × DCHIGH) is odd; both 32-bit and 16-bit accesses are enabled when the value of $DCWIDH \times DCHIGH$ is even.

26.3.10 Source E Read-In Area Setting Register for Output Block (MGR_SESET)

The register MGR_SESET sets the SE area. In a DMA transfer, the total number of pixels to be transferred from the external memory space is obtained by $SEWIDH \times SEHIGH$. The register value is applied to the 2DG in synchronization with the VSYNC signal.

- Only 16-bit access is enabled when the total number of pixels to be transferred (SEWIDH \times SEHIGH) is odd; both 32-bit and 16-bit accesses are enabled when the value of (SEWIDH \times SEHIGH) is even.
- The settings of this register should be the same as those of the MGR_MIXHS and MGR_MIXVS registers.

SEWIDH bits = VLDPH bits in MGR_MIXHS SEHIGH bits = VLDPV bits in MGR_MIXVS

- So that the display does not break up, make settings such that DMA transfer to the SE buffer is efficient. Some recommendations follow as examples.
	- ⎯ SEWIDH bits: Set these bits for a number of pixels that is a multiple of eight (320, 480, and so on).
	- ⎯ Output block DMA transfer size bit: Set this to one (32 bits).
	- Amount transferred per operand: Set this to a large value.
	- DMAC transfer mode: Select pipeline transfer.

26.3.11 Pixel Format Setting Register for Graphics (GR_PIXLFMT)

The register GR_PIXLFMT sets the pixel formats which are used for the input and output buffers. The SE_FMT bit is applied to the 2DG in synchronization with the VSYNC signal. Bits other than the SE_FMT bit are applied to the 2DG immediately.

- If the SB_FMT bits are set to 10, the SSWIDH bits in GR_SABSET should be set as follows: In 16-bit access: For the SSWIDH bits in GR_SABSET, the minimum number of pixels to be transferred and the pixel transfer unit should be set to 4 and $4 \times n$ (n: arbitrary integer), respectively. Setting 2 or 3 pixels to these bits is prohibited.
	- In 32-bit access: For the SSWIDH bits in GR_SABSET, the minimum number of pixels to be transferred and the pixel transfer unit should be set to 8 and $8 \times n$ (n: arbitrary integer), respectively.

26.3.12 Operation Mode Setting Register for Blitter (GR_BLTMODE)

The register GR_BLTMODE sets the blitter operation mode.

- The CRKEY bits are valid only when the SBSEL bits are set to 01 and the BTYPE bits are set to 10.
- When the chromakey is enabled (when the SBSEL bits are set to 01 and the BTYPE bits are set to 10), operation is the same as when the GCOLR bit in GR_BRD1CNT is set to 1. However, the value of the GCOLR bit does not change.
- The LGTYPE bit is valid only when the SBSEL bits are set to 10.
- The SBSEL bits can be set to 00, 01, or 10 when the filling operation is selected (BTYPE = 10), and they can be set to 00, 10, or 11 when the blitting operation is selected (BTYPE = 00).
- If the BTYPE bits are set to 00, setting chromakey is prohibited.

26.3.13 Resize Display Setting Register for Graphics (GR_RISZSET)

The register GR_RISZSET sets the resizing function.

Bit: R/W: Bit: -------------- --- -- ------ 1 1

• When the resizing function is used for the blitter, the EDGE bits should be set as follows: The EDGE bits should be set to 11 if full resizing which resizes the entire source area is performed.

The EDGE bits should be set according to the bit description above if partial resizing which resizes a part of the source area is performed.

• The interference stripes can be reduced by setting the PREON bit to 1 and performing prefiltering.

26.3.14 Resize Mode Select Register for Blitter (GR_RISZMOD)

The register GR_RISZMOD sets the resizing function for the blitter.

Note: The bits in GR_RISZMOD must be set when resizing is performed. These bits need not to be set when resizing is not performed.

26.3.15 Resize Delta Setting Register for Blitter (GR_DELT)

The register GR_DELT sets delta computation results for resizing on the blitter.

Initial

Note: This register must be set before resizing. If resizing is not performed, this register need not be set.

• If one line of data to be transferred by the CPU is 65 pixels or more when an enlargement resizing in the vertical direction is performed, there may be a case that the pixels at the same lines should be transferred twice by the CPU. For details, see section 26.4.3 (6), Duplicate-line Setting of Enlargement Resizing.

26.3.16 Resize Horizontal Starting Phase Register for Blitter (GR_HSPHAS)

The register GR_HSPHAS sets results of the starting position phase computation in the horizontal direction for the blitter resizing.

Notes: 1. This register must be set before resizing. All bits should be cleared to 0 when resizing is not performed.

2. When the H1PHS INTGR bits are odd (H1PHS INTGR $[0] = 1$), be sure to set the SZEL1 bit in GR_DMAC to 0 (16 bits).

26.3.17 Resize Vertical Starting Phase Register for Blitter (GR_VSPHAS)

The register GR_VSPHAS sets results of the starting position phase computation in the vertical direction for the blitter resizing.

Note: This register must be set before resizing. All bits should be cleared to 0 when resizing is not performed.

26.3.18 Resize Horizontal Delta Setting Register for Output Block (MGR_HDELT)

The register MGR_HDELT sets the delta computation results in the horizontal direction for the output block resizing. The register value is applied to the output block in synchronization with the VSYNC signal.

26.3.19 Resize Horizontal Starting Phase Register for Output Block (MGR_HPHAS)

The register MGR_HPHAS sets results of the starting position phase computation in the horizontal direction for the output block resizing. The register value is applied to the output block in synchronization with the VSYNC signal.

26.3.20 Logical Operation Input Data Register for Blitter (GR_LGDAT)

The register GR_LGDAT sets data for logical operation to be performed on the blitter. Since logical operation is performed after format conversion, data should be specified as 5 bits of LGDAT_R (R data), 5 bits of LGDAT_G (G data) and 5 bits of LGDAT_B (B data).

 GR _BLTMODE = 10)

26.3.21 Chromakey Target Color Data Register for Blitter (GR_DETCOL)

The register GR_DETCOL sets the target color for chromakey (α data is not set). Since chromakey processing is performed after format conversion, data should be specified as 5 bits of DETC_R (R data), 5 bits of DETC_G (G data) and 5 bits of DETC_B (B data).

- This register is valid only when the chromakey function has been selected (SBSEL in GR BLTMODE = 01)
- This register value is used differently depending on the CRKEY bits in GR_BLTMODE:

When $CRKEY = 01$: The target color specified in this register is replaced with the replacement color specified in GR_BRDCOL for the SB input data.

When $C R K EY = 10$: The target color specified in this register is blended with the replacement color specified in GR_BRDCOL for the SB input data.

26.3.22 Replacement Color Data Register for Blitter Blending (GR_BRDCOL)

The register GR_BRDCOL is used for chromakey processing and color gradation processing. In chromakey processing, the color specified in this register replaces or is blended with the target color. In color gradation processing, the color specified in this register is used as the replacement color. Since chromakey processing and color gradation processing are performed after format conversion, the data should be specified as 5 bits of BRDC_R (R data), 5 bits of BRDC_G (G data) and 5 bits of BRDC_B (B data).

26.3.23 Blend 1 Control Register for Blitter (GR_BRD1CNT)

The register GR_BRD1CNT specifies settings for blending on the blitter. For details, see section 26.4.3 (3) (a), Blending.

Note: For symbols above, see figure 26.23.

Table 26.5 FBFA Bit Details

- When the fill operation has been selected (BTYPE bits in BR_BLTMODE = 10), the GCOLR bit should be set to 1 and the FBFA bits are arbitrary (basically cleared to 00 because fill operation with blending is performed.)
- The AFTER_A bits are valid only when $GALFA = 1$.
- FBFA should be set to 01 when chromakey processing has been selected.
- GCOLR should be set to 1 to perform fill operation (BTYPE in GR_BLTMODE = 10).

26.3.24 Mixing Mode Setting Register for Output Block (MGR_MIXMODE)

The register MGR_MIXMODE sets mixing mode for the output block. The register value is applied to the output block in synchronization with the VSYNC signal. For details, see section 26.4.4 (5), Blending in Output Block.

Notes: 1. For symbols, see figure 26.55.

- 2. For details on the CBCR bit, see figure 26.15.
- 3. When a system with the low quality external moving picture input has been selected and displays only the graphic image, the MVON bit should be cleared to 0. By this setting, synchronization errors can be prevented by using internal SYNC signal, and this improves the quality of display. When MVON is set to 1, VIHSYNC or VIVSYNC which is synchronous with external moving picture input should be input.
- 4. When α RGB555 is selected for the pixel format (SE_FMT in GR_PIXFMT = 1) and α data of the pixel data is 1, four bits of α data are replaced with CHG_A.
- 5. MVON should be set to 1 before EXTEN in GR_MIXPLY is set to 1. When MVON = 0, setting EXTEN in GR_MIXPLY to 1 is prohibited.
- 6. Follow the procedure below when using the MVON bit to switch VSYNC between external and internal synchronization.
	- (1) Set the GR_MIXPLY register to disable display.
	- (2) Change the MVON bit to change the synchronization of VSYNC.
	- (3) At least twice, check that VSYNC is being generated after the change.
	- (4) Set the GR_MIXPLY register to enable display.

Table 26.6 FCFD Bit Details

Note: The FCFD bits are set automatically by the hardware according to the GR_MIXPLY settings as follows, however, the FCFD bit is not changed: When only the externally supplied moving picture is selected (OUTEN = 0 , EXTEN = 1):

 $FCFD = 011$

When only the graphic image is selected (OUTEN $= 1$, EXTEN $= 0$): FCFD $= 010$ When display is prohibited (OUTEN = 0 , EXTEN = 0): FCFD = 100

26.3.25 Panel-Output Horizontal Timing Setting Register for Output-Block (MGR_MIXHTMG)

The register MGR_MIXHTMG sets the timing of signal output to the panel in the horizontal direction. The register value is applied in synchronization with the VSYNC signal. For details, see section 26.4.1 (5), Setting of Panel Output.

26.3.26 Panel-Output Mixing Horizontal Valid Area Setting Register for Output Block (MGR_MIXHS)

The register MGR_MIXHS sets a valid area for signal output to the panel in the horizontal direction. The register value is applied in synchronization with the VSYNC signal. For details, see section 26.4.1 (5), Setting of Panel Output.

MGR_SESET register) must be the same

The VLDPH bits are equivalent to the SEWIDH bits in MGR_SESET.

26.3.27 Panel-Output Vertical Timing Setting Register for Output-Block (MGR_MIXVTMG)

The register MGR_MIXVTMG sets the vertical timing of signal output to the panel. The register value is applied in synchronization with the VSYNC signal. For details, see section 26.4.1 (5), Setting of Panel Output.

26.3.28 Panel-Output Mixing Vertical Valid Area Setting Register for Output Block (MGR_MIXVS)

The register MGR_MIXVS sets the vertical area for signal output to the panel. The register value is applied in synchronization with the VSYNC signal. For details, see section 26.4.1 (5), Setting of Panel Output.

Note: The settings in this register and in the source E read-in area for the output block (the MGR_SESET register) must be the same

The VLDPV bits are equivalent to the SEHIGH bits in MGR_SESET.

26.3.29 Graphics Block Output SYNC Position Setting Register (GR_VSDLY)

This register specifies the position of the output VSYNC signal. The vertical direction for the moving pictures can vary with the monitor in use. In situations where this is the case, this register can be adjusted to eliminate fluctuations in the vertical direction.

Notes: 1. When the external image is NTSC (when the NTSC bit in the MGR_MIXMODE register has been set to 0), use the initial value (H'160). When PAL is in use, after setting this register to H'100, set the NTSC bit in the MGR_MIXMODE register to 1.

2. Setting all bits of this register to 0 is prohibited.

26.3.30 Video DAC Timing Setting Register (VDAC_TMC)

The setting in this register is for the timing of output to a monitor.

Note: When writing to this register, stop the 2DG module beforehand.

26.4 Operation

Figure 26.2 Block Diagram of the Blitter

Figure 26.3 Block Diagram of the Output Block

26.4.1 Input and Output Operations

(1) Data Bit Map for the Pixel Format

Data in the α RGB444 (16 bits), α RGB555 (16 bits), and α (4 bits) pixel formats are applied to the inputs and output of the 2DG. The formats are shown below.

\bullet α RGB444

α RGB555

• α (4 bits)

Note: Transfer the α data from memory in 4 blocks unit.

Transfer of the data from α 2 or α 3 block is prohibited.

(2) Assignment of Pixels to the Memory Space

Input and output data are mapped on the memory space as shown in figure 26.4. This example shows the data mapping in the case where the selected format is α RGB444 (16 bits).

(3) Relations between Line Pitches and Memory Plane

The target display panels for the 2DG are QVGA $(320 \text{ pixels} \times 240 \text{ lines})$ and WQVGA (480 s) pixels \times 234 lines). The line pitch that determines the relations between memory space and each work screen (such as the character plane and graphics plane) of the SDRAM must be placed on the 64-byte boundaries. Thus, the start addresses of each planes become the following.

XXXX_XX $[4n]$ $[0]$ (H) $(X =$ arbitrary number, n = integer)

So, the starting address must be one of these: XXXX_XX00, XXXX_XX40, XXXX_XX80, or XXXX XXC0.

Figure 26.5 shows an example of the relations between the arrangement of planes A, B, and C, (WQVGA size) and the pitch of display lines in memory for the α RGB444 (16-bits) pixel format.

Figure 26.5 Relations between the Planes and Display-Line Pitch in Memory for the α**RGB444 (16 Bits) Pixel Format**

As shown in figure 26.5, the start address for each of planes A, B, and C is a 64-byte boundary. The readout area from each planes can be set to desired start address and area. However, the access unit is restricted by the byte number per pixel. For example, when the α RGB444 format has been set, access is in 2-byte units and only access to even addresses is allowed, access to odd addresses is prohibited.

(4) Input and Output Buffers

Since transfer to and from buffers SA, SB, DC, and SE in the 2DG must be handled by the DMAC under CPU control, these buffers are mapped on memory spaces (SRAM spaces) of the CPU. Table 26.7 shows the address map of the input and output buffers.

Table 26.7 Address Map of the Input and Output Buffers

Since the 2DG has fixed-size input and output buffers, an image processing is performed with repeated DMA data transfer from these areas. Table 26.8 shows the specifications of the buffers. Each buffer is configured as the double buffer structure.

Table 26.8 Specifications of Input and Output Buffers

(5) Setting of Panel Output

Figure 26.6 shows the relations between the sync signals and the registers setting for display-panel output. These panel-output settings are made in the following registers.

- MGR_MIXHTMG register: WPH bits and PDPH bits
- MGR_MIXHS register: ALLPH bits and VLDPH bits
- MGR_MIXVTMG register: WPV bits and PDPV bits
- MGR_MIXVS register: ALLPV bits and VLDPV bits

Figure 26.6 Relationship between Panel Output and Register Settings

(6) Relations between the Sync Signals for the Output Block and the Individual Clock Signals

(a) With an externally mounted video-decoder LSI

In cases where a moving-picture input is being supplied to the system, set the MVON bit in the MGR_MIXMODE register to 1. Then, the free-running HSYNC_dck (internal signal) is generated from DCLKIN and VSYNC_dck (internal signal) is created from the VIVSYNC signal. After that, the CSYNC signal is composed from the two signals, i.e. HSYNC_dck and VSYNC_dck.

When the system includes an external video decoder LSI but this is not supplying the movingpicture input with a signal, set the MVON bit in MGR_MIXMODE register to 0. In this case, the free-running HSYNC_dck generated from DCLKIN is counted and the hardware itself produces the VSYNC_dck: the CSYNC signal is again composed from the HSYNC_dck and the VSYNC_dck.

(b) Without externally mounted video decoder LSI

In cases where there is no external video-decoder LSI to supply the system with a moving picture input, set the MVON bit in MGR_MIXMODE to 0. In this case, since only the free-running HSYNC_dck generated from DCLKIN is available, the hardware automatically produces the VSYNC_dck signal by counting cycles of HSYNC_dck, and then composes the CSYNC signal from HSYNC_dck and VSYNC_dck.

References:

- The image data from the output block, (the image data composed from the moving-picture and graphics data), and the CSYNC signal are output in synchronization with rising edges of the externally input DCLKIN signal.
- Timing with which the graphics data are read from the SE buffer is controlled by the HSYNC_dck and VSYNC_dck signals, and the MGR_SESET register (timing is not controlled by the VICLK system).
- For externally input moving pictures, the valid number of pixels horizontally is controlled by the VICLKENB signal and the valid number of lines is set by the VLDPV bits in the MGR_MIXVS register.
- Externally input moving pictures specified valid area are resized and then written to the data buffer. Reading of the data from the buffer is controlled by signals HSYNC_dck, and VSYNC_dck, and register MGR_MIXxx.

Figure 26.7 Relations between Externally Input Sync Signal and VICLK

Figure 26.8 Relations between Internal Generated Sync Signals and DCLKIN

VIHSYNC HS_ext VICLKENB Moving picture data Moving picture data (After resizing) Enable signal for moving picture Writing resized data to the data buffer HSYNC_dck (Internal) Reading moving picture data from the data buffer Graphic data Synthesized image data Gr: Abbreviation of graphics 720×74 ns = 53.3us Number of DCLKIN close to HS_ext WPH bits This period varies because of asynchronous (within 1 line). PDPH bits VLDPH bits When moving picture is not input, the timing chart for the VICLK system shown in (1) is not applicable. Only graphic data is output in synchronization with HSYNC_dck. (1) VICLK = 27 MHz synchronous system 480×125 ns = 60us 720 pixels **Telecom** 720 pixels **720 pixels** 720 pixels 480 pixels (moving picture) 480 pixels (moving picture) 480 pixels (moving picture) 480 pixels (moving picture) 480 pixels (moving picture 480 pixels (Gr) 480 pixels (Gr) 480 pixels (Gr + moving picture) 480 pixels (Gr + moving picture) 480 pixels (moving picture) 480 pixels (Gr) 480 pixels (Gr + moving picture) 480 pixels (moving picture) **480 pixels** (moving picture) 480 pixels (moving picture) 480 pixels (moving picture (2) DCLKIN = approx. 8 MHz synchronous system (WQVGA)

(7) Timing for Composition of Moving Pictures and Graphics by the Output Block

Figure 26.12 Timing in the Horizontal Direction (with Moving Pictures Supplied)

Figure 26.13 Timing in the Vertical Direction (with Moving Pictures Supplied)

Figure 26.14 2DG Output Signals and DCLKIN

(8) RGB Conversion of External Input Moving Picture

The timing relations between the output data and clock signal from the external video decoder, i.e. the VICLK, VICLKENB, and VIDATA[7:0] are shown in figure 26.15.

Since the data is in YCbCr422 pixel format and the output data is to be composed with graphics (RGB) data, conversion is initially from YCbCr422 to YCbCr444, and then from YCbCr to RGB.

Figure 26.15 Timing of Conversion from YCbCr422 to YCbCr444

The formulae for converting from YCbCr to RGB are given below.

 $R = 1.164(Y-16) + 1.596(Cr - 128)$

 $G = 1.164$ (Y-16) – 0.391 (Cb – 128) – 0.813 (Cr – 128)

B = 1.164 (Y-16) + 2.018 (Cb-128)

26.4.2 How to Use the DMA

• The 2DG DMAC has four channels. DMAC allocate the channels according to the priority level shown below:

SE buffer DMA (2DG output) > DC buffer DMA (2DG BLT output C) > SB buffer DMA (2DG BLT input B) > SA buffer DMA (2DG BLT input A)

Examples of allocation are as follows:

Channel 1: SE buffer DMA

Channel 2: DC buffer DMA

Channel 3: SB buffer DMA

Channel 4: SA buffer DMA

- When performing a data transfer using the normal CPU transfer instead of the DMAC, set the applicable DMx_MSEL bits of the GR_DMAC register to 11. Note that changing the transfer method to a CPU transfer during a DMA transfer is prohibited.
- For the SA/SB areas, the same DMA setting can be effected by setting the DM34 DSEL and DM34 MSEL bits of the GR_DMAC register.
- The number of bits (32 or 16) per data item is specified in terms of SZSEL1 (blitter)/SZSEL2 (output block). When the lower 2-bit addresses of the starting pixel for the DMA transfer are not "00", clear SZSEL1 and SZSEL2 to 0 (16 bits).
- When $SZSEL1 = 32$ bits, SSWIDH is always set to an even number; for $SZSEL1 = 16$ bits, SSWIDH can be either odd or even. Similarly, for $SZSEL2 = 32$ bits, SEWIDH is always even, and when SZSEL2 = 16 bits, SEWIDH can be either odd or even. Even when the data width of the SDRAM, which is external memory, is 16 bits, it is recommended to set the size to 32 bits when the efficiency of data transfer is important. Basically, use 32 bits if one transfer line = even pixels, and use 16 bits if one transfer line = odd pixels.
- Set the number of data items per operand so that the size of the buffer to be accessed is evenly divisible.

For example, if data transfer size $= 16$ or 32 bits and one operand $= 1, 2, 4$, or 16 data items, full image transfer (480 pixels) can be made to all buffers. However, if data transfer size $= 32$ bits and one operand = 32 data items, a full pixel (480 pixels) transfer to the SE buffer results in $480/(32 \times 2) = 7.5$, the DMA transfer cannot be stopped at 240th word after performing operand transfers seven times (access up to $32 \times 7 = 224$ words), which means that transfers must be made up to the 256th word. In this case, if one operand = 16 data items, a full image transfer can be made by performing operand transfers 15 times.

• When performing DMA access to the 2DG, set the first address of the buffer as the access-start address. Starting DMA access beginning with a middle address is prohibited.

In a DMA transfer to a buffer, if a continuous operand transfer is specified, the maximum transfer size per operation should be set equal to the buffer size to be accessed (for example, 128 bytes for an SB buffer, and 960 bytes for an SE buffer).

If a continuous operand setting is selected, the transfer is continued until all transfer data that is specified is transferred; therefore, the DMA transfer cannot be suspended on account of SB, SA, DC, or SE buffer being full. For this reason, if a continuous operand transfer is to be executed to the SB or SA buffer, line-by-line transfers should be performed solely on the basis of one horizontal pixel setting $= 32$ or 64 pixels. In the case of a DC buffer, line-by-line transfers should be performed solely on the basis of one horizontal pixel setting = 32, 64, 128, or 256 pixels. In the case of an SE buffer, line-by-line transfers should be performed based on the number of pixels that is set on the one horizontal pixel setting.

- If a single-operand transfer is specified in a DMA-based data transfer to a buffer, and if another access equal to the buffer size is possible after completion of access equal to the buffer size, the hardware automatically continues issuing DMA requests in order to align buffer boundaries with operand boundaries. If another access equal to the buffer size is not possible, the DMAC is held in standby until a condition that permits the issuance of a DMA request is satisfied after the execution of internal processing.
- Examples of DMA transfer settings are given below:

The number of pixels transferred to the 2DG is set as horizontal width 64 pixels/3 vertical lines, and each data item transferred consists of 1 pixel = 16 bits. In this case, the total number of pixels transferred will be 16 bits \times 64 pixels \times 3 lines = 384 bytes. If such data is transferred by a DMA transfer, the following two settings are possible:

- Approach 1: If one data item $= 16$ bits/operand $=$ eight data items/transfer mode $=$ pipeline (or cycle stealing)/transfer condition = single-operand transfer is set, and if the number of bytes transferred per DMA transfer is 384 bytes, for each buffer capacity the 2DG performs DMA request control to request a resumption of DMA transfer each time the buffer becomes empty and performs DMA transfers.
- Approach 2: If one data item $= 16$ bits/operand $=$ eight data items/transfer mode $=$ pipeline (or cycle stealing)/transfer condition = continuous-operand transfer is set, and if the number of bytes transferred per DMA transfer is 128 bytes and DMA is executed three times by reloading, because the buffer becomes full in each DMA transfer, for each buffer capacity the 2DG performs DMA request control to request a resumption of DMA transfer each time the buffer becomes empty and performs DMA transfers.

The following are examples of CPU-side settings for using a DMA transfer:

The number of pixels transferred to the 2DG is 52 pixels (horizontal width) \times 20 vertical lines, and the size to be transferred is 32 bits.

During normal DMA use: The amount of DMA transfer per operation is 52 pixels; if transfers are made by reloading 20 times, the OPSEL bits of the DMA mode register must be set to an integer multiple of 52 pixels. For this purpose, a transfer of two data items/operand must be set.

During use of a two-dimensional DMAC: The following settings must be specified: number of blocks per line $DBN = 1$; number of lines per block, DRN = 20; number of column data items per block, $DCDN = 26$ data items: $OPSEL = 2$ data items/operand.

- The number of data items transferred per operand, which is a transfer parameter on the DMA side, can be any value; however, on some transfer areas the situation can arise where one line of data transfer cannot be finished within an HSYNC period (causing an underflow). If this problem occurs, it can be verified by checking the INT_UDRFL bit of the GR_IRSTAT register. If an underflow occurs, increase the number of data items transferred per operand on the DMA side.
- During a data write transfer to the SE, SB, or SA buffer or during a data read transfer from the DC buffer, in either case, DMA transfer or CPU transfer, internally the 2DG processes data by re-assigning internally-generated address. For this reason, even when the CPU side performs a data transfer from any address in the applicable memory space, the 2DG side performs access beginning with the first address in the memory space.
- The direction control bit of the DMAC during a destination transfer from the DMAC to the SE, SB, or SA buffer of the 2DG or during a source transfer from the DC buffer, targets on increment (memory-to-memory image) (memory-to-I/O transfers are prohibited).
- Operation during a CPU transfer

If data access to the 2DG is to be made using the CPU transfer instead of a DMA transfer, the transfer must be performed by taking the maximum capacity of the applicable buffer into consideration (during a DMA transfer, the buffer size is controlled by internal hardware and, therefore, it need not be considered).

For example, if an image consisting of 60 pixels \times 4 lines (for a total of 240 pixels) is to be transferred to the SB buffer by means of CPU transfer,

- 1. Write access to the SB buffer is performed for $128 (= 64 \times 2)$ pixels (generating an INT_SHFUL)
- 2. If the GR_DOSTAT register is read and the SB buffer is not in a two-banks-full condition, the data transfer is resumed in units of 64 pixels.
- 3. If an INT_SHFUL is not generated, the system continues to transfer the remaining 48 pixels, and completes the data transfer process.

In the case of the DC buffer, there will be 256 pixels per bank (or 512 pixels for the two banks). Since access to the SE buffer can adversely affect display, access by means of a CPU transfer is not recommended.

Figure 26.16 shows an example of data transfer to the SB buffer using the DMAC.

Figure 26.16 96-item Data Transfer Timing in Dual/Single Address Transfer Mode (Single-Operand Transfer)

26.4.3 Blitter Operation

(1) Pixel Format Conversion in Blitter

On the blitter, as I/O pixel formats three formats can be set: α RGB444, α RGB555, and α (4 bits). For this reason, internally in the blitter, various operations are performed after converting a given format into a standard format. The size of the standard format is α (4 bits) + RGB (5 bits each) for a total of 19 bits. The formula below shows rules for the conversion of a given format to the standard format:

• α RGB444 (AF83(H)) converted into a standard format α : A(H) \rightarrow A(H) R:F(H) \rightarrow 1E(H) G:8(H) \rightarrow 10(H) B:3(H) \rightarrow 06(H)

Figure 26.17 Pixel Format Conversion 1 in Blitter

 α RGB555 (F599(H)) converted into a standard format $\alpha:1(H) \rightarrow F(H)$ R:1D(H) \rightarrow 1D(H) G:0C(H) \rightarrow 0C(H) B:19(H) \rightarrow 19(H)

Figure 26.18 Pixel Format Conversion 2 in Blitter

• α (4 bits) (B(H)) converted into a standard format $\alpha: B(H) \to B(H)$ R:Csasg_r G:Csasg_g B:Csasg_b

The native value is assigned to the α value. To the RGB value, the BRDC R, BRDC G, and BRDC B bits of the GR_BRDCOL register are assigned.

(2) Summary of Operations between the Blitter and External Memory

The following is a summary of operations between the blitter and external memory:

- 1. The blitter negates the DMA request signal and accepts a DMA transfer from the external memory.
- 2. The SA/SB buffers ((SA1, SA2), (SB1, SB2)), alternately buffer the data received through the DMA transfer.
- 3. When the INT SHFUL and INT ASHFUL are asserted (if there is source buffer half agreement), blitter operations are started.
- 4. Upon completion of blitter operations, a DMA transfer from the DC buffer is performed.
- 5. The above steps 1 to 4 are repeated until all data processing is completed.

As described above, processing such as image synthesis is possible by reading the external memory area first, performing blitter processing on the area that has been read next, and writing back to the same memory area last.

Figure 26.19 Summary of Operations between Blitter and External Memory

(3) Summary of Blitter Operations

The text below provides a summary of blitter operations. Table 26.9 shows allowable combinations of operations. Chromakey processing, color gradation processing, and logical operations cannot be performed simultaneously. Chromakey processing is allowed only during a fill operation. Color gradation processing can be performed only during blitting. Logical operations, blending, and resizing can be performed irrespective of blit/filling (subject to some restrictions).

Notes: Resizing function: During a blitting operation, resizing can be turned on/off on each items; it cannot be used during a filling operation.

- 1. The resizing function can be applied only to full resizing; it cannot be used for partial resizing.
- 2. Only SA-input data is eligible for blending; blending with register-stored data cannot be performed.

As a summary of blitter operations, the text below provides an example where target planes P1 and P2 are blended for full-plane synthesis, and the results are written back to an arbitrary memory space PX on the SDRAM.

The areas for planes P1 and P2 are set as follows: the number of lines in the SSHIGH bits of the GR_SABSET register, and the number of pixels in the SSWIDH bits of the GR_SABSET register. For the PX, the number of lines is set in the DCHIGH bits of the GR_DCSET register, and the number of pixels in the DCWIDH bits of the GR_DCSET register.

The SA and SB buffers each have a 128-byte double-buffer structure ((SA1, SA2), (SB1, SB2)). For example, if one input is selected (i.e., SB_STEN bit = 1 and SA_STEN bit = 0 in the GR BLTPLY register), and if the following values are assigned: SSWIDH bits $= 40$ (pixels) in the GR_SABSET register and SSHIGH bits = 4 (lines) in the GR_SABSET register, blitter operations work as follows:

- 1. Transfers the first 64 pixels to SB1 (SBHF_STAT $(0) = 1$), followed by blitter processing and output.
- 2. Transfers the next 64 pixels to SB2 (SBHF_STAT $(1) = 1$), followed by blitter processing and output.
- 3. Transfers the remaining 32 pixels to SB1 (SBHF_STAT $(0) = 1$), followed by blitter processing and output.
- 4. This completes the blitter operation (SB_STEN=0).

The DC buffer has a 256-byte double-buffer structure (DC1, DC2). For example, if DCWIDH bits $= 60$ (pixels) in the GR_DCSET register and DCHIGH bits $= 5$ (lines) in the GR_DCSET register, blitter operations work as follows:

- 1. Transfers the first 128 pixels to DC1 (DCHF_STAT $(0) = 1$), followed by a DMA transfer.
- 2. Transfers the next 128 pixels to DC2 (DCHF_STAT $(1) = 1$), followed by a DMA transfer.
- 3. Transfers the remaining 44 pixels to DC1 (DCHF_STAT $(0) = 1$), followed by a DMA transfer.

Figure 26.20 Summary of Blitter Operations

(a) Blending

The blending for the synthesis using two planes is performed according to a general formula, based on the following algorithm irrespective of the order in which the color value Cb and α value α b of the background plane and the color value Cf and α value α f of the front plane: (C(i): an offset value if Cpout: color value after blending, and Apout: α value after blending):

Cpout = Mf * (Cf – C(i)) + (1 – Af) * Mb * (Cb – C(i)) + C(i)

Apout = $1 - (1 - Af) * (1 - Ab)$

provided

if the front plane is premultiplied, $Mf = 1$, $Cf = Cpf$ if the front plane is non-premultiplied, $Mf = Af$, $Cf = Cf$ if the background plane is premultiplied, $Mb = 1$, $Cb = Cpb$ if the background plane is non-premultiplied, $Mb = Ab$, $Cb = Cb$

The diagram below shows the relationship between graphic synthesis planes:

Figure 26.22 Relationship between Graphic Synthesis Planes

For the output plane, since its source character and graphics planes are both non-premultiplied, it follows that Mf = α m, Cf = Cm, Mb = α z, and Cb = Cz. Given α RGB444 as the image format, it follows that $C(i) = 0$. Therefore, by applying the algorithms for Cpout and Apout to the output plane, we obtain:

Cgout = α m * Cm + (1 – α m) * α z * Cz α gout = 1 – (1 – α m) * (1 – α z)

Also, for the synthetic output, which is the final output, given that the source output image plane is premultiplied and the moving picture plane is non-premultiplied, it follows that Mb = $\alpha v = 1$. Therefore, by applying the algorithm for Cpout to the synthetic output, we obtain:

 $Cp = 1 * Cgout + (1 - \alpha gout) * 1 * Cv$

The following is an example where priority is given to character display (the case where the α value assigned to the character = 1). In this case, since α m = 1, the blending of the color value of the character part and the α value will result in the following:

Cgout = Cm, α gout = 1

In the final output, the character part will be:

 $Cp = Cm$ (no moving pictures blended)

Consequently, only characters are displayed in the results of output synthesis.

If 1 is assigned as α value ($\alpha z = 1$) of the graphics plane to give priority to graphics plane display, since the front most plane is a character plane, the results of the blending will be as follows, consisting of a blending between characters and graphics, unless the α value of the character plane $= 0$ (α m = 0):

Cgout = α mCm + (1 – α m) Cz, α gout = 1

The final output will be:

 $Cp = \alpha mCm + (1 - \alpha m) Cz$ (no moving pictures blended)

The blending processor in the blitter, which performs pseudo-anti-aliasing with characters, blends characters having an arbitrary α value with an α value = 1 graphic in a rectangular region, and it can transfer the results to the output plane by converting the blending results to an arbitrary α value (register value) for blending with the moving pictures. In this case, the blending results between characters and graphics will be:

Cgout = α mCm + (1 – α m) Cz, α gout = 1

Further, if global α (the BRDC_A bit of the GR_BRDCOL register) is selected using the subsequent stage global α selection bit (the GALFA bit of the GR_BRD1CNT register), the results that are output from the blending processor in the blitter will be:

Cgout = α mCm + (1 – α m) Cz, α gout = α dc = BRDC A bit of GR BRDCOL register

Thus, pixels having αdc values are created in the output plane.

In terms of the 2DG, if synthesis using two planes is performed by blending in the blitter and all color values of the output plane are α -value-computed (i.e., weighted) data, the blending between the moving picture and the output plane in the output block will be performed by selecting premultiplied in the blending processor in the output block. On the other hand, if the color values on the output plane are not α -weighted, the blending between the moving picture and the output plane in the output block will be performed by selecting non-premultiplied in the blending processor in the output block.

Figure 26.23 shows a summary of operations during blending, as follows:

- (1) If GCOLR bit in GR_BRD1CNT register $= 0$, selects data from the SA buffer. If GCOLR bit in GR_BRD1CNT register = 1, selects the value of the GR_BRDCOL register. In fill mode, the value in the GR_BRDCOL register is selected in the same way as if the GCOLR bit of the GR_BRD1CNT register was set (1). However, the value of the GCOLR bit does not change at this time.
- (2) Blends the data selected in Step 1 above with the data from the SB buffer selected on the basis of the SBSEL bit of the GR_BLTMODE register.
- (3) Of the data blended in Step 2 above, the color value Cout is output to the DC buffer as is. With regard to the α value α out, the blended α value is output to the DC buffer as is, provided the GALFA bit of the GR_BRD1CNT register $= 0$. If the GALFA bit of the GR_BRD1CNT register = 1, the AFTER_A bit value of the GR_BRD1CNT register is output to the DC buffer.

Figure 26.23 Summary of Blitter Operations during Blending

(b) Summary of Fill Function

This section shows examples of using the fill function. Set the registers as follows:

- Set the GR_BRDCOL register to specify a fill color.
- Set the GR_SABSET register to specify the number of pixels to be transferred to the SB buffer.
- To specify the number of pixels to be transferred to the DC buffer, set the GR_DCSET register.
- Set the BTYPE bits of the GR_BLTMODE register to 10 in order to specify filling operation for blitter operating mode.
- To specify the blitter function, set the SB_SETEN bit of the GR_BLTPLY register to 1, and the SA_SETEN bit to 0.

With the above settings, the fill function operates as follows:

- 1. By the GR_BLTPLY register settings, image data equal to the pixel count set in the GR_SABSET register is transferred to the SB buffer.
- 2. The pixel data transferred to the SB buffer is blended with the pixel data set in the GR_BRDCOL register.
- 3. Image data equal to the number of pixels set in the GR_DCSET register is output to the DC buffer.

The DMAC first transfers, memory-to-memory, pixels equal to the pixel count set in the register from an SDRAM area specified by the CPU to the SB buffer. After that, the DMAC transfers, memory-to-memory, the pixel data processed in the 2DG from the DC buffer to the SDRAM area specified by the CPU (the original image area whose contents were transferred to the SB buffer). As a result, the specified area on the SDRAM can be replaced with blended pixel data. In this case, the 2DG performs input to the SB buffer and output from the DC buffer.

Figure 26.24 Example of Fill Function

(c) Summary of Blit Function

Blit function without blending

This section shows an example of using the blit function without blending. Set the registers as follows:

- Set the GR_SABSET register to specify the number of pixels to be transferred to the SB buffer.
- Set the GR_DCSET register to specify the number of pixels to be transferred to the DC buffer.
- Set the BTYPE bits of the GR_BLTMODE register to 00 in order to specify blit operation for blitter operating mode.
- Set the FBFA bits of the GR_BRD1CNT register to 01 in order to specify 1-input processing.
- To enable the blit function, set the SB_SETEN bit of the GR_BLTPLY register to 1, and the SA_SETEN bit to 0.

With the above settings, the blit function operates as follows:

- 1. By the GR_BLTPLY register settings, image data equal to the pixel count set in the GR_SABSET register is transferred to the SB buffer.
- 2. The pixel data transferred to the SB buffer undergoes various processings set by the relevant registers.
- 3. Image data equal to the number of pixels set in the GR_DCSET register is output to the DC buffer.

The DMAC first transfers, memory-to-memory, pixels equal to the pixel count set in the register from an SDRAM area specified by the CPU to the SB buffer. After that, the DMAC transfers, memory-to-memory, the pixel data processed in the 2DG from the DC buffer to the SDRAM area specified by the CPU. As a result, the image on any area on the SDRAM can be replaced with pixel data having undergone various image processings by the 2DG. In this case, the 2DG performs input to the SB buffer and output from the DC buffer.

Figure 26.25 Example of Blit Function without Blending

• Blit function with blending (1)

This section shows an example of blending graphic images. Set the registers as follows:

- Set the GR_SABSET register to specify the number of pixels to be transferred to the SA/SB buffers.
- Set the GR_DCSET register to specify the number of pixels to be transferred to the DC buffer.
- Set the BTYPE bits of the GR_BLTMODE register to 00 in order to specify blit operation for blitter operating mode.
- To enable the blit function, set the SB SETEN and SA SETEN bits of the GR BLTPLY register to 1.

With the above settings, the blit function operates as follows:

- 1. By the GR_BLTPLY register settings, image data equal to the pixel count set in the GR_SABSET register is transferred to the SA/SB buffers.
- 2. The pixel data transferred to the SA/SB buffers undergoes various processings set by the relevant registers.
- 3. Image data equal to the number of pixels set in the GR_DCSET register is output to the DC buffer.

The DMAC first transfers, memory-to-memory, pixels equal to the pixel count set in the register from an SDRAM area specified by the CPU to the SA/SB buffers. After that, the DMAC transfers, memory-to-memory, the pixel data processed in the 2DG from the DC buffer to the SDRAM area specified by the CPU. As a result, the image on any area on the SDRAM can be replaced with pixel data having undergone various image processings by the 2DG. In this case, the 2DG performs input to the SA/SB buffers and output from the DC buffer.

Figure 26.26 Example of Blit Function with Blending (1)

Blit function with blending (2)

This section shows an example of blending graphic and character images. Set the registers as follows:

- Set the GR_SABSET register to specify the number of pixels to be transferred to the SA/SB buffers.
- Set the GR_DCSET register to specify the number of pixels to be transferred to the DC buffer.
- Set the BTYPE bits of the GR_BLTMODE register to 00 in order to specify blit operation for blitter operating mode.
- To enable the blit function, set the SB_SETEN and SA_SETEN bits to 1.

In this case, if the α – value-weighted character pixels set in a rectangular area are input into the SB buffer and graphic pixels (α value = F(H)=1) set in the same rectangular area as SB are input into the SA buffer, the output image will be an α -blending of the graphics and characters in the rectangular area followed by resizing. In this case, in order to place character information in the foremost part of the image, it should be input into the SB buffer.

Figure 26.27 Example of Blit Function with Blending (2)

(d) Summary of Chromakey Processing

Since the chromakey processing function is installed only on the input B path, chromakey processing cannot be applied to the input A path. Furthermore, neither logical operations nor color gradation processing can be executed at the same time as chromakey processing. Part of the chromakey specification is that chromakey settings are only effective when fill operations are running. For this reason, select fill mode by setting the BTYPE bits of the GR_BLTMODE register to 10 if chromakey processing is to be used. Additionally, since chromakey processing only operates with fill mode, it is not available in conjunction with the resizing function.

Examples of the effects of executing chromakey processing are as follows:

Figure 26.28 Example of Chromakey Processing (1)

• Chromakey processing example (2) Blit mode setting: Fill mode (GR_BLTMODE register, BTYPE bits = 10) Chromakey type selection: Chromakey blending (GR BLTMODE register, CRKEY bits $= 10$) Blending coefficient selection:

Blending coefficient = $(GR$ BRD1CNT register, FBFA bits = 00) SB output data selection: SB data selection = (GR BLTMODE register, SBSEL bits = 01) Target color setting: Setting for green as the target color in the GR DETCOL register Replacement color setting: Setting for blue as the replacement color in the GR_BRDCOL

Figure 26.29 Example of Chromakey Processing (2)

(e) Summary of Logical Operations

Since the logical operation function is installed only on the input B path, it cannot be performed on the input A path. Furthermore, chromakey processing and color gradation processing cannot proceed simultaneously. As logical operation specifications, logical operation settings are enabled on all operating modes that are set by the BTYPE bits of the GR_BLTMODE register.

Examples of the effects of using logical operations are given below:

Figure 26.30 Example of Logical Operations

(f) Summary of Color Gradation Processing

Since the color gradation processing function is installed only on the input B path, it cannot be performed on the input A path. Irrespective of blending, color gradation processing is available only during blit operations. During such operations, chromakey processing or logical operations cannot be executed at the same time. Although resizing is possible, the use of it is not recommended in situations where character quality is important. Where possible, resizing should be avoided. The color gradation processing is not designed to accommodate partial resizing.

The basic operations of the color gradation processing are as follows:

- 1. DMA-transfers only a specified area from any font area, created with α (4 bits) on the SDRAM, to the 2DG.
- 2. Expands the color that is set in the GR_BRDCOL register.
- 3. Creates a color gradation according to the α value by weighting (multiplying) with the transferred α value and a specified color.

Depending on how a font is created using an α value, pseudo-anti-alias effects can also be produced.

• Color gradation processing example (1) Blit mode setting: Blit mode (GR_BLTMODE register, BTYPE bits = 00) Blending operation: Without blending (GR_BLTPLY register, SB_STEN bit $= 1$ and SA STEN bit $= 0$)

Figure 26.31 Example of Color Gradation Processing (1)

• Color gradation processing example (2)

The text below describes the processing where data in a specified area from the α font region and data from a specified area on the output plane is entered into the 2DG, a color gradation processing is performed, and the results are written back to the specified area on the output plane:

Blit mode setting: Blit mode (GR_BLTMODE register, BTYPE bits = 00) Blending operation: With blending (GR_BLTPLY register, SB_STEN = 1 and SA_STEN bit $= 1$

Figure 26.32 Example of Color Gradation Processing (2)

(4) Resizing

(a) Resize mode setting

Use of the resizing function in the blitter can be specified by the BRISZ bit of the GR_RISZSET register. Resize mode can be specified with respect to the following four factors using the relevant bits of the GR_RISZMOD register.

- A1_H bit: Selects the horizontal α -resizing method. 0: Bilinear method, 1: nearest-neighbor method.
- H1 MTHD bit: Selects the horizontal resizing method. 0: Bilinear method, 1: nearest-neighbor method.
- A1 V bit: Selects the vertical α-resizing method. 0: Bilinear method, 1: nearest-neighbor method.
- V1 MTHD bit: Selects the vertical resizing method. 0: Bilinear method, 1: nearest-neighbor method.

(b) Setting EDGE processing

When full or partial resizing is to be performed, it is necessary to determine whether the right or lower edge of the area to be resized coincides with the right or lower edge of the full resize plane. Therefore, when performing full or partial resizing, the EDGE bits of the GR_RISZSET register must be set accordingly.

An example of a setting for the EDGE bits of the GR_RISZSET register shows the case of the horizontal direction (right edge). In figure 26.33, the decision regarding the EDGE bits of the GR_RISZSET register is depicted in terms of the whole source area, the source update area, and the source-setting area. The mass with an attached color is the region to be updated within the whole source area (the updated source area). Furthermore, when the region indicated by the bold arrow is partially resized, this region is set in registers as the source area (source-setting area). For details, see section 26.4.3 (5), Partial Resizing.

(c) Setting the delta value for use in resizing

The delta value for use in resizing (Ch) can be calculated by using the following equation:

Ch = (source resolution / destination resolution) \times 4096

Note: Ch = 1 / resizing ratio must be calculated from (number of source pixels/number of destination pixels). Also, truncate any fractional part of the result.

For example, if the number of source pixels is 479 and the number of destination pixels is 240, Ch will be:

 $Ch = (479/240) \times 4096 \approx 8174.933 = 1FEE (H)$

Based on these results, the delta value for resizing will be as follows:

Since HDLT_INTGR is the integer part of Ch, HDLT_INTGR = 1 (H) Since HDLT DCML is the fractional part of Ch, HDLT DCML = FEE (H)

If the BRSIZ bit of the GR_RISZSET register is 0 , i.e., the resizing function is not to be used, the operation will be the same as if the HDLT_INTGR and HDLT_DCML bits were set to H'1 and H'000, respectively. However, the values of the HDLT_INTGR and HDLT_DCML bits do not change.

The delta value for use in resizing (Cv) value can be calculated by using the following equation:

 $Cv =$ (source resolution / destination resolution) \times 4096

In this case, through calculation similar to that for Ch, VDLT_INTGR and VDLT_DCML are derived as the integer and fractional parts of Cv (2 and 12 bits, respectively). If the BRSIZ bit of the GR_RISZSET register is 0, i.e., the resizing function is not to be used, the operation will be the same as if the VDLT_INTGR and VDLT_DCML bits were set to H'1 and H'000, respectively. However, the values of the VDLT_INTGR and VDLT_DCML bits do not change.

• Range of settings for full resizing

Since this is full resizing, the integer component (GR_HSPHAS register, H1PHS_INTGR bit) is always 0.

Enlargement: H1PHS_INTGR = H'000, H1PHS_DCML = H'000 to H'(HDLT_DCML – 1) Reduction: H1PHS_INTGR = H'000, H1PHS_DCML = H'000 to H'FFF

• Range of settings for partial resizing Set the integer and fractional parts according to the left-edge pixels of the area for partial resizing.

(d) Setting the source-side starting phase

The source-side starting phase (Psh) can be calculated by using the following equation:

Psh = $Ch \times$ (number of pixels to that where resizing starts) + (initial phase at starting pixel \times 4096)

In this case, the integer part (10 bits) of Psh will be in H1PHS_INTGR and the fractional part (12 bits) of Psh will be in H1PHS_DCML. If there is to be no resizing or the magnification is to remain the same (resizing by a factor of one), set $Psh = 0$.

The source-side starting phase (Psv) can be calculated by using the following equation:

 $Psv = Cv \times$ (number of pixels to that where resizing starts) + (initial phase at starting pixel \times 4096)

In this case, the integer part (9 bits) of Psv will be in V1PHS_INTGR and the fractional part (12) bits) of Psv will be in V1PHS_DCML. If there is to be no resizing or the magnification is to remain the same (resizing by a factor of one), set $Psv = 0$.

• Range of settings for full resizing

Since this is full resizing, the integer part (V1PHS_INTGR) is always 0. Enlargement: V1PHS_INTGR = H'000, V1PHS_DCML = H'000 to H'(VDLT_DCML – 1) Reduction: V1PHS_INTGR = H'000, V1PHS_DCML = H'000 to H'FFF

• Range of settings for partial resizing Set the integer and fractional parts for correspondence with the pixels at the top-edge of the area for partial resizing.

The starting phase is used to vary the proportion of mixing for the two source pixels used as a reference for bilinear filtering in resizing. This can be used to eliminate the omission of pixels in halving of the size (reduction by a factor of two) and so on. However, too large an initial phase can lead to a mismatch of colors at the left edge of the destination. For these reasons, the following limitations apply to settings for the starting phase in resizing.

Limitations:

Always set the integer components of the starting phase (the H1PHS_INTGR and V1PHS_INTGR bits) to zero.

Set the fractional components of the starting phase (the H1PHS_DCML and V1PHS_DCML bits) within the ranges given below:

- Resizing for enlargement: H1PHS DCML and V1PHS DCML bits = H'000 to H' $(VDLT_DCML - 1)$
- Resizing for reduction: H1PHS_DCML and V1PHS_DCML bits = H'000 to H'FFF

Examples of the settings for starting phase in resizing for reduction (magnification by half) are given below (the explanation only applied to the horizontal direction). In figure 26.34, S0 to S4 are source pixels and D0 to D2 are the destination pixels interpolated for the given magnification factor. Destination pixels (1) and (2) indicate the phases for interpolation of the destination pixels for starting phase settings of H'0 and H'800, respectively.

Figure 26.34 Settings for Starting Phase in Resizing for Reduction (Magnification by Half)

The fractional components of the starting phase (the H1PHS_DCML and V1PHS_DCML bits) are asset for an interval of one between source pixels. In figure 26.34, destination pixels (1) shows the positions of interpolation for the destination pixels when the starting phase (H1PHS $DCML$) = H'000, where the position of the leading pixel (D0) is generated with a phase of zero. Destination pixels (2) shows the positions of interpolation for the destination pixels when the starting phase $(H1PHS_DCML) = H'800$, where the position of the leading pixel (D0) is advanced by only the starting phase (H1PHS DCML), and the starting phase is maintained in subsequent interpolation. The proportions of the two source pixels for reference can be varied by attaching or not attaching a starting phase in this way.

The method can thus be used to alleviate the effects of the omission of pixels. Figure 26.35 shows examples of the results when a starting phase is and is not attached in resizing for reduction (magnification by half).

Figure 26.35 Examples of the Results when Settings for Starting Phase in Resizing for Reduction (Magnification by Half)

As shown in figure 26.35, when resizing for reduction of the upper third with one dot omitted is obtained by the bilinear method with a starting phase of H'000, the proportion between the color values of the white and colored pixels becomes 1:0, so the color value of the colored pixels is simply eliminated. When the starting phase is H'800, the proportion between the color values of the white and colored pixels becomes 1:1 and, since the color values of the white and colored pixels are blended in equal portions, the color of the colored pixels is reproduced although the vertical lines disappear.

Since pixels are not omitted in the case of resizing for enlargement, the kind of strong effect seen in the case of resizing for reduction is absent. However, the setting for starting phase can still be used to control the proportions of the two source pixels used as reference in resizing for enlargement (through the bilinear method).

(5) Partial Resizing

When modifying a part of the area on a resized source plane, a partial extraction should be performed on the source area before resizing instead of after. The reason is that the correct boundaries can be maintained by partially extracting the source area before resizing and by pasting the area to the part of the area on the resized plane. An example of this is shown in figure 26.36. Partial resizing, however, cannot be performed when color gradation processing is selected.

The partially modified area (the updated area in figure 26.36) is defined as follows:

```
Vertical offset = Va, horizontal offset = Ha, height = Vb, width = Hb
```
The register settings for the horizontal width and vertical height of the source area are specified by determining a setting area appropriate for partial resizing, instead of specifying the updated area itself, according to the following equations, and by assigning the results to the GR_HSPHAS and GR_VSPHAS registers:

Vertical offset = Va1 = Va − Vx1, horizontal offset: Ha1 = Ha − Hx1

Vertical height: Vb1 = Vb + Vx1 + Vx2, horizontal width: Hb1 = Hb1 + Hx1 + Hx2

Figure 26.36 Setting Areas for Partial Extraction Resizing

Figure 26.37 Summary of Partial Extraction Resizing

(a) What is partial resizing?

Given an original source image area (for example, a 500×300 pixel area) for a fully resized destination image, if only a part (for example, 50×50 pixels) of the source image area is modified, partial resizing refers to the method by which only the modified area (approximately 50 \times 50 pixels) is resized, instead of resizing the entire source image area (500 \times 300).

(b) Defining the area to be partially resized

• EDGE bits of GR_RISZSET register

When performing full or partial resizing, it is necessary to determine whether the right edge or the lower edge of the source area to be resized coincides with the right or lower edge of the full-resize area. For this reason, when performing a full or partial resizing, the EDGE bits of the GR_RISZSET register must be set.

Figure 26.38 explains the values of the EDGE bits to be set in the GR_RISZSET register for areas 1 (TL) to 5 (MM) for partial resizing relative to the full resizing area (area in the bold frame). In reference to the pointers given below, set the EDGE bits according to the position of the partial resizing area.

Figure 26.38 Examples of Partial Resizing Areas

Partial resizing source set area (SSWIDH and SSHIGH bits of GR_SABSET register)

For partial resizing, an area a little larger than the actually updated area is assigned to the register as a partial resizing source area. In this manner, the boundary between the partially updated area and the area that is not updated is eliminated. An example of this given shown in figure 26.39. The area to be set can vary according to the particular resizing ratio (enlargement/reduction) employed.

As shown in figure 26.39, an area expanded by $+\alpha$ vertically and horizontally from the actually updated area (yellow) is assigned to the partial resizing set area (SSWIDH bits/SSHIGH bits). Also, as shown in figure 26.39, if the partially updated area (yellow) is close to the right or lower edge of the fully resized area, in some cases the partially updated area is defined on the right or lower edge as a set area, different from the actual updated area. In such a case, the right and lower edges must be specified to the EDGE bits. (The factor α varies with the resizing ratio (enlargement/reduction)).

(c) Horizontal/vertical starting phase for the source (H1PHS_DCML bits of GR_HSPHAS register /V1PHS_DCML bits of GR_VSPHAS register)

When performing partial resizing, it is necessary to consider the starting phase of the starting pixel for the full resizing in addition to the set area for the partial resizing, in order to ensure that the complete match of the boundary between the output image that is fully resized in advance and the output image that is partially resized later. For example, if the horizontal starting phase (H1PHS_DCML) for the source for full resizing is defined as H'FFF, the same value as full resizing (H'FFF) must also be added to the starting phase (H1PHS_DCML) when partial resizing is performed.

If the parameter H1PHS_DCML or V1PHS_DCML is to be used during the full resizing process, constraint conditions vary between enlargement resizing and reduction resizing; therefore, each parameter should be set in the following ranges:

[Enlargement] H1PHS DCML or V1PHS $DCML = H'000$ to $H'(\text{HDLT} DCML - 1)$ [Reduction] H1PHS_DCML or V1PHS_DCML = H'000 to H'FFF

(d) Determining the area to set up for partial resizing

In figure 26.40, ST, SL, SSHIGH, and SSWIDH are values to be obtained by the user. The other values are known.

Figure 26.40 Partial Resizing Set Area in Source Data Area

In figure 26.41, DT, DL, DCHIGH, and DCWIDH are values to be obtained by the user. The other values are known.

• Definitions: Explanation of codes provided in figures 26.40 and 26.41

- ST': The number of lines from the starting line for the vertical source read-out area for full resizing to the top-edge pixel in the vertical source read-out area in the updated area
- SB': The number of lines from the starting line for the vertical source read-out area for full resizing to the lower-edge pixel in the vertical source read-out area in the updated area
- SL': The number of pixels from the left-edge pixel for the horizontal source read-out area for full resizing to the left-edge (starting) pixel in the horizontal source read-out area in the updated area
- SR': The number of pixels from the left-edge pixel for the horizontal source read-out area for full resizing to the right-edge (end) pixel in the horizontal source read-out area in the updated area
- ST: The number of lines from the starting line for the vertical source read-out area for full resizing to the top-edge pixel in the vertical source read-out area in the partial resize set area
- SL: The number of pixels from the left-edge pixel for the horizontal source read-out area for full resizing to the left-edge (starting) pixel in the horizontal source read-out area in the partial resize set area
- DT: The number of lines from the starting line for the vertical destination area for full resizing to the top-edge pixel in the vertical destination area in the partial resize set area
- DL: The number of pixels from the left-edge pixel for the horizontal destination area for full resizing to the left-edge (starting) pixel in the horizontal destination area in the partial resize set area
- Definitions: Explanation of codes not provided in figures 26.40 or 26.41

GR_HSPHAS register, H1PHS_DCML bit:

For full resizing, the fractional part of the results of calculation of the phase in the horizontal starting position on the source side

GR_VSPHAS register, V1PHS_DCML bit: For full resizing, the fractional part of the results of calculation of the phase in the vertical starting position on the source side

1. Formula for calculating the number of left-edge pixels (DL) and the number of top-edge lines (DT) in the destination set area

Formula for judging DL:

When $SL' > 1$, INT (DH/ $SH \times (SL' - 1) \times INT$ (SH/ DH \times 4096) + H1PHS DCML \geq (SL' -1) × 4096

When $SL \leq 1$, $DL = 0$ for exception handling, regardless of formula for calculation

Formula for calculating DL:

When the result of formula for judging is true: $DL = INT (DH/SH \times (SL' - 1))$

When the result of formula for judging is false: $DL = INT (DH/SH \times (SL'-1)) + 1$ Formula for judging DT:

When $ST' > 1$, INT (DV/ $SV \times (ST' - 1) \times INT$ (SV/ DV \times 4096) + V1PHS_DCML \geq (ST' -1) × 4096

When $ST \le 1$, $DT = 0$ for exception handling, regardless of formula for calculation Formula for calculating DT:

When the result of formula for judging is true: $DT = INT (DV/SV \times (ST' - 1))$

When the result of formula for judging is false: $DT = INT (DV/SV \times (ST - 1)) + 1$

2. Formula for calculating the number of left-edge pixels (SL) and the number of top-edge lines (ST) in the source set area

Formula for calculating SL:

When $SL' > 1$, $SL = INT$ (($DL \times INT$ ($SH/DH \times 4096$) + H1PHS DCML/ 4096) = (H1PHS_INTGR)

When $SL \leq 1$, $SL = 0$ as exception handling

Formula for calculating ST:

When $ST' > 1$, $ST = INT (DT \times INT (SV/DV \times 4096) + V1PHS DCML/ 4096) =$ (V1PHS_INTGR) When $ST \leq 1$, $ST = 0$ as exception handling

3. Formula for calculating the horizontal width (DCWIDH) and the vertical height (DCHIGH) in

the destination set area

Formula for judging DCWIDH:

When $SR' < SH - 2$, INT ((INT (DH/ $SH \times (SR' + 1)$) \times INT (SH/ DH \times 4096) + H1PHS_DCML \geq (SR' + 1) \times 4096

When $SR \geq SH - 2$, DCWIDH = DH – DL for exception handling, regardless of formula for calculation

Formula for calculating DCWIDH:

When the result of formula for judging is true: DCWIDH = INT (DH/ SH \times (SR' + 1)) – DI .

When the result of formula for judging is false: DCWIDH = INT (DH/ SH \times (SR' + 1)) – $DL + 1$

Formula for judging DCHIGH:

When SB' < SV – 2, INT ((INT (DV/ SV \times (SB' + 1)) \times INT (SV/ DV \times 4096) + V1PHS DCML \geq (SB' + 1) \times 4096

When $SB' \geq SV - 2$, DCHIGH = DV – DT for exception handling, regardless of formula for calculation

Formula for calculating DCHIGH:

When the result of formula for judging is true: DCHIGH = INT (DV/ SV \times (SB' + 1)) – **DT**

When the result of formula for judging is false: $DCHIGH = INT (DV/SV \times (SB' + 1))$ – $DT + 1$

4. Formula for calculating the horizontal width (SSWIDH) and the vertical height (SSHIGH) in the source set area

Formula for judging SSWIDH:

When SR' < SH− 2, INT ((INT \times (DH/ SH (SR' + 1)) \times INT (SH/ DH \times 4096) + H1PHS_DCML \geq (SR' + 1) \times 4096

When $SR' \geq SH - 2$, SSWIDH = SH – SL for exception handling, regardless of formula for calculation

Formula for calculating SSWIDH:

When the result of formula for judging is true: SSWIDH = INT (((INT (DH/ SH \times (SR' + 1)) − 1) × INT (SH/ DH × 4096) + H1PHS_DCML)/ 4096) − SL + 2

When the result of formula for judging is false: SSWIDH = INT (((INT (DH/ SH \times (SR' + $1)$) – 0) × INT (SH/ DH × 4096) + H1PHS DCML)/ 4096) – SL + 2

Formula for judging SSHIGH:

When SB' < SV− 2, INT (DV/ SV × (SB' + 1)) × INT (SV/ DV × 4096) + V1PHS_DCML \geq (SB' + 1) \times 4096

When $SB' \geq SV - 2$, SSHIGH = SV – ST for exception handling, regardless of formula for calculation

Formula for calculating SSHIGH:

When the result of formula for judging is true: SSHIGH = INT (((INT (DV/ SV \times (SB' + 1)) − 1) × INT (SV/ DV × 4096) + V1PHS_DCML)/ 4096) − ST + 2

When the result of formula for judging is false: SSHIGH = INT (((INT (DH/ SH \times (SR' + 1)) − 0) × INT (SV/ DV × 4096) + V1PHS_DCML)/ 4096) − ST + 2

5. Calculation of the phase in the starting position on the source side (PHS_H/ PHS_V)

 $PHS_H = INT (SH/DH \times 4096) \times DL + H1PHS_DCML$

Integer part of the results of calculation of the phase in the starting position on the source $side$ (H1PHS_INTGR) = PHS_H (upper 10 bits) = H'

Fractional part of the results of calculation of the phase in the starting position on the source side $(H1PHS_DCML) = PHS_V (lower 12 bits) = H'$

PHS $V = INT (SV/ DV \times 4096) \times DT + V1PHS DCML$

Integer part of the results of calculation of the phase in the starting position on the source side (V1PHS_INTGR) = PHS_H (upper 9 bits) = H'

Fractional part of the results of calculation of the phase in the starting position on the source side (V1PHS_DCML) = PHS_V (lower 12 bits) = H'

Note: In the above formulas, the letters INT mean that the fractional part is rounded to an integer.

(e) Calculation examples of the values in the register setting (only for horizontal direction)

• Conditions

Width of the source area (SH): 280 pixels

Width of the destination area (DH): 350 pixels

Number of the left-edge pixel of the actually updated partial area of the source area (SL'): 55

Number of the right-edge pixel of the actually updated partial area of the source area (SR'): 133

Fractional part of the results of calculation of the phase for the starting position on the source side (H1PHS_DCML (D)): 819 = H'333

• Parameters to be calculated by the user

In the case of partial resizing, the number of the left-edge pixel of the area set as the destination (DL)

In the case of partial resizing, the number of the left-edge pixel of the area set as the source (SL)

In the case of partial resizing, the width of the area set as the destination (DCWIDH)

In the case of partial resizing, the width of the area set as the source (SSWIDH)

Integer part of the result of calculating the phase of the starting position on the source side (H1PHS_INTGR)

Fractional part of the result of calculating the phase of the starting position on the source side (H1PHS_DCML)

1. Calculating DL, the left-edge pixel of the area set as the destination Since $SL' = 55$ pixels, judgment proceeds as follows. Conditional expression:

INT (DH/SH \times (SL' – 1)) \times INT (SH/DH \times 4096) + H1PHS DCML \geq (SL' – 1) \times 4096 INT $(350/280 \times (55 - 1)) \times INT (280/350 \times 4096) + 819 \ge (55 - 1) \times 4096$ Since the result of the conditional expression is "false", the value is obtained as shown below. False: DL = INT (DH/SH \times (SL' – 1)) + 1 = INT (350/280 \times (55 – 1)) + 1 = 68

2. Calculating SL, the left-edge pixel of the area set as the source Since $SL' = 55$ pixels, the value is obtained as follows. $SL = INT (DL \times INT (SH/DH \times 4096) + H1PHS DCML) / 4096$ $=$ INT ((68 \times INT (280/350 \times 4096) + 819)/ 4096) = 54

From the above, the result for the integer part of the phase for the source-side starting position is obtained as H1PHS $INTGR = 54$.

3. Calculating DCWIDH, the width of the destination area Since $SR' = 133$ and $SH = 280$ pixels, judgment proceeds as follows. Conditional expression:

INT (DH/SH \times (SR' + 1)) \times INT (SH/DH \times 4096) + H1PHS_DCML \geq (SR' + 1) \times 4096

INT $(350/280 \times (133 + 1)) \times \text{INT } (280/350 \times 4096) + 819 \ge (133 + 1) \times 4096$

Since the result of the conditional expression is "false", the value is obtained as shown below. False: DCWIDH

 $=$ INT (DH/SH \times (SR' + 1)) – DL + 1 = INT (350/280 \times (133 + 1)) – 68 + 1 = 100

4. Calculating SSWIDH, the width of the source area

Since $SR' = 133$ and $SH = 280$ pixels, judgment proceeds as follows.

Conditional expression:

INT (DH/SH \times (SR' + 1)) \times INT (SH/DH \times 4096) + H1PHS_DCML \geq (SR' + 1) \times 4096

INT $(350/280 \times (133 + 1)) \times \text{INT } (280/350 \times 4096) + 819 \ge (133 + 1) \times 4096$

Since the result of the conditional expression is "false", the value is obtained as shown below. False: SSWIDH

 $=$ INT (((INT (DH/SH \times (SR' + 1)) – 0) \times INT (SH/DH \times 4096) + H1PHS_DCML)/ 4096) $-$ SL+ 2

 $=$ INT (((INT (350/280 × (133 + 1)) – 0) × INT (280/350 × 4096) + 819)/ 4096) – 54 + 2 $= 81$

- Note: Since $SL + SSWIDH \neq SH$ at this time, the horizontal edge is not at the right edge. Accordingly, set the EDGE (0) bit to 0.
- 5. Calculating PHS_H, the result of phase calculation for the source-side starting position PHS = INT (INT (SH/DH \times 4096) \times DL + H1PHS_DCML) = INT (INT (280/350 \times 4096) \times $68 + 819$) = 223587 (H'036963)

From the above result of phase calculation for the source-side starting position, the following result is obtained.

H1PHS_INTGR, integer part of the result of phase calculation for the source-side starting position = PHS H (upper 10 bits) = H'036

H1PHS_DCML, fractional part of the result of phase calculation for the source-side starting position = PHS H (lower 12 bits) = H'963

(f) Determining the starting address of the partial resize area to be assigned to the DMAC by the CPU

• Determining the starting address (Sa) of the source set area

[Definition]

- Sa: Starting address for reading the source data for partial resizing
- SA: Starting address for reading the source data for full resizing
- SGP: Line pitch in the source plane (a 64-byte boundary)

 $SGP = \text{ROUNDUP}$ (byte count per pixel) \times SH/64 bytes) \times 64 bytes

- Note: ROUNDUP in the formula indicates rounding up, i.e. taking the nearest integer above the fractional component.
- ST: The number of lines from the starting line for the vertical source read-out area for full resizing to the top-edge pixel in the vertical source read-out area for partial resizing
- SL: The number of pixels from the left-edge pixel for the horizontal source read-out area for full resizing to the left-edge (starting) pixel in the horizontal source read-out area for partial resizing

Formula for calculating Sa:

 $Sa = SA + SGP \times ST + SL \times (byte count per pixel)$

• Determining the starting address (Da) of the destination set area

Figure 26.43 Determination of Starting Address (Da) of Destination Set Area

[Definition]

- Da: Starting address for writing to the destination area for partial resizing
- DA: Starting address for writing to the destination area for full resizing
- DGP: Line pitch in the destination area (a 64-byte boundary) $DGP = \text{ROUNDUP}$ (byte count per pixel) \times DH/64 bytes) \times 64 bytes
- Note: ROUNDUP in the formula indicates rounding up, i.e. taking the nearest integer above the fractional component.
- DT: The number of lines from the starting line for the vertical destination area for full resizing, to the top-edge pixel in the vertical destination area for partial resizing
- DL: The number of pixels from the left-edge pixel for the horizontal destination area for full resizing, to the left-edge (starting) pixel in the horizontal destination area for partial resizing

Formula for calculating Da:

 $Da = DA + DGP \times DT + DL \times (byte count per pixel)$

• Allocation within memory space

Figure 26.44 depicts the allocation within memory space of a source image that is 100 pixels wide and three lines high.

Figure 26.44 Allocation of images within memory space

- Examples of calculation
- (1) Obtaining Sa, the first address of the area set as the source.
	- (a) Conditions

SA, the first address of the source area in the ca+se of full resizing: $H'10000 (= 65536)$

SH, the number of pixels across the source area in the case of full resizing: 120

ST, the number of lines to the top of the source area to be read out in the case of partial resizing: 15

SL, the number of pixels horizontally to the left edge of the source area (first point) to be read out in the case of partial resizing: 5

Number of bytes per pixel: 2

(b) Obtaining the pitch of rows in the source plane

 $SGP = \text{ROUNDUP}$ ((number of bytes for one pixel) \times SH/64 bytes) \times 64 bytes $=$ ROUNDUP ((2 bytes) \times 120/64 bytes) \times 64 bytes $=$ 256 bytes

(c) Obtaining Sa, the first address of the area set as the source

 $Sa = SA + SGP \times ST + SL \times (number of bytes for one pixel)$ $=65536 + 256$ bytes $\times 15 + 5 \times 2$ bytes $= 69386$ ($= H10F0A$)

(2) Obtaining Da, the first address of the area set as the destination.

(a) Conditions

DA, the first address of the destination area in the case of full resizing: $H'11000 (= 69632)$

DH, the number of pixels across the destination area in the case of full resizing: 277

DT, the number of lines to the top of the destination area in the case of partial resizing: 33

DL, the number of pixels horizontally to the left edge of the destination area (first point) to be read out in the case of partial resizing: 15

Number of bytes per pixel: 2

(b) Obtaining the pitch of rows in the source plane

 $DGP = \text{ROUNDUP}$ ((number of bytes for one pixel) \times DH/64 bytes) \times 64 bytes $=$ ROUNDUP ((2 bytes) \times 277/64 bytes) \times 64 bytes = 576 bytes

- (c) Obtaining Da, the first address of the area set as the destination
- $Da = DA + DGP \times DT + DL \times (number of bytes for one pixel)$ $=69632 + 576$ bytes \times 33 + 15 \times 2 bytes = 888670 (= H'15A5E)

(6) Duplicate-Line Setting of Enlargement Resizing

When performing an enlargement resizing in the vertical direction, in some cases image data of the source image on the memory is used twice in succession by line. For this reason, the pixel data (in units of lines) of the lines used twice needs to be transferred by the CPU to the 2DG twice in succession.

(a) Subject conditions

Duplicate setting applies to the case where a source image, whose horizontal pixel count is larger than the bank size (64 pixels) of the buffer, is to be resized vertically for enlargement (the horizontal resizing ratio is not applicable). In other words, duplicate setting applies to the cases where

horizontal pixel count of source image ≥ 65 and vertical enlargement resizing.

Figures 26.45 and 26.46 illustrate examples.

Figure 26.45 Example of Duplicate-line Setting of Enlargement Resizing (1)

Figure 26.46 Example of Duplicate-line Setting of Enlargement Resizing (2)

(b) Determining a subject line (duplicate-line)

The line to be used twice in succession can be determined according to the following test formula:

[Test formulae]

 $b0 = INT$ ((VDLT_DCML $\times INT$ (SL \times DV/SL) + V1PHS_DCML + (VDLT_DCML \times (– 2)))/4096) $b1 = INT$ ((VDLT_DCML $\times INT$ (SL $\times DV/SL$) + V1PHS_DCML + (VDLT_DCML \times $(-1))$ /4096) $b2 = INT$ ((VDLT_DCML × INT (SL × DV/SL) + V1PHS_DCML + (VDLT_DCML × (0)))/4096) $b3 = INT$ ((VDLT_DCML $\times INT$ (SL $\times DV/SL$) + V1PHS_DCML + (VDLT_DCML \times (1)))/4096) IF (OR (AND ($b1 = (SL - 1)$, $b1 = b0$), AND $(b1 = (SL - 1), b1 = b2)$,

AND $(b2 = (SL - 1), b2 = b3)$, 2, 1)

[Legend]

- SV: Vertical line count of the source image (before resizing)
- DV: Vertical line count of the image after resizing
- SL: Line number of the source image (before resizing) to be determined
- VIPHS_DCML: Fractional part of the initial phase in the vertical direction (allows for settings less than delta)
- VDLT DCML: Fractional part of delta in the vertical direction
- Note: The line is transferred twice in succession if the computational result of the test formula is true (2) ; if it is false (1) , the line is transferred once.

Table 26.10 Examples of Determining Subject Lines (duplicate-lines)

Note: Since SL = 0 lines leads to exception handling, always set the transfer of a line at least once regardless of the multiplier.

An example shows to judge whether or not duplicate-line settings are required in the following case.

• Condition

Number of source pixels (SV): 100

Number of destination pixels (DV): 150

Fractional part of the result of calculating phase of vertical starting position (V1PHS_DCML): $H'500 = 1365$

Line number (SL): 78

• Items calculated by the user

Fractional part of the result of calculating the vertical delta (VDLT_DCML) Cv is calculated for use in obtaining VDLT_DCML. $CV = INT (SV/DV) \times 4096 = INT (100/150) \times 4096 = 2730.$ Since VDLT_DCML becomes the lower-order 12 bits of Cv, VDLT_DCML = 2730.

When the line number (SL) is 78, use the formula given earlier to discern whether duplicate-line settings are or are not necessary.

b0 = INT (2730 × INT (78 × 150/100) + 1365 + (2730 × (-2)))/4096 = 76 b1 = INT (2730 × INT (78 × 150/100) + 1365 + (2730 × (-1)))/4096 = 77 b2 = INT (2730 \times INT (78 \times 150/100) + 1365 + (2730 \times (0)))/4096 = 78 b3 = INT (2730 \times INT (78 \times 150/100) + 1365 + (2730 \times (1)))/4096 = 78 IF (OR (AND (77 = (78 – 1), 77 = 76),

AND
$$
(77 = (78 - 1), 77 = 78)
$$
,

AND
$$
(78 = (78 - 1), 78 = 78)
$$
, 2, 1)

Since all terms of the conditional expression are false, the transfer of line 78 is only to proceed once.

[Example of settings for pixel-data transfer]

As an example, the text below explains Example 1 in table 26.10.

In this case, make settings such that pixel data (in line units) is transferred from the CPU to the 2DG module in the following order.

Data for transfer (in line units) = 0, 1, 1, 2, 3, 3, 4, 5, 5, 6, 7, 7, 8, ...

In this way, lines 1, 3, 5, 7 etc. of pixel data (in line units) must each be transferred twice consecutively.

[Obtaining the number of lines for transfer in the case of duplicate-line settings]

Regardless of whether resizing is full or partial, the number of lines for transfer in the case of duplicate-line settings is always the total obtained from the conditional expression for duplicatelines. Examples for the cases of full or partial resizing are explained below. Furthermore, duplicate-line settings are made in the examples for the number of all horizontal pixels being 65 or more.

Example 1: Duplicate-line settings in the case of full resizing

• Conditions

Number of source lines: 20 Number of destination lines: 40 Multiplication: 2 times

In this case, when the conditional expression given earlier is applied to all lines, the results are as shown in figure 26.47. The number of lines for transfer as the total obtained from the conditional expression for line duplication is thus 39.

Figure 26.47 Duplicate-Line Settings for Full Resizing

Example 2: Duplicate-line settings in the case of partial resizing

Conditions

Number of source lines: 9 Number of destination lines: 18 Multiplication: 2 times Source starting position: Line 5

In this case, when the conditional expression given earlier is applied to all lines, the results are as shown in figure 26.48. The number of lines for transfer as the total obtained from the conditional expression for line duplication is thus 17.

Figure 26.48 Duplicate-line Setting for partial resizing

(c) How to set the register for duplicate-line setting

When lines are to be set in duplicate, the resize register in the 2DG must be set in the same manner as normal resizing settings. The source image size should be set in terms of the number of horizontal pixels and the number of lines that are actually transferred by the CPU to the 2DG. For the resize image size, set the number of horizontal pixels and the number of lines that are actually written to the memory by the 2DG. Examples are shown in figures 26.49 and 26.50.

Figure 26.49 Enlarging an image, both horizontally and vertically,

when the horizontal image has 64 pixels or less (normal setting (not duplicate setting))

Figure 26.50 Enlarging an image, both horizontally and vertically, when the horizontal image has 65 pixels or more (duplicate-line setting)

26.4.4 Output Operations

(1) Summary of Operations between the Output Block and External Memory

The following is a summary of operations between the output block and external memory.

- 1. The output block negates the DMA request signal and accepts a DMA transfer from the external memory.
- 2. The SE buffers (SE1, SE2), alternately buffer the data received through the DMA transfer.
- 3. Triggered by the VSYNC signal, the data undergoes various processings and then output to the panel unit.
- 4. Steps 1 to 3 are repeated until all data processing is completed.

Figure 26.51 Summary of Operations between Output Block and External Memory

(2) Pixel Format Conversion in Output Block

For the output block, either αRGB444 or αRGB555 can be set as the pixel format. The output block uses 6 bits for each color in blending involving the moving pictures. For this reason, a given format is converted into a standard format: α (4 bits) + RGB (6 bits each) for a total of 22 bits. The formula below shows rules for the conversion of a given format to the standard format:

 α RGB444 (AF83 (H)) converted into a standard format α : A (H) \rightarrow A (H) R: F (H) \rightarrow 3C (H) G: 8 (H) \rightarrow 20 (H) B: 3 (H) \rightarrow 0C (H)

Figure 26.52 Pixel Format Conversion in Output Block (1)

 α RGB555 (F599 (H)) converted into a standard format α : 1 (H) \rightarrow CHG A bit in register MGR MIXMODE R: 1D (H) \rightarrow 3A (H) G: 0C (H) \rightarrow 18 (H) B: 19 (H) \rightarrow 32 (H)

Figure 26.53 Pixel Format Conversion in Output Block (2)

(3) Summary of Output Block Operations

As a summary of output block operations, the text below provides an example where data equal to the specified number of pixels is transferred from output plane PX, which is written in an arbitrary memory space on the SDRAM, to the SE buffer using the DMAC, and the data is blended with moving pictures before being output.

The area for memory plane PX is set as follows: the number of lines in the SEHIGH bits of the MGR SESET register, and the number of pixels in the SEWIDH bits of the GR SESET register.

Figure 26.54 Summary of Output Block Operations

The SE buffer has a 960-byte double-buffer structure (SE1, SE2). For example, if the following values are assigned: SEWIDH bits $= 480$ (pixels) in the MGR_SESET register and SEHIGH bits $=$ 240 (lines) in the MGR_SESET register, and if the OUTEN bit $= 1$ in the GR_MIXPLY register to enable graphics display, output block operations work as follows:

- 1. Transfers the first 480 pixels to SE1 (SEHF_STAT $(0) = 1$), followed by VIVSYNC input, constant-rate output processing, and output to the panel unit.
- 2. Transfers the next 480 pixels to SE2 (SEHF_STAT $(1) = 1$), followed by VIVSYNC input, constant-rate output processing, and output to the panel unit.
- 3. Repeats transferring data to SE1 and SE2 up to the 240th line, and outputs interrupt signal INT_FILD at the 240th line.
- 4. The above steps 1 to 3 are repeated until the OUTEN bit of the GR_MIXPLY register is updated to 0.
- **(4) Resizing**

(a) How to set the delta value for use in resizing

The value of a resize delta (Ch) can be determined according to the following equation:

Ch = (source resolution/destination resolution) \times 4096

Note: Ch = 1/resizing ratio must be calculated from source pixel count/destination pixel count.

For example, if the source pixel count is 720 and the destination pixel count is 480, Ch will be:

 $Ch = (720 / 480) \times 4096 = 6144 = 1800$ (H)

Based on these results, the resize delta value will be as follows:

Since MHDLT INTGR is the integer part of Ch (4 bits), MHDLT INTGR = 1 (H) Since MHDLT_DCML is the fractional part of Ch (12 bits), MHDLT_DCML = 800 (H) If the resizing function is not to be used, set the MHDLT_INTGR bit and MHDLT_DCML bit to

H'1 and H'000, respectively.

(b) Howe to set the source-side starting phase for use in resizing

The source-side starting phase (Psh) can be determined according to the following equation:

Psh = Ch \times (starting pixel count) + (starting initial phase \times 4096)

In this case, MH1PHS_DCML will be the fractional part (12 bits) of Psh. If resizing is not performed or a resizing ratio $= 1$, $Psh = 0$ should be set. Note that the starting pixel count is always 0 (first pixel) for the output block, since the output block exclusively resizes moving pictures. Therefore, only the fractional part is necessary for Psh unlike for the blitter.

(5) Blending in Output Block

The following describes the blending in the output block. The blending processor in the output block supplies the output according to the following formula.

 $Cp = (Fc \times Cdc) + (Fd \times Cv)$

Here, Fc and Fd are set using the FCFD bits of the MGR_MIXMODE register. Table 26.11 shows the FCFD bit settings and the corresponding Fc and Fd values.

FCFD (Bit Values) Fc Fd Remarks 000 (Initial value) 1 $1 - \alpha$ dc When SE input image is premultiplied. 001 α dc 1 − α dc When SE input image is non-premultiplied. 010 1 0 Only graphics are output. 011 0 0 1 Only moving pictures are output. 100 0 0 0 Nothing is output. (Black screen) Other than the above \qquad \qquad \qquad Reserved

Table 26.11 Details of FCFD Bits of MGR_MIXMODE Register

Figure 26.55 Summary of Output Block Operations during Blending

26.4.5 Interrupts

2DG interrupt signals are classified into two types: the interrupts related to the blitter (BLT interrupts) and the interrupts related to the output block (output interrupts). Table 26.12 shows the interrupt sources and the conditions on which each source is generated and can be cleared. Figure 26.56 shows the structure of the interrupts.

Figure 26.56 Structure of Interrupts

The 2DG handles the interrupts differently for the following two cases.

(1) When the interrupt source is a DC buffer full, an SA buffer full, an SB buffer full, or an SE buffer full

- (1-1) An interrupt event occurs in the 2DG.
- (1-2) The INT_∗∗∗ and IRQ_∗∗∗ bits in the interrupt status register for graphics (GR_IRSTAT) are set accordingly (interrupt signal = negative logic).
- (1-3) The CPU recognizes the interrupt and reads the interrupt status register for graphics.
- (1-4) The CPU writes 1 to the interrupt reset control register for graphics (GR_INTDIS).
- (1-5) On reception of the value written in the above step, the IRQ_∗∗∗ bit in the GR_IRSTAT register is cleared (thus deasserting the interrupt signal).
- (1-6) Resetting of the response to the interrupt event proceeds within the 2DG.
- (1-7) The INT_∗∗∗ bit in the GR_IRSTAT register is cleared in response to the above step.

Figure 26.57 shows the flow of processing when SB buffer full.

Note that, within the interrupt status register for graphics (GR_IRSTAT), the CPU writing a 1 to a bit of the interrupt reset control register for graphics (in step (1-4) above) clears the bit in the case of IRQ_∗∗∗ bits but does not clear the bit in the case of INT_∗∗∗ bits.

If an interrupt event occurs and is the response within the 2DG is cleared (step (1-6) above) before the CPU has read the GR_ISTAT register (step (1-3) above), reading the value of the GR_ISTAT register will return the value of the register with the corresponding INT_∗∗∗ bit cleared.

Figure 26.57 Interrupt Handling (1)

- **(2) When the interrupt source is a completion of a blit operation, input of a VSYNC signal for the output block, and output underflow for the output block, or capture of a last line by the output block**
- (2-1) An interrupt event occurs in the 2DG.
- (2-2) The INT_∗∗∗ bit in the interrupt status register for graphics (GR_IRSTST) is set accordingly $(int$ errupt signal = negative logic).
- (2-3) The CPU recognizes the interrupt and reads the GR_IRSTAT register.
- (2-4) The CPU writes 1 to the interrupt reset control register for graphics (GR_INTDIS).
- (2-5) The IRQ_∗∗∗ bit in the GR_IRSTAT register is cleared in response to the above step (thus deasserting the interrupt signal).

Figure 26.58 shows the interrupt processing flow.

Figure 26.58 Interrupt Handling (2)

The 2DG interrupt signals are level-sensitive and more than one status bit is assigned to each interrupt signal. Because of this, it is necessary for the CPU to read the interrupt status register until the relevant interrupt signal is reset so that the CPU should recognize all the corresponding status bits and handle them according to the specified priority.

26.5 Appendix. VIDEO OUT (D/A Converter)

VIDEO OUT contains a built-in D/A converter (DAC) of the current output method using a current cell matrix D/A converter; it outputs a current corresponding to 6-bit digital input signals to analog output pins R, G, and B.

26.5.1 Analog Output Current

(1) Determination of the DAC Output Resistor (RL)

Since this DAC is a current output type unit, in order to convert a current into voltage, it requires a resistor (RL) external to the LSI. The RL can be set between 75 and 180 Ω based on an allowable current level, the external load capacity (CL), and a desired settling time (tset). The formula for calculating RL is given below:

Here, we define the time required from the beginning of a change in output to the time it converges to the level within $\pm 1.0\%$ of the final attainable level as settling time.

tset = $4.6 \times \text{RL} \times \text{C}$ tset: $\pm 1\%$ settling time [ns], C: total load capacity [pF] $C = C$ in + CL Cin: internal capacity of module (approximately 5 pF) [pF],

CL: I/O + PKG + on-board capacity [pF]

The condition that must be satisfied by RL in order to obtain the desired settling time will be:

 $RL < tset / {4.6 \times (Cin + CL)}$

Example: (if tset = 18 [ns], CL = 15 [pF])

RL < 18ns / $\{4.6 \times (5 \text{ pF} + 15 \text{ pF})\} = 195.6 \Omega$

Therefore, to settle it at 18 ns, set RL to 195.6 Ω or less (for example, 180 Ω).

(2) Determination of the DAC Output Current (I/O Max.)

This DAC is designed for a maximum output voltage of 1.0 V. Therefore, I/O max. can be calculated once RL is determined.

I/O max. = 1.0 / RL (if RL = 180 Ω, I/O max. = 5.6 mA)

Note: I/O max. should be designed for 13.4 mA or less. Exceeding this limit can cause decreases in performance and reliability.

(3) Determination of the Rext Resistance

As illustrated in figure 26.59, internally in this module the current cells are driven by means of a circuit that uses op amp.

When VCCA = 3.3 V, the "+" pin of the op amp is set so that it will be approximately 0.91 V. Since negative feedback is applied to the op amp, approximately 0.91 V also appears on the REXT pin.

Since each current cell is designed to act as a current mirror with respect to a load circuit for the op amp, reducing the external resistance Rext increases the output current from each current cell, and increases the output current from analog outputs R, G, and B. For full-scale output, the relationship between the output current I/O max. and the Rext resistance is given by the following equation:

Rext = {VCCA \times (4 / 15) + 0.03} \times (1023 / 32) / I/O max.

Figure 26.59 Current Cells and Analog Outputs

This module has a maximum analog output voltage of 1.0 V, as shown in figure 26.60.

The analog output voltage should be adjusted in terms of the value of Rext, to a level that fits in the range shown in figure 26.60. Exceeding this range can cause decreases in precision and reliability.

Figure 26.60 Allowable Analog Ouptut Voltage Range

26.5.2 Notes on Usage

(1) Power Supply Pins

Because the D/A output voltage is a micro-voltage with approximately 1 mV per step, any noise penetrating from the outside of the LSI must be minimized. For this purpose, a ceramic capacitor approximately 0.01 to 0.1 μ F must be placed between each power supply pin and the VSS as close to the LSI pins as possible. It is recommended to place at least one 10-μF capacitor between each power supply and VSSA (figure 26.61).

(2) Power Supply Separation

In order to supply as a noise-free voltage as possible to the analog power supply, the system's analog power should be supplied to this DAC's power supply (figure 26.61). Also, VSSA should be connected to the system's analog ground.

(3) CBU Pin

The CBU pin is used to connect the phase compensation capacitor for op amps. A capacitor should be connected between the CBU pin and VCCA. No other elements or circuits should be connected to CBU pin.

(4) REXT Pin

The REXT pin serves to connect an external resistance element that determines the module's reference current. A resistance element appropriate for the RL should be connected between the REXT pin and VSSA. Notice that any noise on this pin can significantly affect D/A conversion results.

When reviewing pin layout or designing the board, exercise care that the REXT pin does not cross, or run parallel to, DAC output or other digital signals.

26.5.3 Application Example

Figure 26.61 shows a DAC application example.

Figure 26.61 Application Example of DAC

Section 27 Pin Function Controller (PFC)

The pin function controller (PFC) is composed of registers that are used to select the functions of multiplexed pins and assign pins to be inputs or outputs. Tables 27.1 to 27.10 list the multiplexed pins of this LSI.

Table 27.1 Multiplexed Pins (Port A)

Setting of Mode Bits (PAnMD[3:0])

Table 27.2 Multiplexed Pins (Port B)

Table 27.3 Multiplexed Pins (Port C)

					0101
Function 1 (General I/O)	Function 2 (Related Module)	Function 3 (Related Module)	Function 4 (Related Module)	Function 5 (Related Module)	Function 6 (Related Module)
PC10 I/O (port)	DIRECTION output (ATAPI)	TEND1 output (DMAC)	FCE output (FLCTL)	PINT6 input (INTC)	
PC9 I/O (port)	IDERST# output (ATAPI)	DACK1 output (DMAC)	NAF7 I/O (FLCTL)	PINT5 input (INTC)	DACT1 output (DMAC)
PC8 I/O (port)	IDEINT input (ATAPI)	DREQ1 input (DMAC)	NAF6 I/O (FLCTL)	PINT4 input (INTC)	
PC7 I/O (port)	IDEIORDY input (ATAPI)		NAF5 I/O (FLCTL)	PINT3 input (INTC)	
PC6 I/O (port)	IDEIORD# output (ATAPI)		NAF4 I/O (FLCTL)	PINT2 input (INTC)	
PC5 I/O (port)	IDEIOWR# output (ATAPI)		NAF3 I/O (FLCTL)	PINT1 input (INTC)	$\hspace{0.1mm}-\hspace{0.1mm}$
PC4 I/O (port)	IODREQ input (ATAPI)	TIOC4A I/O (MTU2)	NAF2 I/O (FLCTL)	PINT0 input (INTC)	
PC3 I/O (port)	IODACK# output (ATAPI)	TCLKD input (MTU2)	NAF1 I/O (FLCTL)	IRQ3 input (INTC)	
PC2 I/O (port)	IDEA2 output (ATAPI)	TCLKC input (MTU2)	NAF0 I/O (FLCTL)	IRQ2 input (INTC)	
PC1 I/O (port)	IDEA1 output (ATAPI)	TCLKB input (MTU2)			
PC0 I/O (port)	IDEA0 output (ATAPI)	TCLKA input (MTU2)		IRQ0 input (INTC)	
	0000	0001	0010	0011 TIOC4D I/O (MTU2) TIOC4C I/O (MTU2) TIOC4B I/O (MTU2)	0100 FSC output (FLCTL) IRQ1 input (INTC) FOE output (FLCTL)

Setting of Mode Bits (PCnMD[3:0])

Setting of Mode Bits (PDnMD[3:0])

Table 27.5 Multiplexed Pins (Port E)

Setting of Mode Bits (PEnMD[3:0])

Table 27.6 Multiplexed Pins (Port F)

Setting of Mode Bits (PFnMD[3:0])

Table 27.7 Multiplexed Pins (Port G)

Setting of Mode Bits (PGnMD[3:0])

Table 27.8 Multiplexed Pins (Port H)

Table 27.9 Multiplexed Pins (Port J)

Setting of Mode Bits (PJnMD[3:0])

Table 27.10 Multiplexed Pins (Port K)

27.1 Features

- Functions for the multiplexed pins can be selected by setting the control registers.
- When the general I/O function or TIOC I/O function of the MTU2 is specified, the I/O direction should be selected by setting the corresponding I/O register.

27.2 Register Descriptions

The PFC has the following registers.

Table 27.11 Register Configuration

Notes: 1. In 8-bit access, the register can be read but cannot be written to.

2. The initial value depends on the operating mode of the LSI.

27.2.1 Port A I/O Register L (PAIORL)

PAIORL is a 16-bit readable/writable register that is used to set the pins on port A as inputs or outputs. The PA15IOR to PA0IOR bits correspond to the PA15 to PA0 pins, respectively. The setting of PAIORL is valid for the pins for which general I/O function or TIOC I/O function of the MTU2 is selected. PAIORL has no effect on the pins for which other function is selected. If a bit in PAIORL is set to 1, the corresponding pin on port A functions as output. If it is cleared to 0, the corresponding pin functions as input.

27.2.2 Port A Control Registers L1 to L4 (PACRL1 to PACRL4)

PACRL1 to PACRL4 are 16-bit readable/writable registers that are used to select the function of the multiplexed pins on port A. See table 27.1 for the multiplexed functions.

(1) Port A Control Register L4 (PACRL4)

Note: $*$ The initial value is 0000 in 16-bit mode (MD = 0), and 0001 in 32-bit mode (MD= 1).

(2) Port A Control Register L3 (PACRL3)

Note: $*$ The initial value is 0000 in 16-bit mode (MD = 0), and 0001 in 32-bit mode (MD= 1).

(3) Port A Control Register L2 (PACRL2)

Note: $*$ The initial value is 0000 in 16-bit mode (MD = 0), and 0001 in 32-bit mode (MD= 1).

(4) Port A Control Register L1 (PACRL1)

Note: $*$ The initial value is 0000 in 16-bit mode (MD = 0), and 0001 in 32-bit mode (MD= 1).

27.2.3 Port B I/O Register H (PBIORH)

PBIORH is a 16-bit readable/writable register that is used to set the pins on port B as inputs or outputs. The PB18IOR to PB16IOR bits correspond to the PB18 to PB16 pins, respectively. The setting of PBIORH is valid for the pins for which general I/O function is selected and has no effect on the pins for which other function is selected. If a bit in PBIORH is set to 1, the corresponding pin on port B functions as output. If it is cleared to 0, the corresponding pin functions as input.

Bits 15 to 3 in PBIORH are reserved. These bits are always read as 0. The write value should always be 0.

27.2.4 Port B I/O Register L (PBIORL)

PBIORL is a 16-bit readable/writable register that is used to set the pins on port B as inputs or outputs. The PB15IOR to PB0IOR bits correspond to the PB15 to PB0 pins, respectively. The setting of PBIORL is valid for the pins for which general I/O function is selected and has no effect on the pins for which other function is selected. If a bit in PBIORL is set to 1, the corresponding pin on port B functions as output. If it is cleared to 0, the corresponding pin functions as input.

27.2.5 Port B Control Registers H1, H2, L1 to L4 (PBCRH1, PBCRH2, PBCRL1 to PBCRL4)

PBCRH1, PBCRH2, and PBCRL1 to PBCRL4 are 16-bit readable/writable registers that are used to select the function of the multiplexed pins on port B. See table 27.2 for the multiplexed functions.

(1) Port B Control Register H2 (PBCRH2)

Notes: 1. Not initialized by a reset triggered by WDT overflow.

 2. To write to PBCRH2, write by 16-bit or 32-bit access such that the write value for bits 15 to 8 is H'A5 and that for bits 7 to 4 is 0.

(2) Port B Control Register H1 (PBCRH1)

(3) Port B Control Register L4 (PBCRL4)

(4) Port B Control Register L3 (PBCRL3)

(5) Port B Control Register L2 (PBCRL2)

(6) Port B Control Register L1 (PBCRL1)

Note: $*$ The initial value is 0001 in 16-bit mode (MD = 0), and 0000 in 32-bit mode (MD= 1).

27.2.6 Port C I/O Register L (PCIORL)

PCIORL is a 16-bit readable/writable register that is used to set the pins on port C as inputs or outputs. The PC10IOR to PC0IOR bits correspond to the PC10 to PC0 pins, respectively. The setting of PCIORL is valid for the pins for which general I/O function or TIOC I/O function of the MTU2 is selected. PCIORL has no effect on the pins for which other function is selected. If a bit in PCIORL is set to 1, the corresponding pin on port C functions as an output pin. If it is cleared to 0, the corresponding pin functions as an input pin.

Bits 15 to 11 in PCIORL are reserved. These bits are always read as 0. The write value should always be 0.

27.2.7 Port C Control Registers L1 to L3 (PCCRL1 to PCCRL3)

PCCRL1 to PCCRL3 are 16-bit readable/writable registers that are used to select the functions of the multiplexed pins on port C. See table 27.3 for the multiplexed functions.

(1) Port C Control Register L3 (PCCRL3)

(2) Port C Control Register L2 (PCCRL2)

(3) Port C Control Register L1 (PCCRL1)

27.2.8 Port D I/O Register L (PDIORL)

PDIORL is a 16-bit readable/writable register that is used to set the pins on port D as inputs or outputs. The PD2IOR to PD0IOR bits correspond to the PD2 to PD0 pins, respectively. The setting of PDIORL is valid for the pins for which general I/O function is selected and has no effect on the pins for which other function is selected. If a bit in PDIORL is set to 1, the corresponding pin on port D functions as an output. If it is cleared to 0, the corresponding pin functions as an input.

Bits 15 to 3 in PDIORL are reserved. These bits are always read as 0. The write value should always be 0.

27.2.9 Port D Control Register L1 (PDCRL1)

PDCRL1 is 16-bit readable/writable register that is used to select the functions of the multiplexed pins on port D. See table 27.4 for the multiplexed functions.

27.2.10 Port E I/O Register L (PEIORL)

PEIORL is 16-bit readable/writable register that is used to set the pins on port E as inputs or outputs. The PE13IOR, PE11IOR, PE9IOR, and PE7IOR to PE0IOR bits correspond to the PE13, PE11, PE9, and PE7 to PE0 pins respectively. The setting of PEIORL is valid for the pins for which general I/O function or TIOC I/O function of the MTU2 is selected. PEIORL has no effect on the pins for which other function is selected. If a bit in PEIORL is set to 1, the corresponding pin on port E functions as an output pin. If it is cleared to 0, the corresponding pin functions as an input pin.

Bits 15, 14, 12, 10, and 8 in PEIORL are reserved. These bits are always read as 0. The write value should always be 0.

27.2.11 Port E Control Registers L1 to L4 (PECRL1 to PECRL4)

PECRL1 to PECRL4 are 16-bit readable/writable registers that are used to select the functions of the multiplexed pins on port E. See table 27.5 for the multiplexed functions.

(1) Port E Control Register L4 (PECRL4)

(2) Port E Control Register L3 (PECRL3)

(3) Port E Control Register L2 (PECRL2)

(4) Port E Control Register L1 (PECRL1)

27.2.12 Port F I/O Register L (PFIORL)

PFIORL is a 16-bit readable/writable register that is used to set the pins on port F as inputs or outputs. The PF4IOR to PF0IOR bits correspond to the PF4 to PF0 pins, respectively. PFIORL is enabled when the port F pins are functioning as general-purpose inputs/outputs. In other states, PFIORL is disabled. If a bit in PFIORL is set to 1, the corresponding pin on port F functions as an output. If it is cleared to 0, the corresponding pin functions as an input.

Bits 15 to 5 in PFIORL are reserved. These bits are always read as 0. The write value should always be 0.

27.2.13 Port F Control Registers L1, L2 (PFCRL1, PFCRL2)

PFCRL1 and PFCRL2 are 16-bit readable/writable registers that are used to select the functions of the multiplexed pins on port F. See table 27.6 for the multiplexed functions.

(1) Port F Control Register L2 (PFCRL2)

(2) Port F Control Register L1 (PFCRL1)

27.2.14 Port G Control Registers L1, L2 (PGCRL1, PGCRL2)

PGCRL1 and PGCRL2 are 16-bit readable/writable registers that are used to select the functions of the multiplexed pins on port G. See table 27.7 for the multiplexed functions.

(1) Port G Control Register L2 (PGCRL2)

(2) Port G Control Register L1 (PGCRL1)

27.2.15 Port H I/O Register L (PHIORL)

PHIORL is 16-bit readable/writable register that is used to set the pins on port H as inputs or outputs. The PH15IOR to PH0IOR bits correspond to the PH15 to PH0 pins respectively. The setting of PHIORL is valid for the pins for which general I/O function is selected and has no effect on the pins for which other function is selected. If a bit in PHIORL is set to 1, the corresponding pin on port H functions as an output pin. If it is cleared to 0, the corresponding pin functions as an input pin.

27.2.16 Port H Control Registers L1 to L4 (PHCRL1 to PHCRL4)

PHCRL1 to PHCRL4 are 16-bit readable/writable registers that are used to select the functions of the multiplexed pins on port H. See table 27.8 for the multiplexed functions.

(1) Port H Control Register L4 (PHCRL4)

(2) Port H Control Register L3 (PHCRL3)

(3) Port H Control Register L2 (PHCRL2)

(4) Port H Control Register L1 (PHCRL1)

27.2.17 Port J I/O Register L (PJIORL)

PJIORL is 16-bit readable/writable register that is used to set the pins on port J as inputs or outputs. The PJ12IOR to PJ0IOR bits correspond to the PJ12 to PJ0 pins respectively. The setting of PJIORL is valid for the pins for which general I/O function or TIOC I/O function of the MTU2 is selected. PJIORL has no effect on the pins for which other function is selected. If a bit in PJIORL is set to 1, the corresponding pin on port J functions as an output pin. If it is cleared to 0, the corresponding pin functions as an input pin.

Bits 15 to 13 in PJIORL are reserved. These bits are always read as 0. The write value should always be 0.

27.2.18 Port J Control Registers L1 to L4 (PJCRL1 to PJCRL4)

PJCRL1 to PJCRL4 are 16-bit readable/writable registers that are used to select the functions of the multiplexed pins on port J. See table 27.9 for the multiplexed functions.

(1) Port J Control Register L4 (PJCRL4)

(2) Port J Control Register L3 (PJCRL3)

(3) Port J Control Register L2 (PJCRL2)

(4) Port J Control Register L1 (PJCRL1)

27.2.19 Port K I/O Register L (PKIORL)

PKIORL is 16-bit readable/writable register that is used to set the pins on port K as inputs or outputs. The PK11IOR to PK0IOR bits correspond to the PK11 to PK0 pins respectively. The setting of PKIORL is valid for the pins for which general I/O function is selected and has no effect on the pins for which other function is selected. If a bit in PKIORL is set to 1, the corresponding pin on port K functions as an output pin. If it is cleared to 0, the corresponding pin functions as an input pin.

Bits 15 to 2 in PKIORL are reserved. These bits are always read as 0. The write value should always be 0.

27.2.20 Port K Control Register L1 (PKCRL1)

PKCRL1 is 16-bit readable/writable register that is used to select the functions of the multiplexed pins on port K. See table 27.10 for the multiplexed functions.

27.3 Usage Notes

The multiplexed pins listed in tables 27.1 to 27.10 except for pins PE8 to PE13, PF0, PF1, and PG0 to PG7 are provided with weak keepers or pull-up circuit (PB18) in their I/O buffers to prevent the pins from floating into intermediate voltage levels. However, note that the voltage retained in the high-impedance state may fluctuate due to noise.

Section 28 I/O Ports

This LSI has ten ports: A to H, J, and K.

All port pins are multiplexed with other pin functions. The functions of the multiplex pins are selected by means of the pin function controller (PFC).

Each port is provided with data registers for storing the pin data and port registers for reading the states of the pins.

28.1 Features

- 1. Total number of port pins: 107 pins (I/O: 96 pins, Input: 11 pins)
	- \rightarrow Port A: (Input: 16 pins)
	- $-$ Port B: (I/O: 19 pins)
	- $-$ Port C: (I/O: 11 pins)
	- $-$ Port D: (I/O: 3 pins)
	- \rightarrow Port E: (I/O: 11 pins, Input: 3 pins)
	- $-$ Port F: (I/O: 5 pins)
	- ⎯ Port G: (Input: 8 pins)
	- $-$ Port H: (I/O: 16 pins)
	- $-$ Port J: (I/O: 13 pins)
	- $-$ Port K: (I/O: 2 pins)
- 2. Pins with a weak keeper

The following pins of this LSI have a weak keeper circuit or pull-up circuit (PB18) that prevents the pin from floating into intermediate voltage levels.

- $-$ Port A: PA0 to PA15
- ⎯ Port B: PB0 to PB18
- Port C: PC0 to PC10
- Port D: PD0 to PD2
- Port E: PE0 to PE7
- Port F: PF2 to PF4
- $-$ Port H: PH0 to PH15
- Port J: PJ0 to PJ12
- Port K: PK0 and PK1

The I/O pins include a weak keeper circuit or pull-up circuit that fixes the input level high or low when the I/O pins are not driven from outside. Generally in the CMOS products, input levels on unused input pins must be fixed by way of external pull-up or pull-down resistors. However, the I/O pins having a weak keeper circuit or pull-up circuit of this LSI can eliminate these outer circuits and reduce the number of parts in the system. If the pull-up or pull-down resistors is necessary to fix the pin level, use the resistor of $10 \text{ k}\Omega$ or smaller.

28.2 Register Descriptions

The port has the following registers.

Table 28.1 Register Configuration

28.2.1 Port A Data Register L (PADRL)

PADRL is a 16-bit readable/writable register that stores port A data. The PA15DR to PA0DR bits correspond to the PA15 to PA0 pins, respectively.

When a pin function is general output, if a value is written to PADRL, that value is output from the pin, and if PADRL is read, the register value is returned regardless of the pin state.

When a pin function is general input, if PADRL is read, the pin state, not the register value, is returned directly. If a value is written to PADRL, although that value is written into PADRL, it does not affect the pin state. Table 28.2 summarizes PADRL read/write operations.

Table 28.2 Port A Data Register L (PADRL) Read/Write Operation

• Bits 15 to 0 of PADRL

28.2.2 Port A Port Register L (PAPRL)

PAPRL is a 16-bit read-only register, in which the PA15PR to PA0PR bits correspond to the PA15 to PA0 pins, respectively. PAPRL always returns the states of the pins regardless of the PFC setting.

28.2.3 Port B Data Registers H, L (PBDRH, PBDRL)

PBDRH and PBDRL are 16-bit readable/writable registers that store port B data. The PB18DR to PB0DR bits correspond to the PB18 to PB0 pins, respectively.

When a pin function is general output, if a value is written to PBDRH or PBDRL, that value is output from the pin, and if PBDRH or PBDRL is read, the register value is returned regardless of the pin state.

When a pin function is general input, if PBDRH or PBDRL is read, the pin state, not the register value, is returned directly. If a value is written to PBDRH or PBDRL, although that value is written into PBDRH or PBDRL, it does not affect the pin state. Table 28.3 summarizes PBDRH/PBDRL read/write operation.

(1) Port B Data Register H (PBDRH)

(2) Port B Data Register L (PBDRL)

Table 28.3 Port B Data Registers H, L (PBDRH, PBDRL) Read/Write Operations

• Bits 2 to 0 of PBDRH and Bits 15 to 0 of PBDRL

28.2.4 Port B Port Registers H, L (PBPRH, PBPRL)

PBPRH and PBPRL are 16-bit read-only registers, in which the PB18PR to PB0PR bits correspond to the PB18 to PB0 pins, respectively. PBPRH and PBPRL always return the states of the pins regardless of the PFC setting.

(1) Port B Port Register H (PBPRH)

(2) Port B Port Register L (PBPRL)

28.2.5 Port C Data Register L (PCDRL)

PCDRL is a 16-bit readable/writable register that stores port C data. The PC10DR to PC0DR bits correspond to the PC10 to PC0 pins, respectively.

When a pin function is general output, if a value is written to PCDRL, that value is output from the pin, and if PCDRL is read, the register value is returned regardless of the pin state.

When a pin function is general input, if PCDRL is read, the pin state, not the register value, is returned directly. If a value is written to PCDRL, although that value is written into PCDRL, it does not affect the pin state. Table 28.4 summarizes PCDRL read/write operations.

Table 28.4 Port C Data Register L (PCDRL) Read/Write Operations

• Bits 10 to 0 of PCDRL

28.2.6 Port C Port Register L (PCPRL)

PCPRL is a 16-bit read-only register, in which the PC10PR to PC0PR bits correspond to the PC10 to PC0 pins, respectively. PCPRL always returns the states of the pins regardless of the PFC setting.

28.2.7 Port D Data Register L (PDDRL)

PDDRL is a 16-bit readable/writable register that stores port D data. The PD2DR to PD0DR bits correspond to the PD2 to PD0 pins, respectively.

When a pin function is general output, if a value is written to PDDRL, that value is output from the pin, and if PDDRL is read, the register value is returned regardless of the pin state.

When a pin function is general input, if PDDRL is read, the pin state, not the register value, is returned directly. If a value is written to PDDRL, although that value is written into PDDRL, it does not affect the pin state. Table 28.5 summarizes PDDRL read/write operation.

Table 28.5 Port D Data Register L (PDDRL) Read/Write Operation

• Bits 2 to 0 of PDDRL

28.2.8 Port D Port Register L (PDPRL)

PDPRL is a 16-bit read-only register, in which the PD2PR to PD0PR bits correspond to the PD2 to PD0 pins, respectively. PDPRL always returns the states of the pins regardless of the PFC setting.

28.2.9 Port E Data Register L (PEDRL)

PEDRL is a 16-bit readable/writable register that stores port E data. The PE13DR to PE0DR bits correspond to the PE13 to PE0 pins, respectively.

When a pin function is general output, if a value is written to PEDRL, that value is output from the pin, and if PEDRL is read, the register value is returned regardless of the pin state.

When a pin function is general input, if PEDRL is read, the pin state, not the register value, is returned directly. If a value is written to PEDRL, although that value is written into PEDRL, it does not affect the pin state. Table 28.6 summarizes PEDRL read/write operation.

Note: * Depends on the state of the external pin.

Table 28.6 Port E Data Register L (PEDRL) Read/Write Operation

• Bits 13, 11, 9, 7 to 0 of PEDRL

• Bits 12, 10, 8 of PEDRL

28.2.10 Port E Port Register L (PEPRL)

PEPRL is a 16-bit read-only register, in which the PE13PR to PE0PR bits correspond to the PE13 to PE0 pins, respectively. PEPRL always returns the states of the pins regardless of the PFC setting.

28.2.11 Port F Data Register L (PFDRL)

PFDRL is a 16-bit readable/writable register that stores port F data. The PF4DR to PF0DR bits correspond to the PF4 to PF0 pins, respectively.

When a pin function is general output, if a value is written to PFDRL, that value is output from the pin, and if PFDRL is read, the register value is returned regardless of the pin state.

When a pin function is general input, if PFDRL is read, the pin state, not the register value, is returned directly. If a value is written to PFDRL, although that value is written into PFDRL, it does not affect the pin state. Table 28.7 summarizes PFDRL read/write operations.

Table 28.7 Port F Data Register L (PFDRL) Read/Write Operation

• Bits 4 to 0 of PFDRL

28.2.12 Port F Port Register L (PFPRL)

PFPRL is a 16-bit read-only register, in which the PF4PR to PF0PR bits correspond to the PF4 to PF0 pins, respectively. PFPRL always returns the states of the pins regardless of the PFC setting.

28.2.13 Port G Data Register L (PGDRL)

PGDRL is a 16-bit read-only register that stores port G data. The PG7DR to PG0DR bits correspond to the PG7 to PG0 pins, respectively. The general input function of the PG7 to PG0 pins is enabled only when the A/D and D/A converters are halted.

Writing to these bits is ignored, and therefore does not affect the pin state. If these bits are read, the pin state, not the bit value, is directly returned. However, a fixed value is returned for pins for which the A/D or D/A converter function is selected. Table 28.8 summarizes PGDRL read/write operation.

Note: * Depends on the state of the external pin.

Table 28.8 Port G Data Register L (PGDRL) Read/Write Operation

• Bits 7 to 0 of PGDRL

 $m = 1, 0$

28.2.14 Port H Data Register L (PHDRL)

PHDRL is a 16-bit readable/writable register that stores port H data. The PH15DR to PH0DR bits correspond to the PH15 to PH0 pins, respectively.

When a pin function is general output, if a value is written to PHDRL, that value is output from the pin, and if PHDRL is read, the register value is returned regardless of the pin state.

When a pin function is general input, if PHDRL is read, the pin state, not the register value, is returned directly. If a value is written to PHDRL, although that value is written into PHDRL, it does not affect the pin state. Table 28.9 summarizes PHDRL read/write operations.

Table 28.9 Port H Data Register L (PHDRL) Read/Write Operation

• Bits 15 to 0 of PHDRL

28.2.15 Port H Port Register L (PHPRL)

PHPRL is a 16-bit read-only register, in which the PH15PR to PH0PR bits correspond to the PH15 to PH0 pins, respectively. PHPRL always returns the states of the pins regardless of the PFC setting.

28.2.16 Port J Data Register L (PJDRL)

PJDRL is a 16-bit readable/writable register that stores port J data. The PJ12DR to PJ0DR bits correspond to the PJ12 to PJ0 pins, respectively.

When a pin function is general output, if a value is written to PJDRL, that value is output from the pin, and if PJDRL is read, the register value is returned regardless of the pin state.

When a pin function is general input, if PJDRL is read, the pin state, not the register value, is returned directly. If a value is written to PJDRL, although that value is written into PJDRL, it does not affect the pin state. Table 28.10 summarizes PJDRL read/write operations.

Table 28.10 Port J Data Register L (PJDRL) Read/Write Operations

• Bits 12 to 0 of PJDRL

28.2.17 Port J Port Register L (PJPRL)

PJPRL is a 16-bit read-only register, in which the PJ12PR to PJ0PR bits correspond to the PJ12 to PJ0 pins, respectively. PJPRL always returns the states of the pins regardless of the PFC setting.

28.2.18 Port K Data Register L (PKDRL)

PKDRL is a 16-bit readable/writable register that stores port K data. The PK1DR and PK0DR bits correspond to the PK1 and PK0 pins, respectively.

When a pin function is general output, if a value is written to PKDRL, that value is output from the pin, and if PKDRL is read, the register value is returned regardless of the pin state.

When a pin function is general input, if PKDRL is read, the pin state, not the register value, is returned directly. If a value is written to PKDRL, although that value is written into PKDRL, it does not affect the pin state. Table 28.11 summarizes PKDRL read/write operations.

Table 28.11 Port K Data Register L (PKDRL) Read/Write Operations

• Bits 1 and 0 of PKDRL

28.2.19 Port K Port Register L (PKPRL)

PKPRL is a 16-bit read-only register, in which the PK1PR and PK0PR bits correspond to the PK1 and PK0 pins, respectively. PKPRL always returns the states of the pins regardless of the PFC setting.

28.3 Usage Notes

When the PFC has been configured to select the following pin functions, the pin state cannot be read from the data registers or port registers.

• D31 to D16 (data bus)

Section 29 On-Chip RAM

This LSI has an on-chip high-speed RAM, which achieves fast access, and an on-chip RAM for data retention, which can retain data even in deep standby mode. These memory units can store instructions or data.

The memory operation and writing operation of the on-chip high-speed RAM can be enabled or disabled through the RAM enable bits and RAM write enable bits.

For the on-chip RAM for data retention, it is possible to specify whether to retain data in deep standby mode on a page-by-page basis.

29.1 Features

• Pages

On-chip high-speed RAM0 consists of four pages (pages 0, 1, 2, and 3) and on-chip high-speed RAM1 consists of two pages (pages 0 and 1) and each of these pages has a size of 16 Kbytes. The on-chip RAM for data retention consists of four pages and each of these pages has a size of 4 Kbytes.

• Memory map

The on-chip high-speed RAMs are allocated in the address space shown in tables 29.1 and 29.2. The on-chip RAM for data retention is allocated in the address space shown in table 29.3. When a common area for CPU0 and CPU1 is placed on the on-chip high-speed RAMs and the area is exclusively accessed by the TAS.B instruction, the on-chip high-speed RAMs should be accessed through the address space shown in table 29.2.

Table 29.1 Address Spaces of On-Chip High-Speed RAM

Table 29.2 Address Spaces of On-Chip High-Speed RAM (Shadow Spaces)

Table 29.3 Address Spaces of On-Chip RAM for Data Retention

• Port

On-chip high-speed RAM0 is connected to the CPU0 instruction fetch bus, CPU0 memory access bus, and on-chip high-speed RAM0 access bus. When CPU0 accesses on-chip highspeed RAM0 through the address spaces shown in table 29.1, the CPU0 instruction fetch bus or CPU0 memory access bus is used. When CPU0 accesses on-chip high-speed RAM0 through the address spaces shown in table 29.2, the on-chip high-speed RAM0 access bus is used. When CPU1 or DMAC accesses on-chip high-speed RAM0, the on-chip high-speed RAM0 access bus is used in access through the both address spaces shown in table 29.1 and table 29.2.

On-chip high-speed RAM1 is connected to the CPU1 instruction fetch bus, CPU1 memory access bus, and on-chip high-speed RAM1 access bus. When CPU1 accesses on-chip highspeed RAM1 through the address space shown in table 29.1, the CPU1 instruction fetch bus or CPU1 memory access bus is used. When CPU1 accesses on-chip high-speed RAM1 through the address space shown in table 29.2, the on-chip high-speed RAM1 access bus is used. When CPU0 or DMAC accesses on-chip high-speed RAM1, the on-chip high-speed RAM1 access bus is used in access through the both address spaces shown in table 29.1 and table 29.2.

Each page of the on-chip RAM for data retention has one read and write port and is connected to the peripheral bus.

• Priority

When the same page of the on-chip high-speed RAM is simultaneously accessed from different buses, the access is controlled based on the priority. The priority order is on-chip high-speed RAM access bus > memory access bus > instruction fetch bus.

29.2 Usage Notes

29.2.1 Page Conflict

When the same page of the on-chip high-speed RAM is simultaneously accessed from different buses, a page conflict occurs. Although each access is completed correctly, such a conflict degrades the memory access speed. Therefore, it is advisable to provide software measures to prevent such a conflict. Different pages, instead of the same page, can be simultaneously accessed from different buses.

29.2.2 RAME Bit and RAMWE Bit

Before clearing the RAME bits of SYSCR1, SYSCR3, SYSCR5, SYSCR7, SYSCR9, and SYSCR11 to 0 and the RAMWE bits of SYSCR2, SYSCR4, SYSCR6, SYSCR8, SYSCR10, and SYSCR12 to 0, be sure to read and write the arbitrarily-selected same address in each page. Otherwise, the data last written to the corresponding page may not be actually written to the RAM.

Figure 29.1 Examples of Read/Write

29.2.3 Data Retention

Data in the on-chip high-speed RAM and including on-chip data retention RAM are retained in the states other than power-on reset and deep standby mode. In power-on reset and deep standby mode, these RAMs operate as described below.

(1) Power-on Reset

(a) On-Chip High-Speed RAM

Data are retained on a power-on reset by disabling the setting of either the RAME or RAMWE bit.

Data are not retained when the setting of the RAME and RAMWE bits are both enabled.

(b) On-Chip Data Retention RAM

Data are not retained.

(2) Deep Standby Mode

(a) On-Chip High-Speed RAM

Data are not retained.

(b) On-Chip Data Retention RAM

Data are retained in deep standby mode by enabling the setting of the RRAMKP bit, excluding the case that deep standby mode is canceled by power-on reset. In the case that deep standby mode is canceled by interrupt or manual reset pins, power-on reset exception handling is executed, but the data are retained.

Section 30 Power-Down Modes

This LSI supports single-processor mode, dual-sleep mode, software standby mode, deep standby mode, and module standby function as power-down modes. In power-down modes, functions of the CPU, clocks, on-chip memory, or part of on-chip peripheral modules are halted or the power supply is turned off, through which power consumption is reduced. These modes can be exited by a reset or interrupt.

30.1 Power-Down Modes

This LSI has the following power-down modes and function:

- 1. Dual-processor mode
- 2. Single-processor mode (single-processor 0 mode, single-processor 1 mode)
- 3. Dual-sleep mode
- 4. Software standby mode
- 5. Deep standby mode
- 6. Module standby function

Table 30.1 shows the transition conditions for entering the modes from the program execution state, as well as the CPU and peripheral module states in each mode and the procedures for canceling each mode.

Notes: 1. The pin state is retained or set to high impedance. For details, see appendix A, Pin States.

- 2. By specifying the module standby function, individual on-chip peripheral modules (including the RTC) can be halted. To specify the module standby function, set the corresponding MSTP bit in STBCR2 to STBCR7 to 1. To cancel the module standby function, clear the MSTP bit to 0. For the H-UDI and UBC, the module standby function can also be canceled by power-on reset.
- 3. RTC operates when the START bit in the RCR2 register is set to 1. For details, see section 15, Realtime Clock (RTC). When deep standby mode is canceled by a poweron reset, the running state cannot be retained. Make the initial setting for the realtime clock again.
- 4. Setting the bits RRAMKP3 to RRAMKP0 in the RRAMKP register to 1 enables to retain the data in the corresponding area on the on-chip RAM (for data retention) during the transition to deep standby mode. However, the stored contents are initialized when deep standby mode is canceled by a power-on reset.
- 5. Deep standby mode can be canceled by an interrupt (NMI or IRQ) or a reset (manual reset or power-on reset). However, when deep standby mode is canceled by the NMI interrupt or IRQ interrupt, power-on reset exception handling is executed instead of interrupt exception handling. The power-on reset exception handling is executed also in the cancellation of deep standby mode by manual reset.

- 6. The stored contents are initialized when software standby mode is canceled by a power-on reset.
- 7. Clearing the RAME bits in registers SYSCR1, SYSCR3, and SYSCR5 or the RAMWE bits in registers SYSCR2, SYSCR4, and SYSCR6 to 0 (disabled) enables retention of the data in high-speed on-chip RAM0 when software standby mode is canceled by a power-on reset.
- 8. Clearing the RAME bits in registers SYSCR7, SYSCR9, and SYSCR11 or the RAMWE bits in registers SYSCR8, SYSCR10, and SYSCR12 to 0 (disabled) enables retention of the data in high-speed on-chip RAM1 when software standby mode is canceled by a power-on reset.

30.2 Register Descriptions

The following registers are used in power-down modes.

Table 30.2 Register Configuration

30.2.1 Standby Control Register 1 (STBCR1)

STBCR1 is an 8-bit readable/writable register that specifies the states of power-down modes.

Note: When writing to this register, see section 30.4, Usage Notes.

x: Don't care

30.2.2 Standby Control Register 2 (STBCR2)

STBCR2 is an 8-bit readable/writable register that controls the operation of individual modules.

30.2.3 Standby Control Register 3 (STBCR3)

STBCR3 is an 8-bit readable/writable register that controls the operation of modules.

30.2.4 Standby Control Register 4 (STBCR4)

STBCR4 is an 8-bit readable/writable register that controls the operation of modules.

30.2.5 Standby Control Register 5 (STBCR5)

STBCR5 is an 8-bit readable/writable register that controls the operation of modules.

30.2.6 Standby Control Register 6 (STBCR6)

STBCR6 is an 8-bit readable/writable register that controls the operation of modules.

30.2.7 Standby Control Register 7 (STBCR7)

STBCR7 is an 8-bit readable/writable register that controls the operation of modules.

30.2.8 System Control Register 1 (SYSCR1)

SYSCR1 is an 8-bit readable/writable register that enables or disables access (read/write) from CPU0 to each page of the high-speed on-chip RAM0.

Setting the RAMEn $(n = 0 \text{ to } 3)$ bit in SYSCR1 to 1 enables access to page n. Clearing the RAMEn bit to 0 disables access to page n. In this case, an undefined value is returned when reading data or fetching an instruction from page n, and writing to page n is ignored. The initial value of the RAMEn bit is 1.

When clearing the RAMEn bit to 0, be sure to execute instructions to read from and write to the same arbitrary address in page n before clearing the RAMEn bit. If not executed, the data last written to page n may not be actually written to the high-speed on-chip RAM.

SYSCR1 should be set by the program that is placed in a space other than the high-speed on-chip RAM space. Furthermore, an instruction to read SYSCR1 should be located immediately after the instruction to write to SYSCR1. Otherwise, normal access to the high-speed on-chip RAM is not guaranteed.

Note: * For the addresses of each page, see section 29, On-Chip RAM.

30.2.9 System Control Register 2 (SYSCR2)

SYSCR2 is an 8-bit readable/writable register that enables or disables writing from CPU0 to each page of the high-speed on-chip RAM0.

Setting the RAMWEn $(n = 0 to 3)$ bit in SYSCR2 to 1 enables writing to page n. When the RAMWEn bit is cleared to 0, writing to page n is ignored. The initial value of the RAMWEn bit is 1.

When clearing the RAMWEn bit to 0, be sure to execute an instruction to read from and write to the same arbitrary address in page n before setting the RAMWEn bit. If not executed, the data last written to page n may not be written to the high-speed on-chip RAM.

Set SYSCR2 using the program that is placed in a space other than the high-speed on-chip RAM space. Furthermore, an instruction to read SYSCR2 should be located immediately after the instruction to write to SYSCR2. Otherwise, normal access to the high-speed on-chip RAM is not guaranteed.

Note: When writing to this register, see section 30.4, Usage Notes.

Note: * For the addresses of each page, see section 29, On-Chip RAM.

30.2.10 System Control Register 3 (SYSCR3)

SYSCR3 is an 8-bit readable/writable register that enables or disables access (read/write) from CPU1 to each page of the high-speed on-chip RAM0. Other descriptions on this register are the same as SYSCR1.

Note: When writing to this register, see section 30.4, Usage Notes.

30.2.11 System Control Register 4 (SYSCR4)

SYSCR4 is an 8-bit readable/writable register that enables or disables writing from CPU1 to each page of the high-speed on-chip RAM0. Other descriptions on this register are the same as SYSCR2.

Note: When writing to this register, see section 30.4, Usage Notes.

30.2.12 System Control Register 5 (SYSCR5)

SYSCR5 is an 8-bit readable/writable register that enables or disables access (read/write) from the DMAC to each page of the high-speed on-chip RAM0. Other descriptions on this register are the same as SYSCR1.

Note: When writing to this register, see section 30.4, Usage Notes.

30.2.13 System Control Register 6 (SYSCR6)

SYSCR6 is an 8-bit readable/writable register that enables or disables writing from the DMAC to each page of the high-speed on-chip RAM0. Other descriptions on this register are the same as SYSCR2.

30.2.14 System Control Register 7 (SYSCR7)

SYSCR7 is an 8-bit readable/writable register that enables or disables access (read/write) from CPU0 to each page of the high-speed on-chip RAM1. Other descriptions on this register are the same as SYSCR1.

Note: When writing to this register, see section 30.4, Usage Notes.

Note: * For the addresses of each page, see section 29, On-Chip RAM.

30.2.15 System Control Register 8 (SYSCR8)

SYSCR8 is an 8-bit readable/writable register that enables or disables writing from CPU0 to each page of the high-speed on-chip RAM1. Other descriptions on this register are the same as SYSCR2.

Note: When writing to this register, see section 30.4, Usage Notes.

Note: * For the addresses of each page, see section 29, On-Chip RAM.

30.2.16 System Control Register 9 (SYSCR9)

SYSCR9 is an 8-bit readable/writable register that enables or disables access (read/write) from CPU1 to each page of the high-speed on-chip RAM1. Other descriptions on this register are the same as SYSCR1.

Note: When writing to this register, see section 30.4, Usage Notes.

30.2.17 System Control Register 10 (SYSCR10)

SYSCR10 is an 8-bit readable/writable register that enables or disables writing from CPU1 to each page of the high-speed on-chip RAM1. Other descriptions on this register are the same as SYSCR2.

Note: When writing to this register, see section 30.4, Usage Notes.

30.2.18 System Control Register 11 (SYSCR11)

SYSCR11 is an 8-bit readable/writable register that enables or disables access (read/write) from the DMAC to each page of the high-speed on-chip RAM1. Other descriptions on this register are the same as SYSCR1.

Note: When writing to this register, see section 30.4, Usage Notes.

30.2.19 System Control Register 12 (SYSCR12)

SYSCR12 is an 8-bit readable/writable register that enables or disables writing from the DMAC to each page of the high-speed on-chip RAM1. Other descriptions on this register are the same as SYSCR2.

30.2.20 Software Reset Control Register (SWRSTCR)

SWRSTCR is an 8-bit readable/writable register that controls the software reset of the SSIF0 to SSIF5 and the IEB.

30.2.21 High-Impedance Control Register (HIZCR)

HIZCR is an 8-bit readable/writable register that selects whether to retain the pin state or set it to high impedance in software standby mode or deep standby mode.

30.2.22 CPU0/CPU1 Mode Status Registers (C0MSR, C1MSR)

C0MSR and C1MSR are 8-bit read-only registers that indicate current operating mode of respective processors. Writing to these registers is ignored.

• C0MSR

• C1MSR

30.2.23 Data Retention On-Chip RAM Area Specification Register (RRAMKP)

RRAMKP is an 8-bit readable/writable register specifies whether to retain the contents of the corresponding area of the on-chip RAM (for data retention) in deep standby mode.

When the RRAMKP bit is set to 1, the contents of the corresponding area of the on-chip RAM are retained in deep standby mode. When the bit is cleared to 0, the contents of the corresponding area of the on-chip RAM are not retained in deep standby mode.

Note: * For the addresses of each page, see section 29, On-Chip RAM.

30.2.24 Deep Standby Control Register (DSCTR)

DSCTR is an 8-bit readable/writable register that control the states of the external bus control pins and specifies the startup method when deep standby mode is canceled.

30.2.25 Deep Standby Cancel Source Select Register (DSSSR)

DSSSR is a 16-bit readable/writable register that consists of the bits for selecting the interrupt to cancel deep standby mode. IRQ7 to IRQ0 bits are valid only for pins allocated to PJ3 to PJ0 and PC3 to PC0.

30.2.26 Deep Standby Cancel Source Flag Register (DSFR)

DSFR is a 16-bit readable/writable register composed of two types of bits. One is the flags that are used to confirm which interrupt has canceled deep standby mode. The other is the bit that releases the retention of pin state after deep standby mode is canceled. When deep standby mode is canceled by an interrupt (NMI or IRQ) or a manual reset, this register retains the value before the cancellation although power-on reset exception handling is executed. When deep standby mode is canceled by a power-on reset, this register is initialized to H'0000.

All flags must be cleared immediately before transition to deep standby mode.

always be 0.

0: No interrupt on MRES pin 1: Interrupt on MRES pin

Note: When writing to this register, see section 30.4, Usage Notes.

9 MRESF 0 R/W MRES Flag

Initial

8 NMIF 0 R/W NMI Flag

R/W Description

0: No interrupt on NMI pin

Bit Bit Name

30.3 Operation

30.3.1 Transitions in Power-Down Modes

Figure 30.1 illustrates the state transitions between power-down modes.

Figure 30.1 Transitions of States in Power-Down Modes

30.3.2 Dual-Processor Mode

After the reset exception handling, both CPU0 and CPU1 are activated and the LSI enters dualprocessor mode. For details of the sequence after canceling reset, see section 4, Multi-Core Processor.

30.3.3 Single-Processor Mode

(1) Transition to Single-Processor Mode

In dual-processor mode where CPU0 and CPU1 are running, the LSI can enter single-processor mode where either CPU0 or CPU1 works.

When CPU1 executes the SLEEP instruction in dual-processor mode, CPU1 switches from a program execution state to a sleep state, and the LSI enters single-processor 0 mode irrespective of the value of the STBY bit in STBCR1.

On the other hand, when CPU0 executes the SLEEP instruction while the STBY bit in STBCR1 is 0 in dual-processor mode, CPU0 switches from a program execution state to a sleep state, and the LSI enters single-processor 1 mode. However, if CPU0 executes the SLEEP instruction when the SLPERE bit in STBCR1 is 1, the LSI does not enter single-processor 1 mode and a sleep error exception occurs.

The CPU halts after executing the SLEEP instruction, but the contents of its registers remain unchanged. The on-chip peripheral modules continue to operate. The CKIO pin continues to output the clock.

(2) Canceling Single-Processor Mode

Single-processor mode is canceled by an interrupt (NMI, IRQ, PINT, on-chip peripheral module interrupt, and inter-processor interrupt) or a reset (manual reset or power-on reset).

(a) Canceling by an interrupt

When an NMI, IRQ, on-chip peripheral module interrupt, or inter-processor interrupt occurs, single-processor mode is canceled and interrupt exception handling is executed, and then the LSI enters dual-processor mode. At this time, set the interrupt enable bit for the occurring interrupt so that the CPU in the sleep state is enabled to accept the interrupt. For details of the interrupt enable bit, see section 7, Interrupt Controller (INTC). If the priority level of the interrupt occurred is not higher than the interrupt mask level that is set in SR of the CPU, the interrupt request is not accepted and single-processor mode is not canceled.

(b) Canceling by a reset

Single-processor mode is canceled and the reset exception handling is executed by a power-on reset or manual reset, and then the LSI enters dual-processor mode.

30.3.4 Dual-Sleep Mode

(1) Transition to Dual-Sleep Mode

In single-processor mode where CPU0 or CPU1 is running, the LSI can enter dual-sleep mode.

Executing the SLEEP instruction by CPU0 when the STBY bit in STBCR1 is 0 in singleprocessor 0 mode causes a transition from the program execution state to dual-sleep mode. However, if CPU0 executes the SLEEP instruction when the SLPERE bit in STBCR1 is 1, the LSI does not enter dual-sleep mode and a sleep error exception occurs.

On the other hand, when CPU1 executes the SLEEP instruction in single-processor 1 mode, CPU1 switches from the program execution state to dual-sleep mode. The LSI enters dual-sleep mode irrespective of the value of the STBY bit in STBCR1.

The running CPU halts after executing the SLEEP instruction, but the contents of its registers remain unchanged. The on-chip peripheral modules continue to operate. The CKIO pin continues to output the clock.

(2) Canceling Dual-Sleep Mode

Dual-sleep mode is canceled by an interrupt (NMI, IRQ, PINT, on-chip peripheral module interrupt, and inter-processor interrupt) or a reset (manual reset or power-on reset).

(a) Canceling by an interrupt

When an NMI, IRQ, on-chip peripheral module interrupt, or inter-processor interrupt occurs, single-processor mode is canceled and interrupt exception handling is executed. The following transition depends on the setting of the interrupt enable bit that controls enabling/disabling of acceptance of interrupts by each CPU. When both CPUs are enabled to accept interrupts, the LSI enters dual-processor mode. When only CPU0 (CPU1) is enabled to accept interrupts, the LSI enters single-processor 0 (single-processor 1) mode. For details of the interrupt enable bit, see section 7, Interrupt Controller (INTC). If the priority level of the interrupt occurred is not higher than the interrupt mask level that is set in SR of the CPU, the interrupt request is not accepted and single-processor mode is not canceled.

(b) Canceling by a reset

Dual-sleep mode is canceled and the reset exception handling is executed by a power-on reset or manual reset, and then the LSI enters dual-processor mode.

30.3.5 Software Standby Mode

(1) Transition to Software Standby Mode

In single-processor 0 mode where only CPU0 is running, the LSI can enter software standby mode.

After setting to disable an interrupt to CPU1 and then confirming that the SLEEP bit in C1MSR is 1, when CPU0 executes the SLEEP instruction with the STBY bit set to 1 and the DEEP bit cleared to 0 in STBCR1, the LSI switches from a program execution state to software standby mode. However, if CPU0 executes the SLEEP instruction when the SLPERE bit in STBCR1 is 1, the LSI does not enter software standby mode and a sleep error exception occurs.

In software standby mode, not only CPU0 and CPU1 but also the clock and on-chip peripheral modules halt. The clock output from the CKIO pin also stops.

The contents of the CPU0/CPU1 registers and the cache registers remain unchanged. Some registers of on-chip peripheral modules are, however, initialized. As for the states of on-chip peripheral module registers in software standby mode, see section 32.3, Register States in Each Operating Mode.

The CPU takes one cycle to finish writing to STBCR1, and then executes processing for the next instruction. However, it takes one or more cycles to actually write. Therefore, execute a SLEEP instruction after reading STBCR1 to have the values written to STBCR1 by the CPU to be definitely reflected in the SLEEP instruction.

The procedure for switching to software standby mode is as follows:

- 1. Clear the TME bit in the WDT0 timer control/status register (WTCSR0) to 0 to stop the WDT.
- 2. Set the WDT0 timer counter (WTCNT0) to 0 and set the clock select bits CKS[2:0] in WTCSR0 to appropriate values to secure the specified oscillation settling time.
- 3. After setting the STBY bit to 1 and DEEP bit to 0 in STBCR1, read STBCR1.
- 4. Setting to disable an interrupt to CPU1 and confirming that the SLEEP bit in C1MSR is 1, and then make CPU0 to execute the SLEEP instruction.

(2) Canceling Software Standby Mode

Software standby mode is canceled by an interrupt (NMI and IRQ) or a reset (power-on reset or manual reset).

(a) Canceling by an interrupt

When the falling edge or rising edge of the NMI pin (selected by the NMI edge select bit (NMIE) in interrupt control register 0 (C0ICR0, C1ICR0) of the interrupt controller (INTC)) or the falling edge or rising edge of an IRQ pin (IRQ7 to IRQ0) (selected by the IRQn sense select bits (IRQn1S and IRQn0S) in interrupt control register 1 (C0ICR1, C1ICR1) of INTC) is detected, clock oscillation is started. This clock pulse is supplied only to the oscillation settling counter (WDT) used to count the oscillation settling time.

After the elapse of the time set in the clock select bits (CKS[2:0]) in the watchdog timer control/status register (WTCSR0) of WDT0 before the transition to software standby mode, the WDT overflow occurs. Since this overflow indicates that the clock has been stabilized, the clock pulse will be supplied to the entire chip after this overflow. After software standby mode is thus canceled and NMI interrupt exception handling (IRQ interrupt exception handling in case of IRQ) is executed, the LSI enters dual-processor mode. If the priority level of the generated interrupt is equal to or lower than the interrupt mask level specified in the status register (SR) of the CPU, the interrupt request is not accepted and software standby mode is not canceled.

When canceling software standby mode by the NMI interrupt or IRQ interrupt, set the CKS[2:0] bits so that the WDT overflow period will be equal to or longer than the oscillation settling time.

The clock output phase of the CKIO pin may be unstable or fixed to low level immediately after an interrupt is detected and until software standby mode in canceled. When software standby mode is canceled by the falling edge of the NMI pin, the NMI pin should be high when the CPU enters software standby mode (when the clock pulse stops) and should be low when software standby mode is canceled (when the clock is initiated after the oscillation settling). When software standby mode is canceled by the rising edge of the NMI pin, the NMI pin should be low when the CPU enters software standby mode (when the clock pulse stops) and should be high when software standby mode is canceled (when the clock is initiated after the oscillation settling). (This is the same with the IRQ pin.)

(b) Canceling by a reset

When the RES pin or MRES pin is driven low, the LSI enters the power-on reset or manual reset state and software standby mode is canceled. After that, the reset exception handling is executed and then the LSI enters dual-processor mode.

Keep the RES pin or MRES pin low until the clock oscillation settles.

The CKIO pin continues to output the internal clock.

(3) Note on Release from Software Standby Mode

Release from software standby mode is triggered by interrupts (NMI and IRQ) or resets (manual reset and power-on reset). If, however, a SLEEP instruction and an interrupt other than NMI and IRQ are generated at the same time, cancellation of software standby mode may occur due to acceptance of the interrupt.

When initiating a transition to software standby mode, make settings so that interrupts are not generated before execution of the SLEEP instruction.

30.3.6 Software Standby Mode Application Example

This example describes a transition to software standby mode on the falling edge of the NMI signal and the cancellation of software standby mode on the rising edge of the NMI signal. The timing is shown in figure 30.2.

When the NMI pin is changed from high to low level while the NMI edge select bit (NMIE) in the interrupt control register (ICR) is set to 0 (falling edge detection), the NMI interrupt is accepted. When the NMIE bit is set to 1 (rising edge detection) by the NMI exception service routine, and then the SLEEP instruction is executed with the STBY bit set to 1 and the DEEP bit set to 0 in STBCR1, the LSI enters software standby mode. Thereafter, when the NMI pin is changed from low to high level, software standby mode is canceled.

Figure 30.2 NMI Timing in Software Standby Mode (Application Example)
30.3.7 Deep Standby Mode

(1) Transition to Deep Standby Mode

In single-processor 0 mode where only CPU0 is running, the LSI can enter deep standby mode.

After setting to disable an interrupt to CPU1 and confirming that the SLEEP bit in C1MSR is 1, when CPU0 executes the SLEEP instruction with the STBY and DEEP bits in STBCR1 set to 1. the LSI switches from a program execution state to deep standby mode. However, if CPU0 executes the SLEEP instruction when the SLPERE bit in STBCR1 is 1, the LSI does not enter deep standby mode and a sleep error exception occurs.

In deep standby mode, not only CPU0, CPU1, the clock, and on-chip peripheral modules halt, but also power supply is turned off except for that supplied to the RTC and the on-chip RAM (for data retention) area specified by the RRAMKP3 to RRAMKP0 bits in RRAMKP, which can significantly reduce power consumption. Therefore, data in the registers of the CPU0, CPU1, cache, and on-chip peripheral modules are not retained. However, the pin state values immediately before the transition to deep standby mode are retained.

The CPU takes one cycle to finish writing to DSFR, and then executes processing for the next instruction. However, it actually takes one or more cycles to write. Therefore, execute a SLEEP instruction after reading DSFR to definitely reflect the values written to DSFR by the CPU in the SLEEP instruction.

The procedure for switching to deep standby mode is shown below. Figure 30.3 shows its flowchart.

- 1. Set the RRAMKP3 to RRAMKP0 bits in RRAMKP for the corresponding on-chip RAM (for data retention) area that must be retained. Transfer the programs to be retained to the specified areas of the on-chip RAM (for data retention).
- 2. To cancel deep standby mode by means of an interrupt or manual reset, set the appropriate bits in DSSSR to select the pins to be used to cancel deep standby mode. In addition, an appropriate input signal detection mode must be specified (using interrupt control registers 0 and 1 (C0ICR0, C1ICR0, C0ICR1, and C1ICR1) of the interrupt controller (INTC)) for the pins what will be used to cancel deep standby mode. Only rising- or falling-edge detection is effective for canceling deep standby mode. (Deep standby mode cannot be canceled when lowlevel or both-edge IRQ detection is selected.)
- 3. Execute read and write accesses to an arbitrary but the same address for each page in the onchip RAM (for data retention) area. If this is not executed, data last written may not be written to the on-chip RAM. If there is any write to the on-chip RAM (for data retention) hereafter, execute this processing after the last write to the on-chip RAM.
- 4. Set the STBY and DEEP bits in STBCR1 to 1.
- 5. Clear the flag in DSFR and then read DSFR.
- 6. Set to disable an interrupt to CPU1 and confirm that the SLEEP bit in C1MSR is 1, and then execute the SLEEP instruction by CPU0.

Figure 30.3 Flowchart of Transition to Deep Standby Mode

(2) Canceling Deep Standby Mode

Deep standby mode is canceled by interrupts (NMI or IRQ allocated to PJ3 to PJ0 and PC3 to PC0) or a reset (manual reset or power-on reset). When canceling deep standby mode by the NMI or IRQ interrupt, a power-on reset exception handling is executed instead of an interrupt exception handling. When canceling deep standby mode by manual reset, a power-on reset exception handling is also executed. After executing the power-on reset exception handling, the LSI enters dual-processor mode. Figure 30.4 shows the flowchart of canceling deep standby mode.

Figure 30.4 Flowchart of Canceling Deep Standby Mode

(a) Canceling by an interrupt

When the falling edge or rising edge of the NMI pin (selected by the NMI edge select bit (NMIE) in interrupt control register 0 (C0ICR0, C1ICR0) of the interrupt controller (INTC)) or the falling edge or rising edge of an IRQ pin (IRQ7 to IRQ0 allocated to PJ3 to PJ0 and PC3 to PC0) (selected by the IRQn sense select bits (IRQn1S and IRQn0S) in interrupt control register 1 (C0ICR1, C1ICR1) of INTC) is detected, clock oscillation is started after waiting for the power supply stabilization time. After the oscillation settling time has elapsed, deep standby mode is cancelled and the power-on reset exception handling is executed. If the priority level of the generated interrupt is equal to or lower than the interrupt mask level specified in the status register (SR) of the CPU, the interrupt request is not accepted and deep standby mode is not canceled.

The clock output phase of the CKIO pin may be unstable immediately after an interrupt is detected and until deep standby mode is canceled. When deep standby mode is canceled on the falling edge of the NMI pin, the NMI pin should be high when the CPU enters deep standby mode (when the clock pulse stops) and should be low when deep standby mode is canceled (when the clock is initiated after the oscillation settling). When deep standby mode is canceled by the rising edge of the NMI pin, the NMI pin should be low when the CPU enters deep standby mode (when the clock pulse stops) and should be high when deep standby mode is canceled (when the clock is initiated after the oscillation settling). (This is the same with the IRQ pin.)

In addition, the pin levels of the NMI pin and all interrupt pins (IRQ) selected to cancel deep standby mode (by settings in the deep standby mode cancellation source select register) should be as follows during the transition to deep standby mode, regardless of whether or not those pins are actually used to cancel deep standby mode:

- Pins set to cancel deep standby mode at their rising edge should be low during the transition.
- Pins set to cancel deep standby mode at their falling edge should be high during the transition.

(b) Canceling by a reset

When the RES pin is driven low, this LSI enters the power-on reset state and deep standby mode is canceled. Then, the RES pin is driven high and the power-on reset exception handling is executed. When the RES pin is driven low in clock mode 0, 1, or 3, the internal clock output from the CKIO is executed.

When the MRES pin is driven low, this LSI enters the power-on reset state and deep standby mode is canceled. Then, the MRES pin is driven high and the power-on reset exception handling is executed. When the MRES pin is driven high in clock mode 0, 1, or 3, the internal clock output from the CKIO is executed.

Keep the RES or MRES pin low until the clock oscillation settles.

(3) Operation after Canceling Deep Standby Mode

After exiting from deep standby mode, the LSI can be booted up either through the external bus or from the on-chip RAM (for data retention), which can be selected by setting the RAMBOOT bit in DSCTR. By setting the CS0KEEPE bit, the states of the external bus control pins can be retained even after cancellation of deep standby mode. Table 30.3 shows the pin states after cancellation of deep standby mode according to the setting of each bit. Table 30.4 lists the external bus control pins.

Table 30.3 Pin States and Boot-Up Method after Exit form Deep Standby Mode according to DSCTR Register Setting

Table 30.4 External Bus Control Pins in Different Modes

When deep standby mode is canceled by interrupts (NMI or IRO) or a manual reset, the deep standby cancel source flag register (DSFR) can be used to confirm which interrupt has canceled the mode.

Pins retain the state immediately before the transition to deep standby mode. However, in system activation through the external bus, the retention of the states of the external bus control pins is canceled so that programs can be fetched after cancellation of deep standby mode. The pin states are retained after cancellation of deep standby mode until writing 0 to the IOKEEP bit in DSFR from the same bit. Reconfiguration of peripheral functions is required to return to the previous state of deep standby mode. Peripheral functions include all functions such as CPG, INTC, BSC, I/O ports, PFC, and peripheral modules. After the reconfiguration, the retention of the pin state can be canceled by writing 0 to the IOKEEP bit in DSFR.

(4) Note on Transition to Deep Standby Mode

When multiple cancellation sources have been set and more than one cancellation source is input, multiple cancellation source flags are set.

30.3.8 Module Standby Function

(1) Transition to Module Standby Function

Setting the standby control register MSTP bits to 1 halts the supply of clocks to the corresponding on-chip peripheral modules. This function can be used to reduce the power consumption in dualprocessor mode, single-processor mode, and dual-sleep mode. Disable a module before placing it in module standby mode. In addition, do not access the module's registers while it is in the module standby state.

For details on the states of registers, see section 32.3, Register States in Each Operating Mode.

(2) Canceling Module Standby Function

The module standby function can be canceled by clearing each MSTP bit to 0, or by a power-on reset (only possible for RTC, H-UDI, and UBC). When canceling the module standby function by clearing the corresponding MSTP bit to 0, read the MSTP bit to confirm that it has been cleared to 0.

30.4 Usage Notes

30.4.1 Notes on Writing to Registers

When writing to the registers related to power-down modes, note the following.

When writing to the register related to power-down modes, the CPU, after executing a write instruction, executes the next instruction without waiting for the write operation to complete.

Therefore, to reflect the change specified by writing to the register while the next instruction is executed, insert a dummy read of the same register between the register write instruction and the next instruction.

30.4.2 Notice about Deep Standby Control Register (DSCTR)

After (1) power-on reset by \overline{RES} pin is released, and (2) the LSI transits to deep standby mode in case that bit 7 (CS0KEEPE) and bit 6 (RAMBOOT) of deep standby control register (DSCTR) are set to "1", these bits become unable to be written as "0" since then.

To write these as "0", it is necessary to assert RES pin to low.

30.4.3 Notice about Power-On Reset Exception Handling

After (1) power-on reset by RES pin is released, (2) the LSI transit to deep standby mode in case that bit 6 (RAMBOOT) of deep standby control register (DSCTR) is set to $"1", (3)$ the deep standby mode is cancelled, and (4) power-on reset by WDT or H-UDI reset is occurred before power-on reset by RES pin is executed again, then the behavior of the power-on reset exception handling is as table 30.5. So if applicable as above case, PC and SP are necessary to be retained in the area of on-chip RAM for data retention.

Table 30.5 Power-On Reset Exception Handling

After (1) power-on reset by \overline{RES} pin is released, (2) the LSI transit to deep standby mode, and (3) the deep standby mode is cancelled, if there is a possibility that power-on reset by WDT or H-UDI reset is occurred before power-on reset by RES pin is executed again, the settings of WDT or H-UDI should be done in the condition that bit 15 (IOKEEP) and bits 9 to 0 of deep standby cancel source flag register (DSFR) are all cleared after canceling deep standby mode (if some bits are 1, please write these as "0").

If (1) the setting of WDT or H-UDI is done in the condition that IOKEEP bit is not 0, and (2) power-on reset by WDT or H-UDI reset is occurred before power-on reset by RES pin is executed again, the pin status of the pins, whose pin status are retained in deep standby mode and which are not in table 30.4, are kept retained. Additionally, in the case that bit 7 (CS0KEEPE) of deep standby control register (DSCTR) are set to "1", the pin status of the pins in table 30.4 are also keep retained.

If (1) the settings of WDT or H-UDI is done in the condition that bits 9 to 0 are not all 0, and (2) power-on reset by WDT or H-UDI reset is occurred before power-on reset by RES pin is executed again, the internal information about the deep standby canceling source is not cleared, and deep standby mode are cancelled by the wrong canceling source when the LSI attempt to transit to deep standby mode since then.

Section 31 User Debugging Interface (H-UDI)

This LSI incorporates a user debugging interface (H-UDI) for the boundary scan function and emulator support.

31.1 Features

The user debugging interface (H-UDI) is a serial input/output interface that supports JTAG (Joint Test Action Group, IEEE Std.1149.1 and IEEE Standard Test Access Port and Boundary-Scan Architecture).

The H-UDI of this LSI incorporates a boundary scan TAP controller and an emulation TAP controller for controlling the H-UDI interrupt function. When the TRST pin is asserted, including the case of power-on, the boundary scan TAP controller is selected. By inputting the emulation TAP controller switching command, the emulation TAP controller is selected. To switch from the emulation TAP controller to the boundary scan TAP controller, assert the TRST pin.

In ASE mode, the emulation TAP controller is selected. For connection with the emulator, see the manual for the emulator.

Figure 31.1 shows a block diagram of the H-UDI.

Figure 31.1 Block Diagram of H-UDI

31.2 Input/Output Pins

Table 31.1 Pin Configuration

Note: $*$ When the emulator is not in use, fix this pin to the high level.

31.3 Description of the Boundary Scan TAP Controller

The boundary scan TAP controller has the following registers.

Table 31.2 Register Configuration of the Boundary Scan TAP Controller

31.3.1 Bypass Register (BSBPR)

BSBPR is a 1-bit register that cannot be accessed by the CPU. When BSIR is set to BYPASS mode, BSBPR is connected between H-UDI pins TDI and TDO. The initial value is undefined.

31.3.2 Instruction Register (BSIR)

BSIR is a 4-bit register and initialized by TRST assertion or in the TAP test-logic-reset state. This register cannot be accessed by the CPU.

Bits 3 to 0

Table 31.3 Supported Commands for Boundary Scan TAP Controller

31.3.3 Boundary Scan Register (SDBSR)

SDBSR is a shift register located on the PAD to control input/output pins of this LSI. This register cannot be accessed by the CPU. The initial value is undefined.

The EXTEST, SAMPLE/PRELOAD, CLAMP, and HIGHZ commands can be used to perform the boundary scan test that conforms to the JTAG standard. Table 31.4 shows the correspondence between the LSI pins and the bits of the boundary scan register.

Notes: 1. The pin name used for function 1

 2. The pin is open-drain. The pin state is low when driven low, whereas high impedance (Hi-Z) when driven high.

 3. The pin of CONTROL is active-low. When this pin is driven low, the state of the corresponding pin is output.

31.3.4 ID Register (BSID)

BSID is a 32-bit register that cannot be accessed by the CPU. The register can be read from H-UDI pins when the IDCODE command is set, but is not writable.

31.4 Description of the Emulation TAP Controller

To use the emulation TAP controller, enter the emulation TAP controller switching command in the BSIR register of the boundary scan TAP controller. The emulation TAP controller has the following registers.

Table 31.5 Register Configuration of the Emulation TAP Controller

31.4.1 Bypass Register (SDBPR)

SDBPR is a 1-bit register that cannot be accessed by the CPU. When SDIR is set to BYPASS mode, SDBPR is connected between H-UDI pins TDI and TDO. The initial value is undefined.

31.4.2 Instruction Register (SDIR)

SDIR is a 16-bit read-only register and initialized by TRST assertion or in the TAP test-logic-reset state. H-UDI can write to this register regardless of the CPU mode. When a reserved command is set in this register, the operation is not guaranteed. The initial value is H'EFFD.

Note: * The initial value of TI[7:0] is a reserved value, but replace it with a non-reserved value when setting a command.

31.5 Operation

31.5.1 TAP Controller

Figure 31.2 shows the internal states of the TAP controller. This state machine conforms to the state transitions defined by JTAG.

Figure 31.2 TAP Controller State Transitions

Note: The transition condition is the TMS value at the rising edge of TCK. The TDI value is sampled at the rising edge of TCK; shifting occurs at the falling edge of TCK. For details on transition timing of the TDO value, see section 31.5.3, TDO Output Timing. The TDO is at high impedance, except with shift-DR and shift-IR states. During the change to TRST $= 0$, there is a transition to test-logic-reset asynchronously with TCK.

31.5.2 Reset Configuration

Notes: 1. Performs product chip mode and ASE mode settings $\overline{\text{ASEMD}}$ = H, product chip mode ASEMD = L, ASE mode

> 2. In ASE mode, reset hold is entered if the $\overline{\text{TRST}}$ pin is driven low while the $\overline{\text{RES}}$ pin is negated. In this state, the CPU does not start up. When TRST is driven high, H-UDI operation is enabled, but the CPU does not start up. The reset hold state is cancelled by a power-on reset.

31.5.3 TDO Output Timing

When the emulation TAP controller is selected, a transition on the TDO pin is output on the falling edge of TCK with the initial value. However, setting a TDO transition timing switching command in SDIR via the H-UDI pin and passing the Update-IR state synchronizes the TDO transition with the rising edge of TCK. This command does not affect the output timing of the boundary scan TAP controller.

To synchronize the transition of TDO with the falling edge of TCK after setting the TDO transition timing switching command, the TRST pin must be asserted simultaneously with the power-on reset. In the case of power-on reset by the RES pin, the sync reset is still in operation for a certain period in the LSI even after the RES pin is negated. Thus, if the TRST pin is asserted immediately after the negation of the \overline{RES} pin, the TDO transition timing switching command is cleared, resulting in TDO transitions synchronized with the falling edges of TCK. To prevent this, make sure to allow a period of 20 tcyc or longer between the signal transitions of the RES and TRST pins.

31.5.4 H-UDI Reset

An H-UDI reset occurs when an H-UDI reset assert command is set in SDIR. An H-UDI reset is of the same kind as a power-on reset. An H-UDI reset is cleared by setting an H-UDI reset negate command. The required time between the H-UDI reset assert command and H-UDI reset negate command is the same as time for keeping the RES pin low to apply a power-on reset.

31.5.5 H-UDI Interrupt

The H-UDI interrupt function generates an interrupt by setting a command from the H-UDI in SDIR. An H-UDI interrupt is a general exception/interrupt operation, resulting in fetching the exception service routine start address from the exception handling vector table, jumping to that address, and starting program execution from that address. This interrupt request has a fixed priority level of 15.

H-UDI interrupts are accepted in sleep mode, but not in software standby mode.

31.6 Boundary Scan

By setting the commands in BSIR by the H-UDI, the H-UDI pins can be configured for boundary scan mode defined by JTAG.

31.6.1 Supported Instructions

This LSI supports three required instructions (BYPASS, SAMPLE/PRELOAD, and EXTEST) and three optional instructions (IDCODE, CLAMP, and HIGHZ) defined by JTAG.

(1) BYPASS

The BYPASS instruction is a required standard instruction to operate the bypass register. This instruction is used to increase the transfer speed of serial data of other LSIs on the printed circuit board by reducing the shift path. During execution of this instruction, the test circuit does not affect the system circuit.

(2) SAMPLE/PRELOAD

The SAMPLE/PRELOAD instruction inputs a value from the internal circuit of the LSI to the boundary scan register, and output the data from scan path or load the data to the scan path. During execution of the instruction, the value on the input pin of the LSI is transferred to the internal circuit and the value of the internal circuit is output externally from the output pin. Execution of the instruction does not affect the system circuit of the LSI.

In SAMPLE operation, the snapshots of the value transferred from the input pin to the internal circuit and the value transferred from the internal circuit to the output pin are captured in the boundary scan register and then read from the scan path. Capturing of the snapshots is performed in synchronization with the rising edge of TCK in the capture-DR state. The capturing is performed without interfering with normal operation of the LSI.

In PRELOAD operation, an initial value is set in the output latch of the boundary scan register from the scan path before execution of the EXTEST instruction. Without PRELOAD operation, an undefined value is output from the output pin until the first scan sequence is completed (transferred to the output latch) during execution of the EXTEST instruction (the parallel output latch is always output to the output pin with the EXTEST instruction).

(3) EXTEST

The EXTEST instruction tests the external circuit when this LSI is mounted on the printed circuit board. During execution of this instruction, the output pin is used to output the test data (set in advance by the SAMPLE/PRELOAD instruction) from the boundary scan register to the printed circuit board and the input pin is used to capture the test result from the printed circuit board to the boundary scan register. When a test is performed using the EXTEST instruction N times, the N-th test data is scanned-in during (N-1)-th scan-out.

The data loaded in the boundary scan register of the output pin in the capture-DR state of this instruction is not used in testing of the external circuit (an exchange is made in shift operation).

(4) IDCODE

To set the H-UDI pin to IDCODE mode defined by JTAG, set the command in SDIR from the H-UDI pin. When H-UDI is initialized (TRST is asserted or TAP is placed in the test-logic-reset state), IDCODE mode is entered.

(5) CLAMP and HIGHZ

To set the H-UDI pin to CLAMP or HIGHZ mode defined by JTAG, set the command in SDIR from the H-UDI pin.

31.6.2 Notes

- 1. The clock related signals (EXTAL, XTAL, CKIO, AUDIO_X1, AUDIO_X2, USB_X1, USB_X2, RTC_X1, and RTC_X2) are inapplicable to the boundary scan.
- 2. The reset-related signal (RES) is inapplicable to the boundary scan.
- 3. The H-UDI related signals (TCK, TDI, TDO, TMS, TRST, and ASEMD) are inapplicable to the boundary scan.
- 4. The USB related signals (DP0, DM0, DP1, DM1, VBUS, and REFRIN) are inapplicable to the boundary scan.
- 5. 2DG VIDEO OUT related signals (R, G, B, REXT, CBU) are inapplicable to the boundary scan.
- 6. Execute the boundary scan in product chip mode and input the ASEMD pin to high during the RES pin assertion period. And make sure to fix the ASEMD pin at high while executing the boundary scan.

31.7 Usage Notes

- 1. An H-UDI command, once set, will not be modified as long as another command is not set again from the H-UDI. If the same command is to be set continuously, the command must be set after a command (BYPASS mode, etc.) that does not affect chip operations is once set.
- 2. In software standby mode and H-UDI module standby state, none of the functions in the H-UDI can be used. To retain the TAP status before and after standby mode, keep TCK high before entering standby mode.
- 3. Regardless of whether the H-UDI is used, make sure to keep the TRST pin low to initialize the H-UDI at power-on or in recovery from deep standby by the $\overline{\text{RES}}$ pin assertion.
- 4. If the TRST pin is asserted immediately after the setting of the TDO transition timing switching command and the negation of the RES pin, the TDO transition timing switching command is cleared. To avoid this case, make sure to put 20 tcyc or longer between the signal transition timing of the RES and TRST pins. For details, see section 31.5.3, TDO Output Timing.
- 5. When starting the TAP controller after the negation of the TRST pin, make sure to allow 200 ns or longer after the negation.
- 6. Please keep TMS pin high for 200 ns from TRST pin negation.

Section 32 List of Registers

This section gives information on the on-chip I/O registers of this LSI in the following structure.

- 1. Register Addresses (by functional module, in order of the corresponding section numbers)
- Registers are described by functional module, in order of the corresponding section numbers.
- Access to reserved addresses which are not described in this register address list is prohibited.
- When registers consist of 16 or 32 bits, the addresses of the MSBs are given assuming big endian.
- 2. Register Bits
- Bit configurations of the registers are described in the same order as the Register Addresses (by functional module, in order of the corresponding section numbers).
- Reserved bits are indicated by $-$ in the bit name.
- No entry in the bit-name column indicates that the whole register is allocated as a counter or for holding data.
- 3. Register States in Each Operating Mode
- Register states are described in the same order as the Register Addresses (by functional module, in order of the corresponding section numbers).
- For the initial state of each bit, refer to the description of the register in the corresponding section.
- The register states described are for the basic operating modes. If there is a specific reset for an on-chip peripheral module, refer to the section on that on-chip peripheral module.
- 4. Notes when Writing to the On-Chip Peripheral Modules
- To access an on-chip module register, two or more peripheral module clock $(P\phi)$ cycles are required. Care must be taken in system design. When the CPU writes data to an on-chip peripheral register, the CPU executes the succeeding instruction without waiting for the completion of writing to the register. For example, a case in which the system enters software standby mode for power saving is described here. To make this transition, the SLEEP instruction must be performed after setting the STBY bit in the STBCR register to 1. However, a dummy read of the STBCR register is required before executing the SLEEP instruction. If a dummy read is omitted, the CPU executes the SLEEP instruction before the STBY bit is set to 1, thus the system enters sleep mode not software standby mode. A dummy read of the STBCR register is indispensable to complete writing to the STBY bit. To reflect the change made to an on-chip peripheral register while performing the succeeding instruction, dummy-read the register to which write instruction is applied and then execute the succeeding instruction.

32.1 Register Addresses (by functional module, in order of the corresponding section numbers)

32.2 Register Bits

32.3 Register States in Each Operating Mode

Notes: 1. Retains the previous value after an internal power-on reset by means of the WDT.

- 2. The BN3 to BN0 bits are initialized.
- 3. Counting up continues.
- 4. Bits RTCEN and START are retained.
- 5. Bits BC2 to BC0 are initialized.
- 6. Since pin states are read out on the port G data register (PGDRL) and the port registers, values in these registers are neither retained nor initialized.
- 7. Initialized by TRST assertion or in the Test-Logic-Reset state of the TAP controller.

Section 33 Electrical Characteristics

33.1 Absolute Maximum Ratings

Table 33.1 Absolute Maximum Ratings

Caution: Permanent damage to the LSI may result if absolute maximum ratings are exceeded.

33.2 Power-on/Power-off Sequence

Figure 33.1 Power-on/Power-off Sequence

Table 33.2 Time for Power-on/Power-off Sequence

Note: It is recommended that the 1.2-V power supply $(V_{\text{cor}}, PLLV_{\text{cor}}, USBAV_{\text{cor}}, and USBDV_{\text{cor}})$ and the 3.3-V power supply (PV_{cc}, AV_{cc}, USBAPV_{cc}, 2DGAPV_{cc}0, and 2DGAPV_{cc}1) are turned on and off nearly simultaneously.

 An indefinite period of time appears, from the time that power is turned on to the time that both of the 1.2-V power supply and the 3.3-V power supply rise to the Min. voltage (1.1 V for 1.2-V power supply and 3.0 V for 3.3-V power supply), or from the time that either of the 1.2-V power supply or the 3.3-V power supply is turned off and passes the Min. voltage (1.1 V for 1.2-V power supply and 3.0 V for 3.3-V power supply) to the time that both of the 1.2- V power supply and the 3.3-V power supply fall to GND.

 During these periods, states of output pins and in-out pins and internal states become undefined. So it should be as short as possible. Also design the system so that these undefined states do not cause an overall malfunction.

33.3 DC Characteristics

Table 33.3 DC Characteristics (1) [Common Items]

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Caution: When the A/D converter or D/A converter is not in use, the AV_{cc} and AV_{ss} pins should not be open.

Notes: 1. The supply current values are when all output pins and pins with the pull-up function are unloaded.

2. I_{CC} , I_{sleep} , I_{sstby} , and I_{dstby} represent the total currents supplied in the V_{CC} and PLLV_{CC} systems.

3. Pl_{cc} is the current through PV_{cc} (reference value) when there is no load on any output pin and the input pins are fixed. Since actual currents in operation are strongly dependent on the system (waveforms determined by IO, frequency of toggling, and so on), be sure to measure actual values for systems.

4. I_{deph} for the 1.2-V power supply is the total current drawn through V_{CC}, PLV_{CC}, USBAV_{CC}, and USBDV $_{cc}$.

5. I_{demb} for the 3.3-V power supply is the total current drawn through PV_{CC}, AV_{CC}, USBAPV_{CC}, $2DGAPV_{cc}0$, and $2DGAPV_{cc}1$.

Table 33.3 DC Characteristics (2) [Except I² C and USB-Related Pins]

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = MSBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V,$ $T_a = -20$ to 85 °C

Table 33.3 DC Characteristics (3) [I² C-Related Pins***]**

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $\text{cc} = 1.1$ to 1.3 V, USBAV $\text{cc} = 1.1$ to 1.3 V, $PV_{cc} = 3.0$ to 3.6 V, $AV_{cc} = 3.0$ to 3.6 V, USBAPV $_{cc} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = USBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V,$ $T_a = -20$ to 85 °C.

Note: * The PE13/TxD4/SDA2 to PE8/RxD2/SCL0, PF1/SCL3/CRx0, and PF0/SDA3 pins are open-drain pins.

Table 33.3 DC Characteristics (4) [USB-Related Pins***]**

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C.

Note: * REFRIN, VBUS, USB_X1, and USB_X2 pins

Table 33.3 DC Characteristics (5) [USB-Related Pins* **(Low-Speed, Full-Speed, and High-Speed Common Items)]**

Conditions: $V_{CC} = PLLV_{CC} = 1.1$ to 1.3 V, USBDV_{CC} = 1.1 to 1.3 V, USBAV_{CC} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V. $V_{ss} = PLLV_{ss} = USBAV_{ss} = MSBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C.

Note: * DP1, DP0, DM1 and DM0 pins

Table 33.3 DC Characteristics (6) [USB-Related Pins* (**Low-Speed/Full-Speed)]**

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = MSBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V,$ $T_a = -20$ to 85 °C

Note: * DP1, DP0, DM1 and DM0 pins

Table 33.3 DC Characteristics (7) [USB-Related Pins* (**High-Speed)]**

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = MSBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C.

Note: * DP1, DP0, DM1 and DM0 pins

Table 33.4 Permissible Output Currents

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = AV_{ss} = USBAPV_{ss} = 2DGAPV_{ss}0 = 2DGAPV_{ss}1 = 0 V$, $T_a = -20$ to 85 °C

Caution: To protect the LSI's reliability, do not exceed the output current values in table 33.4.

33.4 AC Characteristics

Signals input to this LSI are basically handled as signals in synchronization with a clock. The setup and hold times for input pins must be followed.

Table 33.5 Operating Frequency

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{cc} = 3.0$ to 3.6 V, $AV_{cc} = 3.0$ to 3.6 V, USBAPV $_{cc} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

33.4.1 Clock Timing

Table 33.6 Clock Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.2 EXTAL, AUDIO_X1, AUDIO_CLK, and USB_X1 Clock Input Timing

Figure 33.3 CKIO Clock Input Timing

Figure 33.6 Oscillation Settling Time on Return from Standby (Return by Reset)

Figure 33.7 Oscillation Settling Time on Return from Standby (Return by NMI or IRQ)

Figure 33.8 RTC Clock Oscillation Settling Time

33.4.2 Control Signal Timing

Table 33.7 Control Signal Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $\text{cc} = 1.1$ to 1.3 V, USBAV $\text{cc} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.10 Interrupt Signal Input Timing

33.4.3 Bus Timing

Table 33.8 Bus Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $\text{cc} = 1.1$ to 1.3 V, USBAV $\text{cc} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Note: $*$ The maximum value (fmax) of B ϕ (bus clock) depends on the number of wait cycles and the system configuration of your board.

Figure 33.11 (1) External Address Space: Basic Bus Timing (Normal Access, Read/Write Cycle Wait = 3, CS Assert Wait = 1, Write Data Output Wait = 1, WR/RD Assert Wait = 2, Write Data Output Delay Cycles = 0, Read/Write CS Delay Cycles = 1)

Figure 33.11 (2) External Address Space: Basic Bus Timing (Normal Access, Data Recovery Cycles = 0, Read/Write Cycle Wait = 1, Read/Write CS Delay Cycles = 1, Other Wait Settings = 0)

Figure 33.11 (3) External Address Space: Basic Bus Timing (Normal Access, Data Recovery Cycles = 2, Read/Write Cycle Wait = 1, Read/Write CS Delay Cycles = 1, Other Wait Settings = 0)

Figure 33.12 External Address Space: Basic Bus Timing (Page Read Access, Normal Access Compatible Mode , Read Cycle Wait = 2, Page Read Cycle Wait = 2, CS Assert Wait = 1, RD Assert Wait = 1, Read CS Delay Cycles = 1)

Figure 33.13 External Address Space: Basic Bus Timing (Page Read Access, External Read Data Continuous Assert Mode, Read Cycle Wait = 2, Page Read Cycle Wait = 1, CS Assert Wait = 1, RD Assert Wait = 1, Read CS Delay Cycles = 1)

Figure 33.14 External Address Space: Basic Bus Timing (Page Write Access, Write Cycle Wait = 2, CS Assert Wait = 1, WR Assert Wait = 1, Write Data Output Delay Cycles = 1, Other Wait Settings = 0)

Figure 33.15 External Address Space: Timing with External Wait (Page Read Access to 16-Bit Width Channel, External Read Data Continuous Assert Mode, Read Cycle Wait = 3, Page Read Cycle Wait = 3, Other Wait Settings = 0, External Wait Cycles = 2)

Figure 33.17 SDRAM Space: Single-Write Bus Timing (DCL = 2 (2 Cycles), DRCD = 1 (2 Cycles), DPCG = 1 (2 Cycles))

Figure 33.18 SDRAM Space: Multiple-Read Bus Timing (4 Data Accesses, DCL = 2 (2 Cycles), DRCD = 1 (2 Cycles), DPCG = 1 (2 Cycles))

Figure 33.19 SDRAM Space: Multiple-Write Bus Timing (4 Data Accesses, DCL = 2 (2 Cycles), DRCD = 1 (2 Cycles), DPCG = 1 (2 Cycles))

Figure 33.20 SDRAM Space: Bus Timing of Multiple-Read Across Rows (8 Data Accesses, DCL = 2 (2 Cycles), DRCD = 1 (2 Cycles), DPCG = 1 (2 Cycles))

Figure 33.22 SDRAM Space: Self-Refresh Bus Timing

33.4.4 UBC Timing

Table 33.9 UBC Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $\text{cc} = 1.1$ to 1.3 V, USBAV $\text{cc} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.23 UBC Timing

33.4.5 DMAC Timing

Table 33.10 DMAC Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss}1 = 0 V$, $T_a = -20$ to 85 °C

Figure 33.24 DREQ Input Timing

Figure 33.25 DACK, DACT, TEND Output Timing

33.4.6 MTU2 Timing

Table 33.11 MTU2 Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{cc} = 3.0$ to 3.6 V, $AV_{cc} = 3.0$ to 3.6 V, USBAPV $_{cc} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Note: t_{new} indicates peripheral clock (P ϕ) cycle.

Figure 33.26 MTU2 Input/Output Timing

33.4.7 WDT Timing

Table 33.12 WDT Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss}1 = 0 V$, $T_a = -20$ to 85 °C

Figure 33.28 WDT Timing

33.4.8 SCIF Timing

Table 33.13 SCIF Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = USBAPV_{ss} = 2DGAPV_{ss}0 = 2DGAPV_{ss}1 = 0 V$, $T_a = -20$ to 85 °C

Note: t_{ocyc} indicates the peripheral clock (P ϕ) cycle.

Figure 33.29 SCK Input Clock Timing

Figure 33.30 SCIF Input/Output Timing in Clocked Synchronous Mode

33.4.9 SSU Timing

Table 33.14 SSU Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $\text{cc} = 1.1$ to 1.3 V, USBAV $\text{cc} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = USBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V,$ $T_a = -20$ to 85 °C

Note: t_{pove} indicates the peripheral clock (P ϕ) cycle.

Figure 33.32 SSU Timing (Master, CPHS = 0)

Figure 33.33 SSU Timing (Slave, CPHS = 1)

Figure 33.34 SSU Timing (Slave, CPHS = 0)

33.4.10 IIC3 Timing

Table 33.15 (1) IIC3 Timing I² C Bus Format

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Notes: 1. t_{new} indicates the peripheral clock (P ϕ) cycle.

- 2. Depends on the value of NF2CYC.
- 3. Indicates the I/O buffer characteristic.

Table 33.15 (2) IIC3 Timing Clock Synchronized Serial Format

Notes: 1. t_{pcyc} indicates the peripheral clock $(P\phi)$ cycle.

2. Depends on the value of NF2CYC.

3. Indicates the I/O buffer characteristic.

Figure 33.35 (3) Transmission and Reception Timing

33.4.11 SSIF Timing

Table 33.16 SSIF Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $\text{cc} = 1.1$ to 1.3 V, USBAV $\text{cc} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.36 Clock Input/Output Timing

Figure 33.38 SSIF Transmit Timing (2)

Figure 33.39 SSIF Receive Timing (1)

Figure 33.40 SSIF Receive Timing (2)

33.4.12 RCAN-TL1 Timing

Table 33.17 RCAN-TL1 Timing

Conditions: $V_{\text{cc}} = PLLV_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = MSBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.42 RCAN-TL1 Input/Output Timing

33.4.13 ADC Timing

Table 33.18 ADC Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.43 A/D Converter External Trigger Input Timing

33.4.14 FLCTL Timing

Table 33.19 AND Type Flash Memory Interface Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Note: t_{fcov} indicates the period of one cycle of the FLCTL clock.

 t_{new} indicates the period of one cycle of the peripheral clock (P ϕ).

Figure 33.44 AND Type Flash Memory Command Issuance Timing

Figure 33.46 AND Type Flash Memory Data Read Timing

Figure 33.48 AND Type Flash Memory Status Read Timing

Table 33.20 NAND Type Flash Memory Interface Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = MSBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Note: t_{row} indicates the period of one cycle of the FLCTL clock.

 $t_{\rm wdyc}$ indicates the period of one cycle of the FLCTL clock when the NANDWF bit is cleared to 0, and indicates the period of two cycles of the FLCTL clock when the NANDWF bit is set to 1.

 t_{row} indicates the period of one cycle of the peripheral clock (P ϕ).

Figure 33.49 NAND Type Flash Memory Command Issuance Timing

Figure 33.50 NAND Type Flash Memory Address Issuance Timing

Figure 33.51 NAND Type Flash Memory Data Read Timing

Figure 33.53 NAND Type Flash Memory Status Read Timing

33.4.15 USB Timing

Table 33.21 USB Transceiver Timing (Low-Speed)

Conditions: $V_{\text{cc}} = PLLV_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = MSBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.54 DP and DM Output Timing (Low-Speed)

Figure 33.55 Measurement Circuit (Low-Speed)

Table 33.22 USB Transceiver Timing (Full-Speed)

Conditions: $V_{\text{cc}} = PLLV_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.57 Measurement Circuit (Full-Speed)

Table 33.23 USB Transceiver Timing (High-Speed)

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.58 DP and DM Output Timing (High-Speed)

Figure 33.59 Measurement Circuit (High-Speed)

33.4.16 ATAPI Timing

Table 33.24 Timing of Register Access by PIO Transfer by ATAPI Interface

Table 33.25 ATAPI Interface PIO Data Transfer Timing

Table 33.26 ATAPI Interface Multi-Word Transfer Timing

Table 33.27 ATAPI Interface Ultrta DMA Transfer Timing

Table 33.28 Symbols for ATAPI Interface Ultrta DMA Transfer Timing

Table 33.29 ATAPI Interface DIRECTION Timing

Figure 33.60 Register Transfer and PIO Data Transfer to/from Device

Figure 33.61 Initiating Multiword DMA Data Burst

Figure 33.62 Sustaining Multiword DMA Data Burst

Figure 33.63 Device Terminating Multiword DMA Data Burst

Figure 33.64 Host Terminating Multiword DMA Data Burst

Figure 33.65 Initiating Ultra DMA Data-In Burst

Figure 33.66 Sustained Ultra DMA Data-In Burst

Figure 33.67 Host Pausing Ultra DMA Data-In Burst

Figure 33.68 Device Terminating Ultra DMA Data-In Burst

Figure 33.69 Host Terminating Ultra DMA Data-In Burst

Figure 33.70 Initiating Ultra DMA Data-Out Burst

Figure 33.73 Host Terminating Ultra DMA Data-Out Burst

Figure 33.74 Device Terminating Ultra DMA Data-Out Burst

Figure 33.75 PIO Data Transfer to Device (DIRECTION)

IDECS#[1:0] IDEA[2:0]

IODREQ

Figure 33.78 Initiating Ultra DMA Data-In Burst (DIRECTION)

Figure 33.79 Device Terminating Ultra DMA Data-In Burst (DIRECTION)

Figure 33.80 Host Terminating Ultra DMA Data-In Burst (DIRECTION)

Figure 33.81 Initiating Ultra DMA Data-Out Burst (DIRECTION)

Figure 33.82 Host Terminating Ultra DMA Data-Out Burst (DIRECTION)

Figure 33.83 Device Terminating Ultra DMA Data-Out Burst (DIRECTION)

33.4.17 2DG Timing

Table 33.30 2DG Video Input Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $\text{cc} = 1.1$ to 1.3 V, USBAV $\text{cc} = 1.1$ to 1.3 V, $PV_{cc} = 3.0$ to 3.6 V, $AV_{cc} = 3.0$ to 3.6 V, USBAPV $_{cc} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.84 Video Input Timing

Table 33.31 2DG Display Output Timing

Conditions: $V_{\text{cc}} = PLLV_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.85 DCLKIN Clock Input Timing

Figure 33.86 Display Output Timing

Table 33.32 VIDEO OUT D/A Converter Characteristics

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, RL = 180 Ω, R_{ext} = 5.23 kΩ, T_a = -20 to 85 °C

33.4.18 I/O Port Timing

Table 33.33 I/O Port Timing

Conditions: $V_{\text{cc}} = PLLV_{\text{cc}} = 1.1$ to 1.3 V, USBDV $_{\text{cc}} = 1.1$ to 1.3 V, USBAV $_{\text{cc}} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Figure 33.87 I/O Port Timing

33.4.19 H-UDI Timing

Table 33.34 H-UDI Timing

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV $\text{cc} = 1.1$ to 1.3 V, USBAV $\text{cc} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = USBAPV_{ss} = 2DGAPV_{ss}0 = 2DGAPV_{ss}1 = 0 V$, $T_a = -20$ to 85 °C

Note: $*$ Should be greater than the peripheral clock (P ϕ) cycle time.

Figure 33.88 TCK Input Timing

Figure 33.90 Boundary Scan Input/Output Timing

33.4.20 AC Characteristics Measurement Conditions

- I/O signal reference level: $PV_{cc}/2$ ($PV_{cc} = 3.0$ to 3.6 V, $V_{cc} = 1.1$ to 1.3 V)
- Input pulse level: V_{ss} to 3.0 V (where \overline{RES} , \overline{MRES} , NMI, MD0, MD_CLK1, MD_CLK0, $\overline{\text{ASEMD}}$, $\overline{\text{TRST}}$, and Schmitt trigger input pins are within V_{ss} to PV_{cc}.)
- Input rise and fall times: 1 ns

Figure 33.91 Output Load Circuit

33.5 A/D Converter Characteristics

Table 33.35 A/D Converter Characteristics

Conditions: $V_{\text{cc}} = \text{PLLV}_{\text{cc}} = 1.1$ to 1.3 V, USBDV_{cc} = 1.1 to 1.3 V, USBAV_{cc} = 1.1 to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = ASBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

Note: * Reference values

33.6 D/A Converter Characteristics

Table 33.36 lists the D/A converter characteristics.

Table 33.36 D/A Converter Characteristics

Conditions: $V_{cc} = PLLV_{cc} = 1.1$ to 1.3 V, USBDV $_{cc} = 1.1$ to 1.3 V, USBAV $_{cc} = 1.1$ to 1.3 V, $PV_{\text{cc}} = 3.0$ to 3.6 V, $AV_{\text{cc}} = 3.0$ to 3.6 V, USBAPV $_{\text{cc}} = 3.0$ to 3.6 V, $2DGAPV_{cc}0 = 3.0$ to 3.6 V, $2DGAPV_{cc}1 = 3.0$ to 3.6 V, $V_{ss} = PLLV_{ss} = USBAV_{ss} = MSBAPV_{ss} = 2DGAPV_{ss} = 2DGAPV_{ss} = 0 V$, $T_a = -20$ to 85 °C

33.7 Usage Notes

Place a multilayer ceramic capacitor as a bypass capacitor per a pair of power supply pins. Place the bypass capacitor as close as possible to the power supply pins in this LSI. The capacity values of the capacitor are 0.1 μF to 0.33 μF (recommended values). For the capacitor related to crystal oscillation, see section 5.8, Notes on Board Design.

Table 33.37 shows the combinations of external capacitor.

Appendix

A. Pin States

Table A.1 shows the state of pin function in each operating mode. For input/output pin function, the input buffer state is shown in the table above and the output buffer state is in the table below.

[Legend]

- I: Input
- O: Output
- H: High-level output
- L: Low-level output
- Z: Input pins retain their state, and output pins become high-impedance,.
- +: Output state is retained.
- Notes: 1. Indicates the power-on reset by low-level input to the RES pin. The pin states after a power-on reset by the H-UDI reset assert command or WDT overflow are the same as the initial pin states at normal operation (see section 27, Pin Function Controller (PFC)).
	- 2. Indicates the pin states that the IOKEEP bit in the deep standby cancel source flag register (DSFR) is cleared if the chip has recovered from deep standby mode by the input on any of pins NMI, MRES, and IRQ7 to IRQ0 (see section 30, Power-Down Modes).
	- 3. This LSI shifts to the power-on reset state for a certain period after the recovery from the deep standby mode (see section 30, Power-Down Modes).
	- 4. The week keeper and pull-up circuits included in the I/O pins are turned off.
	- 5. When pins for the connection with a crystal resonator are not used, the input pins (EXTAL, RTC_X1, AUDIO_X1, and USB_X1) must be fixed (pulled up, pulled down, connected to power supply, or connected to ground) and the output pins (XTAL, RTC_X2, AUDIO_X2, and USB_X2) must be open.
	- 6. Depends on the setting of the CKOEN bit in the frequency control register (FRQCR) of the CPG (see section 5, Clock Pulse Generator (CPG)).
	- 7. Depends on the setting of the HIZ bit in the high impedance control register (HIZCR) (see section 30, Power-Down Modes).
	- 8. Depends on the setting of the HIZBSC bit in the high impedance control register (HIZCR) (see section 30, Power-Down Modes).
	- 9. Depends on the setting of the corresponding bit in the deep standby cancel source select register (DSSSR) (see section 30, Power-Down Modes).
	- 10. Depends on the setting of the RTCEN bit in the RTC control register (RCR2) of the RTC (see section 15, Realtime Clock (RTC)).
	- 11. Depends on the AXTALE bit in the standby control register (STBCR1) (see section 30, Power-Down Modes).
	- 12. Depend on the CS0KEEPE bit in the deep standby control register (DSCTR) (see section 30, Power-Down Modes).
	- 13. Z when the TAP controller of the H-UDI is neither the Shift-DR nor Shift-IR state.
	- 14. These are the pin states in product chip mode (ASEMD=H). See the Emulation Manual (provisional title) for the pin states in ASE mode (ASEMD=L).

B. Treatment of Unused Pins

Table B.1 Handling of Pins that are not in Use (Except for H-UDI and Emulator Interface Pins)

Note: It is recommended that the values of pull-up or pull-down resistors are in the range from 4.7 kΩ to 100 kΩ.

Table B.2 Handling of Pins that are not in Use (When H-UDI is Not Used in Product Chip Mode)

Notes: 1. When using the H-UDI, handle these pins as described in the manual for the emulator.

 2. It is recommended that the values of pull-up or pull-down resistors are in the range from 4.7 kΩ to 100 kΩ.

C. Package Dimensions

The package dimension that is shown in the Renesas Semiconductor Package Data Book has Priority.

Figure C.1 Package Dimensions

Main Revisions for This Edition

Index

Numerics

A

B

C

D

E

F

G

H

I

J

L

M

N

O

P

Q

R

FLADR2...

FLDATAR

FLDTFIFO FLECFIFO

FLTRCR... FPSCR FPH. FRMNUM. FRQCR0... FRQCR1...

GR DELT.

RWKCNT

SAR (IIC3)

S

T

U

V

W

SH7205 Group User's Manual: Hardware

Publication Date: 1st Edition, March, 2008 Rev.3.00, March 27, 2014

Published by: Renesas Electronics Corporation

SALES OFFICES

Renesas Electronics Corporation

http://www.renesas.com

Refer to "http://www.renesas.com/" for the latest and detailed information.

Renesas Electronics America Inc.
2801 Scott Boulevard Santa Clara, CA 95050-2549, U.S.A.
Tel: +1-408-588-6000, Fax: +1-408-588-6130

Renesas Electronics Canada Limited 1101 Nicholson Road, Newmarket, Ontario L3Y 9C3, Canada Tel: +1-905-898-5441, Fax: +1-905-898-3220

Renesas Electronics Europe Limited Dukes Meadow, Millboard Road, Bourne End, Buckinghamshire, SL8 5FH, U.K Tel: +44-1628-585-100, Fax: +44-1628-585-900 **Renesas Electronics Europe GmbH**

Arcadiastrasse 10, 40472 Düsseldorf, Germany Tel: +49-211-6503-0, Fax: +49-211-6503-1327

Renesas Electronics (China) Co., Ltd.
Room 1709, Quantum Plaza, No.27 ZhiChunLu Haidian District, Beijing 100191, P.R.China
Tel: +86-10-8235-1155, Fax: +86-10-8235-7679 **Renesas Electronics (Shanghai) Co., Ltd.**

Unit 301, Tower A, Central Towers, 555 Langao Road, Putuo District, Shanghai, P. R. China 200333 Tel: +86-21-2226-0888, Fax: +86-21-2226-0999

Renesas Electronics Hong Kong Limited
Unit 1601-1613, 16/F., Tower 2, Grand Century Place, 193 Prince Edward Road West, Mongkok, Kowloon, Hong Kong
Tel: +852-2265-6688, Fax: +852 2886-9022/9044

Renesas Electronics Taiwan Co., Ltd.
13F, No. 363, Fu Shing North Road, Taipei 10543, Taiwan
Tel: +886-2-8175-9600, Fax: +886 2-8175-9670

Renesas Electronics Singapore Pte. Ltd. 80 Bendemeer Road, Unit #06-02 Hyflux Innovation Centre, Singapore 339949 Tel: +65-6213-0200, Fax: +65-6213-0300

Renesas Electronics Malaysia Sdn.Bhd.
Unit 906, Block B, Menara Amcorp, Amcorp Trade Centre, No. 18, Jln Persiaran Barat, 46050 Petaling Jaya, Selangor Darul Ehsan, Malaysia
Tel: +60-3-7955-9390, Fax: +60-3-7955-9510 **Renesas Electronics Korea Co., Ltd.**

12F., 234 Teheran-ro, Gangnam-Ku, Seoul, 135-920, Korea Tel: +82-2-558-3737, Fax: +82-2-558-5141

SH7205 Group User's Manual: Hardware

R01UH0473EJ0300 (Previous Number: REJ09B0372-0200)

Mouser Electronics

Authorized Distributor

Click to View Pricing, Inventory, Delivery & Lifecycle Information:

[Renesas Electronics](https://www.mouser.com/renesas): [R5S72050W200BG](https://www.mouser.com/access/?pn=R5S72050W200BG)